

组合数学 Combinatorics

4 Linear Homogeneous Recurrence Relation

4-1 Fibonacci Rabbits





The delta of the n^{th} month and $n-1^{th}$ month is given birth by the rabbits in n-2 month. So

$$F_n = F_{n-1} + F_{n-2}$$

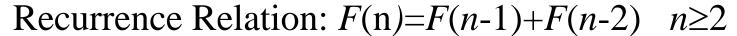
In the first month there's a pair of newly-born rabbits; If a pair of rabbits could give birth to a new pair every month (one male, one female); New rabbits could start giving birth since the third month; The rabbits never die; How many rabbits would there be in the 50th month?

Fibonacci number 1 1 2 3 5 8 13 21 34 55.....

OEIS: A000045

http://oeis.org/A000045





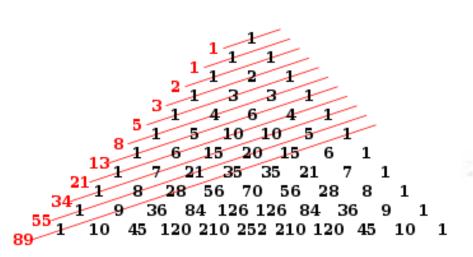
- Initial values: F(0)=0, F(1)=1• In 1150, Indian mathematicians researched the number of arrangements to package items with length 1 and width 2 into boxes. And they described this sequence for the first time.
- In the western world, Fibonacci mentioned a problem about the reproduction of rabbits in Liber Abbaci in 1202.
- Fibonacci, Leonardo 1175-1250
 - Member of the Bonacci family.
 - Travelled to Asia and Africa at 22 with his father and learned to calculate with Indian digits;
 - Played an important role in the recovery of Western Mathematics. And connected Western and Oriental mathematics.
 - G.Cardano: "We could assume that all mathematics we know except the Ancient Greek ones are gotten by Fibonacci.

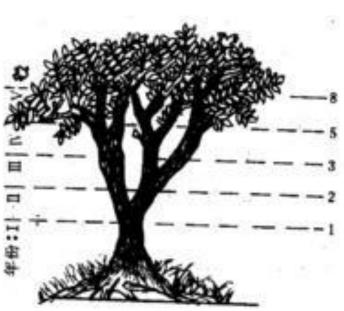


Leonardo of Pisa Fibonacci, Bonacci's son Bonacci: good, natural, simple

Fibonacci number 1 1 2 3 5 8 13 21 34 55.....









Trillium - 3 Petals



Bloodroot — 8 Petals



Devil's Paintbrush - 21 Petals



Sumflower — 55 Petals



St. Johnswort — 5 Petals



Black-eyed Susan — 13 Petals



Ox-eyed Daisy — 34 Petals



Daisy Fleabane — 89 Petals

Fibonacci Numbers

The Fibonacci Quarterly founded in 1963 especially publish the newest researches on this sequence. Which includes:

- -The last digit loops every 60 numbers; the last 2 digits loops every 300 numbers; the last 3 digits loops every 1500 numbers; the last 4 digits loops every 15000 numbers; the last 5 digits loops every 150000 numbers.
- -Every 3rd number could be divided by 2. Every 4th number could be divided by 3. Every 5th number could be divided by 5. Every 6th number could be divided by 8, etc. These divisors can also construct a Fibonacci Sequence.

Fibonacci prime (Sequence <u>A005478</u> in <u>OEIS</u>)

- -In the Fibonacci Sequence, there are primes: 2, 3, 5, 13, 89, 233, 1597, 28657, 514229, 433494437, 2971215073, 99194853094755497,
- -Except n = 4, the indexes of all Fibonacci Primes are primes.
- -However, not all prime index Fibonacci Numbers are primes.
- -Conjecture: Are there infinite primes among Fibonacci Numbers?
- The largest known prime is the 81839th Fibonacci Number, which has 17103 digits.



		1						
3		-2						
		1	1					
						3		
					C	•		
	5							

$$F_1^2 + F_2^2 + \dots + F_n^2 = F_n F_{n+1}$$

Area of the rectangle = Sum of multiple quadrates

Proof without words vs Logic deduction



Recurrence Relation

$$F_0 = 0, F_1 = 1, F_2 = 1 \dots$$

 $F_n = F_{n-1} + F_{n-2}$

Prove the identity: $F_1^2 + F_2^2 + \dots + F_n^2 = F_n F_{n+1}$

Proof: $F_1^2 = F_2 F_1$



Recurrence Relation

Proof:

$$F_0 = 0, F_1 = 1, F_2 = 1 \dots$$

 $F_n = F_{n-1} + F_{n-2}$

$$F_1 + F_2 + \dots + F_n = F_n + 2 - 1$$

 $F_1 = F_3 - F_2$
 $F_2 = F_4 - F_3$
.....

$$F_{n-1} = F_{n+1} - F_n$$
+) $F_n = F_{n+2} - F_{n+1}$

$$F_1 + F_2 + \dots + F_n = F_{n+2-1}$$



Recurrence Relation

$$F_0 = 0, F_1 = 1, F_2 = 1 \dots$$

 $F_n = F_{n-1} + F_{n-2}$

$$F_1 + F_3 + F_5 + \dots + F_{2n-1} = F_{2n}$$

Proof:

$$F_1 = F_2$$

$$F_3 = F_4 - F_2$$

$$F_5 = F_6 - F_4$$

Detailed Expressions?

