

# Bayesian Games

## PhD Microeconomics II

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# Lecture 3 Readings

- Recommended readings

- Fudenberg and Tirole (1991), pp 209-234
- MWG (1995), pp 253-257, pp 282-288

- Further reading

- Gibbons (1992), pp143-163
- Gintis (2009), pp121-131
- Hargreaves-Heap & Varoufakis (2004), pp 60-76
- Myerson (1991), pp 67-83, 127-136, 163-173
- Osborne (2004), pp 271-357
- Osborne & Rubinstein (1994), pp 24-30
- Vega-Redondo (2003), 117-120, 188-204, 217-221

- Seminal contributions

- Aumann (*Ann. of Stat.* 1976) "Agreeing to disagree"
- Harsanyi (*Mang. Sci.* 1967-68) "Games with incomplete informat..."
- Harsanyi (*IJGT* 1973) "Games with randomly perturbed payoffs..."
- Mertens & Zamir (*IJGT* 1985) "Formulation of Bayesian analysis..."

- Recall one of the early assumptions we made in lec 1
  - **Assumption 2:** Players have **complete information**
    - i.e. know all payoff functions
- Is this really true in applications?!
  - Does an employer know how innately productive a worker is?
  - Does the government know people's valuation of a public good?
  - Do you know whether I enjoy being a mean or nice examiner?
  - Does a firm know whether their competitor is low or high cost?
  - Does a car buyer know the true quality of a car?
  - Does an insurance company know how healthy you are?
  - etc...
- Often payoff functions are private information
- **Definition:** A game of **incomplete information** is where at least one payoff function is not known by every player.

# Example: Market Entry with Incomplete Information

- Simple market entry game with incomplete information
  - Player 1 (incumbent):  $S_1 = \{\text{build new factory, don't build}\}$
  - Player 2 (potential entrant):  $S_2 = \{\text{enter, don't enter}\}$
  - Incumbent's cost is private information
  - Sometimes we can use iterative dominance

	<i>E</i>	<i>DE</i>		<i>E</i>	<i>DE</i>
<i>B</i>	0, -1	2, 0	<i>B</i>	3, -1	5, 0
<i>DB</i>	2, 1	3, 0	<i>DB</i>	2, 1	3, 0
	cost = high			cost = low	

- 1 always has dominant strategy: *B* if low cost, *DB* if high cost
- Let  $p_H$  denote 2's prior probability that 1 is high cost
- Considering 2's expected payoff, 2 plays *E* iff  $p_H > 1/2$ , solved.

## Example: Market Entry with Incomplete Information

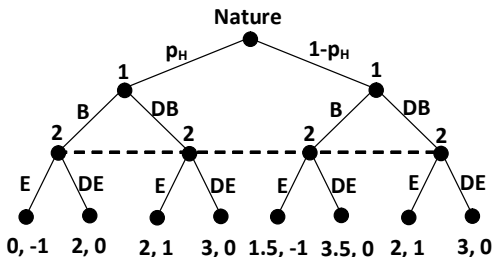
- Typically iterative dominance does not get us far

	<i>E</i>	<i>DE</i>		<i>E</i>	<i>DE</i>
<i>B</i>	0, -1	2, 0	<i>B</i>	1.5, -1	3.5, 0
<i>DB</i>	2, 1	3, 0	<i>DB</i>	2, 1	3, 0
cost = high			cost = low		

- DB* still dom strat if high cost, but no dom strat if low cost
- So 1 has to predict whether or not 2 plays *E*
- Suppose  $y$  is the prob 2 plays *E*
- Examining expected payoffs gives, 1 plays *B* iff  $y < 1/2$
- So 1 has to predict 2's behaviour to BR and 2 cannot infer 1's actions from knowledge of player 1's possible payoffs alone – stuck!

