

Case Study 4: Collaborative Filtering

Graph-Parallel Problems

Synchronous v.
Asynchronous Computation

Machine Learning/Statistics for Big Data
CSE599C1/STAT592, University of Washington

Carlos Guestrin
March 12th, 2013

©Carlos Guestrin 2013

1

Needless to Say, We Need Machine Learning for Big Data



6 Billion
Flickr Photos



28 Million
Wikipedia Pages



1 Billion
Facebook Users



72 Hours a Minute
YouTube

The New York Times
Sunday Review
WORLD U.S. N.Y./REGION BUSINESS TEC

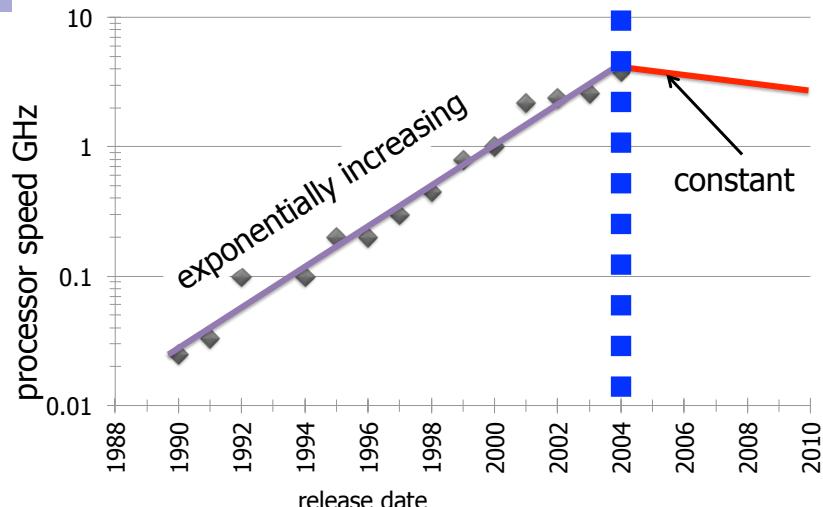
NEWS ANALYSIS
The Age of Big Data
By STEVE LOHR
Published: February 11, 2012

“... data a new class of economic asset, like currency or gold.”

©Carlos Guestrin 2013

2

CPUs Stopped Getting Faster...



3

©Carlos Guestrin 2013

ML in the Context of Parallel Architectures



- But scalable ML in these systems is hard, especially in terms of:
 1. Programmability
 2. Data distribution
 3. Failures

©Carlos Guestrin 2013

4

Move Towards Higher-Level Abstraction

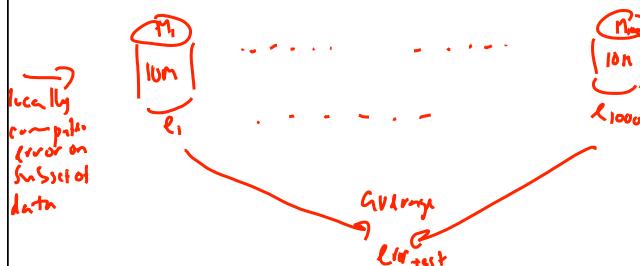
- Distributed computing challenges are hard and annoying!
 1. Programmability
 2. Data distribution
 3. Failures
- High-level abstractions try to simplify distributed programming by hiding challenges:
 - Provide different levels of robustness to failures, optimizing data movement and communication, protect against race conditions...
 - Generally, you are still on your own WRT designing parallel algorithms
- Some common parallel abstractions:
 - Lower-level:
 - Pthreads: abstraction for distributed threads on single machine
 - MPI: abstraction for distributed communication in a cluster of computers
 - Higher-level:
 - Map-Reduce (Hadoop: open-source version): mostly data-parallel problems
 - GraphLab: for graph-structured distributed problems

©Carlos Guestrin 2013

5

Simplest Type of Parallelism: Data Parallel Problems

- You have already learned a classifier
 - What's the test error? $\text{err} = \frac{1}{N_{\text{test}}} \sum_i |y^{(i)} - \text{sign}(w^T x^{(i)})|$
- You have 10B labeled documents and 1000 machines



- Problems that can be broken into independent subproblems are called data-parallel (or embarrassingly parallel)
- Map-Reduce is a great tool for this...
 - Focus of today's lecture
 - but first a simple example

©Carlos Guestrin 2013

6

Data Parallelism (MapReduce)

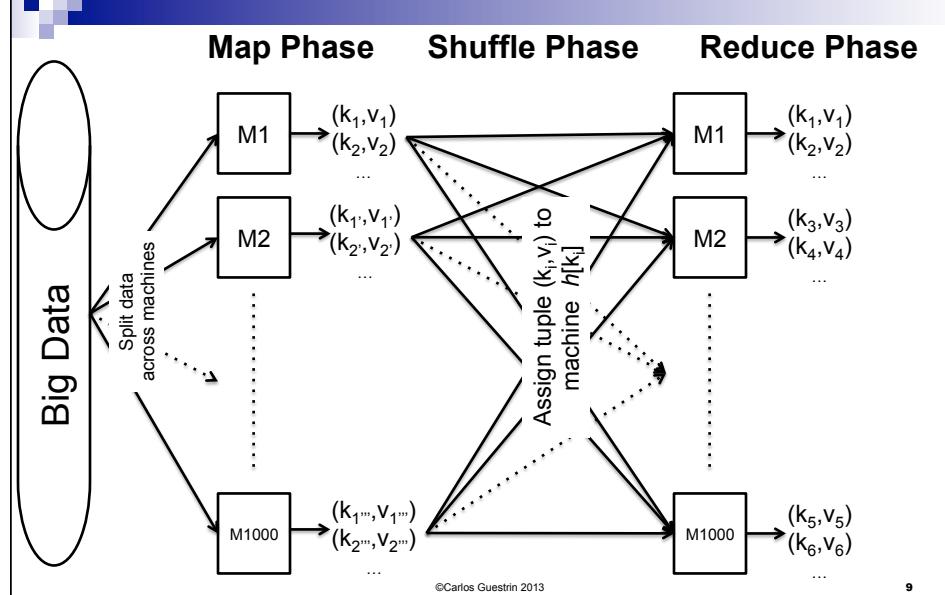


Solve a huge number of **independent** subproblems,
e.g., extract features in images

