

Lemma

$$\min\{7\mathbf{x_1} + \mathbf{x_2} + 5\mathbf{x_3} : \mathbf{x} \; \mathbf{feasible}\} \geq \\ \max\{10\mathbf{y_1} + 6\mathbf{y_2} : \mathbf{y} \; \mathbf{feasible}\}$$

Linear programming duality Theorem

$$\min\{7\mathbf{x_1} + \mathbf{x_2} + 5\mathbf{x_3} : \mathbf{x} \text{ feasible}\} = \max\{10\mathbf{y_1} + 6\mathbf{y_2} : \mathbf{y} \text{ feasible}\}$$

In general

 $\begin{array}{ll} \text{(P)} & \text{(D)} \\ \min \mathbf{c} \cdot \mathbf{x} : & \max \mathbf{b} \cdot \mathbf{y} : \\ \mathbf{A}\mathbf{x} \geq \mathbf{b} & & \mathbf{A}^T \mathbf{y} \leq \mathbf{c} \\ \mathbf{x} > \mathbf{0} & & \mathbf{y} \geq \mathbf{0} \end{array}$

Strong duality Theorem in general

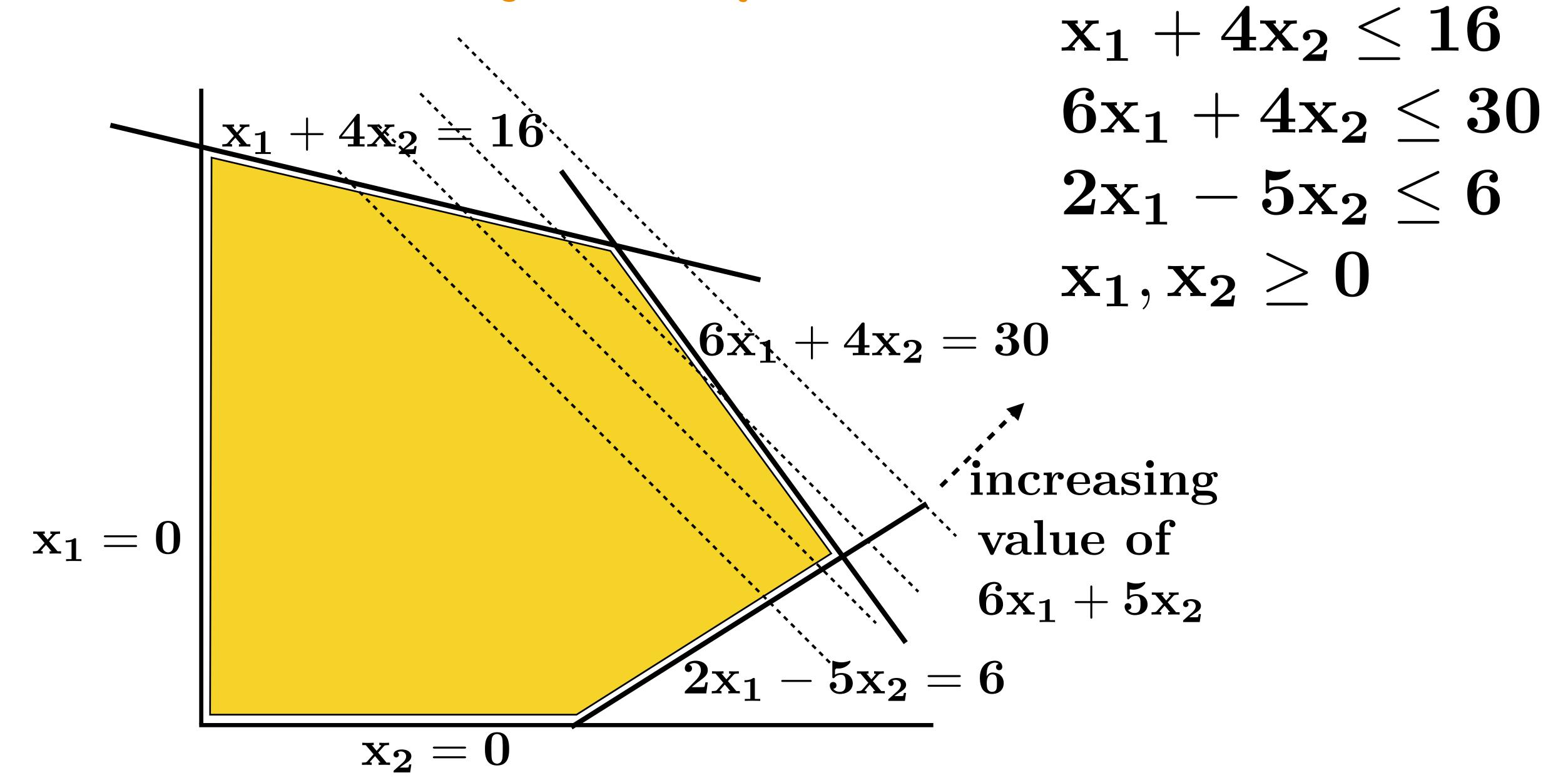
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(P)
min c \cdot x :
Ax \ge b
x \ge 0
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Four possible cases:
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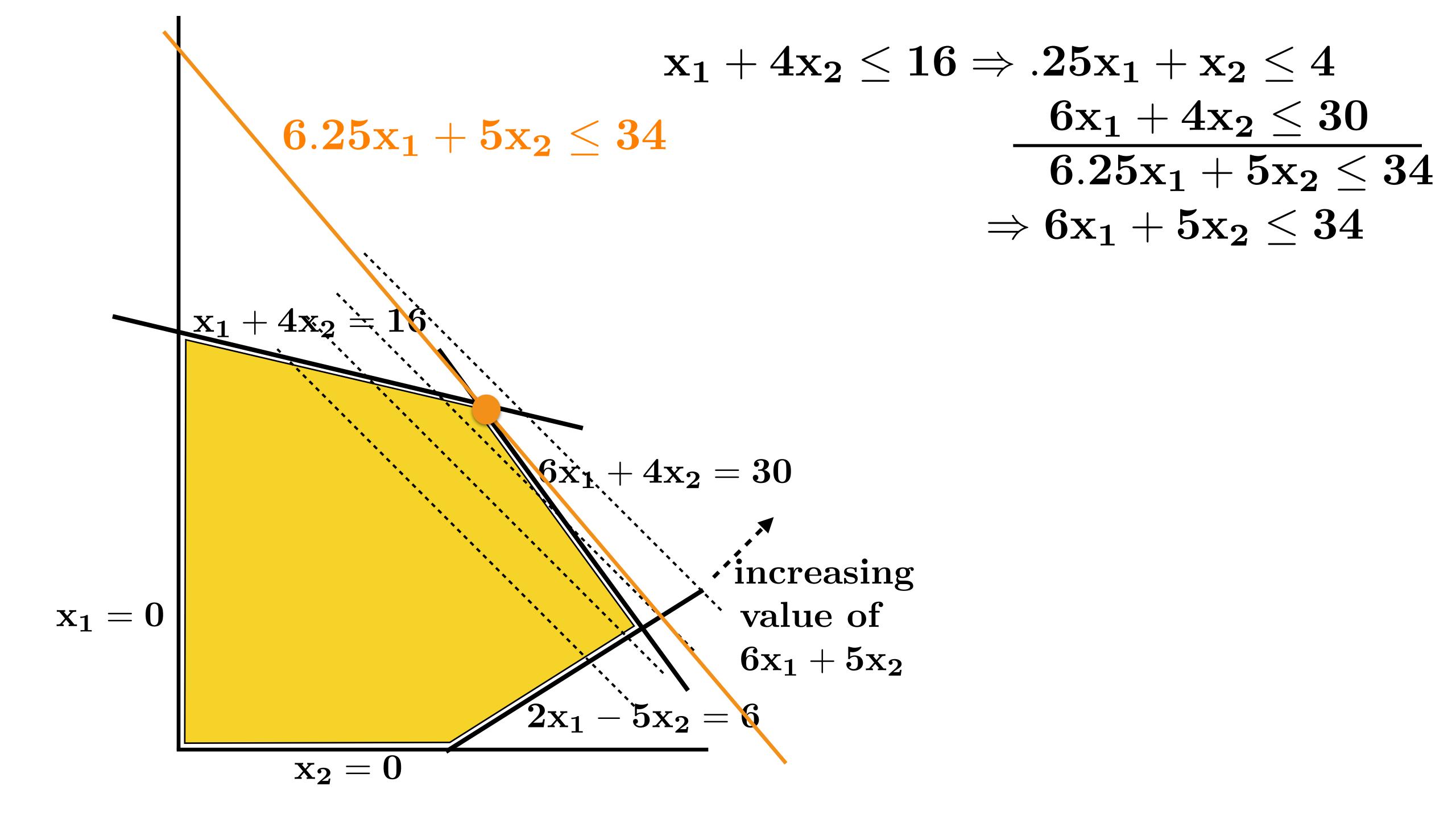
- (P) is empty, (D) has value $+\infty$
- (D) is empty, (P) has value $-\infty$
- value(P)=value(D)
- ((P) and (D) both empty)

