

# Case Study 2: Document Retrieval

## Spectral Clustering

Machine Learning/Statistics for Big Data  
CSE599C1/STAT592, University of Washington

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# Document Retrieval

- Goal: Retrieve documents of interest



# Task 1: Find Similar Documents

## ■ Setup

- **Input:** Query article **X**
- **Output:** Set of k similar articles



# k-Nearest Neighbor

- Articles  $X = \{x^1, \dots, x^N\}, \quad x^i \in \mathbb{R}^d$

- Query:  $x \in \mathbb{R}^d$

- k-NN

- Goal:

Find k articles in  $X$  closest  $x$

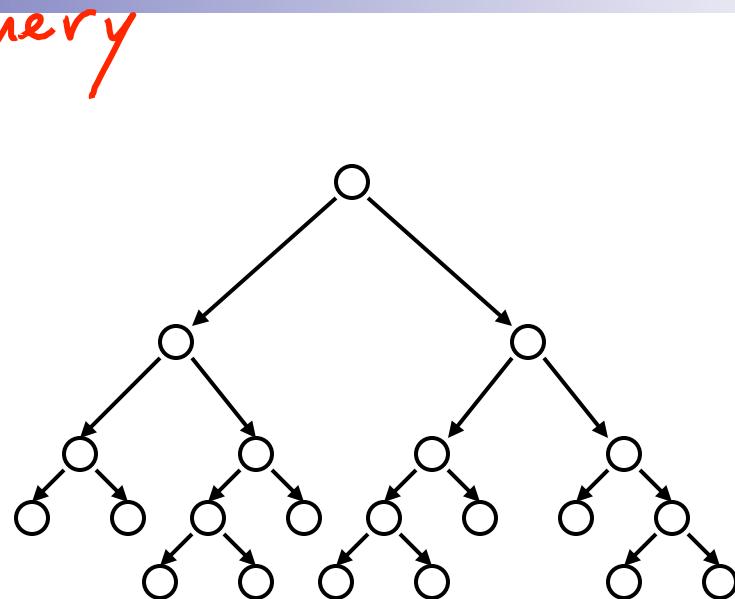
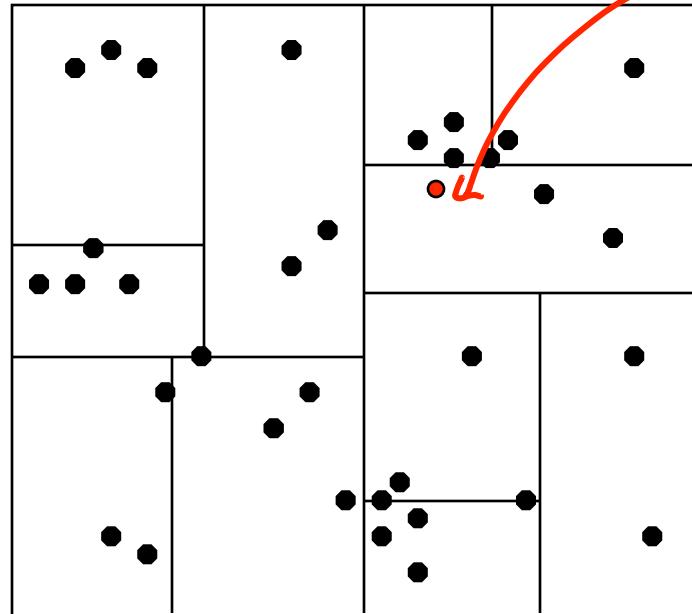
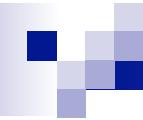
- Formulation:

$$X^{NN} = \{x^{NN_1}, \dots, x^{NN_k}\} \subseteq X$$

$$\text{s.t. } \forall x^i \in X \setminus X^{NN}$$

$$d(x^i, x) \geq \max_{x^{NN_i} \in X^{NN}} d(x^{NN_i}, x)$$

# Nearest Neighbor with KD Trees



- Traverse the tree looking for the nearest neighbor of the query point.

# Task 2: Cluster Documents

## ■ Setup

- **Input:** Corpus of documents
- **Output:** Topic assignment per document



# A Generative Model

- Documents:  $x^1, \dots, x^D$
- Associated topics:  $\underline{z^1, \dots, z^D}$
- Parameters:  $\theta = \{\pi, \beta\}$
- Generative model:

$$z^d \sim \pi \quad d=1, \dots, D$$

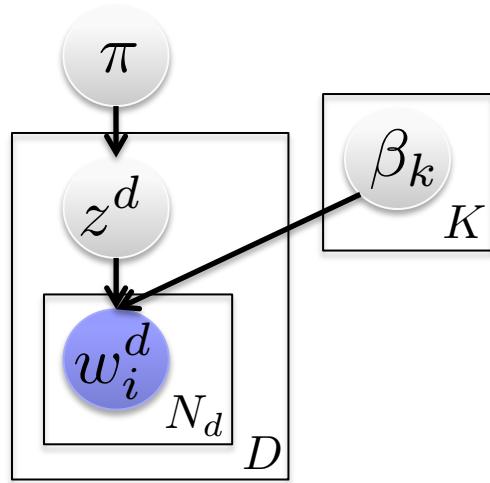
$$w_i^d | z^d \sim \beta_{z^d} \quad i=1, \dots, N$$

↑ word prob. for cluster/topic  $z^d$

Bayesian approach:

$$\pi \sim \text{Dir}(\alpha_1, \dots, \alpha_K)$$

$$\beta_k \sim \text{Dir}(\lambda_1, \dots, \lambda_V)$$



$\beta_k$  is a V-dim pmf  
size of vocab.

# Inference

- Two tasks

- Point estimation:

- Expectation-Maximization (EM)

- Characterize posterior:

- Gibbs sampling
    - Variational methods
    - Stochastic variational inference

$$\hat{\theta}^{\text{ML}}, \text{ or } \hat{\theta}^{\text{MAP}} \leftarrow p(\theta)_{\text{prior}}$$

# EM Algorithm

- Initial guess:  $\hat{\theta}^{(0)}$
- Estimate at iteration  $t$ :  $\hat{\theta}^{(t)}$

- E-Step

Compute  $U(\theta, \hat{\theta}^{(t)}) = E[\log p(y | \theta) | x, \hat{\theta}^{(t)}]$

- M-Step

Compute  $\hat{\theta}^{(t+1)} = \arg \max_{\theta} U(\theta, \hat{\theta}^{(t)}) + \log P(\theta)$

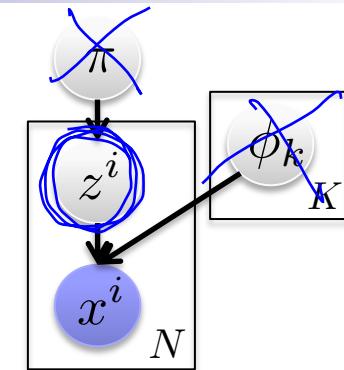
# Collapsed Gibbs Sampling

$$\begin{aligned}\pi &\sim \text{Dir}(\alpha_1, \dots, \alpha_K) \\ \{\mu_k, \Sigma_k\} &\sim F(\phi)\end{aligned}$$

