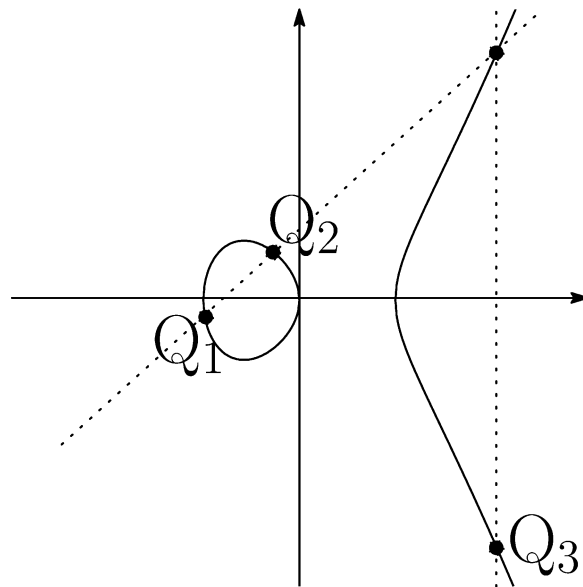


Elliptic Curves and Cryptography (6)

Group Law

- From points Q_1, Q_2 on an elliptic curve, we can create a new point Q_3 .
- We write $Q_3 = Q_1 \oplus Q_2$.
- We define $Q \oplus \infty = Q$.



Elliptic Curves and Cryptography (7)

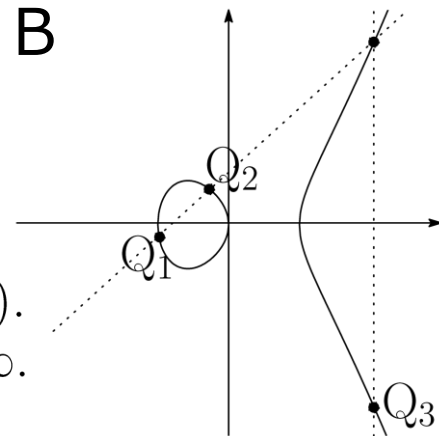
Elliptic Curve $Y^2 = X^3 + AX + B$

- (1) If Q_1 or $Q_2 = \infty$,
 - If $Q_1 = \infty$, put $Q_3 = Q_2$.
 - If $Q_2 = \infty$, put $Q_3 = Q_1$.
- (2) Otherwise, put $Q_1(S_1, T_1)$ and $Q_2(S_2, T_2)$.
 - If $S_1 = S_2$ and $T_1 \neq T_2$, put $Q_3 = \infty$.
 - Otherwise, define $Q_3(S_3, T_3)$ by

$$S_3 = K^2 - S_1 - S_2$$

$$T_3 = K(S_1 - S_3) - T_1$$

$$K = \begin{cases} (T_2 - T_1)/(S_2 - S_1) & \text{if } S_1 \neq S_2 \\ (3S_1^2 + A)/(2T_1) & \text{if } S_1 = S_2. \end{cases}$$



Elliptic Curves and Cryptography (8)

- Given two points Q_1 and Q_2 on the elliptic curve, there is a law (**Group Law**) creating a new point Q_3 .
- Everything is calculated in terms of usual four operations of coordinates.
- We can also define the **Group Law for mod P points**.

Elliptic Curves and Cryptography (9)

- For a mod P point Q , we define

$$[K]Q = Q \oplus \cdots \oplus Q \quad (K-1 \text{ times}).$$

- $[K]Q$ is an analogue of the exponential $A^K \pmod{P}$.
- Using $[K]Q$ instead of A^K , we can construct cryptosystems.

⇒ **Elliptic Curve Cryptography (ECC)**

Summary of Week 4

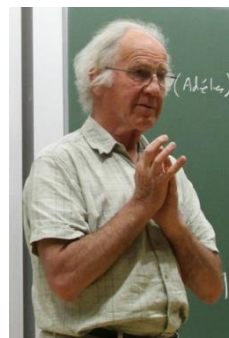
- Basics on Cryptography, Carsar cipher
- The use of Modular Arithmetic in Cryptography.
 - ◆ Discrete Logarithm Problem
 - ◆ Integer Factorization
- The RSA Cryptosystems
- Elliptic Curve Cryptography (ECC)

Plan of Week 5

We will learn more advanced
laws of prime numbers for
elliptic curves.

Let's explore attractive
theorems and conjectures
on elliptic curves!

See you next week!



Bryan John
Birch
(1931-)



Peter
Swinnerton-Dyer
(1927-)

https://en.wikipedia.org/wiki/Bryan_John_Birch

https://en.wikipedia.org/wiki/Peter_Swinnerton-Dyer