

Data Mining Association Rules: Advanced Concepts and Algorithms

Lecture Notes for Chapter 7

Introduction to Data Mining
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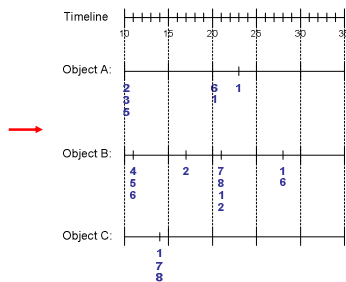
Frequent pattern mining

- It all started with frequent itemsets in supermarket transactions.
- Suppose people have a *bonuscard*, now we can track transactions of a customer through time.
- The pattern becomes more complex: from *itemset* to *sequence of itemsets*.
- More complex patterns: trees (e.g. XML data) and graphs (browsing patterns, chemical structures, etc).

Sequence Data

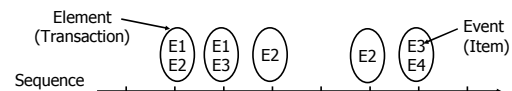
Sequence Database:

Object	Timestamp	Events
A	10	2, 3, 5
A	20	6, 1
A	23	1
B	11	4, 5, 6
B	17	2
B	21	7, 8, 1, 2
B	28	1, 6
C	14	1, 8, 7



Examples of Sequence Data

Sequence Database	Sequence	Element (Transaction)	Event (Item)
Customer	Purchase history of a given customer	A set of items bought by a customer at time t	Books, diary products, CDs, etc
Web Data	Browsing activity of a particular Web visitor	A collection of files viewed by a Web visitor after a single mouse click	Home page, index page, contact info, etc
Event data	History of events generated by a given sensor	Events triggered by a sensor at time t	Types of alarms generated by sensors
Genome sequences	DNA sequence of a particular species	An element of the DNA sequence	Bases A, T, G, C



Formal Definition of a Sequence

- A sequence is an ordered list of elements (transactions)

$$S = \langle e_1, e_2, e_3, \dots \rangle$$
 - Each element contains a collection of events (items)

$$e_i = \{i_1, i_2, \dots, i_k\}$$
 - Each element is attributed to a specific time or location
- Length of a sequence, $|S|$, is given by the number of elements of the sequence
- A k-sequence is a sequence that contains k (not necessarily distinct) events (items)

Examples of Sequence

- Web sequence:

$$\langle \{\text{Homepage}\} \{\text{Electronics}\} \{\text{Digital Cameras}\} \{\text{Canon Digital Camera}\} \{\text{Shopping Cart}\} \{\text{Order Confirmation}\} \{\text{Return to Shopping}\} \rangle$$
- Sequence of initiating events causing the nuclear accident at 3-mile Island:

$$\langle \{\text{clogged resin}\} \{\text{outlet valve closure}\} \{\text{loss of feedwater}\} \{\text{condenser polisher outlet valve shut}\} \{\text{booster pumps trip}\} \{\text{main waterpump trips}\} \{\text{main turbine trips}\} \{\text{reactor pressure increases}\} \rangle$$
- Sequence of books checked out at a library:

$$\langle \{\text{Fellowship of the Ring}\} \{\text{The Two Towers}\} \{\text{Return of the King}\} \rangle$$

Formal Definition of a Subsequence

- A sequence $\langle a_1 a_2 \dots a_n \rangle$ is contained in another sequence $\langle b_1 b_2 \dots b_m \rangle$ ($m \geq n$) if there exist integers $i_1 < i_2 < \dots < i_n$ such that $a_1 \subseteq b_{i_1}, a_2 \subseteq b_{i_2}, \dots, a_n \subseteq b_{i_n}$

Data sequence	Subsequence	Contain?
$\langle \{2,4\} \{3,5,6\} \{8\} \rangle$	$\langle \{2\} \{3,5\} \rangle$	Yes
$\langle \{1,2\} \{3,4\} \rangle$	$\langle \{1\} \{2\} \rangle$	No
$\langle \{2,4\} \{2,4\} \{2,5\} \rangle$	$\langle \{2\} \{4\} \rangle$	Yes