$$z^i \sim \pi$$

$$x^i | z^i \sim N(x^i; \mu_{z^i}, \Sigma_{z^i})$$

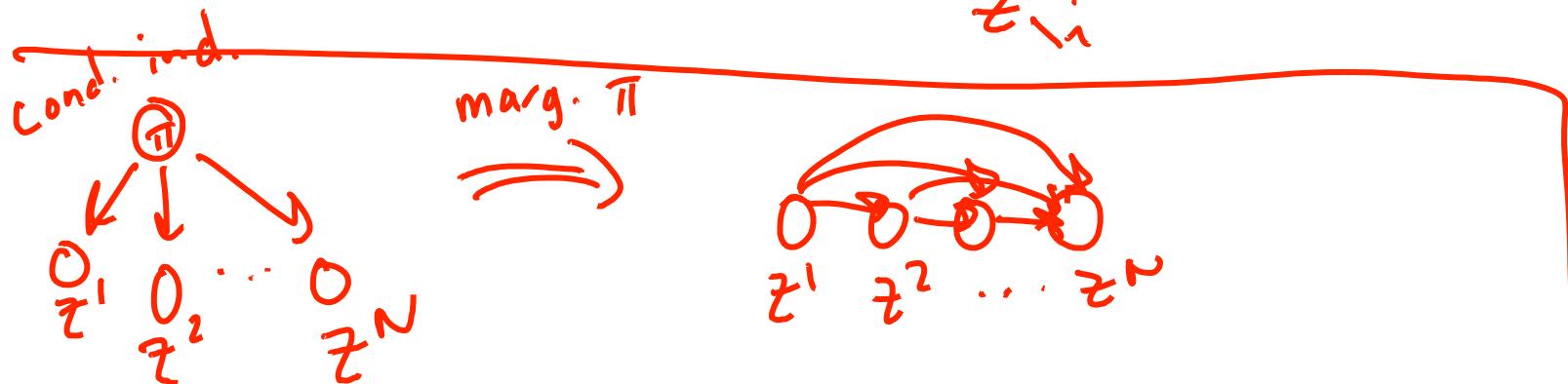
- Collapsed sampler



For  $i=1, \dots, N$

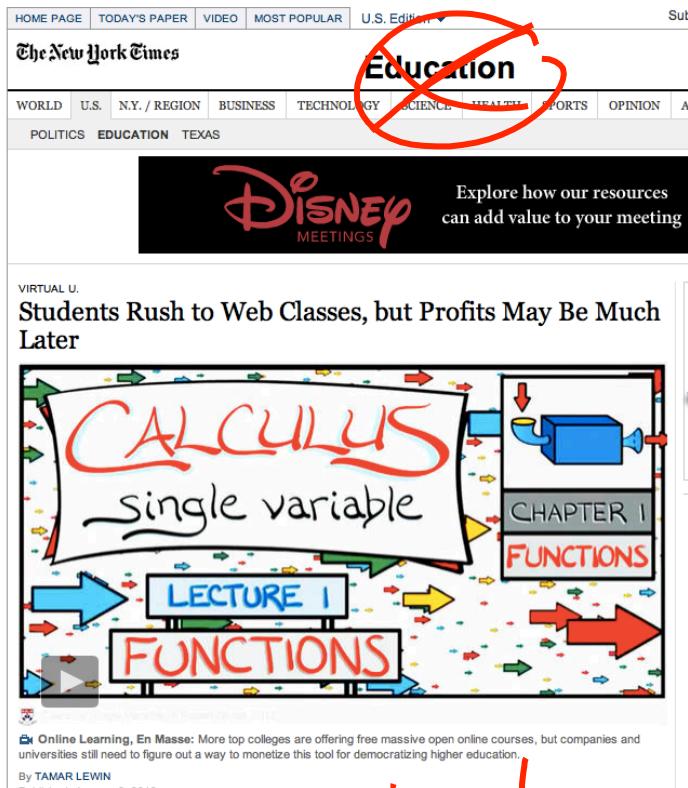
$$z^{i(t)} \sim p(z^i | z^{1(t)}, \dots, z^{i-1(t)}, z^{i+1(t)}, \dots, z^{N(t)}, x_{1:N}, \alpha, \lambda)$$

$\underbrace{\hspace{10em}}$   
 $z_{\setminus i}^{(t)}$

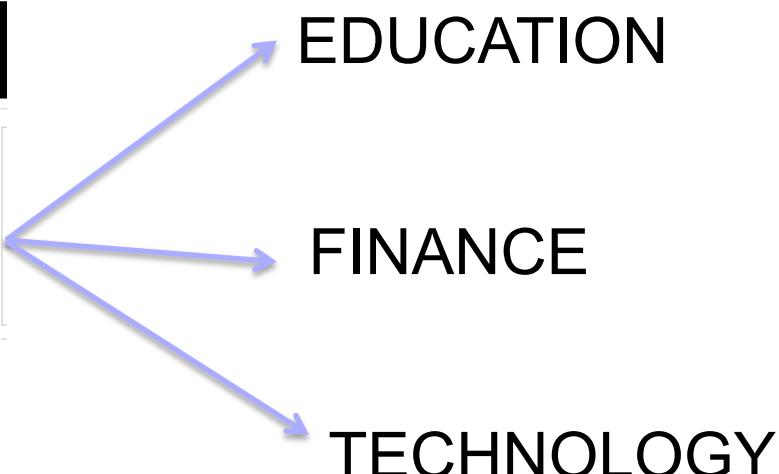


# Task 3: Mixed Membership Model

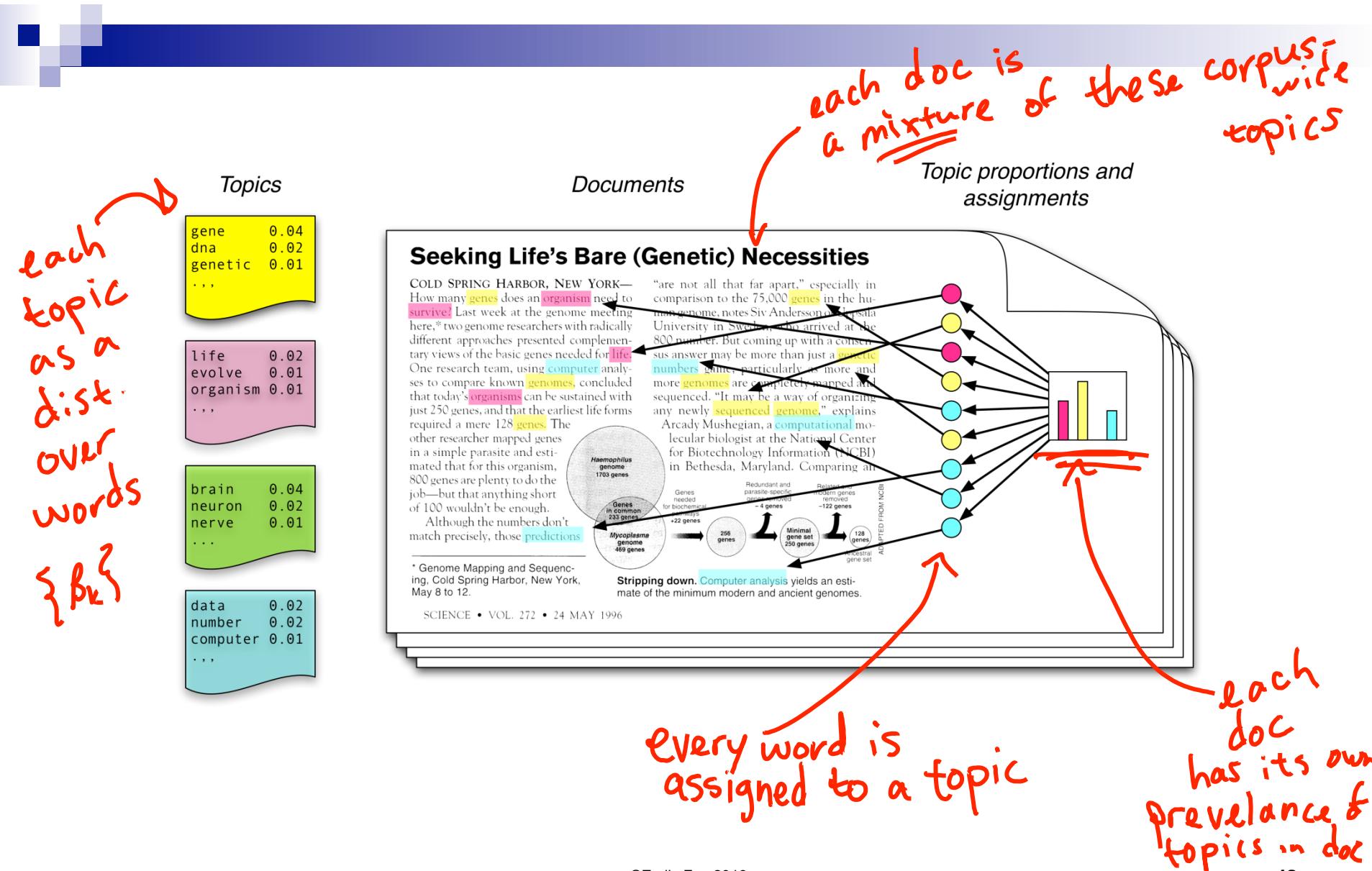
- **Setup:** Document may belong to multiple clusters



mixed membership



# Latent Dirichlet Allocation (LDA)



# Variational Methods

- Recall task: Characterize the posterior  $p(\theta, z | x)$   
    params   latent vars   obs
- Turn posterior inference into an optimization task
- Introduce a “tractable” family of distributions over parameters and latent variables
  - Family is indexed by a set of “free parameters”
  - Find member of the family closest to:  $p(\theta, z | x)$   
*call the family  $Q$  and want  $q \in Q$  that  
is closest to  $p(\theta, z | x)$*
- Questions:
  - How do we measure “closeness”?
  - If the posterior is intractable, how can we approximate something we do not have to begin with?