Map-Reduce Abstraction

- Map: *Transforms a data element*
 - Data-parallel over elements, e.g., documents
 - Generate (key,value) pairs
 - “value” can be any data type
 - Reduce: *Take all values associated with a key and aggregate*
 - Aggregate values for each key
 - Must be commutative-associate operation
 - Data-parallel over keys
 - Generate (key,value) pairs
 - Map-Reduce has long history in functional programming
 - But popularized by Google, and subsequently by open-source Hadoop implementation from Yahoo!
- Word count map (document) for word in doc emit (word, 1)*
- Reduce (word, count: list(int))*
c = 0
for i in count
c += count[i]
emit (word, c)
- map reduce ('Uw', [1, 17, 0, 0, 12])*
emit ('Uw', 30)

Map-Reduce – Execution Overview



Issues with Map-Reduce Abstraction

- Often all data gets moved around cluster
 - Very bad for iterative settings
- Definition of Map & Reduce functions can be unintuitive in many apps
 - Graphs are challenging
- Computation is synchronous

SGD for Matrix Factorization in Map-Reduce?

$$\epsilon_t = L_u^{(t)} \cdot R_v^{(t)} - r_{uv}$$

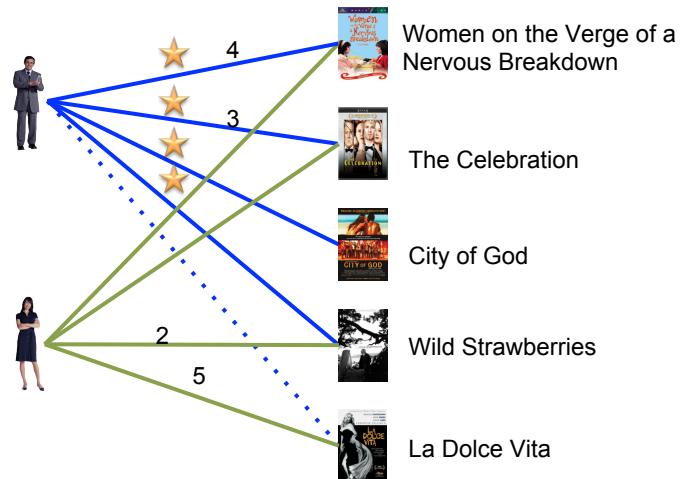
$$\begin{bmatrix} L_u^{(t+1)} \\ R_v^{(t+1)} \end{bmatrix} \leftarrow \begin{bmatrix} (1 - \eta_t \lambda_u) L_u^{(t)} - \eta_t \epsilon_t R_v^{(t)} \\ (1 - \eta_t \lambda_v) R_v^{(t)} - \eta_t \epsilon_t L_u^{(t)} \end{bmatrix}$$

- Map and Reduce functions???
- Map-Reduce:
 - Data-parallel over all mappers
 - Data-parallel over reducers with same key
- Here, one update at a time!

©Carlos Guestrin 2013

11

Matrix Factorization as a Graph



©Carlos Guestrin 2013

12

Flashback to 1998

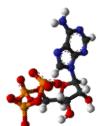


First Google advantage:
a Graph Algorithm & a System to Support it!

Social Media



Science



Advertising



Web



- **Graphs** encode the **relationships** between:

People

Facts

Products

Interests

Ideas

- **Big: 100 billions of vertices and edges** and rich metadata

- Facebook (10/2012): 1B users, 144B friendships
- Twitter (2011): 15B follower edges

