

Lab 04: Demonstrating Static vs Dynamic Scoping – E/20/037

Task 1

Run the following python code, and explain the output.

```
a = 5

def f():
    print(a)

def g():
    a = 10
    f()

g()
```

- Python uses **static (lexical) scoping**.
- When f() executes, it looks for variable a in its lexical environment (where it was defined in the source code), not where it was called from.
- Since f() was defined at the global level, it refers to the global a = 5.
- The local a = 10 inside g() is not visible to f() because f() was not defined inside g().
- In static scoping, variable binding is determined by the physical structure of the code, not the call stack.

Task 2

Convert the Python code from **Task 1** into Perl syntax. You will perform a two-part experiment by modifying the variable declaration inside the function g().

Case A: Using ‘**my**’ Keyword

Case B: Using ‘**local**’ Keyword

Run the code and record the output. Explain the output for two cases.

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HelloWorld.pl

```
1 $a = 5;
2
3 sub f {
4     print "$a\n";
5 }
6
7 sub g {
8     my $a = 10; # Using 'my'
9     f();
10}
11
12 g();|
```

STDIN

Input for the program (Optional)

Output:

5

The ‘**my**’ keyword creates a **lexically scoped** variable.

This means `my $a = 10` creates a NEW variable `$a` that only exists within the `g()` subroutine’s lexical scope.

When `f()` is called, it looks for `$a` in its lexical environment (where it was defined), which is the global scope.

So it prints the global `$a = 5`. This demonstrates **static scoping**.

HelloWorld.pl

```
1 $a = 5;
2
3 - sub f {
4     print "$a\n";
5 }
6
7 - sub g {
8     local $a = 10; # Using 'local'
9     f();
10}
11
12 g();
```

STDIN
Input for the program (Optional)

Output:
10

The local keyword creates a **dynamically scoped** variable.

It temporarily modifies the global \$a for the duration of g()'s execution and any subroutines called from within g().

When f() is called from g(), it sees the temporarily modified value \$a = 10 because dynamic scoping looks at the **call stack** at runtime, not the lexical structure.

This demonstrates **dynamic scoping**.

Task 3

Complete the Perl code below by filling in the blanks with either ‘`my`’ (Lexical) or ‘`local`’ (Dynamic). For each case, run the code in Perl, and explain the output.

```
$x = "Global";  
  
sub level_3 {  
    print "Level 3 sees x as: $x\n";  
}  
  
sub level_2 {  
    # TO DO  
    _____ $x = "Level 2 Value";  
    level_3();  
}
```

```
sub level_1 {  
    # TO DO  
    _____ $x = "Level 1 Value";  
    level_2();  
}  
  
level_1();
```

Cases you have to explain:

Case	Level 1 Key word	Level 2 Keyword
A	<code>my</code>	<code>my</code>
B	<code>local</code>	<code>my</code>
C	<code>my</code>	<code>local</code>
D	<code>local</code>	<code>local</code>

Case A: Level 1 = my, Level 2 = my

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HelloWorld.pl 44a2txbr9 🖊

```
1 $x = "Global";
2
3 sub level_3 {
4     print "Level 3 sees x as: $x\n";
5 }
6
7 sub level_2 {
8     my $x = "Level 2 Value"; # my
9     level_3();
10}
11
12 sub level_1 {
13     my $x = "Level 1 Value"; # my
14     level_2();
15 }
16
17 level_1();|
```

STDIN

Input for the program (Optional)

Output:

Level 3 sees x as: Global

Both my keywords create **lexically scoped variables**. Each creates a new variable \$x that only exists within its respective subroutine. When level_3() executes, it looks for \$x in its lexical scope (where it was defined), which is the global scope. It cannot see the lexically scoped variables in level_1 or level_2 because they are isolated to their own scopes.

Case B: Level 1 = local, Level 2 = my

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HelloWorld.pl 44a2txbr9 🖊

```
1 $x = "Global";
2
3 sub level_3 {
4     print "Level 3 sees x as: $x\n";
5 }
6
7 sub level_2 {
8     my $x = "Level 2 Value"; # my
9     level_3();
10}
11
12 sub level_1 {
13     local $x = "Level 1 Value"; # local
14     level_2();
15 }
16
17 level_1();|
```

STDIN

Input for the program (Optional)

Output:

Level 3 sees x as: Level 1 Value

level_1 uses local, which **temporarily modifies the global \$x** to "Level 1 Value" for the **duration of the call chain**. However, level_2 uses my, which creates a NEW lexically scoped variable that only exists within level_2 itself. When level_3() is called, it looks at the **global \$x which has been temporarily modified by the local in level_1**.

Case C: Level 1 = my, Level 2 = local

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HelloWorld.pl 44a2txbr9 

```
1 $x = "Global";
2
3 sub level_3 {
4     print "Level 3 sees x as: $x\n";
5 }
6
7 sub level_2 {
8     local $x = "Level 2 Value"; # local
9     level_3();
10}
11
12 sub level_1 {
13     my $x = "Level 1 Value"; # my
14     level_2();
15 }
16
17 level_1();
```

STDIN
Input for the program (Optional)

Output:
Level 3 sees x as: Level 2 Value

level_1 uses my, creating a lexically scoped variable that doesn't affect the global \$x. level_2 uses local, which **temporarily modifies the global \$x** to "Level 2 Value". When level_3() executes, it sees this local modification in the call stack. Result: prints "Level 2 Value".

Case D: Level 1 = local, Level 2 = local

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HelloWorld.pl 44a2vqv9q 

```
1 $x = "Global";
2
3 sub level_3 {
4   print "Level 3 sees x as: $x\n";
5 }
6
7 sub level_2 {
8   local $x = "Level 2 Value"; # local
9   level_3();
10}
11
12 sub level_1 {
13   local $x = "Level 1 Value"; # local
14   level_2();
15 }
16
17 level_1();|
```

STDIN
Input for the program (Optional)

Output:
Level 3 sees x as: Level 2 Value

Both use local, creating dynamically scoped temporary modifications to the global \$x. level_1 first sets it to "Level 1 Value", then level_2 sets it to "Level 2 Value". Since **local creates temporary values that are visible in the call stack, and level_2's modification is more recent** (deeper in the call stack), level_3() sees "Level 2 Value".

'my' creates isolated lexical scopes (static scoping), while 'local' temporarily modifies the global variable for the duration of the call chain (dynamic scoping).