

We can construct an arbitrary length hash function $H^s(x)$ from a fixed length hash function $h^s(x)$ with inputs of length $2n$ and output of length n . If (Gen, h) is a fixed length collision resistant hash function, then (Gen, H) is also a collision resistant hash function.

Let (Gen, h) be a fixed length collision resistant hash function for input of length $2l(n)$ and output of length $l(n)$. We can set $l(n) = l$. On input a key s and a string x of any length L such that $L < 2^{l(n)}$, a variable length hash function (Gen, H) can be constructed as follows using Merkle-Demgard transform:

- Set the length $L = \text{length of the string } x$. Let there be $B = \lceil L/l \rceil$ blocks. Pad x with zeros so that its length is a multiple of l . Parse the result as the sequence of l -bit blocks $x_1 \dots x_B$. Set $x_{B+1} = L$ where L is encoded using exactly l bits.
- Set $z_0 = 0^l$ as initialising vector
- For $i=1$ to $B+1$, compute $z_i = h^s(z_{i-1} \parallel x_i)$
- Output z_{B+1}

