

# Steiner forest

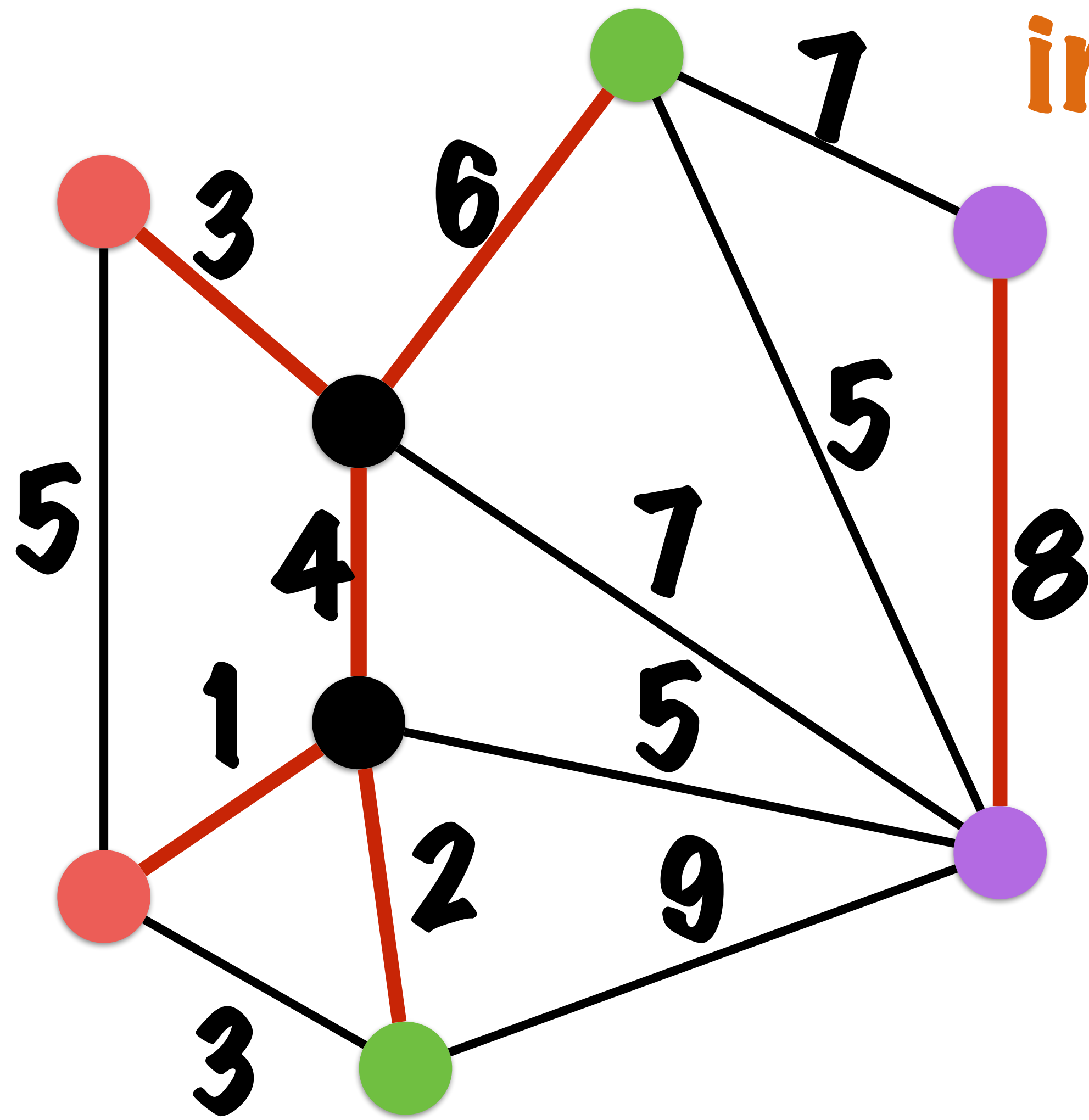




# Integer programming model

**input:** sets of terminals  $S_1, S_2, \dots$

**output:** set of edges



**min output cost**

**s.t.**

**for all  $i$**

**for all  $u, v$  in  $S_i$**

**$u$  and  $v$  are connected**

# Integer programming model

