

# Delegated Anonymous Credentials with Revocation Capability for IoT Service Chains

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**Abstract**—The abstract goes here.

**Index Terms**—

## I. INTRODUCTION

Internet of Things (IoT) enables physical objects also called *Things* to communicate with each other and to their human operators. This opens up a myriad of use cases such as smart homes, smart factories, smart cities, smart healthcare, smart grids etc [1]. It is also expected that such connected devices could reach upto 50 billion by 2020 [2]. The IoT devices (for example, a smart bulb or a temperature sensor) are very constrained in terms of memory, processing power, storage and most often are battery powered. Unlike the traditional computers these devices cannot perform computationally intensive tasks and are intended for minor operations of sensing and actuation. Also most of these devices are out in the open without any physical supervision making them easily susceptible to physical attacks.

Owing to the resource constraints and physical openness, IoT devices have been targets of various attacks [3] at physical, network and application layer. IoT devices also collect lot of personal information like user's location, eating habits, medical history etc because of which there has been a growing concern among consumers of such services. Unlike normal computers, these devices cannot provide an interface where the user can look up what personal information is being shared and with whom. In [4] the authors define privacy in IoT as a guarantee for the subject

- To be aware of the privacy risks imposed by smart things.
- Control over collection and processing of personal information
- Control over subject's personal information being disseminated outside of his control sphere.

They then categorize privacy threats and challenges of IoT into a) Identification b) Localization and tracking c) Profiling d) Privacy-violating interaction and presentation e) Inventory and life cycle tracking and f) linkage.

IoT services do not act in silos. They interact with each other and with external entities to provide a complete package of services to the user. For example, in Home Automation, based on the user who is entering the house (say Owner

vs Guest), a completely different set of services may get invoked. The service interactions and invocations depend on the roles and capabilities of the user invoking them. We call the sequence of services that get invoked as *IoT Service Chain*. In this paper we look at providing security and privacy to users and IoT devices that invoke IoT service chains. The rest of the paper is organized as follows. Section 2..... Section 3... Section 4..

## II. MOTIVATION AND RELATED WORK

1) *IoT Service Chains*: We introduced *IoT Service Chains* in section I to refer to the chain of services invoked when an event occurs. Individual services in the chain interact with each other, on-behalf of the initiator towards a common goal. Initiator could either be a human or an IoT device. Each service in the chain would in-turn validate the user/IoT device's credentials for the desired service functionality. IoT devices can either act on behalf of their operator or can be independent of it. For example, if a smart Heart Monitoring System (HMS) detects a low pulse rate for a patient, it immediately needs to initiate the Advanced Cardiac Life Support (ACLS) by injecting an IV of an antidote, reading and interpreting the ECG, starting CPR, inform the doctor etc. The HMS in this case acts as an independent device and does not impersonate the patient. But in case of a smart fridge, which keeps track of the stock of milk, it automatically places an order for replenishment by using the owners credit card details. In this case, the smart fridge acts on behalf of the owner.

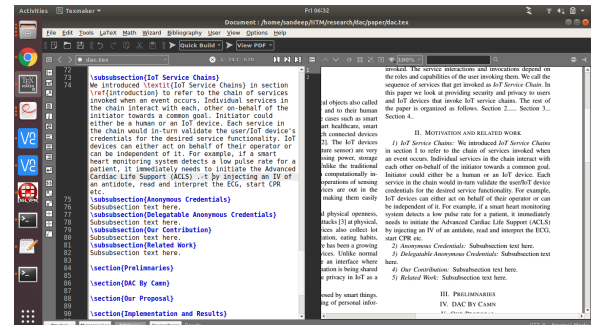


Fig. 1. IoT Service Chain

As can be seen in Fig. 1 there are 5 services in the IoT service chain. xxx explanation of HMS, ACLS etc interactions among services, authentication etc

One thing to notice from the above figure is that while the service chain is being invoked, the user/IoT device has to be online so that it supplies the necessary credentials for authentication and authorization. This isn't a huge problem if the initiating entity is a user (having a smart phone or a tablet). But for a constrained IoT device generating authentication and authorization information for each service in the chain would result in draining of its resources very quickly. Also, most of the IoT devices are duty cycled and may not remain active till the chain completes. This again results in loss of packets or in retransmission. It will be even more challenging if the user/device's privacy needs to be protected through out the chain. In this paper we focus on some of the options on how device privacy can be ensured when invoking the service chain without pushing the device to its resource limits.

