



$$\begin{aligned} \text{rear} &= (\text{rear} + 1) \% \text{SIZE} \\ &= (-1 + 1) \% 5 \\ &= 0 \% 5 = 0 \end{aligned}$$

$$\begin{aligned} \text{rear} &= \text{rear} + 1 \\ &= 4 + 1 \\ &= 5 \end{aligned}$$

$$\begin{aligned} \text{rear} &= (0 + 1) \% 5 \\ &= 1 \% 5 = 1 \end{aligned}$$

$$\text{rear} = (\text{rear} + 1) \% \text{SIZE}$$

$$\begin{aligned} \text{rear} &= (1 + 1) \% 5 \\ &= 2 \% 5 = 2 \end{aligned}$$

$$= (2 + 1) \% 5$$

$$= 3 \% 5 = 3$$

$$\begin{aligned} &= (3 + 1) \% 5 \\ &= 4 \% 5 = 4 \end{aligned}$$

$$\begin{aligned} &= (4 + 1) \% 5 \\ &= 5 \% 5 = 0 \end{aligned}$$

Enqueue:

- 1) Check if queue is not full
- 2) Increment rear as  
 $\text{Rear} = (\text{rear} + 1) + \text{SIZE}$
- 3) Add element at rear position
- 4) If  $\text{front} == -1$ , make  $\text{front} = 0$

Dequeue:

Check if Queue is not empty.

Increment front.

If front is 4 and rear is 0

to delete the rear position, we cannot increment front as  $\text{front}++$  Will be index 5 but we want to delete index 0;

So,

$\text{Front} = \text{front} + 1 \% \text{SIZE}$

If deleting the last element in queue

If ( $\text{front} == \text{rear}$ )

$\text{Front} = \text{rear} = -1$

Queue empty condition

If ( $\text{rear} == -1$ ) queue is empty

Circular Queue full condition

When Queue is full,

Rear = 0 front = 1

Rear = 1, front = 2

Rear = 2 front = 3

Rear = 3 front = 4

**rear = 4, front = 0**

Front == rear + 1

3 == 2+1

But when rear = 4 front = 0

0 == 4+1 does not satisfy the above condition

So the queue full condition can be

Front == ( $\text{rear} + 1$ ) % SIZE

0 == (4+1) % 5

0 == 5%5

0==0

To enqueue, instead of  $\text{rear}++$

We say,

$\text{Rear} = (\text{rear} + 1) \% \text{SIZE}$

**Option 2:**

If  $\text{rear} == \text{SIZE} - 1$

Rear = 0;

Else

Rear++;