

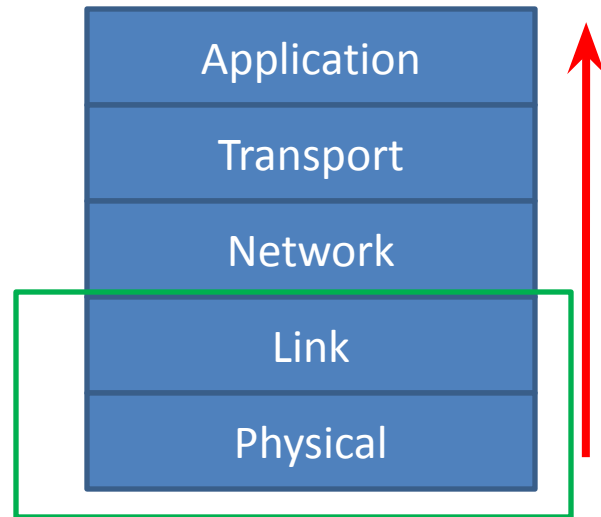
Computer and Network Security: Physical/Link Layer Attacks and Solutions

Kameswari Chebrolu

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Outline

- Attacks at different layers of the protocol stack
 - Background Material
 - Various Attacks
 - Solutions to the same

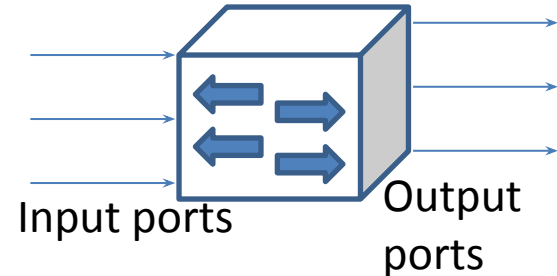
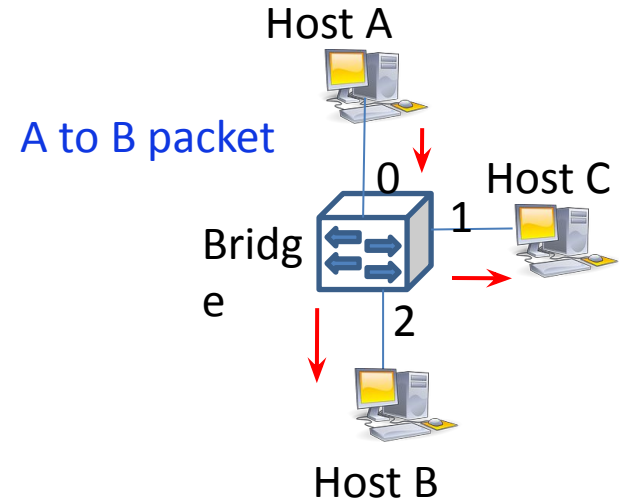


Physical/Link Layer Roles

- What is the role of physical layer?
 - Encode bits into signals (e.g. voltage levels)
- What is the role of link layer?
 - Next hop delivery (framing, error-control, media access and switching)
- Link-layer Switching: Star topology and learning bridges
 - Bus topology outdated, no role for CSMA/CD

Forwarding

- MAC Address: identifies source and destination nodes
- How to forward?
 - Host A sending packet to Host B
 - Host A sending packet to Host C
- Manual configuration: Tedious
- Automatic simple strategy:
Forward on all interfaces except the one on which received



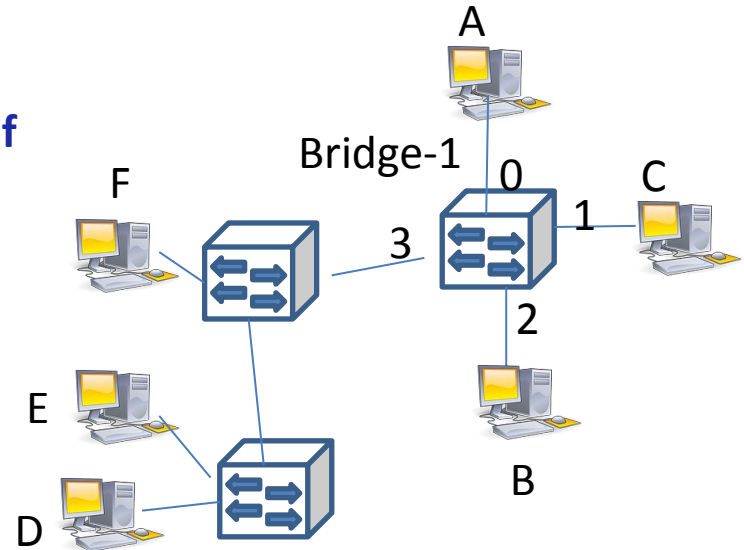
Learning Bridges

- Idea: Inspect source address and map it to port on which the frame was received
 - Each entry purged after some period unless refreshed