- The support of a subsequence w is defined as the fraction of data sequences that contain w
- A *sequential pattern* is a frequent subsequence (i.e., a subsequence whose support is $\geq \text{minsup}$)

Sequential Pattern Mining: Definition

- Given:
 - a database of sequences
 - a user-specified minimum support threshold, *minsup*
- Task:
 - Find all subsequences with support $\geq \text{minsup}$

Sequential Pattern Mining: Challenge

- Given a sequence: $\langle \{a\} \{b\} \{c\} \{d\} \{e\} \{f\} \{g\} \{h\} \{i\} \rangle$
 - Examples of subsequences: $\langle \{a\} \{c\} \{f\} \{g\} \rangle$, $\langle \{c\} \{d\} \{e\} \rangle$, $\langle \{b\} \{g\} \rangle$, etc.
- How many k -subsequences can be extracted from a given n -sequence?

$\langle \{a\} \{b\} \{c\} \{d\} \{e\} \{f\} \{g\} \{h\} \{i\} \rangle \quad n = 9$
 $k=4:$

↓	↓	↓	↓	↓	↓	↓	↓	↓
Y	-	-	Y	Y	-	-	-	Y
⏟			⏟			⏟		
$\langle \{a\} \{b\} \{c\} \{d\} \rangle$			$\langle \{c\} \{d\} \{e\} \{f\} \rangle$			$\langle \{f\} \{g\} \{h\} \{i\} \rangle$		

Answer: $\binom{n}{k} = \binom{9}{4} = 126$

Apriori property for sequences

Let D be a database that contains a collection of data sequences d . The support of a sequence t is the fraction of all data sequences that contain t :

$$s(t) = \frac{|\{d \in D : t \text{ is a subsequence of } d\}|}{|D|}$$

Apriori property:

If a data sequence d contains a sequence t , then it will also contain any subsequence of t .

Therefore:

If w is a subsequence of t , then $s(w) \geq s(t)$.

Sequential Pattern Mining: Example

Object	Timestamp	Events
A	1	1,2,4
A	2	2,3
A	3	5
B	1	1,2
B	2	2,3,4
C	1	1,2
C	2	2,3,4
C	3	2,4,5
D	1	2
D	2	3,4
D	3	4,5
E	1	1,3
E	2	2,4,5

Minsup = 50%

Examples of Frequent Subsequences:

$\langle \{1,2\} \rangle$	$s=60\%$
$\langle \{2,3\} \rangle$	$s=60\%$
$\langle \{2,4\} \rangle$	$s=80\%$
$\langle \{3\} \{5\} \rangle$	$s=80\%$
$\langle \{1\} \{2\} \rangle$	$s=80\%$
$\langle \{2\} \{2\} \rangle$	$s=60\%$
$\langle \{1\} \{2,3\} \rangle$	$s=60\%$
$\langle \{2\} \{2,3\} \rangle$	$s=60\%$
$\langle \{1,2\} \{2,3\} \rangle$	$s=60\%$

Generalized Sequential Pattern (GSP)

- Step 1:**
 - Make the first pass over the sequence database D to yield all the 1-element frequent sequences
- Step 2:**
 - Repeat until no new frequent sequences are found
 - Candidate Generation:**
 - Merge pairs of frequent subsequences found in the $(k-1)$ th pass to generate candidate sequences that contain k items
 - Candidate Pruning:**
 - Prune candidate k -sequences that contain infrequent $(k-1)$ -subsequences
 - Support Counting:**
 - Make a new pass over the sequence database D to find the support for these candidate sequences
 - Candidate Elimination:**
 - Eliminate candidate k -sequences whose actual support is less than *minsup*

Candidate Generation

- Base case ($k=2$):

- Merging two frequent 1-sequences $\langle \{i_1\} \rangle$ and $\langle \{i_2\} \rangle$ will produce three candidate 2-sequences: $\langle \{i_1\} \{i_2\} \rangle$, $\langle \{i_2\} \{i_1\} \rangle$ and $\langle \{i_1 i_2\} \rangle$
- Each frequent 1-sequence $\langle \{i_m\} \rangle$ produces the candidate 2-sequence $\langle \{i_m\} \{i_m\} \rangle$.