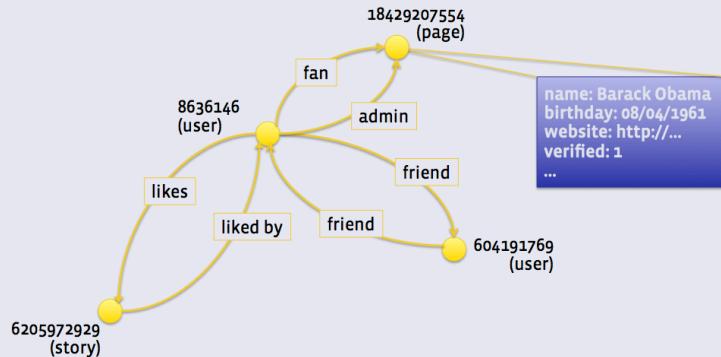
©Carlos Guestrin 2013

14

Facebook Graph

Data model

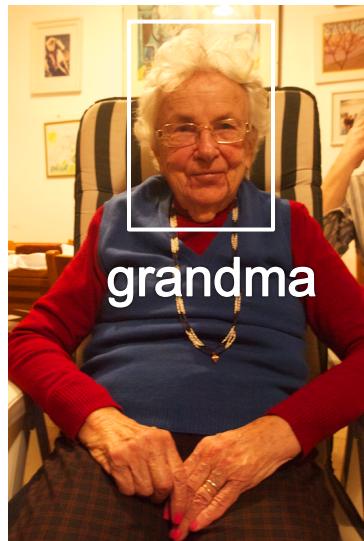
Objects & Associations



Slide from Facebook Engineering presentation 15

©Carlos Guestrin 2013

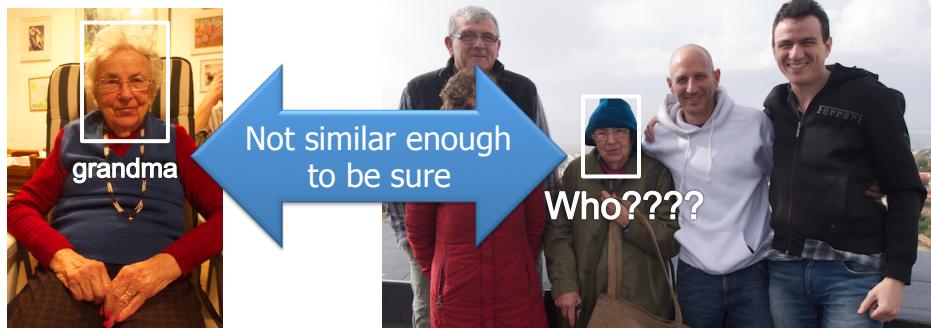
Label a Face and Propagate



©Carlos Guestrin 2013

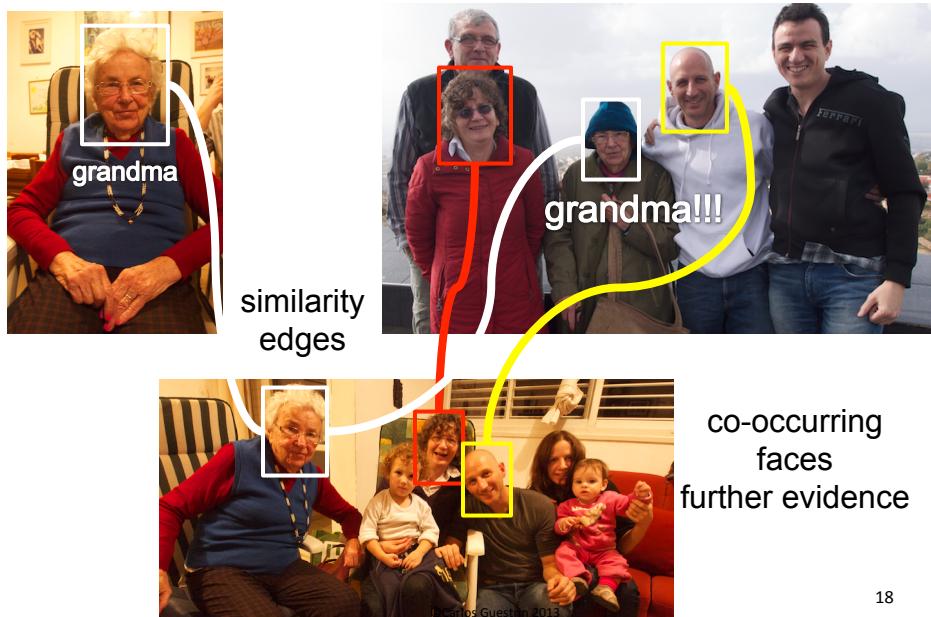
16

Pairwise similarity not enough...



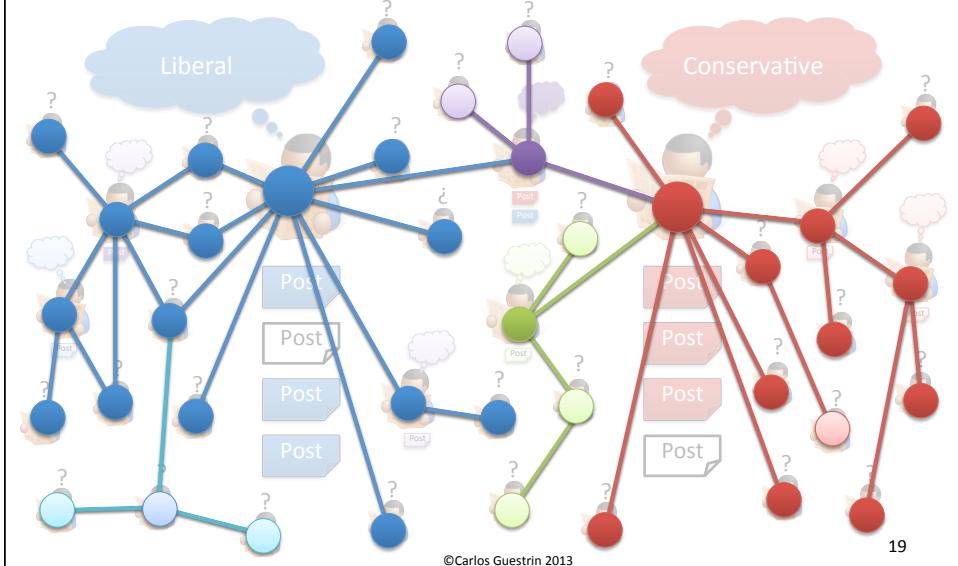
17

Propagate Similarities & Co-occurrences for Accurate Predictions

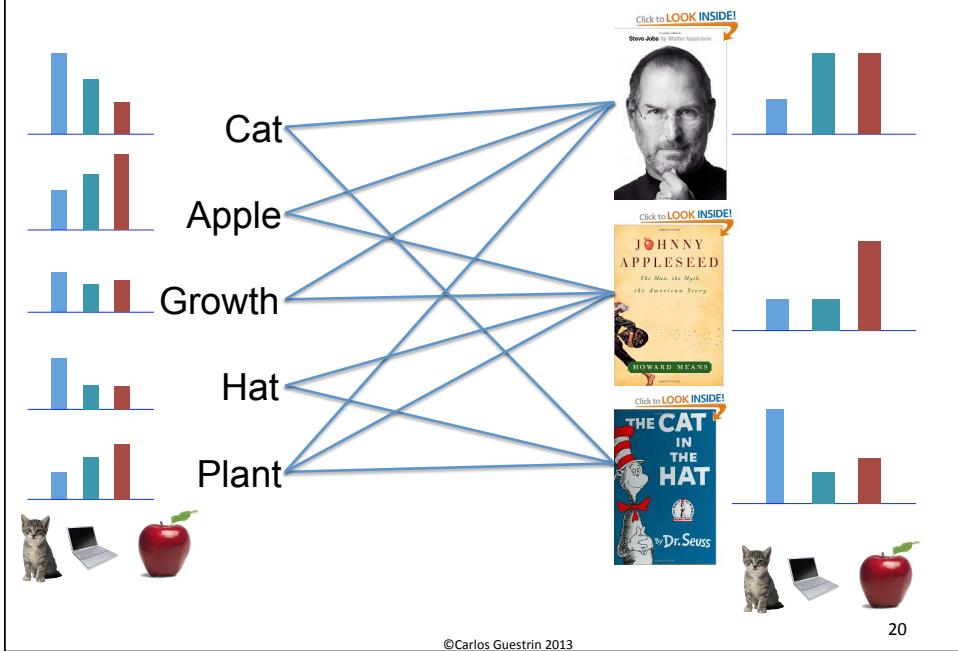


18

Example: Estimate Political Bias



Latent Topic Modeling (LDA)