# Variational Methods

- Similarity measure:

$$\begin{aligned} D(q(z, \theta) \parallel p(z, \theta | x)) &= E_q[\log q(z, \theta)] - E_q[\log p(z, \theta | x)] \\ &= E_q[\log q(z, \theta)] - E_q[\log p(z, \theta | x)] \end{aligned}$$

$\underbrace{-L(q)}_{\text{add to a const}} \quad \underbrace{\neq \log p(x)}_{\text{const.}}$

- Evidence lower bound (ELBO)

$$\underbrace{\log p(x)}_{\text{const.}} = D(q(z, \theta) \parallel p(z, \theta | x)) + L(q) \geq L(q)$$

$\xrightarrow{\text{add to a const}} \quad \equiv$

- Therefore, minimizing KL is equivalent to maximizing a lower bound on the marginal likelihood:

□ Max  $L(q) = \min D(q || p) = \max \text{lower bound of } \log p(x)$   $\xleftarrow{\text{of } q} \text{entropy}$

$$L(q) = E_q[\log p(\theta, z, x)] \neq E_q[\log q(\theta, z)]$$

$\underbrace{\quad}_{\text{const.}}$

# Task 2: Cluster Documents

## ■ Setup

- **Input:** Corpus of documents
- **Output:** Topic assignment per document



Sports



FIFA WORLD CUP  
Brasil

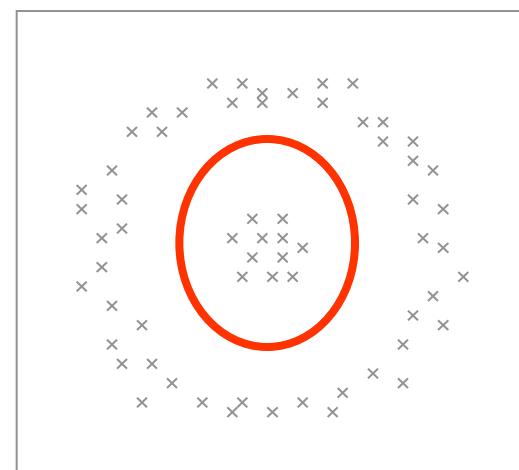
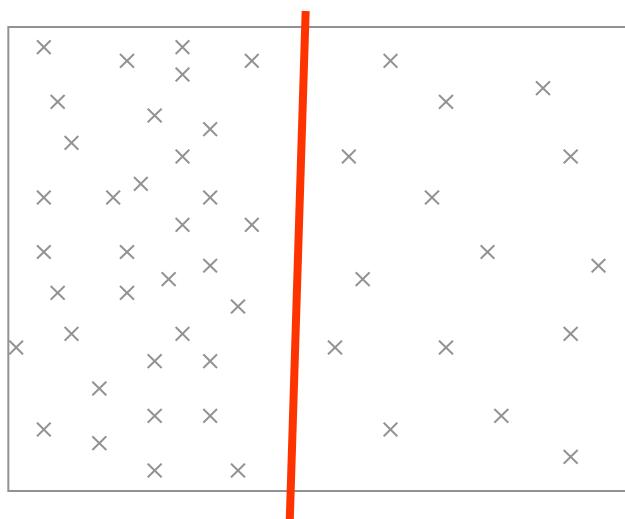


world news



# New Approach: Spectral Clustering

- **Goal:** Cluster observations
- **Method:**
  - Use similarity metric between observations
  - Form a similarity graph
  - Use standard linear algebra and optimization techniques to cut graph into connected components (clusters)



# Setup

- Data:  $x^1, \dots, x^N$   $x^i = \text{doc } i \dots$  maybe use tf-idf
- Similarity metric:

$s_{ij}$  bt  $x^j$  and  $x^i$  (eg cosine similarity)  $x^i \in \mathbb{R}^V$

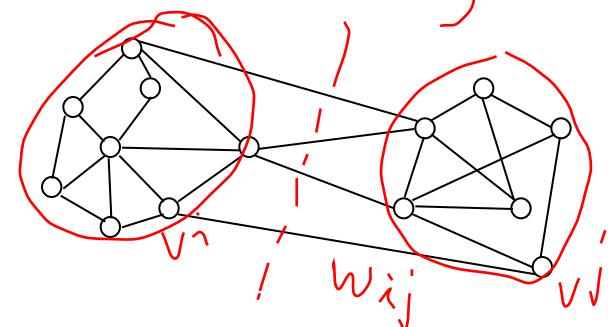
$$s_{ij} = s_{ji}$$

- Similarity graph

□ Nodes  $\checkmark^i$  for each  $x^i$

□ Edge weights

$$w_{ij} = \text{fcn of } s_{ij}$$



$$G = \{V, E\}$$

- Problem: Want to partition graph such that edges between groups have low weights

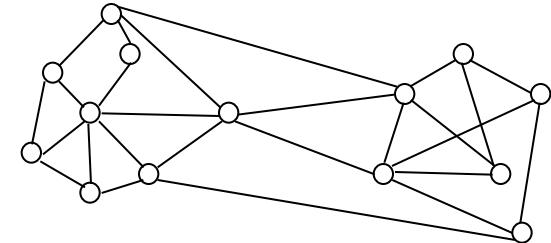
# Types of Graphs

- **$\epsilon$ -neighborhood:**

- Only include edges with distances  $< \epsilon$
- Treat as unweighted  $w_{ij} = \epsilon$

- **k-NN:**

- Connect  $v_i$  and  $v_j$  if  $v_j$  is a k-NN of  $v_i$
- Weighted by similarity  $s_{ij} = w_{ij}$
- Directed  $\rightarrow$  undirected



