

Broadcast authentication

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C. Gunter, S. Khanna, K. Tan, S. S. Venkatesh (2004)

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- ❖ **Randomised selective authentication**: The case for the defence — exploit the mismatch in the needs of the attacker and the client.
 - ❖ **The attacker's need: have the receiver examine the vast majority of the spurious packets.** Defang the attack by having the receiver randomly rejects a fraction $1 - p$ of incoming packets. [Reject rate determined by the spare computational capacity at the receiver and the maximum attack rate. For example, if $p = 0.1$ then only ten percent of the attack gets through.]
 - ❖ **The sender's need: have one signed packet validated by the receiver.** Sender sends n copies of her cryptographically signed packet. The probability that at least one signed packet is verified is approximately $1 - \text{Po}(0; \lambda) = 1 - e^{-\lambda}$ where $\lambda = np$ is the Poisson parameter. [If $p = 0.1$ and $n = 25$ there is a 92% chance that a signed packet makes it through the blockade; if $n = 40$ the chance jumps to 98%.]