For example, if $\langle \{A\} \rangle$ and $\langle \{B\} \rangle$ are frequent, this produces the candidate 2-sequences: $\langle \{A\} \{B\} \rangle$, $\langle \{B\} \{A\} \rangle$, $\langle \{A, B\} \rangle$, $\langle \{A\} \{A\} \rangle$ and $\langle \{B\} \{B\} \rangle$.

Candidate Generation

- General case ($k>2$):

- A frequent $(k-1)$ -sequence w_1 is merged with frequent $(k-1)$ -sequence w_2 to produce a candidate k -sequence if the subsequence obtained by removing the first (using the lexicographic order of events within each element) event in w_1 is the same as the subsequence obtained by removing the last event in w_2 .
 - The resulting candidate after merging is given by the sequence w_1 extended with the last event of w_2 .
 - If the last two events in w_2 belong to the same element, then the last event in w_2 becomes part of the last element in w_1
 - Otherwise, the last event in w_2 becomes a separate element appended to the end of w_1

Candidate Generation Examples

- Merging the sequences $w_1 = \langle \{1\} \{2\} \{3\} \{4\} \rangle$ and $w_2 = \langle \{2\} \{3\} \{4\} \{5\} \rangle$ will produce the candidate sequence $\langle \{1\} \{2\} \{3\} \{4\} \{5\} \rangle$ because the last two events in w_2 (4 and 5) belong to the same element
- Merging the sequences $w_1 = \langle \{1\} \{2\} \{3\} \{4\} \rangle$ and $w_2 = \langle \{2\} \{3\} \{4\} \{5\} \rangle$ will produce the candidate sequence $\langle \{1\} \{2\} \{3\} \{4\} \{5\} \rangle$ because the last two events in w_2 (4 and 5) do not belong to the same element
- We do not have to merge the sequences $w_1 = \langle \{1\} \{2\} \{3\} \{4\} \rangle$ and $w_2 = \langle \{1\} \{2\} \{4\} \{5\} \rangle$ to produce the candidate $\langle \{1\} \{2\} \{3\} \{4\} \{5\} \rangle$ because if the latter is a viable candidate, then it can be obtained by merging w_1 with $\langle \{2\} \{4\} \{5\} \rangle$

GSP Example

Frequent 3-sequences

$\langle \{1\} \{2\} \{3\} \rangle$
 $\langle \{1\} \{2\} \{5\} \rangle$
 $\langle \{1\} \{5\} \{3\} \rangle$
 $\langle \{2\} \{3\} \{4\} \rangle$
 $\langle \{2\} \{5\} \{3\} \rangle$
 $\langle \{3\} \{4\} \{5\} \rangle$
 $\langle \{5\} \{3\} \{4\} \rangle$

Candidate Generation

$\langle \{1\} \{2\} \{3\} \{4\} \rangle$
 $\langle \{1\} \{2\} \{5\} \{3\} \rangle$
 $\langle \{1\} \{5\} \{3\} \{4\} \rangle$
 $\langle \{2\} \{3\} \{4\} \{5\} \rangle$
 $\langle \{2\} \{5\} \{3\} \{4\} \rangle$

Candidate Pruning

$\langle \{1\} \{2\} \{5\} \{3\} \rangle$

For example: $\langle \{2\} \{5\} \{3\} \rangle + \langle \{5\} \{3\} \{4\} \rangle : \langle \{2\} \{5\} \{3\} \{4\} \rangle$
 $\langle \{1\} \{5\} \{3\} \rangle + \langle \{5\} \{3\} \{4\} \rangle : \langle \{1\} \{5\} \{3\} \{4\} \rangle$
 $\langle \{2\} \{5\} \{3\} \{4\} \rangle$ is pruned because 3-subsequence $\langle \{2\} \{3\} \{4\} \rangle$ is not frequent.