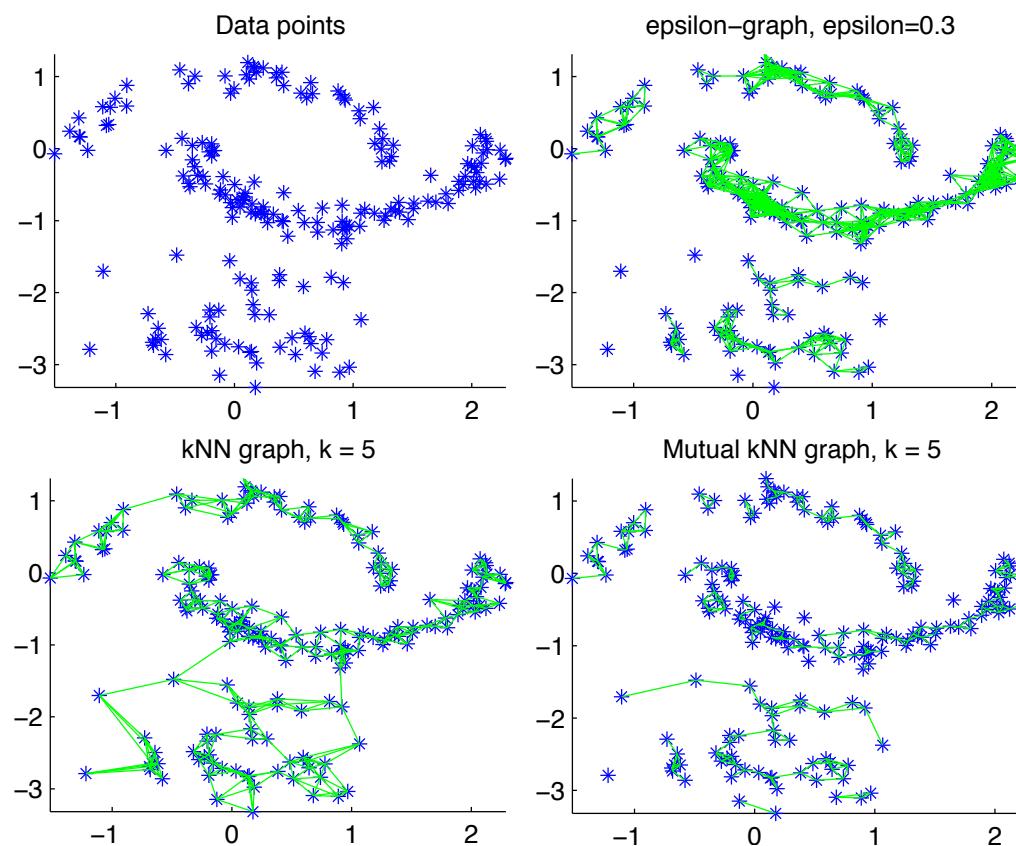
- **Mutual k-NN:**

- Same as k-NN, but only include mutual k-NN

# Issues with Choosing Graph

- Choosing graph construction techniques and parameters is non-trivial

choice  
matters



From  
von Luxburg  
2007

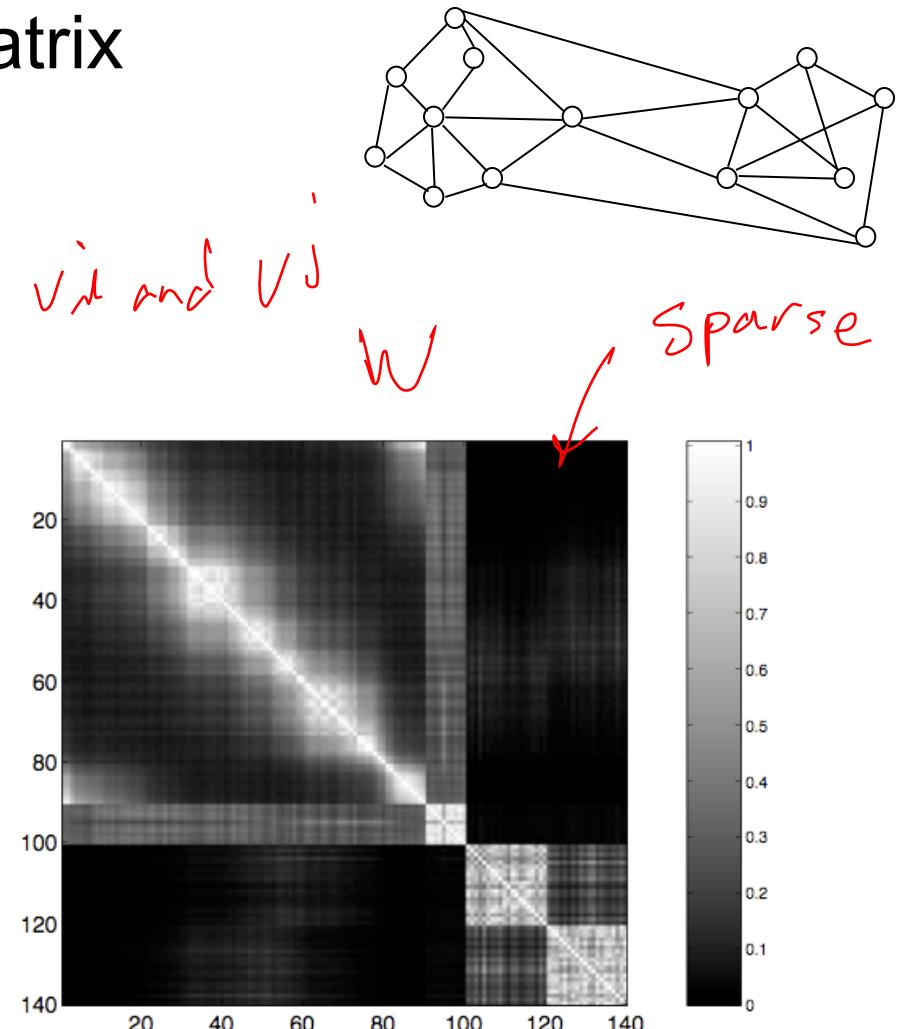
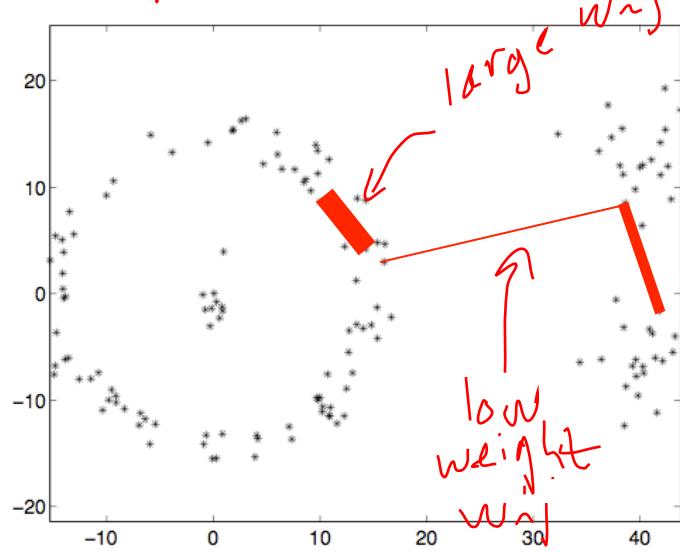
# Graph Terminology I

## ■ Weighted adjacency matrix

$$W = (w_{ij})_{i,j=1,\dots,N}$$

$w_{ij} = 0 \Rightarrow$  no edge bt  $v_i$  and  $v_j$

$$w_{ij} \geq 0$$



# Graph Cuts

- **Problem:** Partition graph such that edges between groups have low weights

- Define:  $\underline{W(A, B)} = \sum_{i \in A, j \in B} w_{ij}$

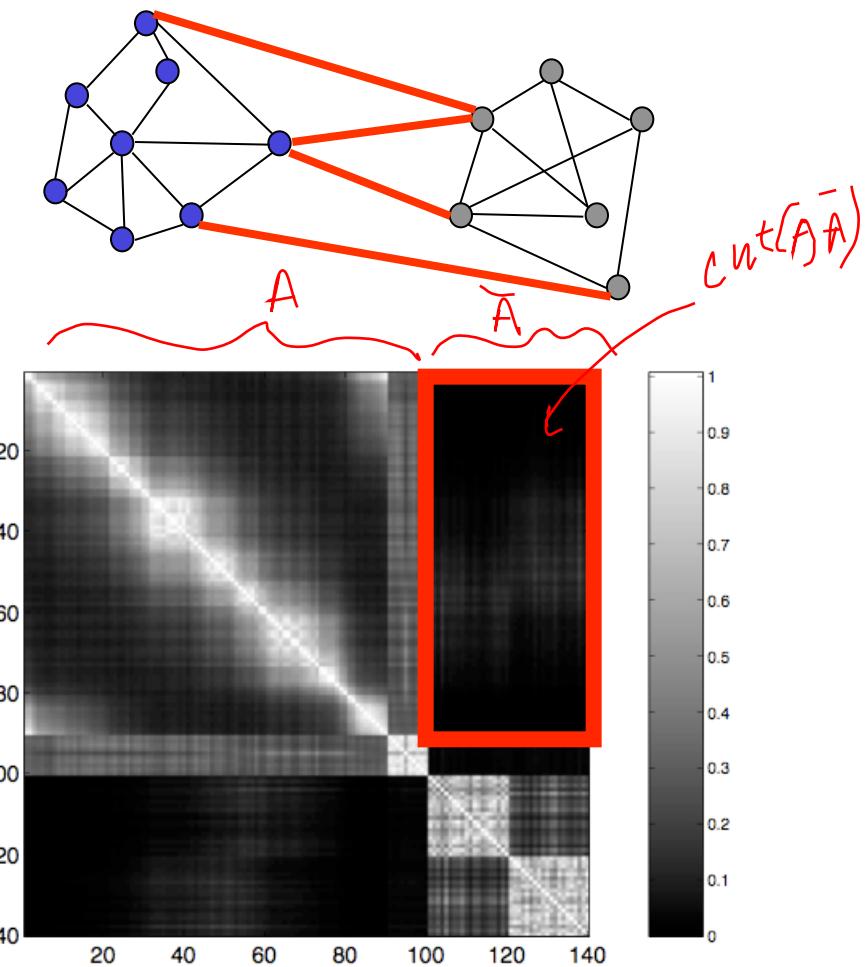
- MinCut problem:

$$\text{Cut}(A_1, \dots, A_k) \stackrel{k}{=} \frac{1}{2} \sum_{i=1}^k W(A_i, \bar{A}_i)$$

Choose

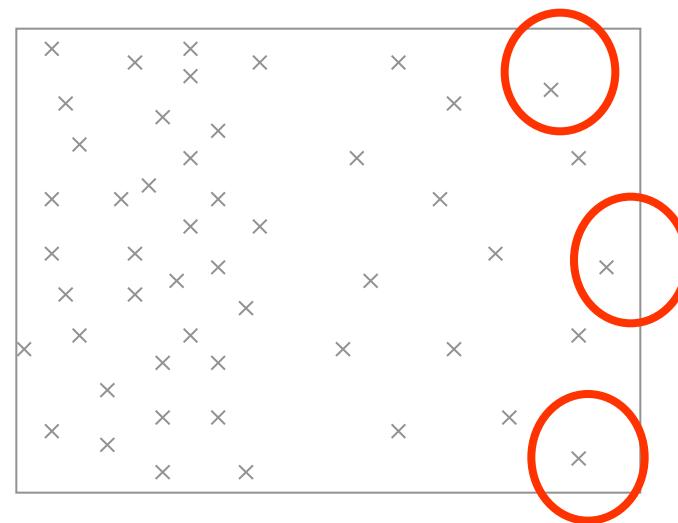
$$A_1, \dots, A_k = \underset{\substack{A_1, \dots, A_k \\ \text{disjoint} \subset V}}{\operatorname{arg\,min}} \text{Cut}(A_1, \dots, A_k)$$

