## Facility location



#### Linear programming relaxation

Primal: minimize 
$$\sum_{i \in F} \sum_{j \in C} c_{ij} x_{ij} + \sum_{i \in F} f_i y_i$$

Dual variables:

subject to

$$\sum x_{ij} \geq 1,$$

$$j \in C$$

$$\alpha_j$$

$$y_i - x_{ij} \geq 0,$$

$$j \in C, i \in F$$

$$\beta_{ij}$$

$$y_i \geq 0$$
,

 $i \in F$ 

$$i \in F$$

$$x_{ij} \geq 0,$$

$$j \in C, i \in F$$

#### Taking the dual

Dual: maximize

$$\sum_{j \in C} \alpha_j$$

Primal variables

subject to

$$\alpha_j - \beta_{ij} \leq c_{ij},$$

$$j \in C, i \in F$$

$$x_{ij}$$

$$\sum \beta_{ij} \leq f_i,$$

$$i \in F$$

$$y_i$$

$$\alpha_j \geq 0$$
,

 $j \in C$ 

$$j \in C$$

$$\beta_{ij} \geq 0,$$

$$j \in C, i \in F$$

#### Recall:

#### Complementary slackness conditions

# If x is optimal for (P) and y optimal for (D) then for every i:

$$\begin{aligned} \mathbf{c_i} &= \sum_{j} \mathbf{a_{ij}y_j} \text{ or } \mathbf{x_i} = \mathbf{0} \\ &\text{ and for every j:} \\ \mathbf{b_j} &= \sum_{i} \mathbf{a_{ij}xi} \text{ or } \mathbf{y_j} = \mathbf{0} \end{aligned}$$

## Complementary slackness conditions:

1. 
$$\forall i \in F, j \in C: x_{ij} > 0 \implies \alpha_j - \beta_{ij} = c_{ij}$$

$$2. \ \forall i \in F: y_i > 0 \implies \sum_{j \in C} \beta_{ij} = f_i$$

3. 
$$\forall j \in C: \alpha_j > 0 \implies \sum_{i \in F} x_{ij} = 1$$

$$4. \ \forall i \in F, j \in C: \beta_{ij} > 0 \implies y_i = x_{ij}$$

 $eta_{ij}$ : Contribution of client j for opening facility i

Complementary slackness condition 2:  $\forall i \in F: y_i > 0 \implies \sum_{j \in C} \beta_{ij} = f_i$ 

If 
$$y_i=1$$
 then  $\sum_{j\in C} \beta_{ij}=f_i$  we say that facility i is fully paid

## Complementary slackness condition 4: $\forall i \in F, j \in C: \beta_{ij} > 0 \implies y_i = x_{ij}$

thus, 
$$\forall i \in F, j \in C: y_i \neq x_{ij} \implies \beta_{ij} = 0$$

Recall 
$$y_i - x_{ij} \ge 0$$
, If  $y_i \ne x_{ij}$  then  $\beta_{ij} = 0$ 

and so, client j does not contribute to opening any facility except the one it is connected to

#### Complementary slackness

condition 1:  $\forall i \in F, j \in C: x_{ij} > 0 \implies \alpha_j - \beta_{ij} = c_{ij}$ 

#### thus, if client j is assigned to facility i we have:

$$\alpha_j - \beta_{ij} = c_{ij}$$

We define  $\alpha_j$  as the total price paid by client j.

The total price is made of the use of edge (i,j)  $c_{ij}$  and the contribution to the opening cost  $\beta_{ij}$ 

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