**output:** set of edges

$$x_e \in \{0, 1\}$$

**min output cost**  
**s.t.**

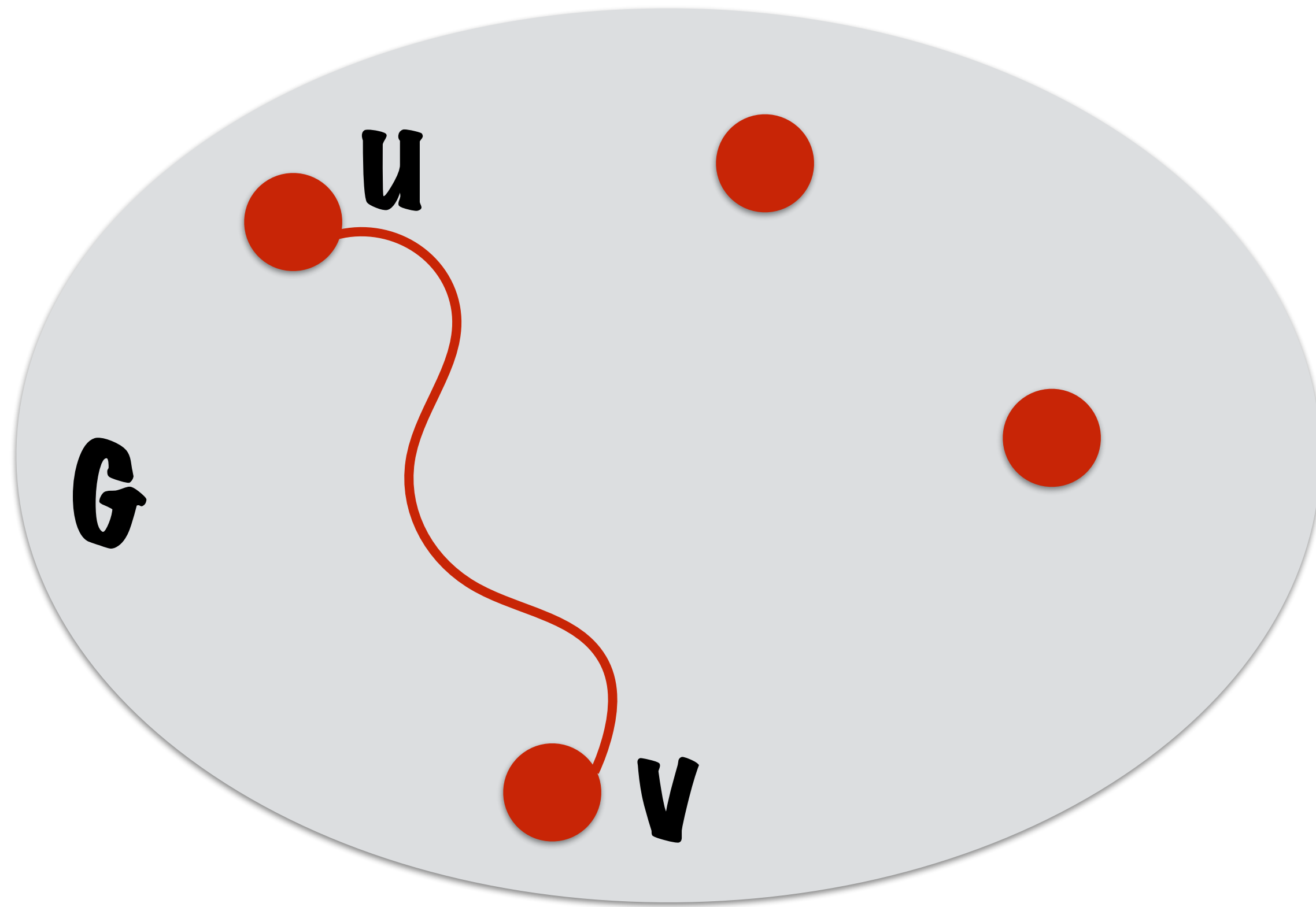
$$\min \sum_e c_e x_e$$

**for all  $i$**   
**for all  $u, v$  in  $S_i$**   
 **$u$  and  $v$  are connected**

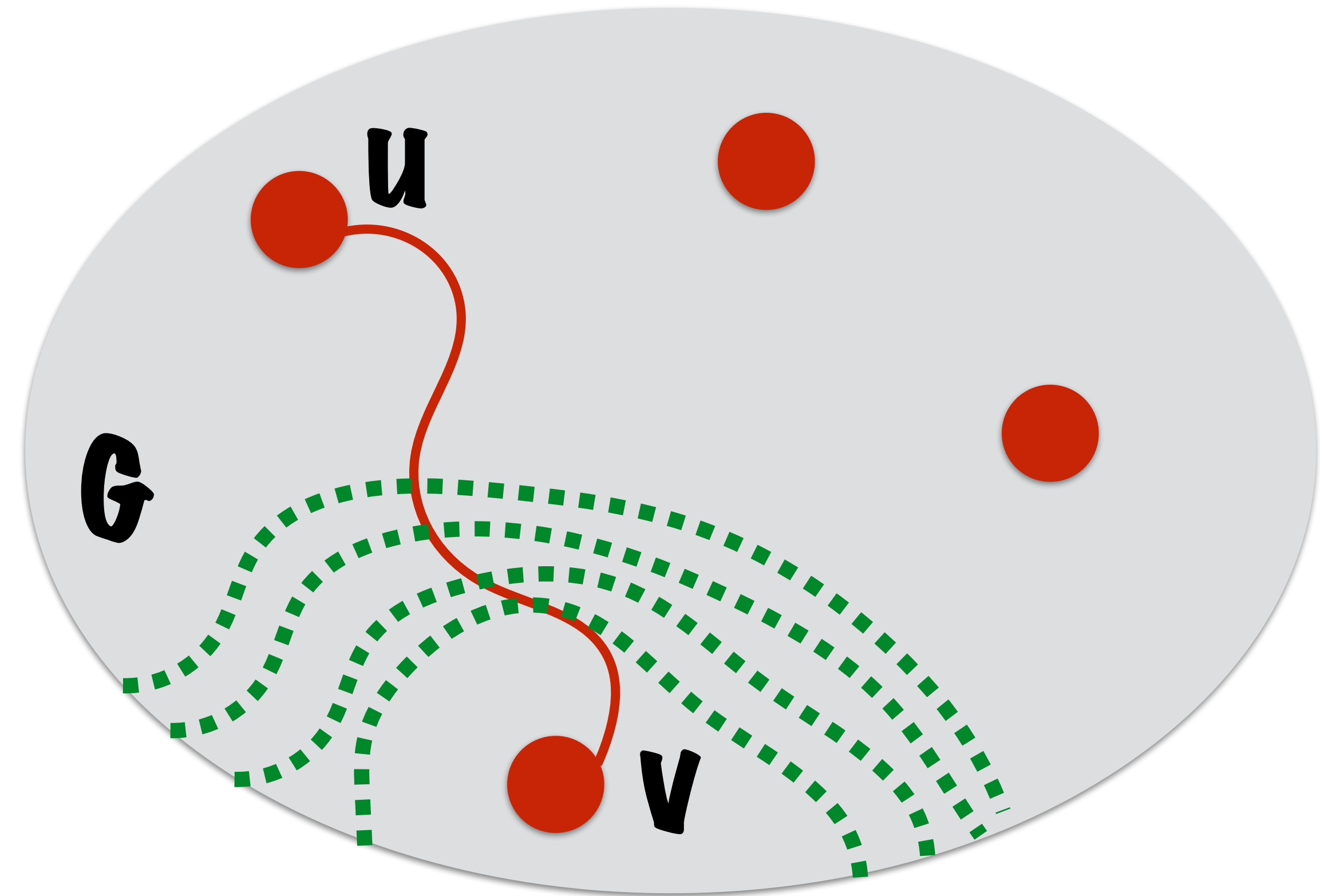
$$\forall i \forall u, v \in S_i$$

?

