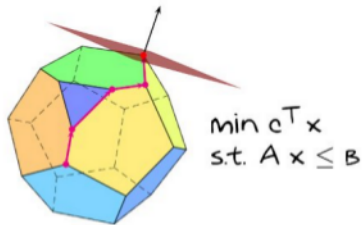
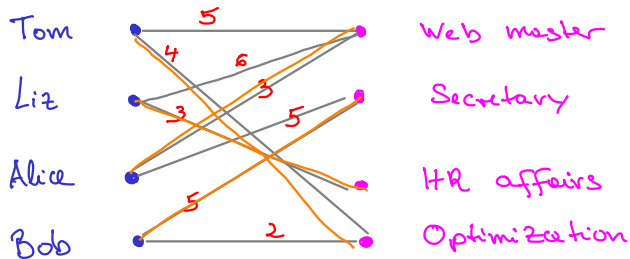


Matchings and vertex covers

- ▶ Example
- ▶ A min-max relation



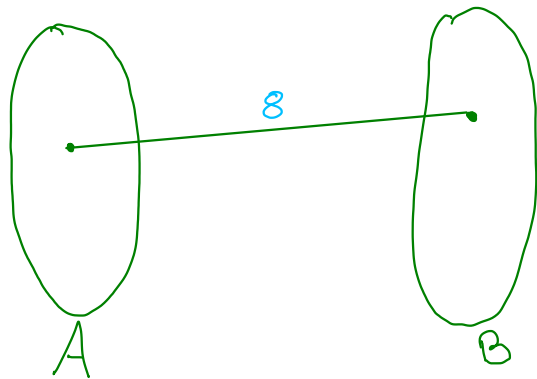
Assigning jobs to employees



Score = 15

Bipartite graphs

A graph $G = (V, E)$ is **bipartite**, if one can partition V into $V = A \cup B$ such that each edge $e \in E$ satisfies $|e \cap A| = |e \cap B| = 1$.



Weights:

$w: E \rightarrow \mathbb{N}_0$

Notation:

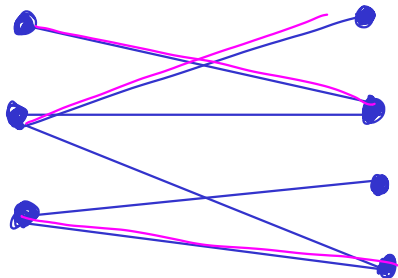
$u \subseteq E:$

$$w(u) = \sum_{e \in u} w_e$$

Matchings

A *matching* is a subset $M \subseteq E$ of the edges such that each $e_1 \neq e_2 \in M$ satisfy $e_1 \cap e_2 = \emptyset$.

The edges in a matching “do not touch”.

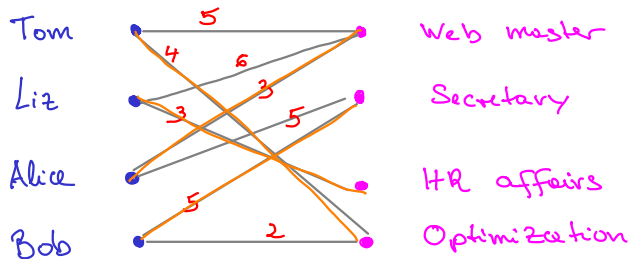


$$|M| = 3$$

The maximum weight (bipartite) matching problem

Given a (bipartite) graph $G = (V, E)$ and **edge weights** $w : E \rightarrow \mathbb{N}_0$, determine a matching $M \subseteq E$ such that

$$w(M) := \sum_{e \in M} w_e \quad \text{is maximal.}$$



$$w(M) = 15$$

Question:

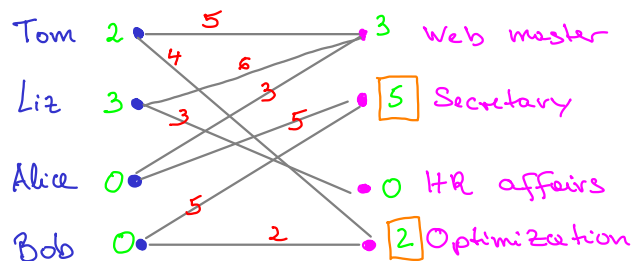
Is M maximal

w-vertex covers

Let $G = (V, E)$ be a graph with edge weights $w : E \rightarrow \mathbb{N}_0$. A **w-vertex cover** is a vector $y \in \mathbb{N}_0^{|V|}$ such that

$$\forall uv \in E : y_u + y_v \geq w_{uv}.$$

The **value** of a w-vertex cover y is $\sum_{v \in V} y_v$.



$$3 + 4 = 7 \geq 5$$

$$\text{VALUE} = 15$$

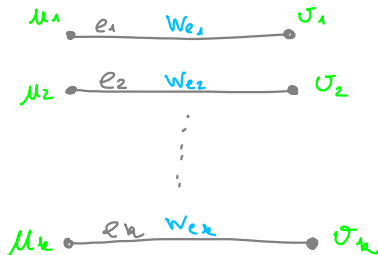
w -vertex cover \geq weight of matching

Lemma (Weak duality)

Let $G = (V, E)$ be a graph and let $w : E \rightarrow \mathbb{N}_0$ be edge-weights. If M is a matching of G and if y is a w -vertex cover of G , then

$$w(M) \leq \sum_{v \in V} y_v.$$

Proof:



Weight of M

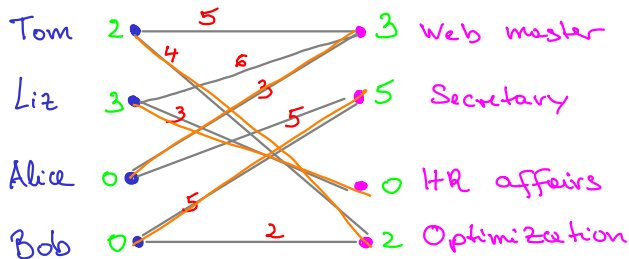
$$M = \{e_1, e_2, \dots, e_k\}$$

$$\sum_{i=1}^k \underbrace{w_{e_i}}_{\leq y_{u_i} + y_{v_i}}$$

$$\leq \sum_{v \in V} y_v$$

VAL. OF
 w -VERTEX
COVER

The job assignment is optimal



$$W(\pi) = 15$$

$$\sum y_v = 15$$

Both!

y and π are optimal