+)
$$F_{2n-1} = F_{2n} - F_{2n-2}$$

$$\therefore F_1 + F_3 + F_5 + \dots + F_{2n-1} = F_{2n}$$



组合数学 Combinatorics

4 Linear Homogeneous Recurrence Relation

4-2 Expressions of Fibonacci Numbers





Magic

• There's a 80cm × 80cm quadrate tablecloth. How to convert it to a 1.3m × 50cm one?

0, 1, 1, 2, 3, 5, 8, 13, 21,....

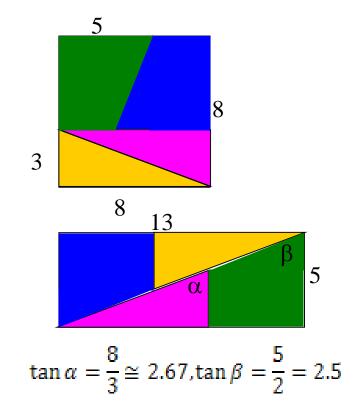
$$F(n)*F(n) - F(n-1)F(n+1) = (-1)^n$$

 $n=0,1,2$

Larger tablecloths?

$$F(100)=?$$

Direct expressions?





Fibonacci Recurrence

$$F_0 = 0$$
, $F_1 = 1$,
 $F_n = F_{n-1} + F_{n-2}$

Assume
$$G(x) = F_1 x + F_2 x^2 + \cdots$$

$$x^{3}: F_{3} = F_{2} + F_{1}$$

$$x^{4}: F_{4} = F_{3} + F_{2}$$
+)

$$G(x) - x^2 - x = x(G(x) - x) + x^2G(x)$$

$$\therefore (1-x-x^2)G(x)=x$$

$$\therefore G(x) = \frac{x}{1 - x - x^2} = \frac{x}{(1 - \frac{1 - \sqrt{5}}{2}x)(1 - \frac{1 + \sqrt{5}}{2}x)} = \frac{A}{1 - \frac{1 + \sqrt{5}}{2}x} + \frac{B}{1 - \frac{1 - \sqrt{5}}{2}x}$$



Fibonacci Recurrence

$$\begin{cases} A+B=0\\ \frac{\sqrt{5}}{2}(A-B)=1 \end{cases} \begin{cases} A+B=0\\ A-B=\frac{2}{\sqrt{5}} \end{cases} A = \frac{1}{\sqrt{5}}, \quad B=-\frac{1}{\sqrt{5}} \end{cases}$$

$$\therefore G(x) = \frac{1}{\sqrt{5}} \left[\frac{1}{1-\frac{1+\sqrt{5}}{2}x} - \frac{1}{1-\frac{1-\sqrt{5}}{2}x} \right] = \frac{1}{\sqrt{5}} \left[(\alpha-\beta)x + (\alpha^2-\beta^2)x^2 + \cdots \right]$$

$$\alpha = \frac{-2}{1-\sqrt{5}} = \frac{1+\sqrt{5}}{2}, \quad \beta = \frac{2}{1+\sqrt{5}} = \frac{1-\sqrt{5}}{2}$$

$$F_{n} = \frac{1}{\sqrt{5}}(\alpha^{n} - \beta^{n}) = \frac{1}{\sqrt{5}}((\frac{1+\sqrt{5}}{2})^{n} - (\frac{1-\sqrt{5}}{2})^{n})$$

$$\frac{F_n}{F_{n-1}} = \frac{1+\sqrt{5}}{2} \approx 1.618$$



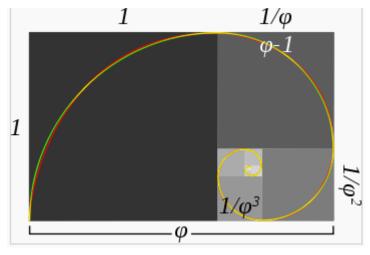
Fibonacci Sequence

$$\boldsymbol{F}_n = \boldsymbol{F}_{n-1} + \boldsymbol{F}_{n-2}$$

$$F_n = F_{n-1} + F_{n-2}$$
 $F_n = \frac{1}{\sqrt{5}} (\alpha^n - \beta^n) = \frac{1}{\sqrt{5}} ((\frac{1+\sqrt{5}}{2})^n - (\frac{1-\sqrt{5}}{2})^n)$

$$\frac{F_n}{F_{n-1}} = \frac{1+\sqrt{5}}{2} \approx 1.618$$

$$\varphi = [1; 1, 1, 1, \dots] = 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \cdots}}}$$



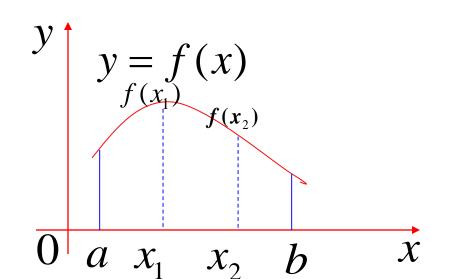


Applications in Optimization Methods

Assume that function f(x) reaches its maximum at $x = \xi$. Design an optimization algorithm to find the extreme point to a certain extent within finite iterations. The simplest way is to trisect the interval (a,b).

$$x_1 = a + \frac{1}{3}(b-a), \quad x_2 = a + \frac{2}{3}(b-a)$$

如下图:

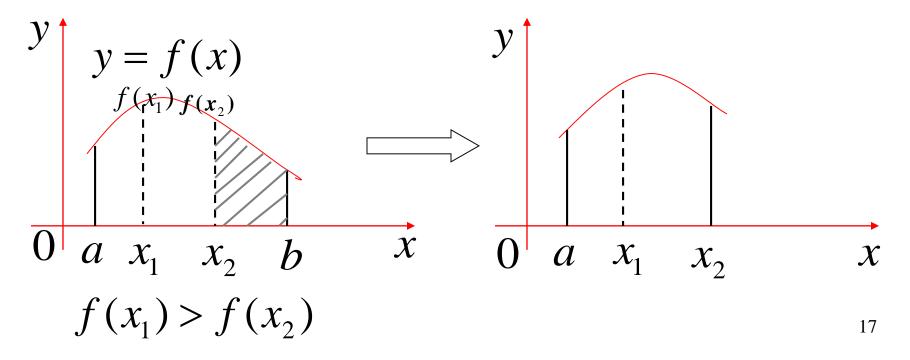




§ 2.4 Applications in Optimization Method

Discuss according to the sizes of f(1), f(2)

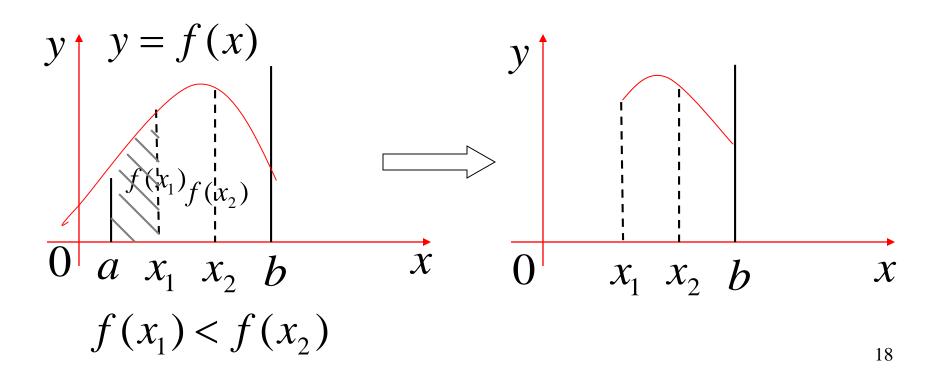
When $f(x_1) > f(x_2)$, the maximum ξ must be in (a, x_2) , interval (x_2, b) could be removed.





§ 2.4 Applications in Optimization Method

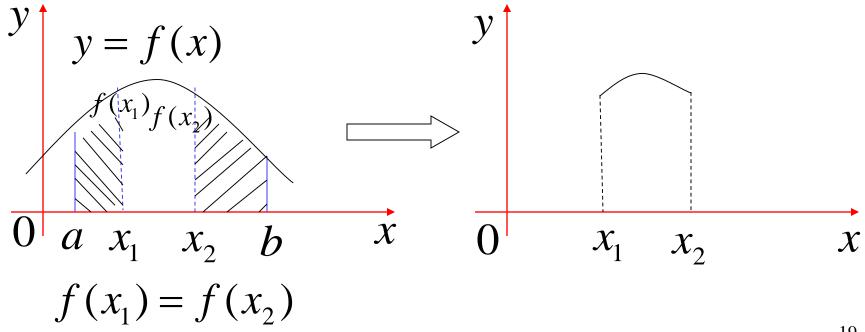
When $f(x_1) < f(x_2)$, the maximum ξ must be in (x_1, b) , the range (a, x_1) could be removed.





§ 2.4 **Application in Optimization Method**

When $f(x_1) = f(x_2)$, the maximum ξ must be in (x_1, x_2) , so both (a, x_1) and (x_2, b) could be removed.





y = f(x) $f(x_1) f(x_2)$

So with 2 tests, at least we could reduce the range to 2/3 of the origin domain. For example $f(x_1) > f(x_2)$ and we find the maximum in (a, x_2) .

If continue using threefold division method, it's a fact that x_1 is not used. So we image to have 2 symmetrical point x, l-x in (0,1) to test.





$0 \quad 1-x \quad x \quad 1$

If keep (0, x), then we keep testing at points x^2 , (1 - x)x in (0, x). If

$$x^2 = (1 - x)$$

Then the test at (1 - x) could be used again. So we save one test. We have:

$$x^{2} + x - 1 = 0$$

$$\therefore x = \frac{-1 + \sqrt{5}}{2} = 0.618$$

$$0.382, 0.618$$

$$0.236, 0.382$$

$$0.146, 0.236$$

$$0 \quad 0.382 = (0.618)^2 \quad 0.618 \quad 1$$



Application in Optimization Method

This is the 0.618 optimization method. When finding unimodal maximums in (0,1), we could test at:

$$x_1 = 0.618$$
, $x_2 = 1 - 0.618 = 0.3832$

points. For example keep (0,0.618), as

 $(0.618)^2 = 0.328$ e only need to test once at

$$0.618 \times 0.328 = 0.236$$



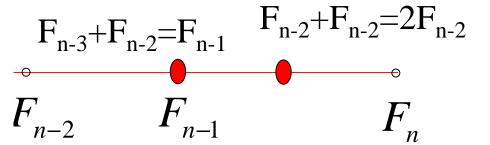
Applications in Optimization Method

We could use Fibonacci Sequence in Optimization Method. Its difference from 0.618 method is to decide the number of tests before testing. We introduce in 2 situations.