# Integer programming model

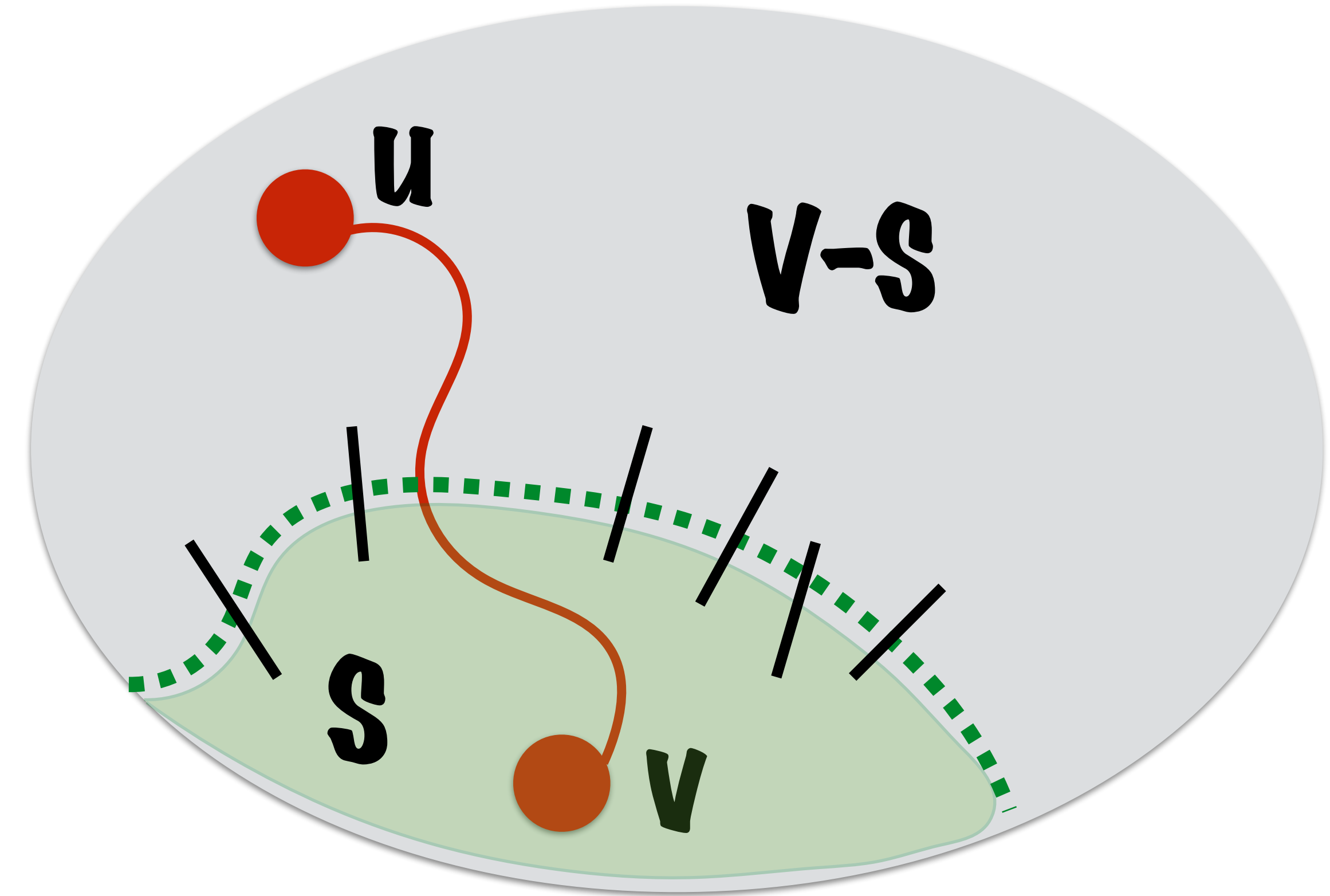