ML Tasks Beyond Data-Parallelism



Map Reduce

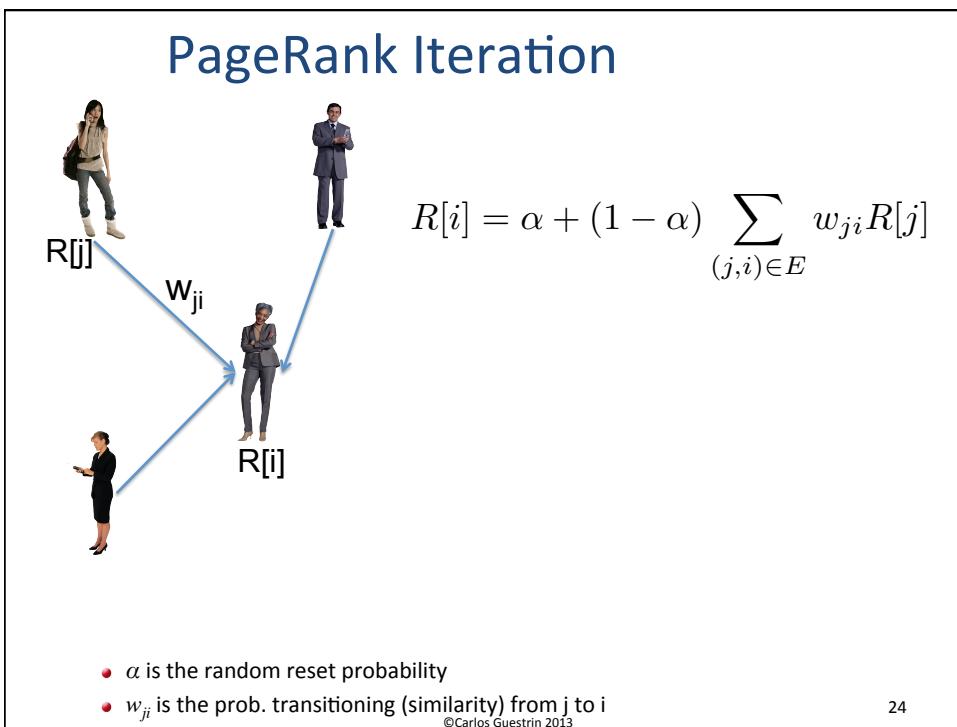
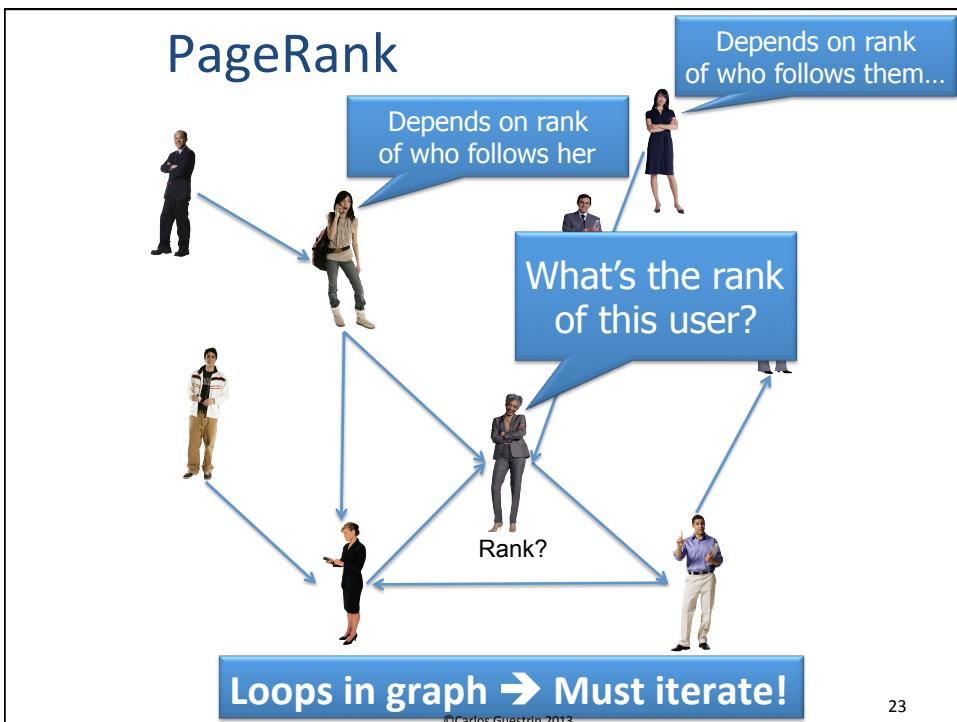
Feature Extraction
Cross Validation
Computing Sufficient Statistics

Graphical Models
Gibbs Sampling
Belief Propagation
Variational Opt.
Semi-Supervised Learning
Label Propagation
CoEM
Collaborative Filtering
Tensor Factorization
Graph Analysis
PageRank
Triangle Counting