- Trivial to solve for  $k=2$



# Issues with MinCut

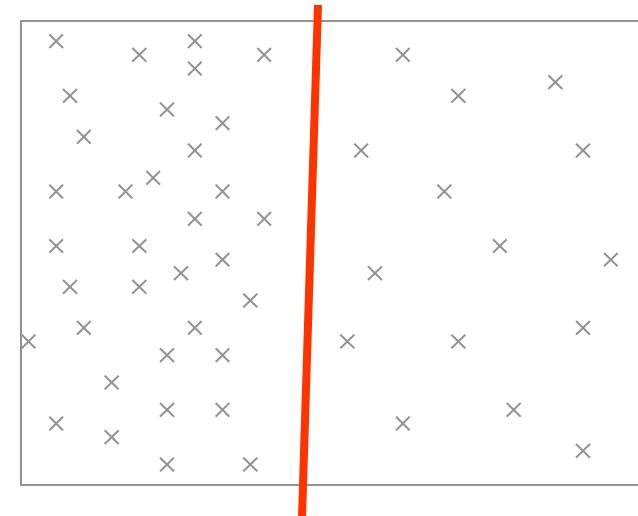
- MinCut favors isolated clusters



nothing working against this

# Cuts Accounting for Size

- Ratio cuts (RatioCut)
- Normalized cuts (Ncut)
- Lead to “balanced” clusters

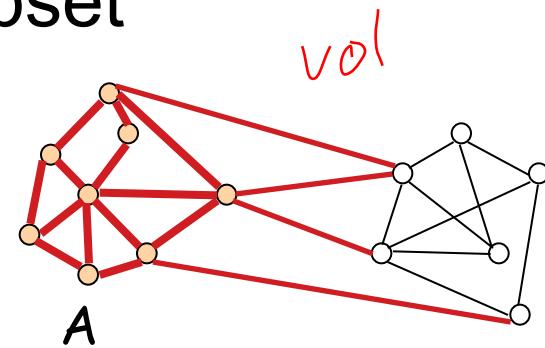


- First need more graph terminology...  
*to measure "size" of the clusters*

# Graph Terminology II

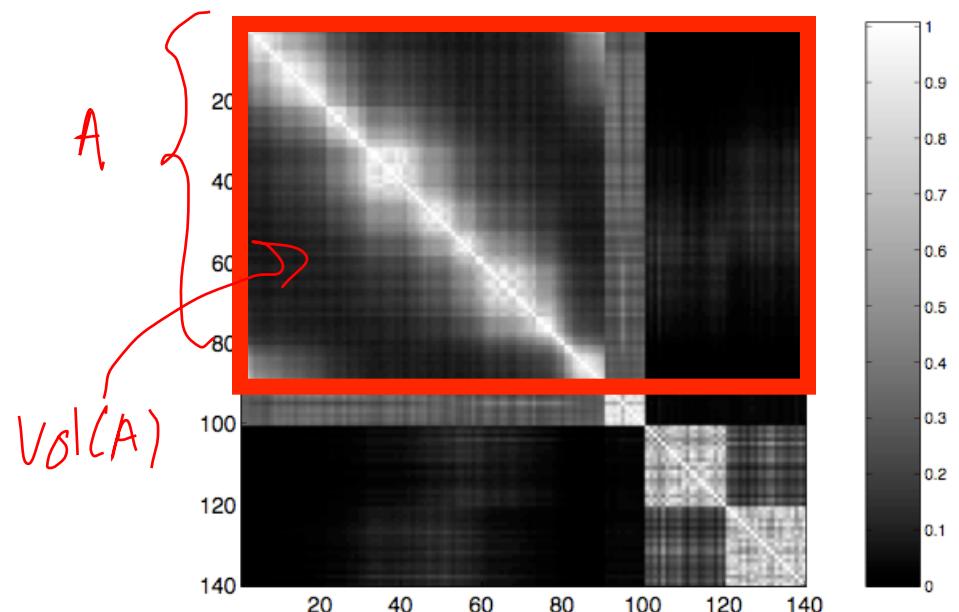
- Two measures of size of a subset
  - Cardinality:

$|A| = \# \text{ of vertices in } A$



- Volume:

$$\text{vol}(A) = \sum_{i \in A} \sum_{j=1}^N w_{ij}$$



# Cuts Accounting for Size

## ■ Ratio cuts (RatioCut)

□  $k=2$

$$\text{RatioCut}(A, \bar{A}) = \underbrace{\text{cut}(A, \bar{A})}_{\sim} \left( \frac{1}{|A|} + \frac{1}{|\bar{A}|} \right)$$

□ General  $k$

$$\text{RatioCut}(A_1, \dots, A_k) = \frac{1}{2} \sum_i \frac{w(A_i, \bar{A}_i)}{|A_i|}$$

*min when  
|A| and  
|Ā|  
coincide*

## ■ Normalized cuts (Ncut)

□  $k=2$

$$\text{Ncut}(A, \bar{A}) = \text{cut}(A, \bar{A}) \left( \frac{1}{\text{vol}(A)} + \frac{1}{\text{vol}(\bar{A})} \right)$$

□ General  $k$

$$\text{Ncut}(A_1, \dots, A_k) = \frac{1}{2} \sum_i \frac{w(A_i, \bar{A}_i)}{\text{vol}(A_i)}$$

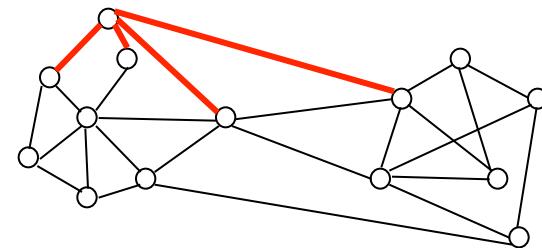
## ■ Problem is NP-hard! Look at relaxation.