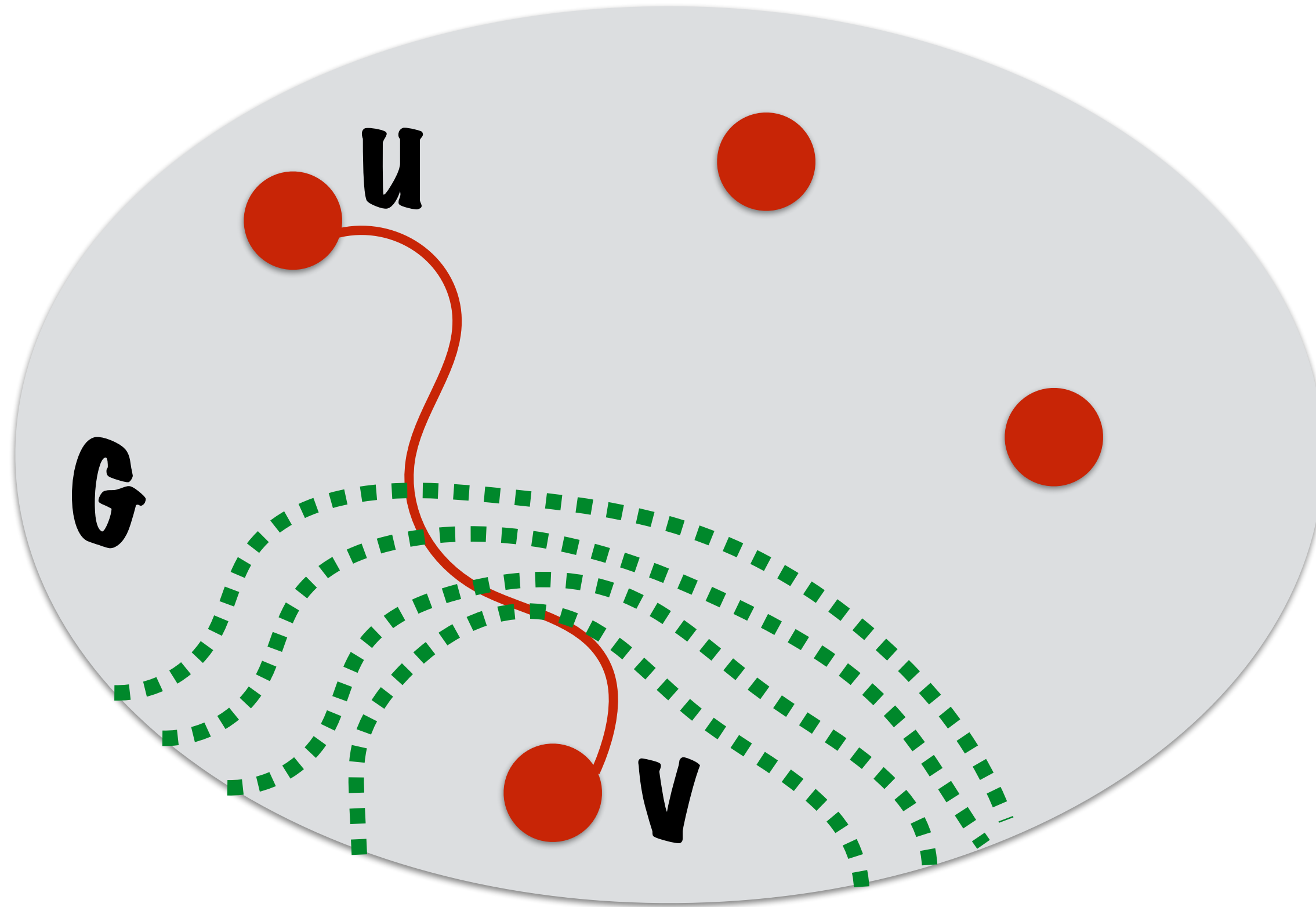