2) *Attribute Based Anonymous Credentials*: In traditional credential based system, a central authority grants credentials to users and systems. These credentials can be password or token (certificate) based. When the user wants to access a resource he presents the credential to it which inturn validates the credential. The resource (also called service) trusts the central authority and thereby the credentials issued by it. Once the credential verification is done, the service checks whether the user has got the right access level to access the resource. One of the primary concern with such a system is that the service gets to know the complete details of the user presenting the credential. for example, in public key certificates issued to users, the service will know how long the credential (certificate in this case) is valid, the country, state, organization, email address etc of the user, who are all the intermediate chain of issuers etc. The only information that matters to the service is whether the certificate is valid and is issued by the central authority and whether the user possesses the required access right. But in this case lot more user details are available to it. With so much personal information available to the service, there is always a chance of misuse of data if it falls in wrong hands. According to European Union General Data Protection Regulation (GDPR) Data Minimization principle, entities should only process adequate, relevent and limited personal data that is necessary in relation to the purposes for thish they are processed. As mentioned in [4], the privacy problem becomes manifold in IoT world. One, because there are huge number of constrained IoT devices without proper security measures in place. And two, there is very little control over what personal information is disseminated from these devices to the outside world.

Anonymous Credentials introduced in [5] provide a way in which the user can prove that he holds a credential without revealing any information about the user. The verifier cannot also forge the user's credential. In Attribute Based Anonymous Credentials, the user can selectively prove that he holds the set of attributes needed by the verifier and not reveal all of his attributes, thereby maintaining privacy. The prover creates a zero-knowledge proof of possession of the credential which the verifier verifies using the public key of the central authority.

Some of the most popular Anonymous Credential Systems (ACS) are [6] and [7].

A Delegatable Anonymus Credential (DAC) System, introduced in [8] not only allows the users to generate anonymous credentials, but also allows them to anonymously delegate their credentials to other entities. For example, in a hierarchical setup, the root issuer may delegate its issuing authority to region wise issuers who in turn may delegate to division wise issuers and so on. Although there can be multiple levels of delegation, the anonymous credential generated by the prover at any level can only be verified by the public key of the root issuer. One of the primary advantages of DAC is that it alleviates the burden on the root issuer to issue credentials but still maintains anonymity of all the issuers in the chain.

3) *Our Contribution*: In section II-1 we talked about IoT service chains and the difficulty of ensuring user/IoT device's privacy during the chain propagation. We address this problem with our proposal on *Delegated Anonymous Credentials in IoT Service Chains*. We describe a mechanism in which the IoT device can delegate its credentials to a controller which in turn generates an anonymous credential based on the attributes needed by the service. Our scheme is based on the DAC system developed by Camenisch *et al.* [9]. The authors developed a scheme where credentials are not delegated anonymously but the prover anonymously proves that the entire chain of delegation is valid. The verifier verifies the anonymous credential just with the root issuer's public key.

The following are major contributions in this work.

- a) Discuss in detail the problem of ensuring privacy to users and IoT devices in IoT service chains.
- b) Implementation of the DAC scheme outlined in [9]. We used the Pairing Based Cryptography Library from [10] for the implementation. We implemented the full L-level credential Delegation, Presentation and Verification.
- c) Credential revocation has been mentioned as "future work" in [9]. We implemented revocation of credentials in our scheme.
- d) Once the verifier verifies the anonymous credential token, the next logical step would be to communicate with the prover to exchange data. In order to do that securely, a common session key needs to be established between the two parties. We demonstrate how the session key can be established.
- e) We then used the above DAC implementation to realize Delegated Anonymous Credentials in IoT service chains. We outline the framework consisting of Root issuer, Controller, User/IoT device and Service as the principal components. We describe the messages exchanged between these components and outline how privacy of users can be ensured.
- f) We implemented the above scheme and discuss the various possibilities on how token verification and policy implementation can be placed across constrained IoT services. Various metrics like time taken, number and size of messages exchanged, memory, CPU etc are evaluated for the various models that we discuss.

4) *Related Work*: Subsubsection text here.



1) *Credential Revocation*: One of the important functionalities of any credential based system is to account for revoked credentials. Revocation allows the administrator to stop rouge elements from entering into the system. The authors in [9] left out credential revocation as a future item. In this section we discuss how credential revocation can be achieved using the author's concrete instantiation.