Host	Port
A	0
B	2
C	1
D	3
E	3
F	3

Bridge-1's table

Extended LAN
(A to F are all
neighbours as if
connected to a
single bus **)

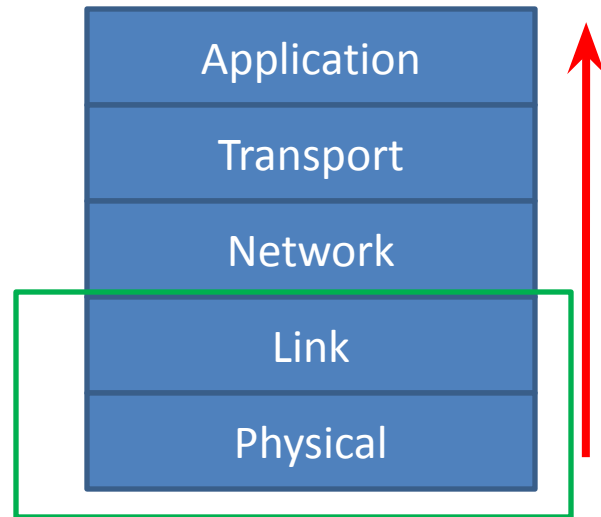


Algorithm

- If a frame received at bridge for destination D on port p
 - No entry for D in the table, forward on all ports except port p
 - If entry for D in forwarding table corresponds to p, drop frame
 - If entry for D in forwarding table corresponds to $i \neq p$, then forward on i

Outline

- Attacks at different layers of the protocol stack
 - ~~Background Material~~
 - **Various Attacks**
 - Solutions to the same



Attacks and Defenses (Outline)

- Eavesdropping
- Disruption
- Spoofing
- Protocol specific attacks

Eavesdropping

- Communication media and topology matters
 - Fiber, Wireless, Ethernet star
- Can splice a cable and sniff
- Eavesdropping without piercing cable also possible (based on electromagnetic radiation)
 - See Operation Ivy Bells:
https://en.wikipedia.org/wiki/Operation_Ivy_Bells

Eavesdropping

- Wireless media
 - Put interface in monitor/promiscuous mode; sniff media using tcpdump/wireshark tools
 - Can sniff network/transport/app layer headers/data
- Ethernet switching based on star topology
 - Learning bridges make it hard to sniff (frame is forwarded only on the relevant port)
 - Possible to work around. How?

Eavesdropping in Ethernet

- **MAC flooding:** Forces a switch to flood unicast traffic instead of route along right path
- How achieved?
 - Bridges set aside limited memory to store forwarding table (MAC address to port)
 - Feed switch many Ethernet frames with different source addresses □ evict valid entries
 - No entry for a destination □ forward on all ports (other than one received) □ scope for eavesdropping

Defense

- MAC Limiting: Limit the number of MAC addresses learnt from a port.
 - Ports connected to end-hosts (low limit) vs other switches (high limit)
- Encryption of link layer payload (or at higher layers, e.g. HTTPS)

Attacks and Defenses (Outline)

- ~~Eavesdropping~~
- **Disruption**
- **Spoofing**
- Protocol specific attacks

Disruption

- Jam ongoing communication
 - Easier in wireless environment; Transmit at same time with high signal strength
 - Packet discarded due to bit corruption (checksum fail)
 - Cut cables (especially ones that serve many customers)
- **Defence:** No easy solution
 - Localizing jammers; Redundancy in the system

Spoofing

- Impersonate some one else (take their address)
 - Blind spoofing: spoofing without eavesdropping
- Used to cover track or access unauthorized resource
- Inject packets with fake source MAC address
 - In linux, can be done via “ifconfig” (e.g. `ifconfig eth0 hw ether 00:84:48:AA:D3:61`) with root permission

Attacks and Defenses (Outline)

- ~~Eavesdropping~~
- ~~Disruption~~
- ~~Spoofing~~
- **Protocol specific attacks: Address Resolution Protocol (ARP)**

ARP Spoofing

- **ARP Spoofing** can allow eavesdropping, denial of service and MITM attacks
- But what is ARP?