# Harsanyi's Transformation

- Harsanyi (1967/68): transform *incomplete* info into *imperfect* info
  - Incomplete info on costs  $\rightarrow$  Imperfect info about "Nature"'s move



- Nature chooses Incumbent's "type" (cost)
- Entrant cannot distinguish between the different types
- **Assumption 8:** Players have **common priors** on moves by Nature
- NE of imperf info game is a *Bayesian Nash Equilibrium*

# Bayesian-Nash Equilibrium of Market Entry Game

- Solving for the BNE

- Know player 1 plays  $DB$  if high cost, so just consider low cost:

- Let  $x$  be player 1's probability playing  $B$  when cost is low
- Let  $y$  be player 2's probability playing  $E$
- We've already worked out 1's BR

$$x = \begin{cases} 0 & \text{if } y > 1/2 \\ \in [0, 1] & \text{if } y = 1/2 \\ 1 & \text{if } y < 1/2 \end{cases}$$

- Compare 2's expected payoffs to get his BR

$$y = \begin{cases} 0 & \text{if } x > 1/[2(1-p_H)] \\ \in [0, 1] & \text{if } x = 1/[2(1-p_H)] \\ 1 & \text{if } x < 1/[2(1-p_H)] \end{cases}$$

- To identify BNE find  $(x, y)$  such that
  - $x$  is optimal for 1 with low cost given 2's strategy
  - $y$  is optimal for 2 given 1's strategy and beliefs  $p_H$
- Two BNE:  $(x = 0, y = 1) \forall p_H$ ;  $(x = 1, y = 0)$  iff  $p_H \leq \frac{1}{2}$

# The Example and Bayesian Games in General

- Reflect on properties of Bayesian games/example
  - Incomplete information game
    - Player 1's cost private info
  - Introduce "Nature" who chooses "type" first → imperf info game
    - Nature decided high or low cost for player 1
  - *Ex ante* probabilities same for all players, common prior assumption
    - Both knew  $p_H$  and  $(1 - p_H)$
  - Nature chooses, endowing players with private info in the *interim*
    - Player 1 found out high or low cost, player 2 didn't
  - $i$ 's strategies in transformed game: action for each of  $i$ 's poss types
    - Player 1 had an action for high cost and an action for low cost
    - Player 2 had 1 action
    - But if player 1 knows his type why a strategy for each of his types?!
    - Similar logic to extensive form's complete contingent plans... For 1 to BR need to think what 2 will do, 2 does not know cost, so thinks probabilistically given  $p_H$  about what 1 would do if high and if low cost.



# The Example and Bayesian Games in General

- Reflect on properties of Bayesian games/example
  - A Bayesian Nash equilibrium consists of two things:
    - 1 A strategy profile
    - 2 Beliefs specified for each player about "types" of others
  - Our two BNE,  $(s_1^*(cost), s_2^*, p_H)$ , were

$$\left( \begin{array}{l} DB \text{ if } cost = low \\ DB \text{ if } cost = high \end{array}, E, \forall p_H \right) \& \left( \begin{array}{l} B \text{ if } cost = low \\ DB \text{ if } cost = high \end{array}, DE, p_H \leq \frac{1}{2} \right)$$

- A BNE requires each player be maximising  $E[\text{payoff}]$  given others' strategies and beliefs (accurate on own type and probabilistic on others')
  - 1 knows his type so maximises his  $E[\text{payoff}]$  given 2's strategy (which depends on  $p_H$ )
  - 2 doesn't know 1's type, so maximises  $E[\text{payoff}]$  given  $p_H$  and 1's strategy
- *Ex post* payoff profile depends on full set of choices (including Nature's)

# Bayesian Game Description

- Augment game description with "types" – the private info
- Bayesian game description

$$\Gamma^B := (N, \{S_i\}, \Theta, p(\cdot), \{u_i(s, \theta)\})$$

- Players –  $N$
- Types –  $\Theta, p(\cdot)$ 
  - $i$ 's type  $\theta_i \in \Theta_i$ , a random variable, realisation known only to  $i$
  - Assume  $|\Theta_i|$  is finite
  - $p(\theta_1, \dots, \theta_n)$  is the objective joint prob distribution of  $\{\theta_i\}_{i \in N}$
  - Assume  $p(\theta_i) > 0$  for all  $\theta_i \in \Theta_i$
  - Let  $p(\theta_{-i} | \theta_i)$  denote  $i$ 's conditional prob of  $-i$ 's type given own type
  - Let  $\theta = \times_{i \in N} \theta_i$  and  $\Theta = \times_{i \in N} \Theta_i$
  - Everything is common knowledge except realised  $\{\theta_i\}_{i \in N}$
- Strategies –  $S_i$ 
  - Choices over actions (contingent plans for extensive form, later)
  - $s \in S$ , usual definitions of strategy profiles
- Payoffs –  $u_i(s, \theta)$