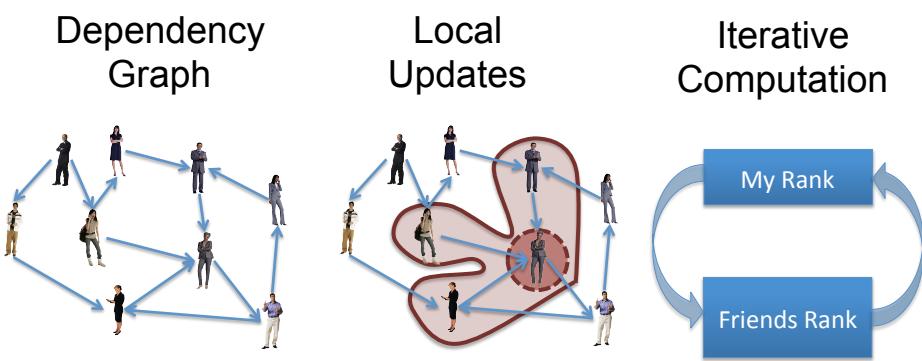
©Carlos Guestrin 2013

21

Example of a Graph-Parallel Algorithm



Properties of Graph Parallel Algorithms



©Carlos Guestrin 2013

25

Addressing Graph-Parallel ML



Map Reduce

Feature Extraction
Cross Validation
Computing Sufficient Statistics

Graph-Parallel Abstraction

Graphical Models **Semi-Supervised Learning**
Gibbs Sampling Label Propagation
Belief Propagation CoEM
Variational Opt.

Collaborative Filtering **Data-Mining**
Tensor Factorization PageRank
Triangle Counting

©Carlos Guestrin 2013

26

Graph Computation:

Synchronous

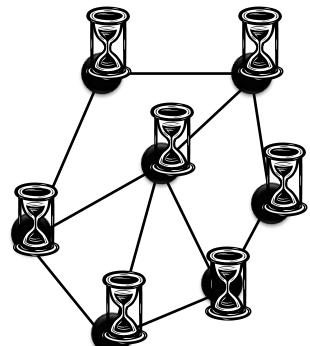
v.

Asynchronous

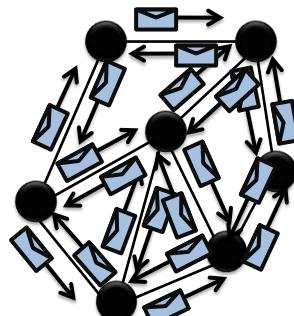
Bulk Synchronous Parallel Model:
Pregel (Giraph)

[Valiant '90]