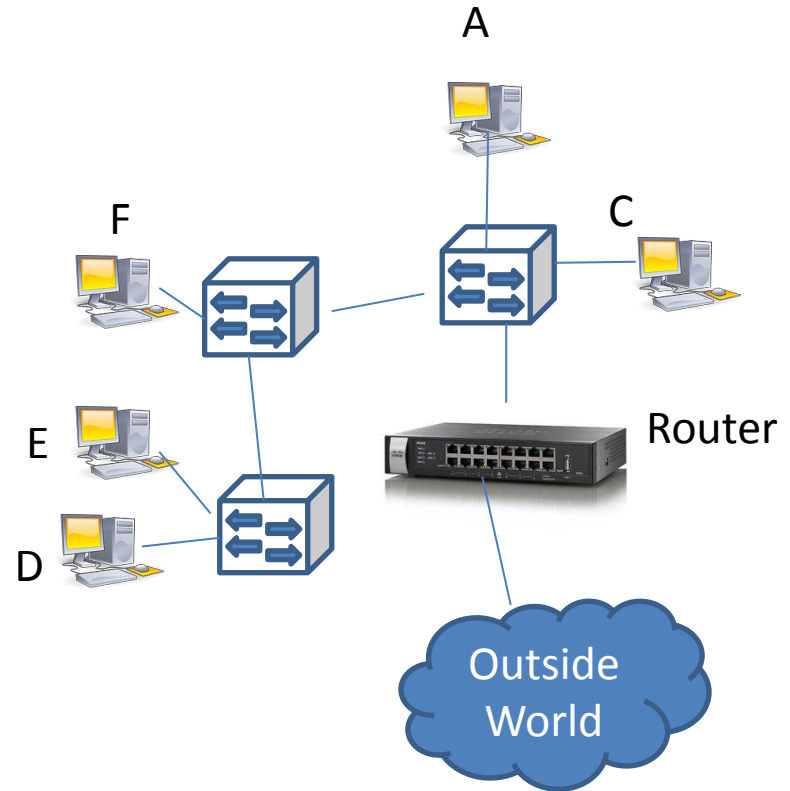
ARP (Address Resolution Protocol)

- At network layer: Routing module at a node decides who the next hop node is (its IP address)
- At link layer: transfer over the next-hop link needs next hop's MAC address
 - Link here can refer to extended LAN

ARP (Address Resolution Protocol)

- ARP helps determine MAC address corresponding to given IP address (Request is broadcast)
 - Host with matching IP address replies (Reply is unicast)
- Each host maintains a cache with IP to MAC translations
 - Entries in cache timed out periodically (say 15 min)

Example



Address Resolution Protocol (ARP)

- When forwarding a datagram, check cache, if no mapping, invoke ARP
- Originator: Add entry to cache corresponding to target (obtained from ARP reply)
- Target: Add entry to cache corresponding to the originator (obtained from ARP request)
- Intermediate hosts: Refresh existing entries (obtained from ARP request)

Gratuitous ARPs

- Generated by a host to inform others of its IP to MAC mapping
- Could be a request or reply
 - If request, no reply will occur
 - If reply, there was no preceding request