GSP Exercise

Number	Frequent 3-sequences
1	$\langle \{1, 2, 3\} \rangle$
2	$\langle \{1, 2\} \{3\} \rangle$
3	$\langle \{1\} \{2, 3\} \rangle$
4	$\langle \{1, 2\} \{4\} \rangle$
5	$\langle \{1, 3\} \{4\} \rangle$
6	$\langle \{1, 2, 4\} \rangle$
7	$\langle \{2, 3\} \{3\} \rangle$
8	$\langle \{2, 3\} \{4\} \rangle$
9	$\langle \{2\} \{3\} \{3\} \rangle$
10	$\langle \{2\} \{3\} \{4\} \rangle$

Generate the candidate 4-sequences.
Which ones are pruned?

GSP Exercise

Number	Frequent 3-sequences
1	$\langle \{1, 2, 3\} \rangle$
2	$\langle \{1, 2\} \{3\} \rangle$
3	$\langle \{1\} \{2, 3\} \rangle$
4	$\langle \{1, 2\} \{4\} \rangle$
5	$\langle \{1, 3\} \{4\} \rangle$
6	$\langle \{1, 2, 4\} \rangle$
7	$\langle \{2, 3\} \{3\} \rangle$
8	$\langle \{2, 3\} \{4\} \rangle$
9	$\langle \{2\} \{3\} \{3\} \rangle$
10	$\langle \{2\} \{3\} \{4\} \rangle$

1 + 7: $\langle \{1, 2, 3\} \{3\} \rangle$

GSP Exercise

Number	Frequent 3-sequences
1	<{1,2,3}>
2	<{1,2} {3}>
3	<{1} {2,3}>
4	<{1,2} {4}>
5	<{1,3} {4}>
6	<{1,2,4}>
7	<{2,3} {3}>
8	<{2,3} {4}>
9	<{2} {3} {3}>
10	<{2} {3} {4}>

1 + 7: <{1,2,3} {3}>
1 + 8: <{1,2,3} {4}>

GSP Exercise

Number	Frequent 3-sequences
1	<{1,2,3}>
2	<{1,2} {3}>
3	<{1} {2,3}>
4	<{1,2} {4}>
5	<{1,3} {4}>
6	<{1,2,4}>
7	<{2,3} {3}>
8	<{2,3} {4}>
9	<{2} {3} {3}>
10	<{2} {3} {4}>

1 + 7: <{1,2,3} {3}>
1 + 8: <{1,2,3} {4}>
2 + 9: <{1,2} {3} {3}>

GSP Exercise

Number	Frequent 3-sequences
1	<{1,2,3}>
2	<{1,2} {3}>
3	<{1} {2,3}>
4	<{1,2} {4}>
5	<{1,3} {4}>
6	<{1,2,4}>
7	<{2,3} {3}>
8	<{2,3} {4}>
9	<{2} {3} {3}>
10	<{2} {3} {4}>

1 + 7: <{1,2,3} {3}>
1 + 8: <{1,2,3} {4}>
2 + 9: <{1,2} {3} {3}>
2 + 10: <{1,2} {3} {4}>

GSP Exercise

Number	Frequent 3-sequences
1	<{1,2,3}>
2	<{1,2} {3}>
3	<{1} {2,3}>
4	<{1,2} {4}>
5	<{1,3} {4}>
6	<{1,2,4}>
7	<{2,3} {3}>
8	<{2,3} {4}>
9	<{2} {3} {3}>
10	<{2} {3} {4}>