# Graph Terminology III

## ■ Degree

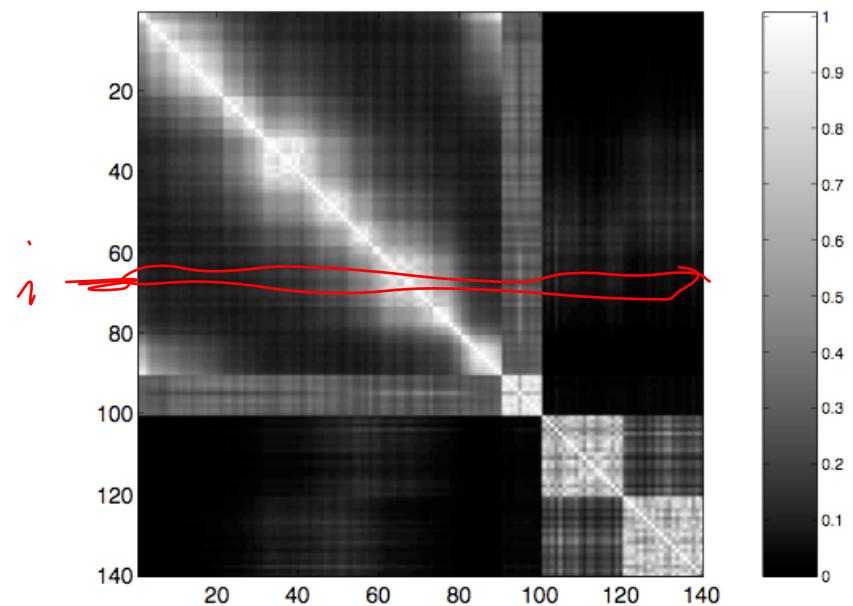
$$d_i = \sum_{j=1}^N w_{i,j}$$

only counts  
neighbors

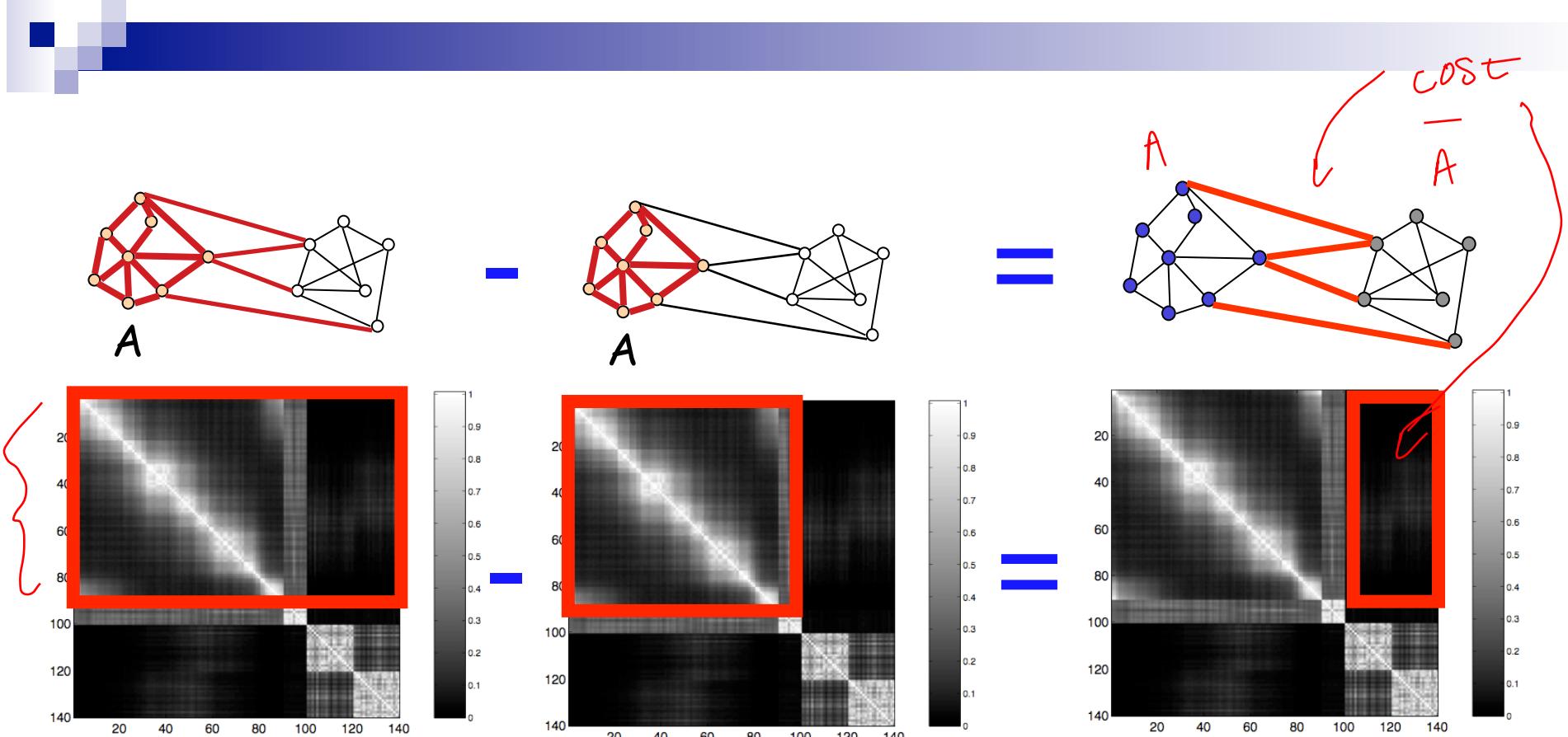


## ■ Degree matrix

$$D = \begin{bmatrix} d_1 & & \\ & \ddots & \\ & & d_N \end{bmatrix}$$



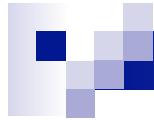
# Restating Cut Metric



$\text{Volume}(A)$  - "association( $A$ )" = cut cost

# Restating Cut Metric

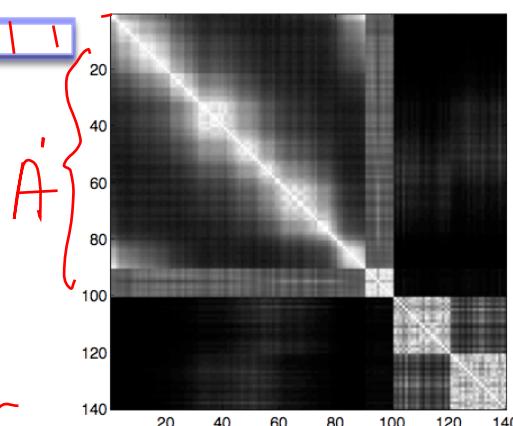
- Assoc.



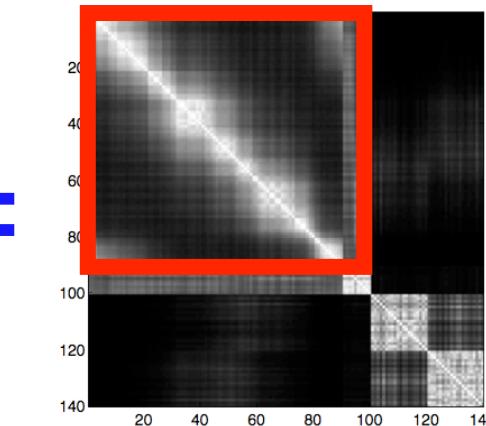
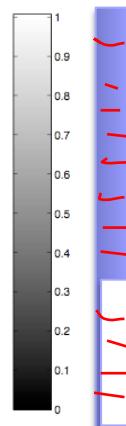
$x^T$



$W$



$x$

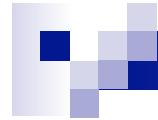


binary vector  
that's an  
indicator on set A

$\mathbb{1}_A$

# Restating Cut Metric

- Volume



$$x^T \quad D \quad x$$

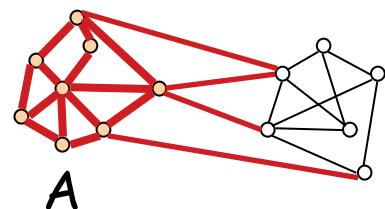
$$\begin{bmatrix} d_1 & 0 & 0 & 0 \\ 0 & d_2 & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & d_N \end{bmatrix}$$

Sum of weights in row 2

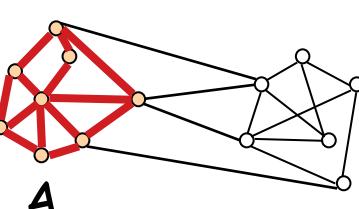
$=$

Vol(A)

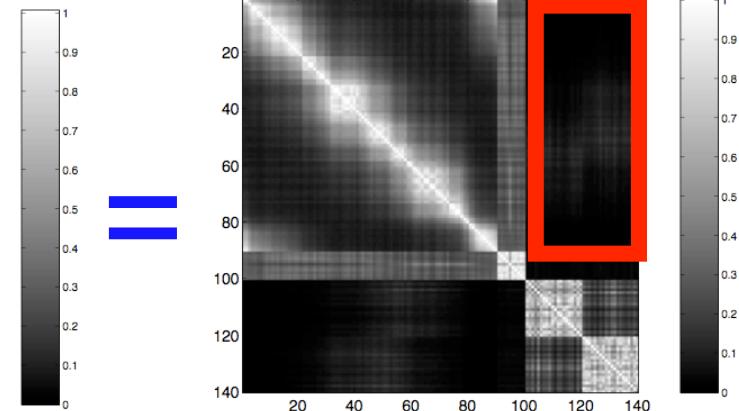
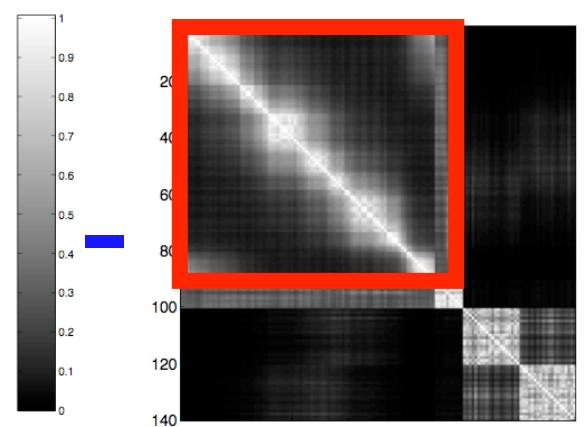
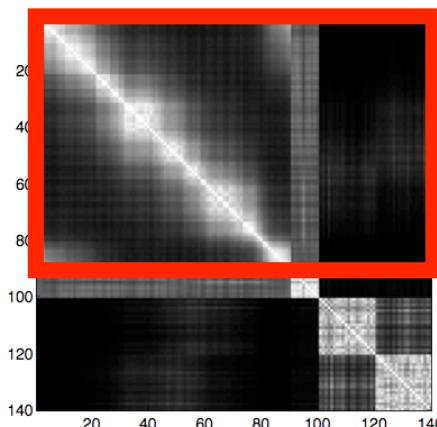
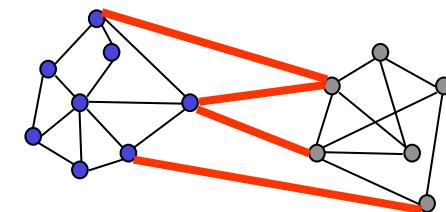
# Restating Cut Metric