Compute



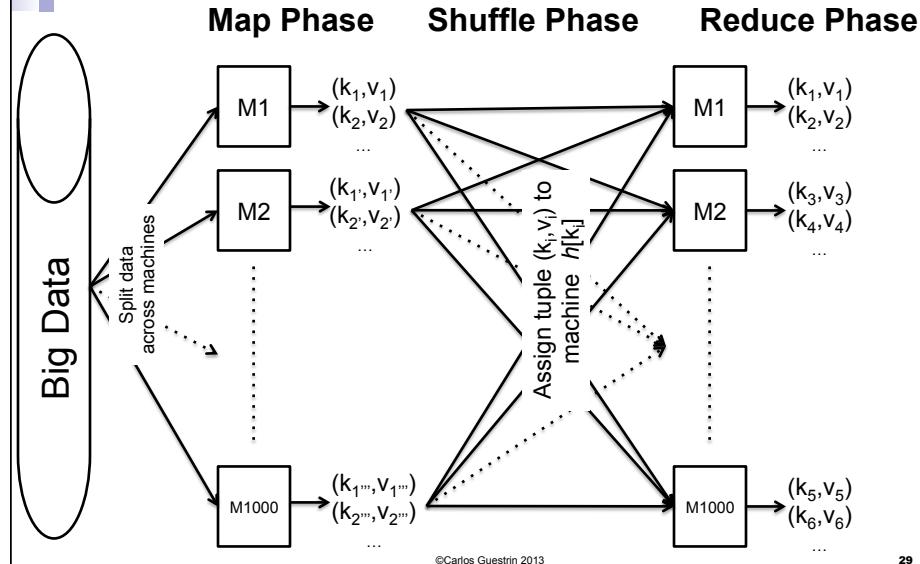
Communicate



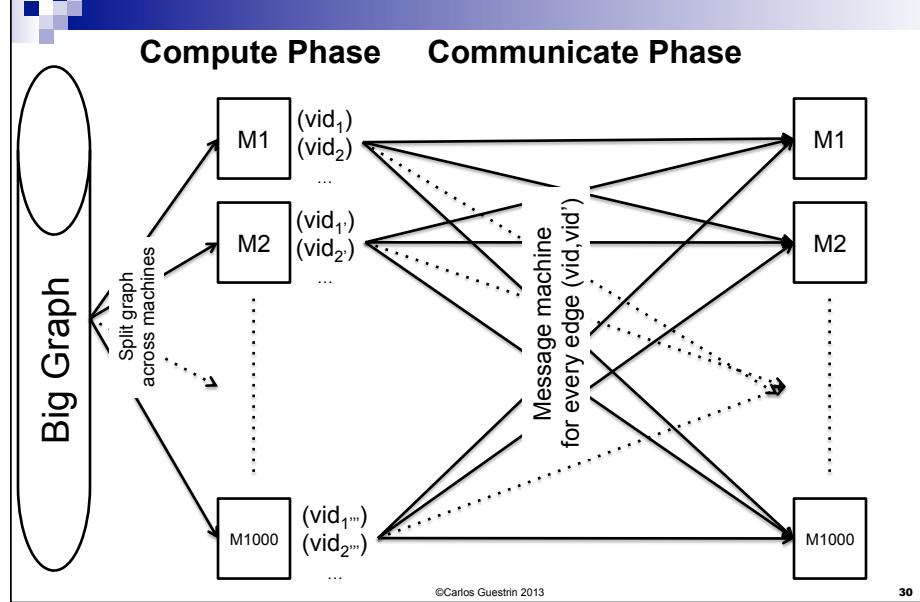
28

©Carlos Guestrin 2013

Map-Reduce – Execution Overview



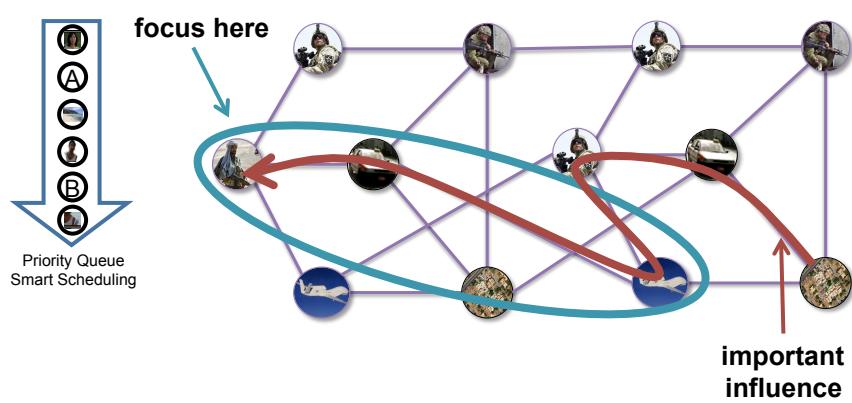
BSP – Execution Overview

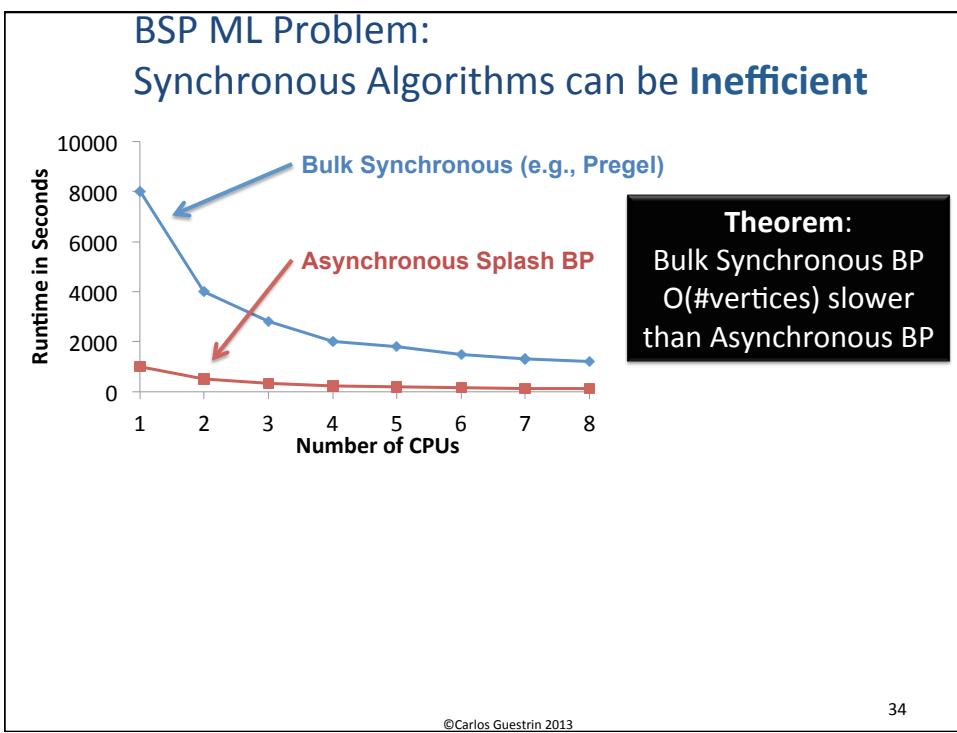
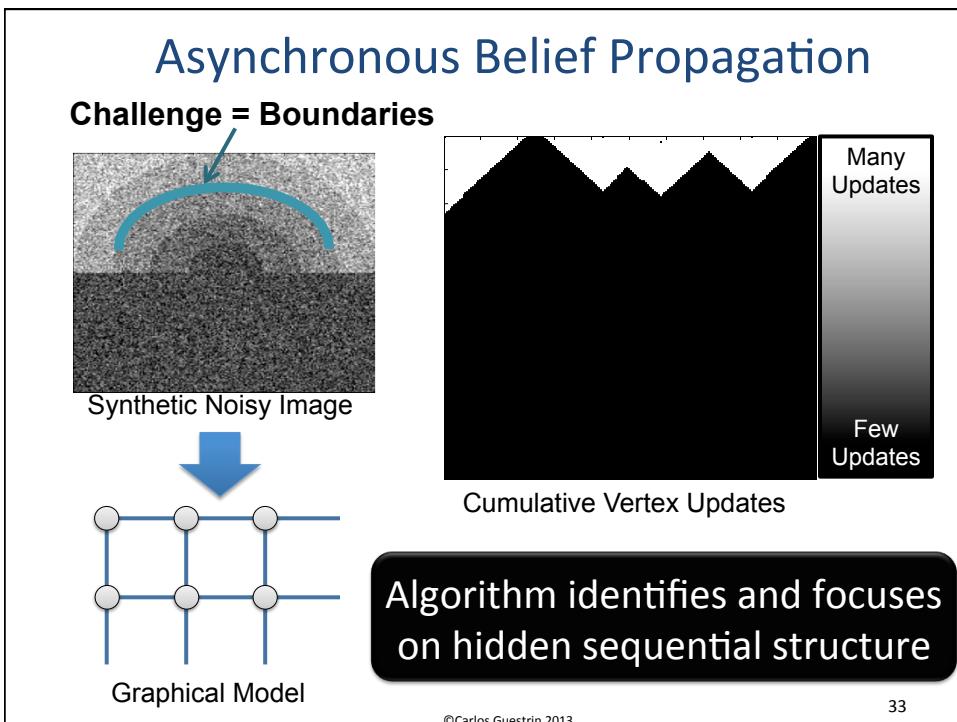


*Bulk synchronous
parallel model
provably inefficient
for some ML tasks*

Analyzing Belief Propagation

[Gonzalez, Low, G. '09]