1 + 7: <{1,2,3} {3}>
1 + 8: <{1,2,3} {4}>
2 + 9: <{1,2} {3} {3}>
2 + 10: <{1,2} {3} {4}>
3 + 7: <{1} {2,3} {3}>

GSP Exercise

Number	Frequent 3-sequences
1	<{1,2,3}>
2	<{1,2} {3}>
3	<{1} {2,3}>
4	<{1,2} {4}>
5	<{1,3} {4}>
6	<{1,2,4}>
7	<{2,3} {3}>
8	<{2,3} {4}>
9	<{2} {3} {3}>
10	<{2} {3} {4}>

1 + 7: <{1,2,3} {3}>
1 + 8: <{1,2,3} {4}>
2 + 9: <{1,2} {3} {3}>
2 + 10: <{1,2} {3} {4}>
3 + 7: <{1} {2,3} {3}>
3 + 8: <{1} {2,3} {4}>

GSP Exercise

Number	Frequent 3-sequences
1	<{1,2,3}>
2	<{1,2} {3}>
3	<{1} {2,3}>
4	<{1,2} {4}>
5	<{1,3} {4}>
6	<{1,2,4}>
7	<{2,3} {3}>
8	<{2,3} {4}>
9	<{2} {3} {3}>
10	<{2} {3} {4}>

1 + 7: <{1,2,3} {3}>
1 + 8: <{1,2,3} {4}>
2 + 9: <{1,2} {3} {3}>
2 + 10: <{1,2} {3} {4}>
3 + 7: <{1} {2,3} {3}>
3 + 8: <{1} {2,3} {4}>

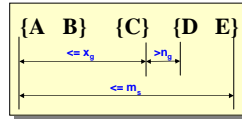
GSP Exercise: candidate pruning

Number	Frequent 3-sequences
1	$\langle \{1,2,3\} \rangle$
2	$\langle \{1,2\} \{3\} \rangle$
3	$\langle \{1\} \{2,3\} \rangle$
4	$\langle \{1,2\} \{4\} \rangle$
5	$\langle \{1,3\} \{4\} \rangle$
6	$\langle \{1,2,4\} \rangle$
7	$\langle \{2,3\} \{4\} \rangle$
8	$\langle \{2\} \{3\} \{4\} \rangle$
9	$\langle \{2\} \{3\} \{3\} \rangle$
10	$\langle \{2\} \{3\} \{4\} \rangle$

$1 + 7: \langle \{1,2,3\} \{3\} \rangle$
 $1 + 8: \langle \{1,2,3\} \{4\} \rangle$
 $2 + 9: \langle \{1,2\} \{3\} \{3\} \rangle$
 $2 + 10: \langle \{1,2\} \{3\} \{4\} \rangle$
 $3 + 7: \langle \{1\} \{2,3\} \{3\} \rangle$
 $3 + 8: \langle \{1\} \{2,3\} \{4\} \rangle$

$\langle \{1,2,3\} \{4\} \rangle$ is not pruned
 because $\langle \{1,2,3\} \rangle$,
 $\langle \{1,2\} \{4\} \rangle$ and $\langle \{1,3\} \{4\} \rangle$ are all frequent.

Timing Constraints (I)



x_g : max-gap
 n_g : min-gap
 m_s : maximum span

$x_g = 2, n_g = 0, m_s = 4$

Data sequence	Subsequence	Contain?
$\langle \{2,4\} \{3,5,6\} \{4,7\} \{4,5\} \{8\} \rangle$	$\langle \{6\} \{5\} \rangle$	Yes
$\langle \{1\} \{2\} \{3\} \{4\} \{5\} \rangle$	$\langle \{1\} \{4\} \rangle$	No
$\langle \{1\} \{2,3\} \{3,4\} \{4,5\} \rangle$	$\langle \{2\} \{3\} \{5\} \rangle$	Yes
$\langle \{1,2\} \{3\} \{2,3\} \{3,4\} \{2,4\} \{4,5\} \rangle$	$\langle \{1,2\} \{5\} \rangle$	No