-



=



$$x^T D x$$

-

$$x^T W x$$

=

$$x^T (D - W) x$$

$\uparrow$   
 $1_A$

$\text{cut}(A, \bar{A})$  ↗ key

# Graph Laplacian

- Definition:  $L = D - W$

- Facts:

- Symmetric, positive semi-definite
  - Eigenvalues

$$0 = \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_N$$

eigvec  $u_1 = \mathbf{1}$

- Invariance to self-edges

$$\begin{aligned} L_{ii} &= d_i - w_{ii} \\ L_{ij} &= -w_{ij} \end{aligned} \quad \left. \begin{array}{l} \text{don't depend on } w_{ii} \\ \text{don't depend on } w_{ij} \end{array} \right\}$$

- Inner product in  $L$  space

$$\forall f \in \mathbb{R}^N \quad f' L f = \frac{1}{2} \sum_{i,j} w_{ij} (f_i - f_j)^2$$

useful later

# Relationship to Identifying Connected Components

- Proposition:

- The multiplicity  $k$  of eigenvalue 0 of  $L$  is equal to the number of connected components  $A_1, \dots, A_k$

Furthermore,  $u_1, \dots, u_k = \sum_{A_1, \dots, A_k} 1_{A_i}$

- Proof: Assume graph is connected ( $k=1$ )

$$0 = u_1^T L u_1 = \sum_{i,j} w_{ij} (u_{1,i} - u_{1,j})^2$$

If  $w_{ij} > 0 \Rightarrow u_{1,i} = u_{1,j}$

Since  $\exists$  a path bt all  $i, j$ , then

$$u_1 = \text{constant} = 1$$

# Relationship to Identifying Connected Components

- Proposition:

- The multiplicity  $k$  of eigenvalue 0 of  $L$  is equal to the number of connected components

- Proof: Assume  $k$  connected components

Assume WLOG that they're ordered  $A_1, \dots, A_k$

$$L = \begin{pmatrix} L_1 & & \\ & \ddots & 0 \\ 0 & & L_k \end{pmatrix}$$

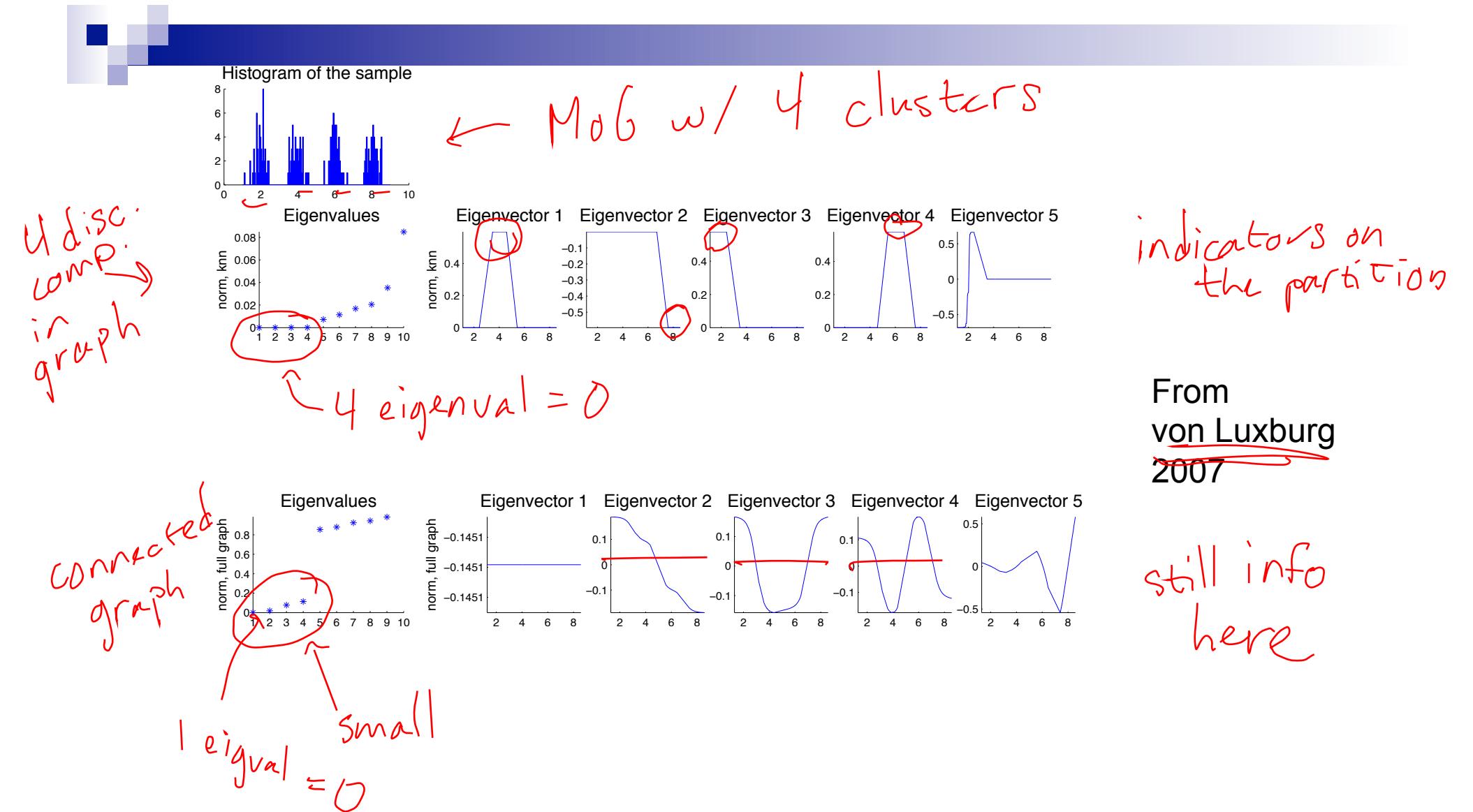
graph Laplacian  
for subgraph  $A_i$

$$\text{eigval}(L) = \bigcup_i \text{eigval}(L_i)$$

each has 1 eigenval  
equal to 0  
+ corr. eigvec  $\mathbf{1}_{A_i}$

$\Rightarrow$  eigvecs are indicators on the partition

# Example – Mixture of Gaussians



# Graph Laplacians and Ratio Cuts

- Ratio cuts for  $k=2$
- Define cluster indicator variables:

$$f_{A_i} = \begin{cases} \sqrt{|A|}/|A| & v_i \in A \\ -\sqrt{|A|}/|\bar{A}| & v_i \in \bar{A} \end{cases}$$

- Properties:

$\sum f_{A_i} = |A| \sqrt{|A|}/|A| - |\bar{A}| \sqrt{|\bar{A}|}/|\bar{A}| = 0$

 $\|f_A\|^2 = N$

- RatioCut

$$\text{RatioCut}(A, \bar{A}) = \frac{f_A' L f_A}{\|f_A\|^2} \quad \text{for } f_A \text{ as}$$

- Reformulating RatioCut problem

$$\min_{A \in V} f_A' L f_A \quad \text{s.t. } f_A \text{ defined as above, } f_A \in \mathbb{R}, \|f_A\| = \sqrt{N}$$

$f_{A_i}$  are in a discrete set

# Relaxation to Formulation

- Let  $f$  be arbitrary continuous vector

$$\min_{f \in \mathbb{R}^N} f' L f \quad \text{s.t. } f \parallel 1 \quad \|f\| = \sqrt{N}$$

graph  
laplacian      1st eigvec  
of  $L$       const

- Rayleigh-Ritz Theorem

- Which vector maximizes objective subject to constraint that the vector is orthogonal to the first eigenvector and has bounded norm?

$$f = u_2(L) = \text{eigvec assoc w/ 2nd smallest eigenval}$$

# Mapping Back to Partition

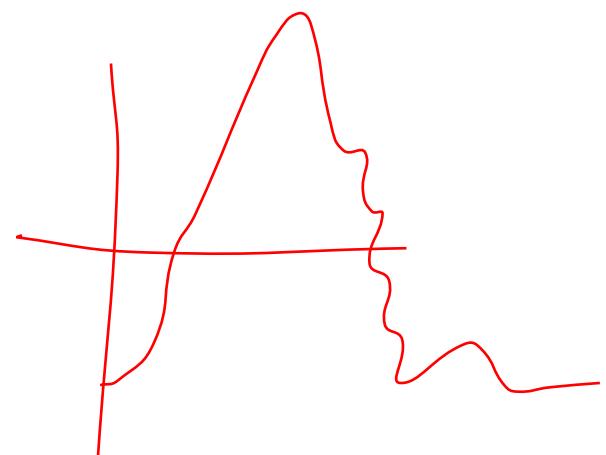
- To obtain partition, transform continuous  $f$  to a discrete indicator