(a) The possible testing number is some Fn_{\circ}

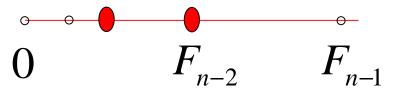


- •At this point testing points are division points F_{n-1} and F_{n-2} .
 - •If F_{n-1} is better, remove the part smaller than F_{n-2} ;
 - •The remained part contains F_n - $F_{n-2} = F_{n-1}$ division points,
 - •In which test pints F_{n-2} and F_{n-3} are correspondent to former index $F_{n-2}+F_{n-2}=2F_{n-2}$ and $F_{n-3}+F_{n-2}=F_{n-1}$. Just right, F_{n-1} has been tested in the previous round.



•If F_{n-2} is better, remove the part larger than F_{n-1} .

•In the remained part there are F_{n-1} division points, in the next step among test points F_{n-2} and F_{n-3} , F_{n-2} has been tested. \circ So among the F_n possible tests, we could find the extreme with at most n-1 tests.



One difference between the Fibonacci Method and 0.618 method is that Fibonacci Method could be used when the parameters are all integers. If the number of possible tests are smaller than F_n but larger than F_{n-1} , we could add several imagine points to make it F_n points. We could assume these image points to be worse than any other points without actually compare them.

Elliott wave

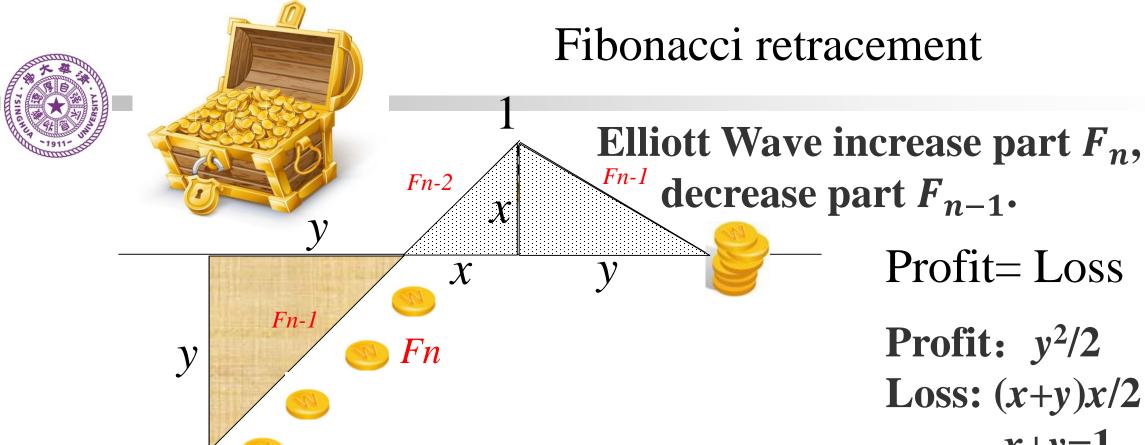


A complete loop includes 8 waves (5 increase. 3 decrease)

A complete period includes 8 waves, in which 5 are increasing, 3 are decreasing. They are all Fibonacci Numbers. In details we could get 34 waves and 144 waves, they are also Fibonacci Numbers.

Common retracement ratio are 0.382, 0.5 and 0.618. It mainly reflects the psychology of investors.





/ If always invest uniformly from 0 to 1, and then it decreases by x The sustentation x=?

Profit= Loss

Profit: $y^2/2$

Loss: (x+y)x/2

$$x+y=1$$

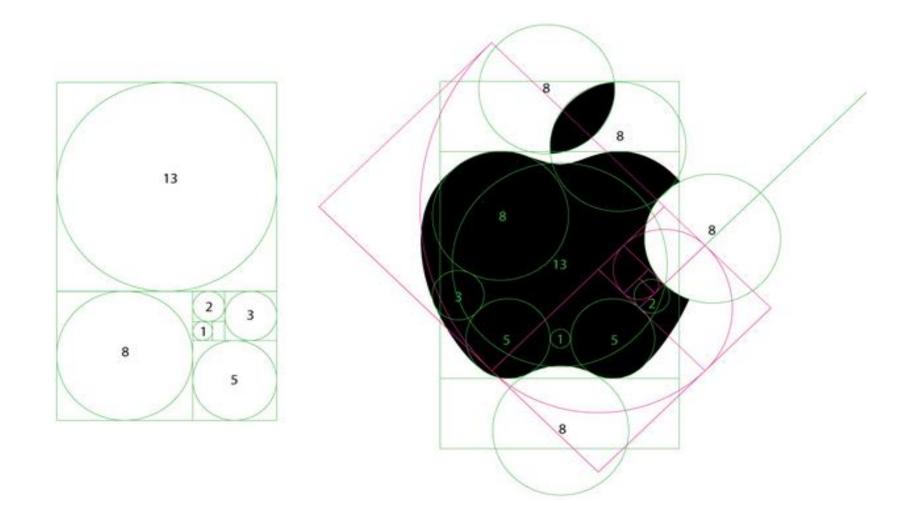
$$x = 0.382$$



Fibonacci retracement









组合数学 Combinatorics

4 Linear Homogeneous Recurrence Relation

4-3 Linear Homogeneous Recurrence Relation





$$F_{n}$$
- F_{n-1} - F_{n-2} =0

$$h(n) - 3h(n-1) + 2h(n-2) = 0$$

Summary

Linear summation

$$-RHS = 0$$

Coefficients are constants

Def If sequence $\{a_n\}$ satisfies:

$$a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0,$$

$$a_0 = d_0, a_1 = d_1, \dots, a_{k-1} = d_{k-1},$$

 C_1, C_2, \cdots, C_k and $d_0, d_1, \cdots, d_{k-1}$ are constants, $C_k \neq$

0, so this expression is called a k^{th} -order linear homogeneous recurrence relation of $\{a_n\}$.