Synchronous v. Asynchronous

- Bulk synchronous processing:
 - Computation in phases
 - All vertices participate in a phase
 - Though OK to say no-op
 - All messages are sent
 - Simpler to build, like Map-Reduce
 - No worries about race conditions, barrier guarantees data consistency
 - Simpler to make fault-tolerant, save data on barrier
 - Slower convergence for many ML problems
 - In matrix-land, called Jacobi Iteration
 - Implemented by Google Pregel 2010
- Asynchronous processing:
 - Vertices see latest information from neighbors
 - Most closely related to sequential execution
 - Harder to build:
 - Race conditions can happen all the time
 - Must protect against this issue
 - More complex fault tolerance
 - When are you done?
 - Must implement scheduler over vertices
 - Faster convergence for many ML problems
 - In matrix-land, called Gauss-Seidel Iteration
 - Implemented by GraphLab 2010, 2012

©Carlos Guestrin 2013

35

Case Study 4: Collaborative Filtering

GraphLab

Machine Learning/Statistics for Big Data
CSE599C1/STAT592, University of Washington
Carlos Guestrin
March 12th, 2013

©Carlos Guestrin 2013

36

The GraphLab Goals

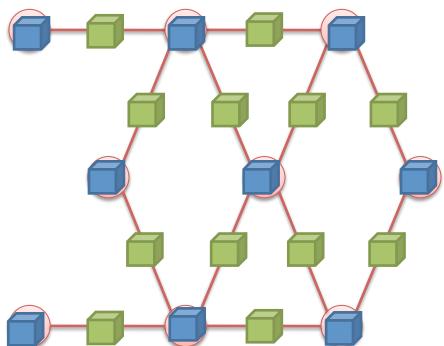


©Carlos Guestrin 2013

37

Data Graph

Data associated with vertices and edges



Graph:

- Social Network

Vertex Data:

- User profile text
- Current interests estimates

Edge Data:

- Similarity weights

©Carlos Guestrin 2013

38

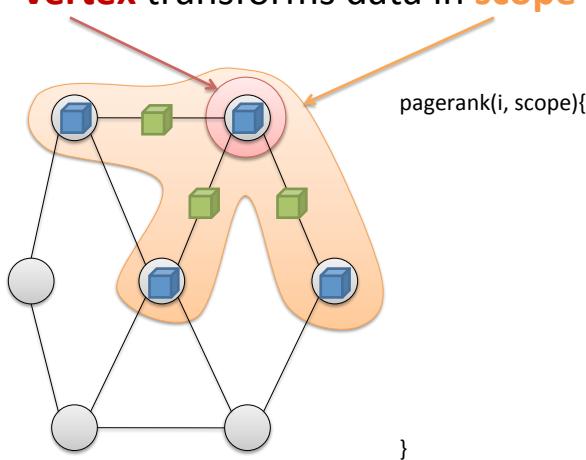
How do we *program* graph computation?

“Think like a Vertex.”

-Malewicz et al. [SIGMOD’10]

Update Functions

User-defined program: applied to
vertex transforms data in **scope** of vertex



Update Function Example: Connected Components

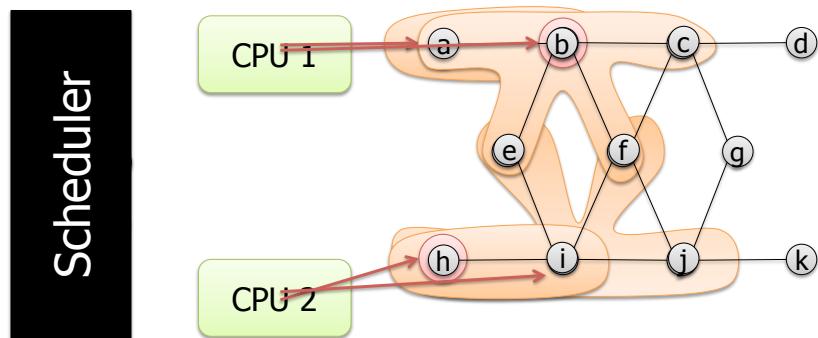


©Carlos Guestrin 2013

41

The Scheduler

The **scheduler** determines order vertices are updated



©Carlos Guestrin 2013

42

Example Schedulers

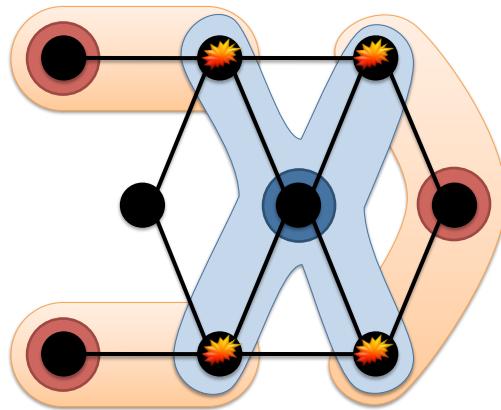
- Round-robin
- Selective scheduling (skipping):
 - round robin but jump over un-scheduled vertex
- FIFO
- Prioritize scheduling
 - Hard to implement in a distributed fashion
 - Approximations used (each machine has its own priority queue)

©Carlos Guestrin 2013

43

Ensuring Race-Free Code

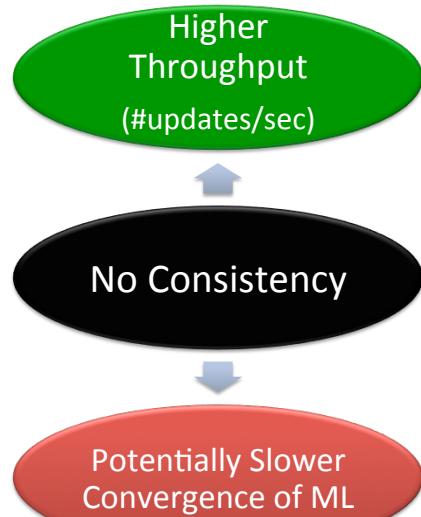
How much can computation **overlap**?