Timing Constraints: Formal Definition

A data sequence $d = \langle d_1 \dots d_m \rangle$ contains a sequence $w = \langle w_1 \dots w_n \rangle$ if there exist integers $i_1 < \dots < i_n$ such that $w_j \subseteq d_{i_j}$ and

1. maxgap: $\text{time}(d_{i_j}) - \text{time}(d_{i_{j-1}}) \leq \text{maxgap}$
2. mingap: $\text{time}(d_{i_j}) - \text{time}(d_{i_{j-1}}) > \text{mingap}$
3. maxspan: $\text{time}(d_{i_n}) - \text{time}(d_{i_1}) \leq \text{maxspan}$

Note: this definition applies if the window size = 0, otherwise things get more complicated (see timing constraints II for definition of window size).

Mining Sequential Patterns with Timing Constraints

- Approach 1:
 - Mine sequential patterns without timing constraints
 - Postprocess the discovered patterns
- Approach 2:
 - Modify GSP to directly prune candidates that violate timing constraints
 - Question:
 - ♦ Does Apriori principle still hold?

Apriori Principle for Sequence Data

Object	Timestamp	Events
A	1	1,2,4
A	2	2,3
A	3	5
B	1	1,2
B	2	2,3,4
C	1	1,2
C	2	2,3,4
C	3	2,4,5
D	1	2
D	2	3,4
D	3	4,5
E	1	1,3
E	2	2,4,5

Suppose:

$x_g = 1$ (max-gap)
 $n_g = 0$ (min-gap)
 $m_s = 5$ (maximum span)
 $\text{minsup} = 60\%$

$\langle \{2\} \{5\} \rangle$ support = 40%
 but
 $\langle \{2\} \{3\} \{5\} \rangle$ support = 60%

Problem exists because of max-gap constraint
 No such problem if max-gap is infinite

Maxgap and the apriori property

- Due to the maxgap constraint, the support of a sequence can be increased by inserting an element between two elements.
- Not by inserting an element at the beginning or the end of the sequence.
- This spoils the apriori property that the support of a sequence is never bigger than the support of any of its subsequences.
- No pruning possible anymore?

Continuous Subsequences

- s is a contiguous subsequence of $w = \langle e_1 \rangle \langle e_2 \rangle \dots \langle e_k \rangle$ if any of the following conditions hold:
 - s is obtained from w by deleting an item from either e_1 or e_k
 - s is obtained from w by deleting an item from any element e_i that contains more than 2 items
 - s is a contiguous subsequence of s' and s' is a contiguous subsequence of w (recursive definition)
- Examples: $s = \langle \{1\} \{2\} \rangle$
 - is a contiguous subsequence of $\langle \{1\} \{2\} \rangle$, $\langle \{1\} \{2\} \{3\} \rangle$, and $\langle \{3\} \{4\} \{1\} \{2\} \{3\} \{4\} \rangle$
 - is not a contiguous subsequence of $\langle \{1\} \{3\} \{2\} \rangle$ and $\langle \{2\} \{1\} \{3\} \{2\} \rangle$

Apriori principle with maxgap constraint

Apriori property:

If a data sequence d contains a sequence t , then it will also contain any *contiguous* subsequence of t .

Therefore:

If w is a *contiguous* subsequence of t , then $s(w) \geq s(t)$.

Apriori principle with maxgap constraint

Example:

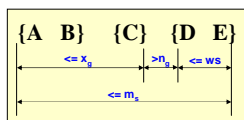
Suppose $t = \langle \{A\}, \{A,B\}, \{C\}, \{B,C\} \rangle$ is contained in a data sequence d (meaning it also satisfies the maxgap constraint), then any contiguous subsequence of t also satisfies the maxgap constraint with respect to d .

We can only violate the maxgap constraint by removing $\{A,B\}$ or $\{C\}$ but that would result in a non-contiguous subsequence.