# Bayesian Nash Equilibrium

- Harsanyi transformation – Nature chooses  $\theta_i$  first

- Strategies can be conditioned on  $\theta_i$ 
  - A pure strategy is a function of type  $s_i(\theta_i)$
  - Where  $s_i(\cdot) \in S_i^{\Theta_i}$ , all maps  $\Theta_i \rightarrow S_i$
  - Strategy profiles,  $s(\cdot) \in S^{\Theta}$ , defined in usual way

- **Definition:**  $s(\cdot)$  is a **Bayes-Nash equilibrium** if for each  $i$

$$s_i(\cdot) \in \arg \max_{s'_i \in S_i^{\Theta_i}} \sum_{\theta_i} \sum_{\theta_{-i}} p(\theta_i, \theta_{-i}) u_i(s'_i(\theta_i), s_{-i}(\theta_{-i}), (\theta_i, \theta_{-i})).$$

- Since  $p(\theta_i) > 0$  for all  $i$ , ex ante and interim conditions equivalent

$$s_i(\theta_i) \in \arg \max_{s'_i \in S_i} \sum_{\theta_{-i}} p(\theta_{-i} | \theta_i) u_i(s'_i, s_{-i}(\theta_{-i}), (\theta_i, \theta_{-i})).$$

- $i$  plays BR to conditional distribution of opponents' strategies for each type that he might end up playing (since doesn't know which)
- BNE existence: appeal to Nash's Theorem, transformed game standard

# Another Example: Public Goods Provision

- Public Goods Provision Game

- $i = \{1, 2\}$ ,  $A_i = \{\text{contribute}, \text{don't}\}$
- Payoffs

	$C$	$D$
$C$	$1 - c_1, 1 - c_2$	$1 - c_1, 1$
$D$	$1, 1 - c_2$	$0, 0$

- Complete information

- (Practice for you: find all the NE of this game), for now note:

- 1 Never efficient to play  $CC$  &  $CC$  is never an NE (assuming  $c_{1,2} \neq 0$ )
- 2 If  $c_1 < 1 < c_2$  then  $CD$  is efficient & this is the NE (analogous for  $DC$  when  $c_2 < 1 < c_1$ )
- 3 If  $\max\{c_1, c_2\} < 1$  then NE is either  $CD$  or  $DC$ , provided in both NE so more efficient than  $DD$ , not most efficient if higher cost provides

# Public Goods Provision with Incomplete Info

- Introduce incomplete information

- $c_i$  private info

- Common knowledge that  $c_i$  independently drawn from same continuous strictly increase cdf  $P(\cdot)$  over  $[\underline{c}, \bar{c}]$ , where  $\underline{c} < 1 < \bar{c}$ .

- Strategies and payoffs

- $i$ 's pure strategy,  $s_i(c_i) : [\underline{c}, \bar{c}] \rightarrow \{0, 1\}$  (contribute or not)
  - $i$ 's payoff,  $u_i(s_i, s_j, c_i) = \max[s_1, s_2] - c_i s_i$

- BNE

- Where  $(s_i^*(\cdot), s_j^*(\cdot))$  such that for each  $i$  and every possible  $c_i$ ,  $s_i^*(c_i)$  maximises  $E_{c_j}[u_i(s_i, s_j^*(c_j), c_i)]$
  - Let  $z_j \equiv \Pr(s_j^*(c_j) = 1)$  i.e. eqm prob of opponent contributing
  - $i$  contributes if cost < benefit  $\cdot (1 - z_j)$ , thus  $s_i^*(c_i) = 1$  iff  $c_i < 1 - z_j$
  - Thus types  $[\underline{c}, c_i^*]$  contribute (empty if  $c_i^* < \underline{c}$ )
  - Similarly,  $j$  contributes iff  $c_j \in [\underline{c}, c_j^*]$
  - Since  $z_j = \Pr(\underline{c} \leq c_j \leq c_j^*) = P(c_j^*)$ , eqm cutoff satisfies  $c_i^* = 1 - P(c_j^*)$
  - Then if unique,  $c_i^*$  and  $c_j^*$  should satisfy  $c^* = 1 - P(1 - P(c^*))$