- Cluster coordinates

$f_i \in \mathbb{R}$  into  $C, \bar{C}$   
using k-means

- Return

$$\begin{cases} v_i \in A & \text{if } f_i \in C \\ v_i \in \bar{A} & \text{if } f_i \in \bar{C} \end{cases}$$



# Ratio Cuts for General k

- Define cluster indicator variables:

$$F_{ij} = \begin{cases} 1/\sqrt{|A_j|} & \text{if } i \in A_j \\ 0 & \text{otherwise (ow)} \end{cases} \quad F_A \in \mathbb{R}^{N \times k} \quad F'_A F_A = I$$

- RatioCut

$$\text{RatioCut}(A_1, \dots, A_k) = \sum_{i=1}^k f'_{Ai} L f_{Ai} = \text{Tr}(F'_A L F_A)$$

- Reformulating RatioCut problem

$$\min_{A_1, \dots, A_k} \text{Tr}(F'_A L F_A) \quad \text{and} \quad F_A \quad \text{w/} \quad F'_A F_A = I$$

- Relaxation

$$\min_{F \in R^{N \times k}} \text{Tr}(F' L F) \quad \text{s.t.} \quad F' F = I$$

# Ratio Cuts for General k

- Relaxation:

$$\min_{F \in R^{N \times k}} \text{Tr}(F' L F) \quad \text{s.t. } F' F = I$$

- Solution:

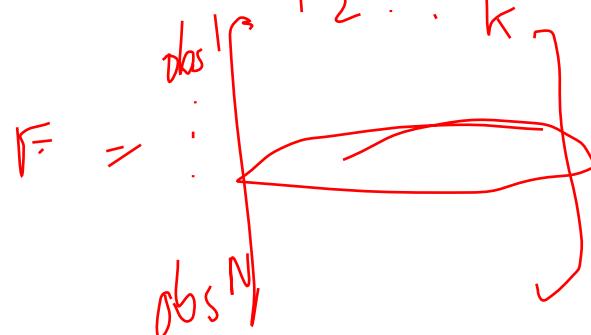
standard trace min problem

→ choose  $F$  containing first  $k$  eigvec(L)

$$\begin{bmatrix} | & | \\ u_1, \dots, u_k & | \\ | & | \end{bmatrix}$$

- To obtain partition:

cluster rows of  $F$  using k-means



if obs i is in cluster w/ obs j  
then the rows are the same

# Graph Laplacians and Norm. Cuts

- Normalized cuts for  $k=2$

$A, \bar{A}$

- Define cluster indicator variables:

$$f_{A,i} = \begin{cases} \sqrt{\text{vol}(\bar{A})} / \text{vol}(A) & v_i \in A \\ -\sqrt{\text{vol}(A)} / \text{vol}(\bar{A}) & v_i \in \bar{A} \end{cases}$$

- Properties:

$$(Df_A)^T \mathbb{1} = 0 \quad \text{and} \quad f_A^T Df_A = \text{vol}(V)$$

- Ncut

$$\text{Ncut}(A, \bar{A}) = \frac{f_A^T L f_A}{\text{vol}(V)}$$

- Reformulating Ncut problem

$$\min_{A \subset V} f_A^T L f_A \quad \text{s.t.} \quad Df_A \perp \mathbb{1} \quad \text{and} \quad f_A^T Df_A = \text{vol}(V)$$

# Relaxation to Formulation

- Let  $f$  be arbitrary continuous vector

$$\min_{f \in \mathbb{R}^N} F' L F \text{ s.t. } Df \perp \underbrace{\mathbb{I}}_{\text{const}} \quad F' D F = \text{vol}(v)$$

$$\Downarrow f = D^{-1/2} g$$

$$\min_{g \in \mathbb{R}^r} g' D^{-1/2} L D^{-1/2} g \text{ s.t. } g \perp \underbrace{D^{1/2} \mathbb{I}}_{\text{const}} \quad \|g\|^2 = \text{vol}(v)$$

$\xrightarrow{\text{1st eigvec of } L_{\text{sym}}}$

- Rayleigh-Ritz Theorem

$$g = u_2(L_{\text{sym}})$$

$$\Rightarrow f = D^{-1/2} u_2(L_{\text{sym}}) = u_2(L_{rw})$$

Equiv to  $f$  soln of  $L_u = \boxed{DU}$        $\boxed{I - D^{-1}W}$

# Normalized Cuts for General k

- Define cluster indicator variables:

$$F_{ij} = \begin{cases} 1/\sqrt{\text{vol}(A_j)} & v_i \in A_j \\ 0 & \text{ow} \end{cases} \quad F'_{\mathcal{A}} F_{\mathcal{A}} = I$$
$$F'_{\mathcal{A}} D F_{\mathcal{A}} = I$$

- Reformulating RatioCut problem

$$\min_{A_1, \dots, A_k} \text{Tr}(F'_{\mathcal{A}} L F_{\mathcal{A}}) \text{ s.t. } F'_{\mathcal{A}} D F_{\mathcal{A}} = I$$

- Relaxation

$$\min_{H \in R^{N \times k}} \text{Tr}(H' D^{-1/2} L D^{-1/2} H) \text{ s.t. } H'H = I$$

- Solution:

- $H$  is matrix of first  $k$  eigenvectors of  $\cancel{L_{sym}}$ , which is equivalent to the approximate  $F$  being the first  $k$  eigenvectors of  $\cancel{L_{rw}}$

# Random Walks on Graphs

- Stochastic process with random jumps from  $v_i$  to  $v_j$  wp:
- Transition matrix:
- Connection to graph Laplacian:
- Intuitively, want to partition graph s.t. random walk stays in cluster for a while and rarely jumps between clusters

# Random Walks on Graphs

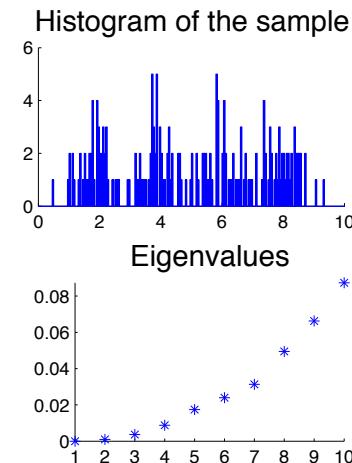
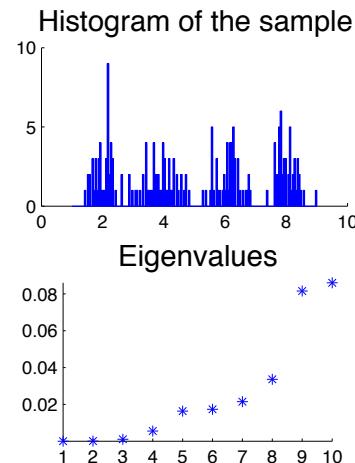
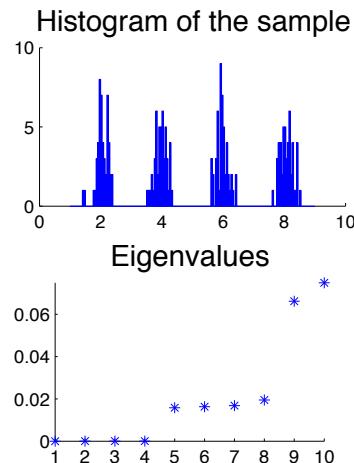
- Assume that stationary distribution exists and is unique. Then,
- Proposition:  $\text{Ncut}(A, \bar{A}) = P(A | \bar{A}) + P(\bar{A} | A)$
- Proof:
- Minimizing normalized cuts is equivalent to minimizing the probability of transitioning between clusters

# Notes

- No guarantee to quality of approximation
- Sensitive to choice of similarity graph (see earlier)
- Which graph Laplacian to use?
  - If degrees in graph vary significantly, then Laplacians are quite different
  - ~~In general,  $L_{rw}$  behaves the best~~
  - ~~Volume gives better measure of within-cluster similarity than cardinality~~
  - Normalized cuts has consistency results, Ratio cuts does not

# Notes

- Choosing the number of clusters  $k$  can be hard
  - Easy when clusters are well-separated



From  
von Luxburg  
2007

- k-means to return partition from solution to relaxation is *an* approach, but not the only