Uses of Gratuitous ARPs

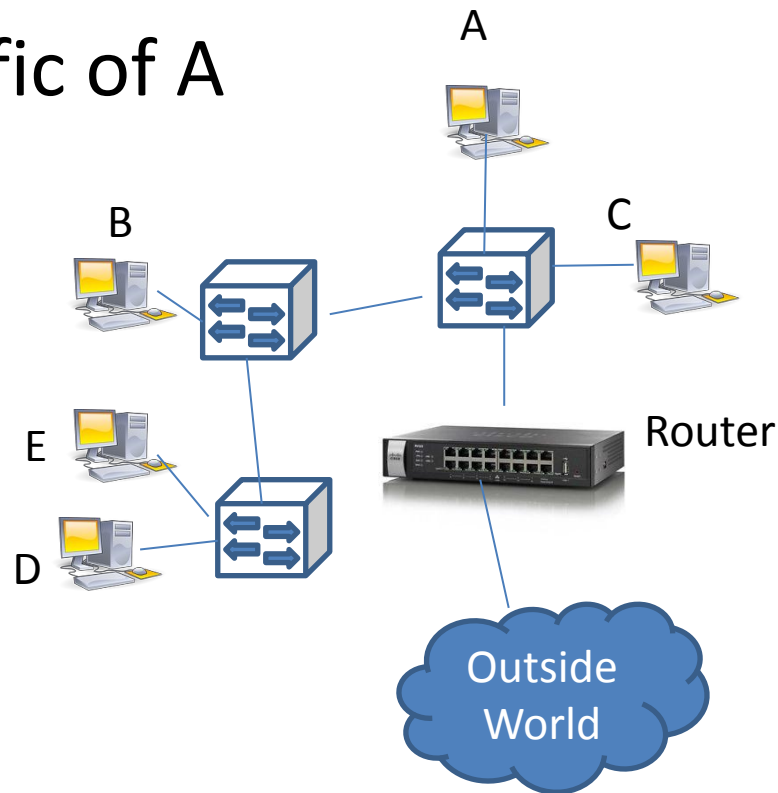
- Issued whenever IP or MAC address of an interface changes or brought up from down state
 - Help rectify cached ARP entries
 - Report IP address conflicts (duplicate IP)
 - Inform bridges of the location of new host

Vulnerability

- ARP is a stateless protocol
 - A host will act on a reply regardless of whether a request precedes it
- ARP has no authentication mechanism to verify identity of sender

ARP Cache Poisoning

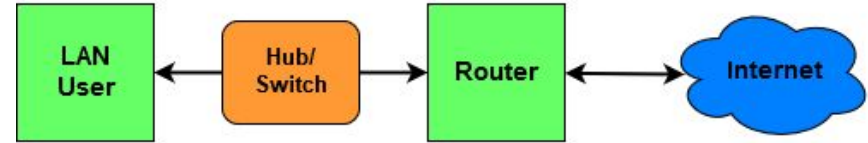
- Eve E wants to examine traffic of A



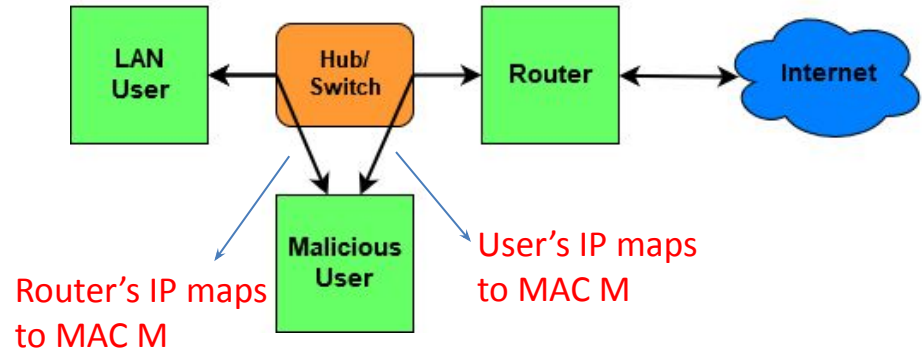
ARP Cache Poisoning

- Associate IP address of target host with Attacker's MAC
 - Traffic meant for target received by attacker
 - Inspect, modify, drop (denial of service) traffic
 - Can also do MITM, where both parties unaware of the attack
- Legitimate uses: backup-servers; debug traffic

Routing under normal operation



Routing subject to ARP cache poisoning



MITM Attack

Solutions

- Hosts maintain 'static read-only' ARP tables for critical services
 - Not very convenient; not scalable
- Ethernet switches can check for same IP address mapped to many MACs and alert sys-ads via email
 - Can also verify mapping from DHCP servers
- OS level: Ignore unsolicited replies (but attacker can work around it)
- Tools: Arpwatch, xARP, ArpStar
 - Analyze all observed ARP traffic and determine spoofing attacks

Summary

- Covered some Link layer background (switching, learning bridges, ARP)
- Different types of attacks: Eavesdropping, Disruption, Spoofing and ARP spoofing
- Each attack has a solution that works in practice