**$u$  and  $v$  are connected**



**every  $u$ - $v$  cut is crossed**

# Integer programming model



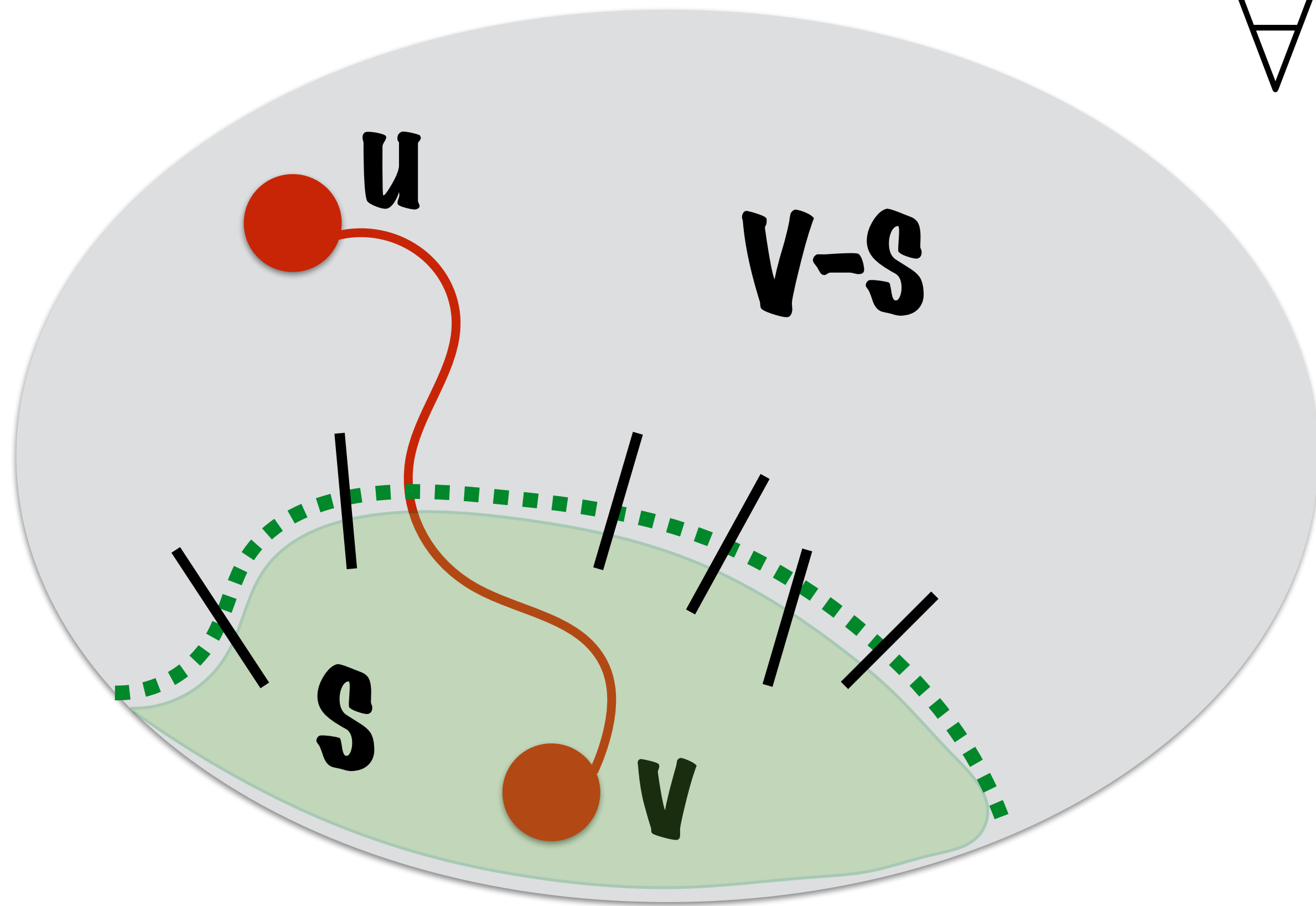
**every  $u$ - $v$  cut is crossed**

$$S \subseteq V : |S \cap \{u, v\}| = 1$$

$$\delta(S) = \{e \in S \times V \setminus S\}$$

for all  $i$   
 for all  $u, v$  in  $S_i$   
 $u$  and  $v$  are connected

