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QQ4

QQ4

0 points possible (ungraded)

Which of the following statements are true concerning the Traveling Salesman Problem? (You may check more than one choice)

- a) The solution will always be different depending on the start node
- b) The nearest neighbor heuristic is very fast, but not guaranteed to provide an optimal tour

 ✓
- lacktriangledown c) The nearest neighbor heuristic is a construction algorithm. lacktriangledown
- d) The nearest neighbor heuristic is an improvement algorithm.
- ullet e) If N is the number of nodes in the network, then there will be N arcs in the tour ullet



Explanation

The correct answers are b, c, e..

- a) False. It may be different, but it will not necessarily be different.
- b) True. This is pretty much the definition of a heuristic!
- c) True. The heuristic creates or builds a tour, so it is a construction algorithm
- d) False. Since this is the opposite of (c)!
- e) True. Simply looking at a tour should convince you of this since every node is visted once and only once.

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You have used 1 of 3 attempts

1 Answers are displayed within the problem

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$ \mathbf{Z} $	Part a) How could 2 different starting points yield same paths? If I'm reading the question correctly, we are asking if each unique starting point will have an optimized tr	4
\leq	Part e) If there are N nodes, won't there N-1 arcs? Do we consider starting and ending points to be nodes with set N? If so, shouldn't there be N-1 arcs for	3

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