$$h(n) = 2h(n-1) + 1, h(1) = 1$$

$$a_n = a_{n-1} + a_{n-2} a_{n-3}$$
 $a_{n-1} = a_{n-2} = a_{n-3} = 1$



Fibonacci Recurrence

$$F_n = F_{n-1} + F_{n-2}$$
 $F_0 = 0, F_1 = 1$

Assume
$$G(x) = F_1 x + F_2 x^2 + \cdots$$

$$\therefore (1-x-x^2)G(x)=x$$

$$\therefore G(x) = \frac{x}{1 - x - x^2} = \frac{x}{(1 - \frac{1 - \sqrt{5}}{2}x)(1 - \frac{1 + \sqrt{5}}{2}x)} = \frac{A}{1 - \frac{1 + \sqrt{5}}{2}x} + \frac{B}{1 - \frac{1 - \sqrt{5}}{2}x}$$

Factoring?

$$(1-ax)^{-1} = 1 + ax + a^2x^2 + \dots$$

$$(1-x-x^2) = (1-\frac{1-\sqrt{5}}{2}x)(1-\frac{1+\sqrt{5}}{2}x)$$

(**Factor Theorem**) If a is a root of linear polynomial f(x), which means f(a) = 0, then polynomial f(x) has a factor x - a.

We need factor (1 - ax),

If a is a root of linear polynomial $f(x^{-1})$, which means f(a) = 0, then polynomial $f(x^{-1})$ has a factor $x^{-1} - a = (1-ax)/x$.

$$(1-ax)^{-1} = 1 + ax + a^2x^2 + \dots$$

$$(1-x-x^2) = (1-\frac{1-\sqrt{5}}{2}x)(1-\frac{1+\sqrt{5}}{2}x)$$

$$(1-x-x^2) = x^2((x^{-1})^2-x^{-1}-1) = x^2((m)^2-m-1)$$

Let $m=x^{-1}$ $C(m) = m^2-m-1 = (m-\alpha)(m-\beta)$
Substitute $m=x^{-1}$ in, get $F_n = F_{n-1} + F_{n-2}$
 $F(x) = x^2(x^{-1}-\alpha)(x^{-1}-\beta) = (1-\alpha x)(1-\beta x)$

$$\alpha = \frac{-2}{1 - \sqrt{5}} = \frac{1 + \sqrt{5}}{2},$$
$$\beta = \frac{2}{1 + \sqrt{5}} = \frac{1 - \sqrt{5}}{2}$$

Linear Homogeneous Recurrence Relation

• The recurrence expression of Fibonacci Sequence

$$F_0 = 0$$
, $F_1 = 1$, $G(x) = \frac{x}{1 - x - x^2} = \frac{1 - \alpha x}{(1 - \alpha x)(1 - \beta x)}$
 $F_n = F_{n-1} + F_{n-2}$

Denominator becomes $F(x) = x^2((x^{-1})^2 - x^{-1} - 1) = x^2((m)^2 - m - 1)$ Let $m = x^{-1}$

$$C(m) = m^{2} - m - 1 = (m - \alpha)(m - \beta)$$

$$\alpha = \frac{1 + \sqrt{5}}{2}, \quad \beta = \frac{1 - \sqrt{5}}{2}$$
Substitute $m = x^{-1}in$, get

 $F(x) = x^2(x^{-1} - \alpha)(x^{-1} - \beta) = (1 - \alpha x)(1 - \beta x)$

$$G(x) = \frac{x}{(1 - \frac{1 - \sqrt{5}}{2}x)(1 - \frac{1 + \sqrt{5}}{2}x)} = \frac{A}{1 - \frac{1 + \sqrt{5}}{2}x} + \frac{B}{1 - \frac{1 - \sqrt{5}}{2}x}$$

• Recurrence expression of Hanoi Tower

$$h(n)-2h(n-1)=1$$
 $H(x) = \frac{x}{(1-x)(1-2x)} = \frac{x}{1-3x+2x^2}$
 $h(n-1)-2h(n-2) = \text{Subtract and get}$
 $C(x) = x^2-3x+2$
 $Roots\ of\ C(x)=0\ are\ 1\ and\ 2$

$$a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0,$$



Linear Homegeneous Recurrence Relation

Assume G(x) is a generating function of $\{a_n\}$:

$$G(x) = a_0 + a_1 x + \dots + a_n x^n + \dots$$

$$x^k (a_k + C_1 a_{k-1} + C_2 a_{k-2} + \dots + C_k a_0) = 0$$

$$x^{k+1} (a_{k+1} + C_1 a_k + C_2 a_{k-1} + \dots + C_k a_1) = 0$$

$$\dots$$

$$x^n (a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k}) = 0$$

Adds up both sides of these equations, get

$$G(x) - \sum_{i=0}^{k-1} a_i x^i + C_1 x \left(G(x) - \sum_{i=0}^{k-2} a_i x^i \right) + \dots + C_k x^k G(x) = 0$$