# The Effect of Incomplete Information

- What do the BNE look like? Effect of incomplete cost info?
- Consider two illustrative cases
  - $P(\cdot)$  uniform on  $[0, 2]$ ,  $P(c) \equiv \frac{c}{2}$ 
    - $c^*$  is unique, sub  $\frac{c}{2}$  into BNE condition, solve  $c^*$ ,  $c^* = \frac{2}{3}$
    - **Result:** Less efficient than complete info
    - Under-contribution: Play  $D$  for  $c \in (2/3, 1)$  despite ex post benefit  $>$  cost, &  $1 - P(c^*) = \frac{2}{3}$  chance not supplied by other
    - Over-contribution:  $c \in [0, 2/3]$  both play  $C$  when only need one to.
  - Some  $P(\cdot)$  on  $[\underline{c}, \bar{c}]$  where  $\underline{c} \geq 1 - P(1)$ 
    - 2 asymmetric BNE
    - One always plays  $D$ , other always plays  $C$  for  $c \leq 1$
    - e.g. BNE with player 1 always  $D$ :  $c_1^* = 1 - P(1) \leq \underline{c}$  and  $c_2^* = 1$
    - i.e. 1 never contributes because min cost greater than gain; 2 always contributes since otherwise zero probability of provision.
    - **Result:** Efficiency vs uniform unclear, less efficient than complete
    - BNE where 1 always  $D$ , if  $c_1 < \frac{2}{3}$  &  $c_2 > 1$ , not provided (is if uniform)
    - BNE where 2 always  $D$ , if  $c_1 < \frac{3}{4}$  &  $c_2 > 1$ , provided (not if uniform)
    - Coordination on inefficient BNE possible

# The Example and Bayesian Games in General

- Prior beliefs are v important in Bayesian games
  - Note how different the BNE looked with the two different cdfs
- Often assume types drawn independently
  - $p(\theta_{-i} | \theta_i) = p(\theta_{-i})$  here, but types correlated in some games
- Bayesian games often have many BNE
  - Complete info game had 2 NE
  - Incomplete info game had 3 BNE that we identified (may be more)
  - NE are supported by consistent strategies
  - BNE supported by consistent beliefs and strategies
    - Many more possible combinations
- Monotonicity and cut-off rules are common in BNE
  - Despite the differences between the equilibria
  - Equilibrium contribution strategies were monotonic functions of type
  - Contribute up to some type, then don't – cut-off rule
- Information often has efficiency consequences
  - In our game incomplete info reduced efficiency (a common result)

# Prior Beliefs

- Where do they come from?
  - Similar to Nash conjectures – e.g. culture, focal, learning etc
  - Could be an objective distribution e.g. male-female ratio in species
- What form "should" they take?
  - Modellers often use uniform or other typical statistical distributions
  - Best to be as general as possible subject to tractability constraints
- Is everyone having the same priors plausible? (Assumption 8)
  - Same info sets and CKR then technically, yes.
  - Aumann (1976): Even if different info sets, still common beliefs:
    - Given CKR, moment players discover holding differing beliefs, incentive to revise beliefs (to incorporate new info)
    - Rational players cannot "agree to disagree"
  - Reasonable?
    - "How many coins in the jar?" vs "Does god exist?"
    - Repeated game vs one-shot game (common *priors* even then?)
    - Costly information transmission
    - Incentives to acquire info – if all free-ride, could agree to disagree!



