

Let $f(n) = n^2$ and the language $A' = \text{pad}(A, f(n))$,

$$\therefore A' = \{x\#^{n^2-n} : x \in A \wedge n = |x|\}.$$

Now, $A \in \text{TIME}(n^6)$

$\Leftrightarrow \exists \text{ TM } M_A \text{ that decides } A \text{ in } O(n^6) \text{ time.}$

A' is f -padded version of A .

From M_A , let's construct a TM $M_{A'}$ that decides A' .

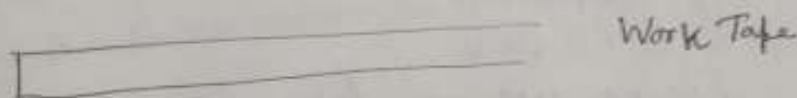
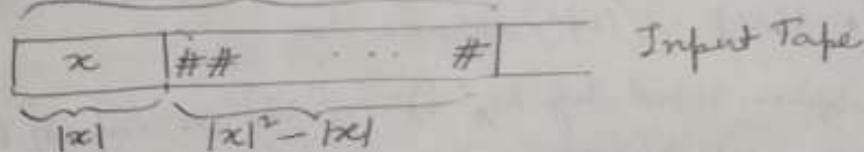
$M_{A'}$ "on input y :

1. Check if $\exists x \in \Sigma^* \mid y = x\#^{n^2-n}$ in $O(|x|^2)$ time,
 /x count backwards the number of #s and
 check if there are $|x|^2 - |x|$ such #s, in
 $O(|x|^2 - |x|) = O(|x|^2)$ time */

1.1. If Not, output REJECT

1.2. Else, run M_A on x for $O(|x|^6)$ time and
 output its answer. /x ignoring the #s */ "

Input length = $|x|^2$



Clearly the running time for $M_{A'} = O(|x|^2) + O(|x|^6) = O(|x|^6)$.

but input length of $M_{A'} = |x|^2 = n$ (let)

$\Rightarrow \text{runtime} = O((|x|^2)^3) = O(n^3)$ w.r.t. the length of input

$\Rightarrow A' \in \text{TIME}(n^3)$

(Also, $M_{A'}$ always halts since
 M_A " " ")

2.

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To prove: $P \neq \text{SPACE}[n^3]$

By contradiction

Let's assume $P = \text{SPACE}[n^3]$, to the contrary.

By Space Hierarchy Theorem \exists a language $A \mid A \in \text{SPACE}[n^6]$
but $A \notin \text{SPACE}[n^3]$, since $n^3 \neq o(n^6)$... ①

Now, the language $A' = \text{pad}(A, n^2) \in \text{SPACE}[n^3] = P$, we can show using the similar technique as in 9.13 as follows:

Construct TM, $M_{A'}$ from M_A (deciding A)

$M_{A'}(y)$

check if y is of the form $\underbrace{\# \# \#}_{|x|} \underbrace{x}_{|x|^2 - |x|} \#$ $|x|^2 = |x|$

/* Count # number of # using $\log(|x|^2 - |x|) = O(\log|x|^2)$ space */

if y no then REJECT

else run M_A on x and output whatever it does

/* uses $O(|x|^6)$ space */

Hence total space used by $M_{A'}$ is $O(|x|^6)$ on input ^{with} length $|x|^2$

$\Rightarrow A' \in \text{SPACE}[n^3] = P$

$\Rightarrow A \in P$ (since we can determine $w \in A$ by padding w with $|w|^2 - |w|$ #'s and then testing whether $w' = w \# |w|^2 - |w| \in A'$ in polynomial time)

But $P = \text{SPACE}[n^3]$, by assumption

$\Rightarrow A \in \text{SPACE}[n^3]$... ②

① & ② $\Rightarrow \Leftarrow$

(Proved)