Modified Candidate Pruning Step

- Without maxgap constraint:
 - A candidate k -sequence is pruned if at least one of its $(k-1)$ -subsequences is infrequent
- With maxgap constraint:
 - A candidate k -sequence is pruned if at least one of its *contiguous* $(k-1)$ -subsequences is infrequent
- So with a maxgap constraint we can do less pruning.

Timing Constraints (II)



x_g : max-gap
 n_g : min-gap
 ws : window size
 m_s : maximum span

$x_g = 2$, $n_g = 0$, $ws = 1$, $m_s = 5$

Data sequence	Subsequence	Contain?
$\langle \{2,4\} \{3,5,6\} \{4,7\} \{4,6\} \{8\} \rangle$	$\langle \{3\} \{5\} \rangle$	No
$\langle \{1\} \{2\} \{3\} \{4\} \{5\} \rangle$	$\langle \{1,2\} \{3\} \rangle$	Yes
$\langle \{1,2\} \{2,3\} \{3,4\} \{4,5\} \rangle$	$\langle \{1,2\} \{3,4\} \rangle$	Yes

Modified Support Counting Step

- Given a candidate pattern: $\langle \{a, c\} \rangle$
 - Any data sequences that contain
 - $\langle \dots \{a\} \{c\} \dots \rangle$,
 - $\langle \dots \{a\} \dots \{c\} \dots \rangle$ (where $\text{time}(\{c\}) - \text{time}(\{a\}) \leq ws$)
 - $\langle \dots \{c\} \dots \{a\} \dots \rangle$ (where $\text{time}(\{a\}) - \text{time}(\{c\}) \leq ws$)
- will contribute to the support count of candidate pattern

Exercise: timing constraints

Data sequence: $d = \langle \{1,2,3\} \{2,4\} \{2,4,5\} \{3,5\} \{6\} \rangle$ where elements occur on consecutive time points 1...5.

Given are the following sequences $w = \langle e_1 e_2 \dots e_i e_{i+1} \dots e_n \rangle$

1. $\langle \{1\} \{2\} \{3\} \rangle$
2. $\langle \{1,2,3,4\} \{5,6\} \rangle$
3. $\langle \{2,4\} \{2,4\} \{6\} \rangle$
4. $\langle \{1\} \{2,4\} \{6\} \rangle$
5. $\langle \{1,2\} \{3,4\} \{5,6\} \rangle$

Are they subsequences of d , subject to the following timing constraints?
mingap=0 (interval between last event in e_i and first event in $e_{i+1} > 0$)

maxgap=3 (interval between first event in e_i and last event in $e_{i+1} \leq 3$)

maxspan=5 (interval between first event in e_1 and last event in $e_n \leq 5$)

ws=1 (time between first and last events in $e_i \leq 1$)

Exercise: timing constraints

$d = \langle \{1,2,3\} \{2,4\} \{2,4,5\} \{3,5\} \{6\} \rangle$

Given are the following sequences $w = \langle e_1 e_2 \dots e_i e_{i+1} \dots e_n \rangle$

1. $\langle \{1\} \{2\} \{3\} \rangle$ Yes
2. $\langle \{1,2,3,4\} \{5,6\} \rangle$
3. $\langle \{2,4\} \{2,4\} \{6\} \rangle$
4. $\langle \{1\} \{2,4\} \{6\} \rangle$
5. $\langle \{1,2\} \{3,4\} \{5,6\} \rangle$

Are they subsequences of d , subject to the following timing constraints?
mingap=0 (interval between last event in e_i and first event in $e_{i+1} > 0$)

maxgap=3 (interval between first event in e_i and last event in $e_{i+1} \leq 3$)

maxspan=5 (interval between first event in e_1 and last event in $e_n \leq 5$)

ws=1 (time between first and last events in $e_i \leq 1$)

Exercise: timing constraints

$d = \langle \{1,2,3\} \{2,4\} \{2,4,5\} \{3,5\} \{6\} \rangle$

Given are the following sequences $w = \langle e_1 e_2 \dots e_i e_{i+1} \dots e_n \rangle$