# Interim vs Ex Ante Strict Dominance

- A BNE is an *equilibrium* in the sense that it's predictable
- So player  $i$  must
  - Predict player  $j$ 's strategy choice – to do...
    - Consider how each  $j \neq i$  *thinks* player  $i$  will play
    - Consider  $j$ 's *beliefs* about  $i$ 's type
- In this prediction process, how should we view types  $\theta_i$  and  $\theta'_i$ ?
  - 1 A single player making type-contingent decisions at ex ante stage
  - 2 Two different "individuals" one of which nature will pick to play
- In 1, ex-ante predictions so all types of  $i$  predict the same
  - Similar to Harsanyi's original formulation
- In 2, interim predictions so predictions may differ between types of  $i$ 
  - e.g. genetically determined preferences – ex ante impossible
- No difference for BNE (players have common beliefs)
- But does it matter for strict dominance?

# Strict Dominance and the Public Good Game

- Interim strict dominance
  - Given  $c_i$ , which of  $i$ 's strategies are not strictly dominated?
    - $D$ : not strictly dominated (play  $D$  if expect  $j$  to contribute,  $\forall c_i$ )
    - If  $c_i > 1$  then  $C$  strictly dominated for  $i$
    - If  $\underline{c} > 1 - P(1)$ , no more dominated strategies
    - So, for example, interim dominance permits  $c_i \in [\underline{c}, c']$  don't contribute and  $c_i \in (c', 1]$  do contribute, former expect  $j$  to contribute if  $c_j < 1$  and latter expect  $j$  never contributes
    - Could not happen in a BNE (cut-off rule/monotonicity in equilibrium)
- Ex ante strict dominance
  - Don't know  $c_i$  yet, which of  $i$ 's strategies is not strictly dominated?
    - Any  $s_i(\cdot)$  that has player contribute with prob  $z > 0$  and is not a cut-off rule is strictly ex ante dominated by a strategy where player contributes iff  $c_i < c'$ , where  $c' = P(z)$
    - For any  $s_j$ ,  $i$  receives public good with same prob, but his expected cost of provision is strictly lower
    - Intuitively, if  $i$  is a single player, then any beliefs of  $j$ 's strategy that make it attractive to contribute at  $c'$  also make it attractive to contribute at  $c_i < c'$

# Interim vs Ex Ante Strict Dominance

- Generally, more strategies dominated ex ante than interim

- For a given type-contingent strategy  $\hat{\sigma}_1(\cdot)$  of player 1
  - Easier to find  $\sigma_1(\cdot)$  satisfying ex ante dominance condition

$$\begin{aligned} & \sum_{\theta_1} p_1(\theta_1) \sum_{\theta_{-1}} p(\theta_{-1} | \theta_1) u_1(\sigma_1(\theta_1), \sigma_{-1}(\theta_{-1}), \theta) \\ & > \sum_{\theta_1} p_1(\theta_1) \sum_{\theta_{-1}} p(\theta_{-1} | \theta_1) u_1(\hat{\sigma}_1(\theta_1), \sigma_{-1}(\theta_{-1}), \theta) \end{aligned}$$

for all  $\sigma_{-1}(\cdot)$ , than to find  $s_1$  and  $\theta_1$  satisfying interim constraints

$$\begin{aligned} & \sum_{\theta_{-1}} p(\theta_{-1} | \theta_1) u_1(s_1, \sigma_{-1}(\theta_{-1}), \theta) \\ & > \sum_{\theta_{-1}} p(\theta_{-1} | \theta_1) u_1(\hat{\sigma}_1(\theta_1), \sigma_{-1}(\theta_{-1}), \theta) \end{aligned}$$

for all  $\sigma_{-1}(\cdot)$ .