Linear Homegeneous Recurrence Realtion

$$(1 + C_1 x + C_2 x^2 + \dots + C_k x^k)G(x) = \sum_{j=0}^{k-1} C_j x^j \sum_{i=0}^{k-1-j} a_i x^i$$

Let $P(x) = \sum_{i=0}^{k-1} C_i x^j \sum_{i=0}^{k-1-j} a_i x^i$, the order of polynomial $P(x) \le k-1$.

$$G(x) = \frac{P(x)}{(1 + C_1 x + \dots + C_k x^k)}$$

$$(1 - ax)^{-1} = 1 + ax + a^2 x^2 + \dots$$

$$a_{n} + C_{1}a_{n-1} + C_{2}a_{n-2} + \dots + C_{k}a_{n-k} = 0,$$

$$C(m) = (m - a_{1})^{k_{1}}(m - a_{2})^{k_{2}} \dots (m - a_{i})^{k_{i}}$$

$$k_{1} + k_{2} + \dots + k_{i} = k$$

$$m = x^{-1}$$

$$P(x)$$

$$= \frac{P(x)}{(1-a_1x)^{k_1}(1-a_2x)^{k_2}\cdots(1-a_ix)^{k_i}}$$

Linear Homogeneous Recurrence Relation

$$F_n$$
- F_{n-1} - F_{n-2} =0 $h(n) - 3h(n-1) + 2h(n-2)$ =0 x^2 - x - 1 = 0 x^2 - $3x$ + 2 = 0

Def if sequence $\{a_n\}$ satisfies:

$$a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0,$$

$$a_0 = d_0, a_1 = d_1, \dots, a_{k-1} = d_{k-1},$$

 $C_1, C_2, \dots C_k$ and $d_0, d_1, \dots d_{n-1}$ are constants, $C_k \neq 0$, then this expression is called a kth-order linear homogeneous recurrence relation of $\{a_n\}$.

$$C(x) = x^{k} + C_{1}x^{k-1} + \dots + C_{k-1}x + C_{k}$$

Characteristic Polynomial

$$F_n - F_{n-1} - F_{n-2} = 0$$
 $h(n) - 3h(n-1) + 2h(n-2) = 0$

Def if sequence
$$\{a_n\}$$
 satisfies: $a_n + C_1 a_{n-1} + C_2 a_{n-2} + \cdots + C_k a_{n-k} = 0$,

$$a_0 = d_0, a_1 = d_1, \dots, a_{k-1} = d_{k-1},$$

 $C_1, C_2, \dots C_k$ and $d_0, d_1, \dots d_{n-1}$ are constants, $C_k \neq 0$, then this expression is called a kth-order linear homogeneous recurrence relation of $\{a_n\}$.

$$C(x) = x^{k} + C_{1}x^{k-1} + \dots + C_{k-1}x + C_{k}$$

Characteristic Polynomial



Linear Homegeneous Recurrence Relation

Now we discuss the calculation by situations

(1) Characteristic Polynomial C(x) has distinct real roots

Assume
$$C(x) = (x - \alpha_1)(x - \alpha_2) \cdots (x - \alpha_k)$$

G(x) could be simplified as l_1, l_2, \dots, l_k could be solved by

$$I_{1} + I_{2} + \dots + I_{k} = d_{0}$$

$$I_{1}\alpha_{1} + I_{2}\alpha_{2} + \dots + I_{k}\alpha_{k} = d_{1}$$

$$\dots$$

$$k-1 + I_{k}\alpha^{k-1} + I_{k}\alpha^{k-1} - d_{k}$$

$$(1-ax)^{-1} = 1 + ax + a^2x^2 + \dots$$

in which $d_0, d_1 \dots d_{k-1}$ are initial values of a_n



Linear Homegeneous Recurrence Relation

• Fibonacci Sequence Linear Homogeneous Recurrence Relation

$$F_0 = 0$$
, $F_1 = 1$, $F_n = F_{n-1} + F_{n-2}$

Def If sequence $\{a_n\}$ satisfies:

$$a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0,$$

 $a_0 = d_0, a_1 = d_1, \dots, a_{k-1} = d_{k-1},$

Characteristic Polynomial

$$C(x) = x^{2}-x-1 = (x-\alpha)(x-\beta)$$

$$\alpha = \frac{1+\sqrt{5}}{2}; \beta = \frac{1-\sqrt{5}}{2}$$

$$F_{n} = A\alpha^{n} + B\beta^{n}$$

$$F_0 = 0$$
, $F_1 = 1$

$$\begin{cases} A+B=0\\ \frac{\sqrt{5}}{2}(A-B)=1 \end{cases}$$

 $C_1, C_2, \dots C_k$ and $d_1, d_2, \dots d_{n-1}$ are constants.