1. $\langle \{1\} \{2\} \{3\} \rangle$ Yes
2. $\langle \{1,2,3,4\} \{5,6\} \rangle$ No (maxgap constraint violated)
3. $\langle \{2,4\} \{2,4\} \{6\} \rangle$
4. $\langle \{1\} \{2,4\} \{6\} \rangle$
5. $\langle \{1,2\} \{3,4\} \{5,6\} \rangle$

Are they subsequences of d , subject to the following timing constraints?
mingap=0 (interval between last event in e_i and first event in $e_{i+1} > 0$)

maxgap=3 (interval between first event in e_i and last event in $e_{i+1} \leq 3$)

maxspan=5 (interval between first event in e_1 and last event in $e_n \leq 5$)

ws=1 (time between first and last events in $e_i \leq 1$)

Exercise: timing constraints

$d = \langle \{1,2,3\} \{2,4\} \{2,4,5\} \{3,5\} \{6\} \rangle$

Given are the following sequences $w = \langle e_1 e_2 \dots e_i e_{i+1} \dots e_n \rangle$

1. $\langle \{1\} \{2\} \{3\} \rangle$ Yes
2. $\langle \{1,2,3,4\} \{5,6\} \rangle$ No
3. $\langle \{2,4\} \{2,4\} \{6\} \rangle$ Yes
4. $\langle \{1\} \{2,4\} \{6\} \rangle$
5. $\langle \{1,2\} \{3,4\} \{5,6\} \rangle$

Are they subsequences of d , subject to the following timing constraints?
mingap=0 (interval between last event in e_i and first event in $e_{i+1} > 0$)

maxgap=3 (interval between first event in e_i and last event in $e_{i+1} \leq 3$)

maxspan=5 (interval between first event in e_1 and last event in $e_n \leq 5$)

ws=1 (time between first and last events in $e_i \leq 1$)

Exercise: timing constraints

$d = \langle \{1,2,3\} \{2,4\} \{2,4,5\} \{3,5\} \{6\} \rangle$

Given are the following sequences $w = \langle e_1 e_2 \dots e_i e_{i+1} \dots e_n \rangle$

1. $\langle \{1\} \{2\} \{3\} \rangle$ Yes
2. $\langle \{1,2,3,4\} \{5,6\} \rangle$ No
3. $\langle \{2,4\} \{2,4\} \{6\} \rangle$ Yes
4. $\langle \{1\} \{2,4\} \{6\} \rangle$ Yes
5. $\langle \{1,2\} \{3,4\} \{5,6\} \rangle$

Are they subsequences of d , subject to the following timing constraints?
mingap=0 (interval between last event in e_i and first event in $e_{i+1} > 0$)

maxgap=3 (interval between first event in e_i and last event in $e_{i+1} \leq 3$)

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Exercise: timing constraints

$d = \langle \{1,2,3\} \{2,4\} \{2,4,5\} \{3,5\} \{6\} \rangle$

Given are the following sequences $w = \langle e_1 e_2 \dots e_i e_{i+1} \dots e_n \rangle$

1. $\langle \{1\} \{2\} \{3\} \rangle$ Yes
2. $\langle \{1,2,3,4\} \{5,6\} \rangle$ No
3. $\langle \{2,4\} \{2,4\} \{6\} \rangle$ No
4. $\langle \{1\} \{2,4\} \{6\} \rangle$ Yes
5. $\langle \{1,2\} \{3,4\} \{5,6\} \rangle$ No (violates mingap and maxgap constraint)

Are they subsequences of d , subject to the following timing constraints?
mingap=0 (interval between last event in e_i and first event in $e_{i+1} > 0$)

maxgap=3 (interval between first event in e_i and last event in $e_{i+1} \leq 3$)

maxspan=5 (interval between first event in e_1 and last event in $e_n \leq 5$)

ws=1 (time between first and last events in $e_i \leq 1$)