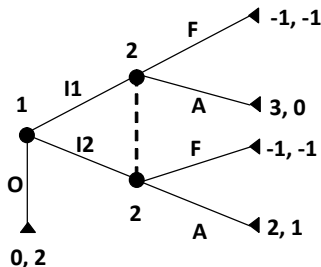
- Intuition:
  - Interim: Many beliefs (differing by type), need domination for all
  - Ex ante: One belief (same for all types), domination easier

# Purification of Mixed Strategy Equilibria

- Critique of using mixed strategies in complete info games
  - People don't flip coins to make decisions!
- Can justify mixed strategies using Bayesian games
  - Complete info game example:
    - $S_i = \{invest, don't\}$ ;  $u_i$ : 1 if only  $i$  invests,  $-1$  if both do,  $0$  if don't
    - Only symmetric NE is playing invest with prob  $\frac{1}{2}$
  - Introduce incomplete information
    - If only  $i$  invests gets  $(1 + \theta_i)$  where  $\theta_i$ , private, is uniform on  $[-\varepsilon, \varepsilon]$
    - BNE: symmetric pure strats  $s_i(\theta_i < 0) = \text{don't}$  and  $s_i(\theta_i \geq 0) = \text{invest}$
    - Each firm expects other to invest with prob  $\frac{1}{2}$ , invest iff  $\frac{1}{2}(1 + \theta_i) + \frac{1}{2}(-1) \geq 0$ , i.e.  $\theta_i \geq 0$
    - As  $\varepsilon \rightarrow 0$ , pure strategy BNE  $\rightarrow$  mixed-strategy NE of complete info
- Known as **purification** of a mixed strategy equilibrium
  - **Result** (Harsanyi, 1973): Any MSNE can be obtained as limit of pure strategy equilibrium in sequence of slightly perturbed games.
    - Intuition: players play pure strats in Bayesian games, don't know opponent's type, so effectively play as if facing a mixed strat

# Extensive Form Bayesian Games

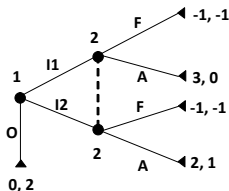
- Only considered incomplete info in strategic form so far
- Extensive form example:



- 2 pure strat NE:  $(O, F \text{ if enter})$  &  $(I1, A \text{ if enter})$
  - First NE doesn't seem reasonable: 2 prefers to accom once enters
  - Subgame Perfection too weak – only 1 subgame, both NE are SPNE!
- Need something stronger for extensive form Bayesian games

# Continuation Games vs Subgames

- Subgame perfection has no bite as only 1 well-defined subgame



- Definition:** A **continuation game**,  $C$  of  $\Gamma^E$ , is a subset of  $\Gamma^E$ 
  - Starting at an information set  $h$  (may or may not be singleton)
  - Containing only all successor nodes of  $x \in h$
  - If  $x' \in C$  and  $x'' \in h(x')$ , then  $x'' \in C$  (i.e. no chopping up info sets)
  - Adopt other properties (payoff functions, player/action labels) from  $\Gamma^E$
- How many continuation games in the example?
- To apply the "spirit" of subgame perfection for continuation games
  - Need a probability distribution over nodes in the first info set
  - Then require BNE in every continuation game

# Beliefs and Sequential Rationality

- **Definition:** A **system of beliefs** in  $\Gamma^E$  is a probability  $\mu(x) \in [0, 1]$  for each  $x$  in  $\Gamma^E$  such that  $\sum_{x \in h} \mu(x) = 1$  for all  $h \in H$ .
  - i.e. player's assessment of the relative probability of being at each node in that information set conditional upon that info set being reached
- **Definition:** A strategy profile  $\sigma$  in  $\Gamma^E$  is **sequentially rational** at  $h$  given a system of beliefs  $\mu$  if

$$E \left[ u_{l(h)} \mid h, \mu, \sigma_{l(h)}, \sigma_{-l(h)} \right] > E \left[ u_{l(h)} \mid h, \mu, \hat{\sigma}_{l(h)}, \sigma_{-l(h)} \right]$$

for all  $\hat{\sigma}_{l(h)} \in \Delta \left( S_{l(h)} \right)$ . If  $\sigma$  satisfies this condition for all  $h \in H$ , then  $\sigma$  is **sequentially rational given belief system  $\mu$** .

