

The RSA Cryptosystems (6)

RSA Key Generation

- **P** \neq **Q** large prime numbers, **N** = PQ
- Choose random $1 < \mathbf{E} < (P-1)(Q-1)$
s.t. E and $(P-1)(Q-1)$ are relatively prime.
 $\Rightarrow 1 < \mathbf{D} < (P-1)(Q-1)$ multiplicative
inverse **ED** \equiv **1 (mod (P-1)(Q-1))**
- After that, P,Q should be discarded safely.
We will **only use N,E,D** in the cryptosystem.

The RSA Cryptosystems (7)

- $N(=PQ)$, E are public. (E is **Public Key**.)

RSA Encryption (E is **Encryption Key**.)

- ◆ The **plaintext** is $0 \leq X \leq N-1$.
- ◆ The **ciphertext** is $Y \equiv X^E \pmod{N}$.



RSA Decryption (D is **Decryption Key**.)

- ◆ From the **ciphertext** Y ,
calculate $Z \equiv Y^D \pmod{N}$. Then $Z=X$.