As mentioned in the previous section, if a user at level  $K-1$  wants to delegate certain attributes ( $\vec{a}_K$ ) to another user, he signs the message ( $cpk_K, \vec{a}_K$ ) with  $csk_{K-1}$ . We propose to extend the signature message to include a Hash of ( $cpk_K, \vec{a}_K$ ). We call this hash as the credential hash at level  $K$ . Therefore,

$$h_K = Hash(cpk_K, \vec{a}_K)$$

$$\sigma_K = Sign_{csk_{K-1}}(cpk_K, h_K, \vec{a}_K)$$

Hash in this case can be one of the *Secure Hash Algorithms* in the SHA-2 family. The credential hash is a unique representation of the public key and the delegated attributes at any level. Since the signature includes the credential hash, signature verification would fail if the hash is tampered. The delegator forwards  $\sigma_K, h_K, \vec{a}_K$  to the delegatee.  $h_K$  forms a part of credential at every level. The attribute token generated for presentation would now look like

$$at \leftarrow NIZK\{\}$$

When generating the attribute token, unlike other attributes, the credential hash value is always revealed to the verifier. There is no danger of someone modifying the hash as that would result in signature verification failure. If the root issuer or any delegator in the chain needs to revoke a credential, he would publish the credential hash in a Black List (BL) of credential hashes. The verifier has to check the credential hash for each level against the BL and if there is a match, verification should fail.

The instantiation of the NIZK for the attribute token would now look like,

xxxxxx  
xxxxxx

2) *Session Key Establishment*: In the IoT scenario, once the token generated by the IoT device/User is verified by the IoT service, the next logical step would be to exchange data. In order to securely perform data exchange, both the endpoints should establish a common secret key. In the earlier section, we talked about how the authors use signature proof of knowledge over message  $m$  to prove the possession of the Level- $L$  private key. The proof then acts as a Digital Signature over  $m$ . Instead of any random message, the prover can generate an ephemeral public, private key pair ( $pk, sk$ ) and use the public key  $pk$  as the message over which the signature proof is performed. This public key is also sent along with the attribute token to the verifier. Once the token verification succeeds, the verifier generates a random secret key and uses the prover's ephemeral public key to encrypt it. It then sends the encrypted secret key over to the prover which then decrypts it using the private key  $sk$ . Any attempt by a Man-In-The-Middle to modify the public key sent to the verifier would

result in the token verification failure. The scheme assumes that the response from the verifier is not tampered. Additional measures will have to be taken if verifier's identity and its response have to be validated at the prover.

## V. THE PROPOSAL

In this section we discuss in detail the architectural components of our proposal and the messages exchanged between them. One of the primary requirements of the system was to provide a light weight solution for IoT devices and users to anonymously access IoT services.

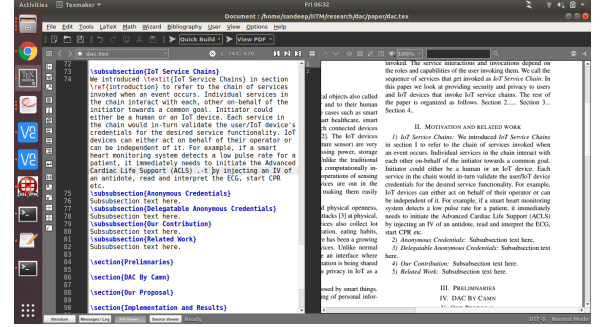


Fig. 3. System Architecture

Fig. 3 illustrates the system architecture with the following major components.

1) *The Root Issuer*: As the name implies, the Root Issuer (RI) issues Level-1 credentials to users/IoT devices. The RI service maintains a public, private key pair ( $ipk, isk$ ) using which it signs the Level-1 credentials. The public key  $ipk$  is distributed across the system so that services can validate credential tokens presented from any delegation level just by using  $ipk$ .

The RI generates the system public parameters  $G1, G2, g1, g2$  and the bilinear map  $e$ .

All users/IoT devices are assumed to hold one long term public, private key pair. Any user/IoT device which wishes to obtain credentials from RI, submits the credential request containing their public key and the attributes that they wish to possess. RI then validates the entity's request and generates the credential. As mentioned in section IV-1, RI signs the combination of user/device's public key, the credential hash and the attributes, with its private key  $isk$ . the credential hash is the hash of the user's public key and all the attributes that the RI wishes to assign to the entity. The entity's credential would then be a combination of the signature, the credential hash, the attributes and its public, private key pair.

If the root issuer wishes to revoke an entity's credential, it would publish the credential hash in a BlackList (BL) that is available for everyone in the system to view. Only the RI can add, delete or modify the entries in the BL. Since the credential hash is hashed over public key and attributes, no one can infer any useful information about the entity just by looking at the hash.

## VI. IMPLEMENTATION AND RESULTS

## VII. CONCLUSION

The conclusion goes here.

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