$$x_e \in \{0, 1\}$$

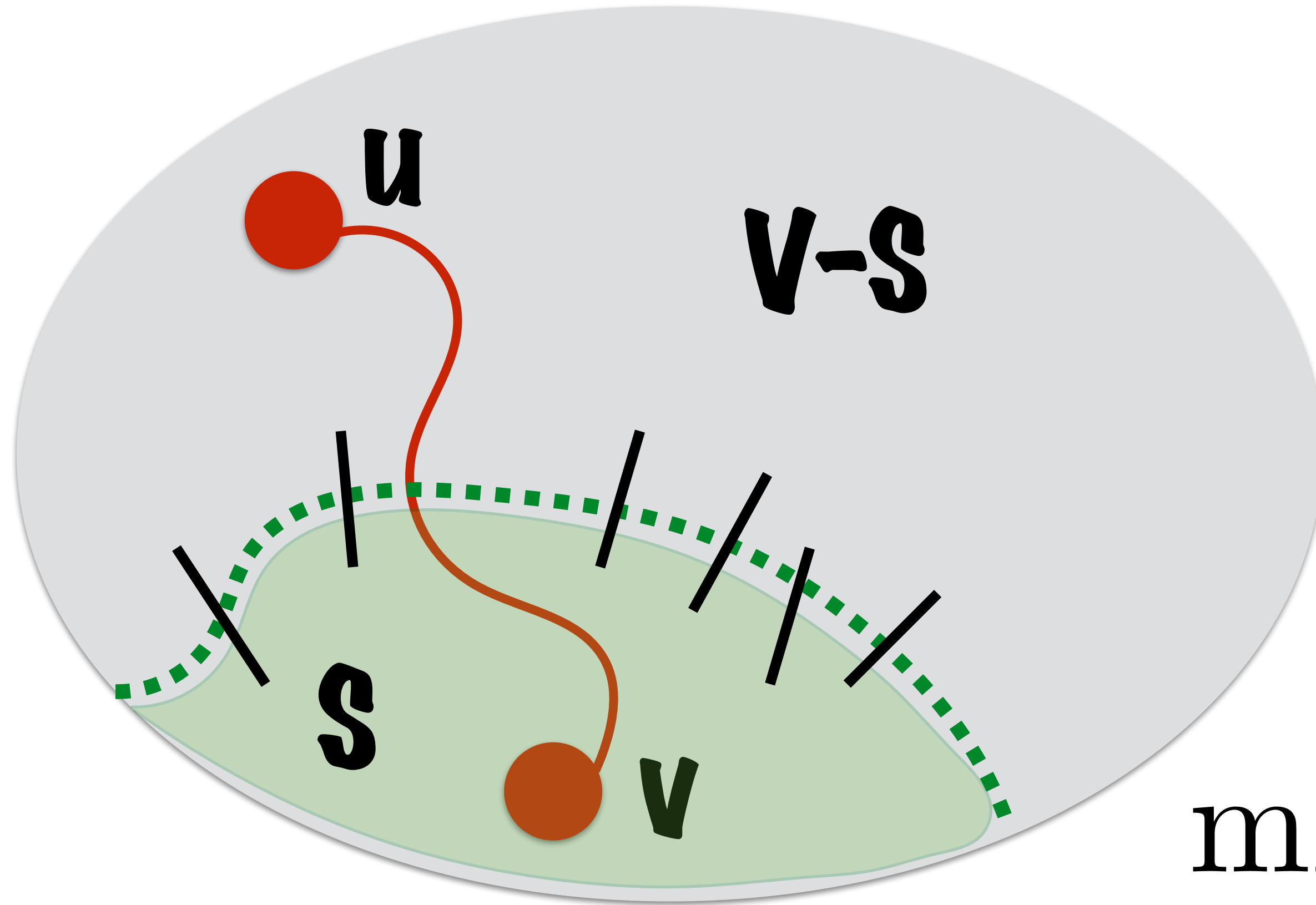


$$\forall i \forall u, v \in S_i$$

$$\forall S \text{ s.t. } |S \cap \{u, v\}| = 1$$

$$\sum_{e \in \delta(S)} x_e \geq 1$$

$$\mathcal{S} = \{\mathbf{S} : \exists i \exists \mathbf{u}, \mathbf{v} \in \mathbf{S}_i : |\mathbf{S} \cap \{\mathbf{u}, \mathbf{v}\}| = 1\}$$