©Carlos Guestrin 2013

44

Need for Consistency?

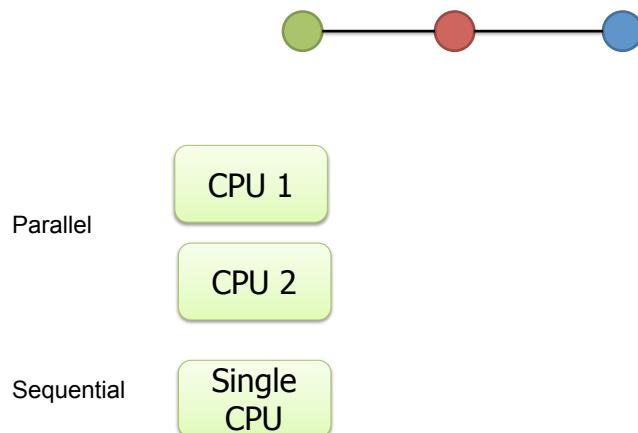


©Carlos Guestrin 2013

45

GraphLab Ensures Sequential Consistency

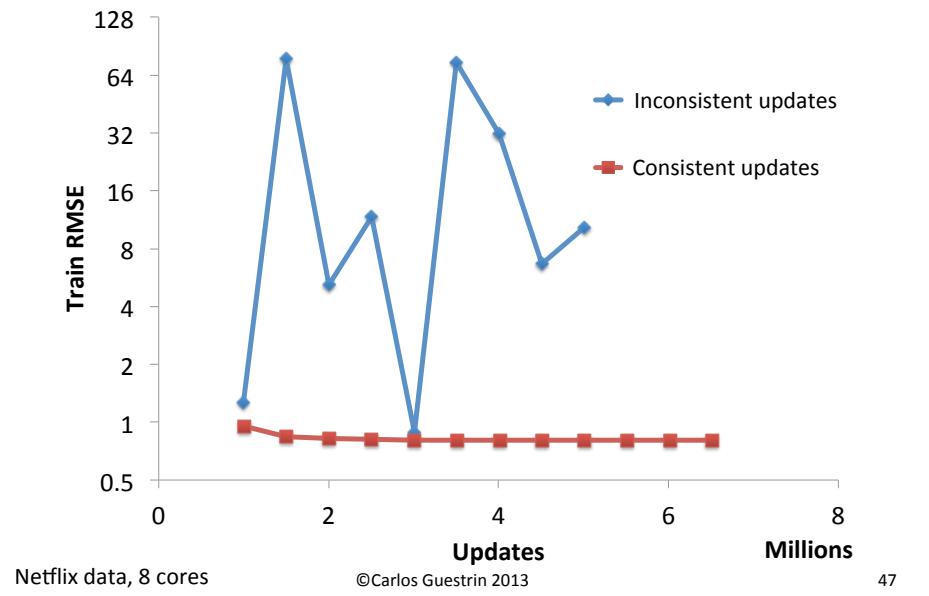
For each parallel execution, there exists a sequential execution of update functions which produces the same result



©Carlos Guestrin 2013

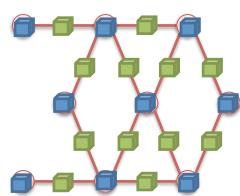
46

Consistency in Collaborative Filtering

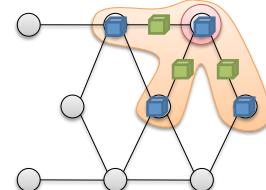


The GraphLab Framework

Graph Based
Data Representation



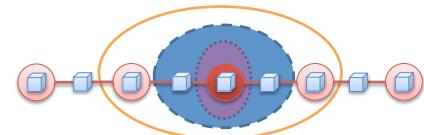
Update Functions
User Computation



Scheduler



Consistency Model



©Carlos Guestrin 2013

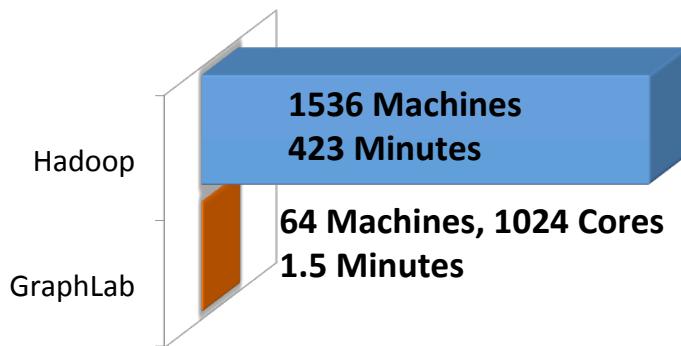
48

Triangle Counting in Twitter Graph



40M Users
1.2B Edges

Total:
34.8 Billion Triangles



Hadoop results from [Suri & Vassilvitskii '11]

49

CoEM (Jones et al., 2005)

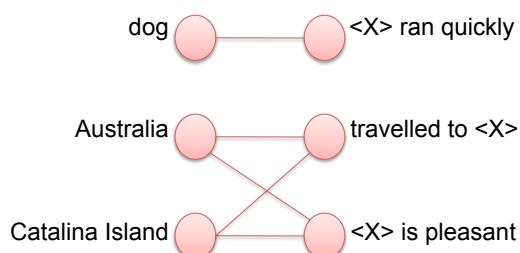
Named Entity Recognition Task

Is "Dog" an animal?

Is "Catalina" a place?

Vertices: 2 Million

Edges: 200 Million



©Carlos Guestrin 2013

50

Never Ending Learner Project (CoEM)

Hadoop	95 Cores	7.5 hrs
Distributed GraphLab	32 EC2 machines	80 secs

©Carlos Guestrin 2013

51

What you need to know...

- Data-parallel versus graph-parallel computation
- Bulk synchronous processing versus asynchronous processing
- GraphLab system for graph-parallel computation
 - Data representation
 - Update functions
 - Scheduling
 - Consistency model

©Carlos Guestrin 2013

52