Characteristic Polynomial

$$C(x) = x^{k} + C_{1}x^{k-1} + \dots + C_{k-1}x + C_{k}$$

1) Characteristic polynomial has distinct real roots, k different real roots

$$C(x) = (x - a_1)(x - a_2) \cdots (x - a_k)$$

$$a_n = l_1 a_1^n + l_2 a_2^n + \cdots + l_k a_k^n$$

$$F_{n} = \frac{1}{\sqrt{5}}(\alpha^{n} - \beta^{n}) = \frac{1}{\sqrt{5}}((\frac{1+\sqrt{5}}{2})^{n} - (\frac{1-\sqrt{5}}{2})^{n})$$

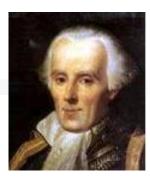
Generating Function



Def 2-1 For sequence a_0 , a_1 , a_2 ..., construct function

$$G(x) = a_0 + a_1 x + a_2 x^2 + \dots,$$

G(x) is the generating function of $a_0, a_1, a_2...$



Laplae 1812 AD



Berelli 1705- AD

 C_0, C_1, \ldots, C_n

Seems to be functions but it's actually a mapping



Euler 1764- AD

Recurrence $a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0$,

Coefficients

$$C(x) = x^{k} + C_{1}x^{k-1} + \dots + C_{k-1}x + C_{k}$$

Integer Segmentation

Generating function G(x) as a bridge $a_n = l_1 a_1^n + l_2 a_2^n + \cdots + l_k a_k^n$

 $= x^2 + 2x^3 + 3x^4 + 4x^5 + 5x^6 + \dots$



Linear Homogeneous Recurrence Relation

Def If sequence $\{a_n\}$ satisfies:

$$a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0, \quad (2-5-1)$$

 $a_0 = d_0, a_1 = d_1, \dots, a_{k-1} = d_{k-1}, \quad (2-5-2)$

 $C_1, C_2, \cdots C_k$ and $d_0, d_1, \cdots d_{k-1}$ are constants

Characteristic Polynomial $C(x) = x^k + C_1 x^{k-1} + \cdots + C_{k-1} x + C_k$

1) Characteristic polynomial has k distinct real roots

$$C(x) = (x - a_1)(x - a_2) \cdots (x - a_k)$$

$$a_n = l_1 a_1^n + l_2 a_2^n + \cdots + l_k a_k^n$$

In which $l_1, l_2, \dots l_k$ are undetermined coefficients.

$$(1+x)^{\alpha} = \sum_{n=0}^{\infty} \frac{\alpha(\alpha-1)...(\alpha-n+1)}{n!} x^n \qquad \alpha \in R \qquad \sum_{k=0}^{\infty} (k+1)x^k = \frac{1}{(1-x)^2}$$

Characteristic Polynomial has multiple roots

• Eg
$$a_n - 4a_{n-1} + 4a_{n-2} = 0$$
, $a_0 = 1, a_1 = 4$

$$x^{2}: a(2) = 4a(1) - 4a(0)$$

$$x^{3}: a(3) = 4a(2) - 4a(1)$$

$$+) \qquad \cdots \cdots$$

$$A(x) = \frac{1}{1 - 4x + 4x^{2}}$$

$$A(x) = \frac{1}{1 - 4x + 4x^{2}} = \frac{1}{(1 - 2x)^{2}}$$

$$= (1 - 2x)^{-2} = \sum_{k=0}^{\infty} C(k+1,k)2^{k} x^{k}$$

$$= \sum_{k=0}^{\infty} (k+1)2^{k} x^{k}$$

$$a_{n} = (n+1)2^{n}$$

Generating Function Method Characteristic Equation Method

Characteristic Equation: $x^2-4x+4=(x-2)^2$

Generating Function Form: $A(x) = \frac{ax + b}{(1 - 2x)^2}$

Partial Fractions:
$$A(x) = \frac{A}{(1-2x)} + \frac{B}{(1-2x)^2}$$

$$a_n = A \times 2^n + B(n+1)2^n = (A' + Bn)2^n$$

$$a_0 = A' = 1$$
, $a_1 = (1+B)2 = 4$

$$A' = 1, B = 1$$

$$a_n = (n+1)2^n$$

$$G(x) = \sum_{i=1}^{t} \sum_{j=1}^{k_i} \frac{A_{ij}}{(1 - \alpha_i x)^j}$$

$$(1+x)^{\alpha} = \sum_{n=0}^{\infty} \frac{\alpha(\alpha-1)...(\alpha-n+1)}{n!} x^{n} \qquad \alpha \in R$$

(2) Characteristic Polynomial C(x) has multiple roots
Assume β is a k-multiple root of C(x), it could be simplified as $\sum_{i=1}^{k} \frac{A_j}{(1-\beta x)^j}$

$$x^n$$
's coefficients $a_n = \sum_{j=1}^k A_j {j+n-1 \choose n} \beta^n$, in which
$${j+n-1 \choose n} = {j+n-1 \choose j-1}$$

is a j-l-order polynomial of n. So a_n is the product of β and a k-l-order polynomial of n. The term related to the solution of recurrence relation is: $(A_0 + A_1 n + \cdots + A_{k-1} n^{k-1})\beta^n$

in which A_0, A_1, \dots, A_{k-1} are k undetermined coefficients.