- i.e. no player wants to change strategy once reach an info set given beliefs about what's happened and opponent's strategy

# Consistency of Beliefs

- In a weak perfect Bayesian equilibrium
  - ① Strategies must be sequentially rational given beliefs
  - ② Beliefs must be consistent with strategies
    - Similar to Nash conjectures, i.e. beliefs are correct
    - Like SPNE & NE, require BNE in each continuation game
- Illustrate belief consistency:
  - Assume all  $i$  play completely mixed strategies, every  $h$  reached with positive prob
  - For each  $x$  in a player's  $h$ , compute  $\Pr(x|\sigma)$  then assign conditional prob of being at each  $x$  given play has reached that  $h$  using Bayes' rule

$$\Pr(x|h, \sigma) = \frac{\Pr(x|\sigma)}{\sum_{x' \in h} \Pr(x'|\sigma)}$$

- E.g. Previous game if 1 plays  $\sigma : \sigma(O) = \frac{1}{4}, \sigma(I1) = \frac{1}{2} \text{ \& } \sigma(I2) = \frac{1}{4}$ 
  - Prob of reaching 2's info set given  $\sigma$  is  $\frac{3}{4}$
  - Bayes' rule: prob of being at  $I1$  given info set reached is  $\frac{2}{3}$  ( $I2 = \frac{1}{3}$ )
  - 2's should have beliefs  $\frac{2}{3}$  and  $\frac{1}{3}$  to be consistent with 1's  $\sigma$



# Consistency of Beliefs Off the Path

- What about if players are not playing completely mixed strategies?
- $\exists$  some  $x$  not reached with positive probability
- Cannot use Bayes' rule to calculate the probability of reaching these
- Never go off the eqm path, even if played repeatedly, so how do you form beliefs about these  $x$ ?
- WPBE = agnostic... have whatever beliefs you like off the path!
  - This is why it is "weak" relative to the solution concepts we'll look at next lecture!

# Weak Perfect Bayesian vs Nash Equilibrium

- **Definition:**  $(\sigma, \mu)$  is a **weak perfect Bayesian equilibrium** in  $\Gamma^E$  if

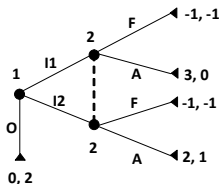
- ①  $\sigma$  is sequentially rational given  $\mu$
- ②  $\mu$  is consistent with  $\sigma$  and Bayes' rule on the path, thus any  $h$  such that  $\Pr(h|\sigma) > 0$  must have

$$\mu(x) = \frac{\Pr(x|\sigma)}{\Pr(h|\sigma)} \text{ for all } x \in h.$$

- Contrast with definition of NE in this context
- **Definition:**  $\sigma$  is an **NE** in  $\Gamma^E$  if  $\exists$  some  $\mu$  such that
  - ①  $\sigma$  is sequentially rational given  $\mu$  for all  $h$  such that  $\Pr(h|\sigma) > 0$
  - ②  $\mu$  is consistent with  $\sigma$  and Bayes' rule on the path
- Sequential rationality
  - Only required on eqm path by NE
  - Required on and off the path by WPBE
  - Thus  $WPBE \subseteq NE$

# Applying WPBE to the Market Entry Example

- Applying WPBE to the example:



- 2 must play "A if enter" in any WPBE
  - This is optimal if get to that info set for any  $\mu$
- Previous NE (O, F if enter) is not WPBE
- Is other NE, (I1, A if enter), WPBE?
  - Need a  $\mu$  satisfying condition 2 that makes these  $\sigma$  sequentially rational
  - To satisfy condition 2, must set  $\mu(I1) = 1$  (info set reached with positive prob given  $\sigma$ )
  - These  $\sigma$  are indeed sequentially rational given that  $\mu$
  - Thus (I1, A if enter) is a WPBE.

# Summary

- Complete information is unrealistic for most applications
- Bayesian games allow us to model incomplete information
- Harsanyi: transform incomplete info game into imperfect info game
  - Then consider the Bayesian-Nash equilibrium
- Beliefs important in Bayesian games, typically assume common priors
- Equilibria characterised by beliefs & strategies so often more equilibria
- Strict dominance: difference between ex ante and interim application
- Mixed strategies can be defended as the limit of perturbed games
- When considering dynamic games, Bayesian-Nash is often too weak
- WPBE requires sequential rationality on and off the path
  - No requirements on beliefs off the path