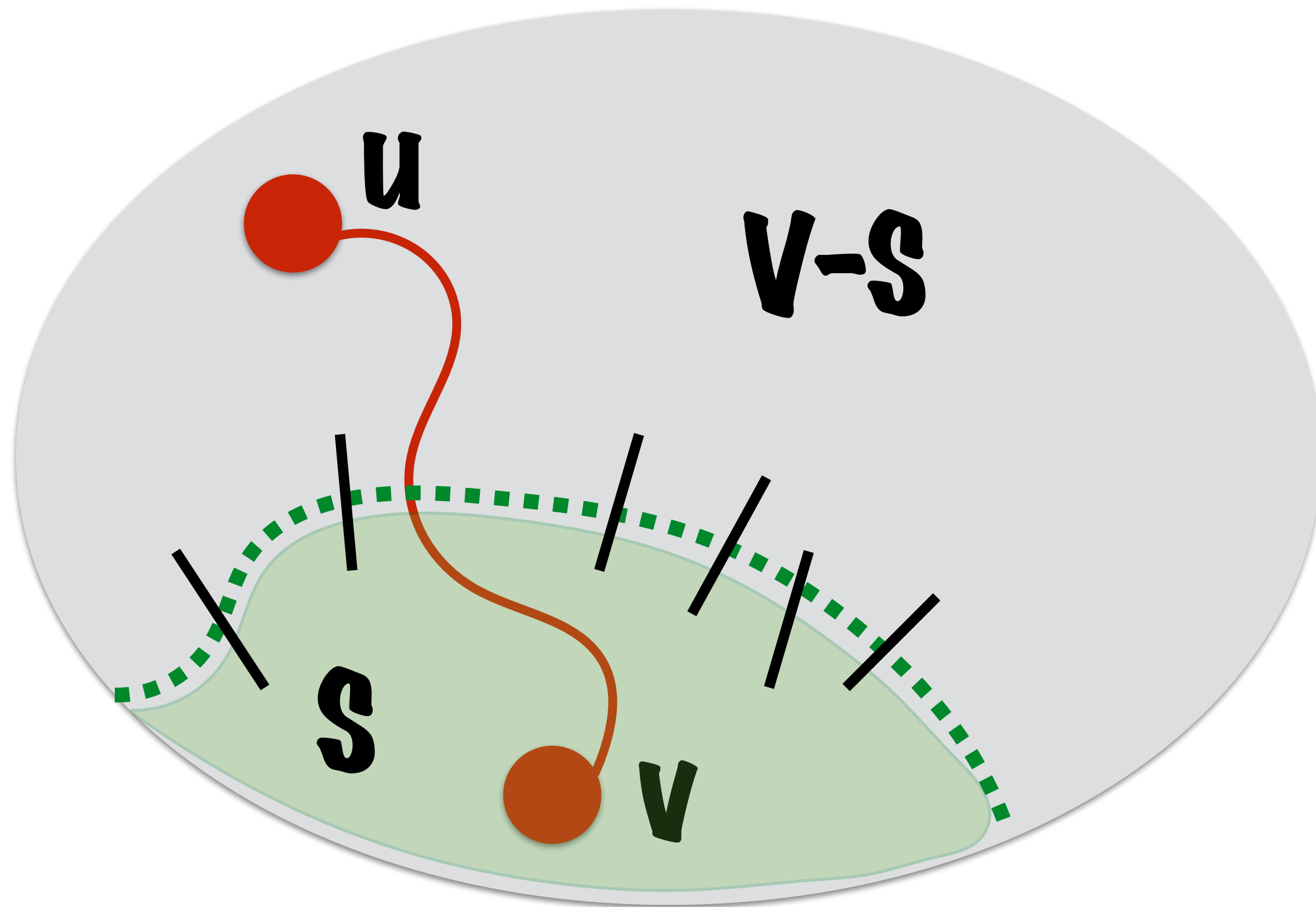
$$\min \sum_e c_e x_e :$$

$$\forall \mathbf{S} \in \mathcal{S} \sum_{e \in \delta(\mathbf{S})} x_e \geq 1$$