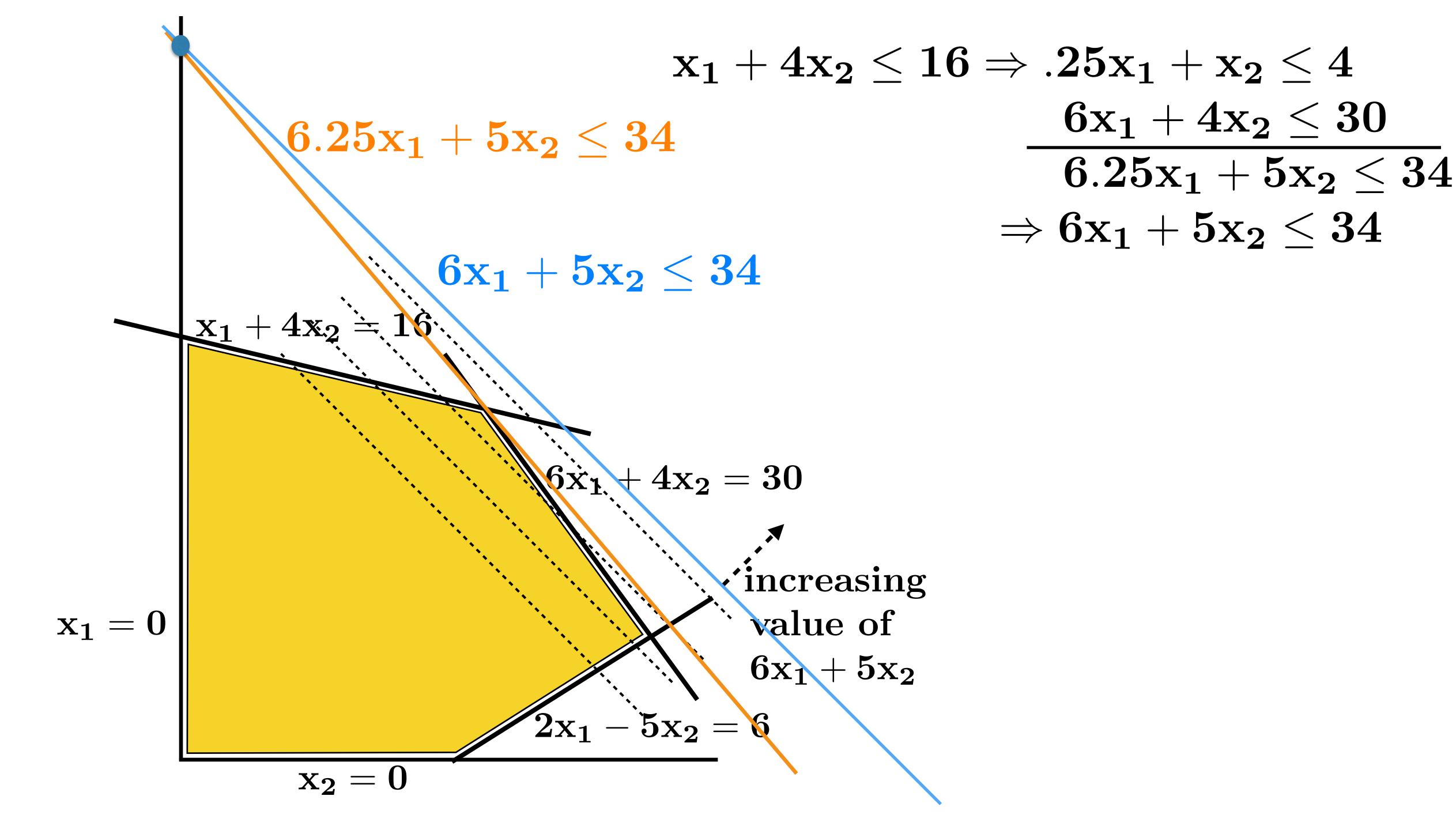
 $maxb \cdot y:$ $A^{T}y \leq c$ $y \geq 0$

What about the geometry?



 $\max 6x_1 + 5x_2$:





Linear programming duality Theorem

$$(P) = (D)$$

$$\max \mathbf{c} \cdot \mathbf{x} : \min \mathbf{b} \cdot \mathbf{y} :$$

$$\mathbf{A}\mathbf{x} \le \mathbf{b}$$

$$\mathbf{x} \ge \mathbf{0}$$

$$\mathbf{y} \ge \mathbf{0}$$

OPT constraints of (P)
and a convex combination
that imply exactly
the right upper bound