$$x_e \in \{0, 1\}$$

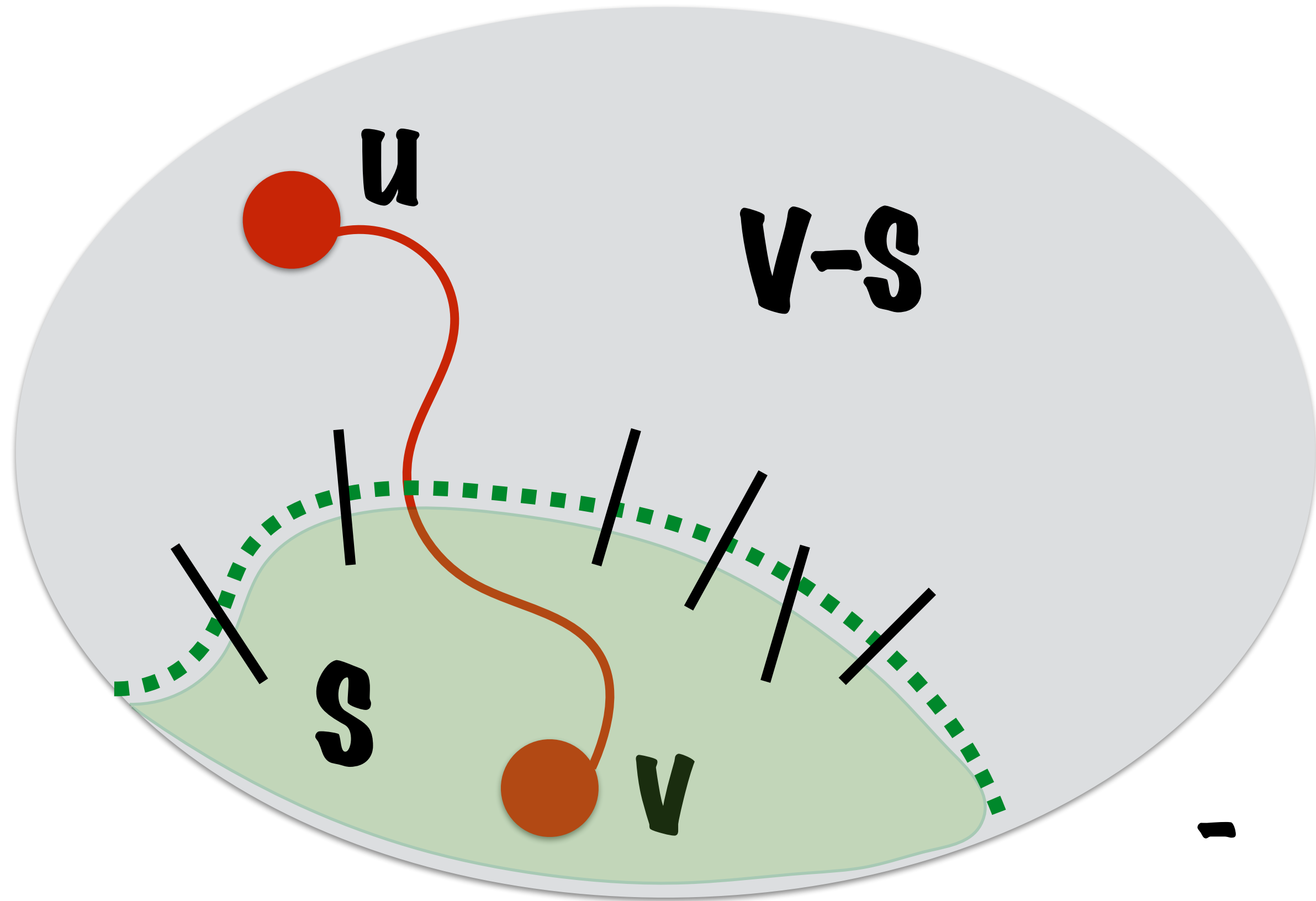
$$\begin{aligned} & \min \sum_e c_e x_e : \\ & \forall S \in \mathcal{S} \sum_{e \in \delta(S)} x_e \geq 1 \\ & x_e \in \{0, 1\} \end{aligned}$$

$$\begin{aligned} & \min \sum_e c_e x_e : \\ & \forall S \in \mathcal{S} \sum_{e \in \delta(S)} x_e \geq 1 \\ & x_e \geq 0 \end{aligned}$$





$$\begin{aligned} & \min \sum_e c_e x_e : \\ & \forall S \in \mathcal{S} \sum_{e \in \delta(S)} x_e \geq 1 \\ & x_e \geq 0 \end{aligned}$$



- Exponential  
number of constraints  
but  
still solvable**
- - **does not need to be solved**



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