• Eg
$$a_n - 4a_{n-1} + 4a_{n-2} = 0$$
, $a_0 = 1, a_1 = 4$

Characteristic Equation is: $x^2 - 4x + 4 = (x - 2)^2$

$$\boldsymbol{a}_n = (\boldsymbol{A}_1 + \boldsymbol{A}_2 \boldsymbol{n}) 2^n$$

$$\boldsymbol{a}_0 = \boldsymbol{A}_1 = 1$$

$$a_1 = (1 + A_2) 2 = 4, \qquad A_2 = 1$$

$$a_n = (1+n)2^n$$

Distinct real roots Multiple real roots Conjugate complex roots?

$$x^2 - x + 1 = 0$$

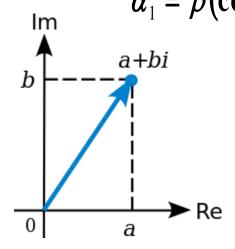
Conjugate Complex Roots

• Quadratic Formula:
$$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

• When b^2 -4ac<0, there's no real root, two complex roots.

$$x_{1,2} = \frac{-b \pm i \times \sqrt{4ac - b^2}}{2a}$$

$$\alpha_1 = \rho(\cos\theta + i\sin\theta), \alpha_2 = \overline{\alpha_1} = \rho(\cos\theta - i\sin\theta)$$



Trigonometrical form of complex number z = a + bi: $z = \rho(\cos\theta + i \sin\theta)$

$$\rho = \sqrt{a^2 + b^2}$$



§ 2.5 Linear Homogeneous Recurrence Relation

(3) Characteristic Polynomial C(x) has conjugate complex roots Assume that a_1 , a_2 are a pair of conjugate complex roots of C(x).

$$\alpha_1 = \rho(\cos\theta + i\sin\theta), \alpha_2 = \overline{\alpha_1} = \rho(\cos\theta - i\sin\theta)$$

In
$$\frac{A_1}{1-\alpha_1 x} + \frac{A_2}{1-\alpha_2 x}$$
 the coefficient of x^n is:

$$A_1\alpha_1^n + A_2\alpha_2^n$$

$$A_1 \alpha_1^n + A_2 \alpha_2^n$$

$$= (A_1 + A_2) \rho^n \cos n\theta + i(A_1 - A_2) \rho^n \sin n\theta$$

$$= A \rho^n \cos n\theta + B \rho^n \sin n\theta$$

In which
$$A = A_1 + A_2$$
, $B = (i)(A_1 - A_2)$

When calculating in reality, we could solve the conjugate complex roots at first, then calculate undetermined coefficients *A*, *B* to avoid the intermediate complex number calculations.



$\underline{A_1\alpha_1^n + A_2\alpha_2^n} = A\rho^n \cos n\theta + B\rho^n \sin n\theta$

• Eg
$$a_n = a_{n-1} - a_{n-2}, a_1 = 1, a_2 = 0$$

Characteristic equation: $x^2 - x + 1 = 0$

$$x = \frac{1 \pm \sqrt{-3}}{2} = \cos\frac{\pi}{3} \pm i \sin\frac{\pi}{3} = e^{\pm\frac{\pi}{3}i}$$

$$a_n = A_1 \cos\frac{n\pi}{3} + A_2 \sin\frac{n\pi}{3}$$

$$a_{1} = \frac{1}{2}A_{1} + \frac{\sqrt{3}}{2}A_{2} = 1$$

$$a_{2} = -\frac{1}{2}A_{1} + \frac{\sqrt{3}}{2}A_{2} = 0$$

$$A_{1} = 1; A_{2} = \frac{\sqrt{3}}{3}$$

$$a_n = \cos\frac{n\pi}{3} + \frac{\sqrt{3}}{3}\sin\frac{n\pi}{3}$$



Summary of Linear Recurrence Relation

According to the non-zero roots of C(x)

1) k distinct non-0 real roots $C(x) = (x - a_1)(x - a_2) \cdots (x - a_k)$

$$a_n = l_1 a_1^n + l_2 a_2^n + \dots + l_k a_k^n$$

In which l_1, l_2, \dots, l_k , are undetermined coefficients.

2) A pair of conjugate complex root $\alpha_1 = \rho e^{i\theta}$ and $\alpha_2 = \rho e^{-i\theta}$:

$$a_n = A\rho^n \cos n\theta + B\rho^n \sin n\theta$$

In which A, B are undetermined coefficients.

3) Has root α_1 with multiplicity of k.

$$(A_0 + A_1 n + \cdots + A_{k-1} n^{k-1})\alpha_1^n$$

In which A_0, A_1, \dots, A_{k-1} are k undetermined coefficients.



组合数学 Combinatorics

4 Linear Homogeneous Recurrence Relation

4-4 Applications





Linear Homogeneous Recurrence Relation

Eg: Solve
$$S_n = \sum_{k=0}^{n} k$$

 $S_n = 1 + 2 + 3 + \dots + (n-1) + n$
 $S_{n-1} = 1 + 2 + 3 + \dots + (n-1)$
 $\therefore S_n - S_{n-1} = n$
Similarly $S_{n-1} - S_{n-2} = n - 1$
Subtract, get $S_n - 2S_{n-1} + S_{n-2} = 1$
Similarly $S_{n-1} - 2S_{n-2} + S_{n-3} = 1$
 $\therefore S_n - 3S_{n-1} + 3S_{n-2} - S_{n-3} = 0$
 $S_0 = 0$, $S_1 = 1$, $S_2 = 3$

$$S_n - 3S_{n-1} + 3S_{n-2} - S_{n-3} = 0$$
$$S_0 = 0, \quad S_1 = 1, \quad S_2 = 3$$

Corresponding Characteristic Equation is

$$m^3 - 3m^2 + 3m - 1 = (m - 1)^3 = 0$$

 $m = 1$ is a 3-multiple root

$$S_{n} = (A + Bn + Cn^{2})(1)^{n} = A + Bn + Cn^{2}$$

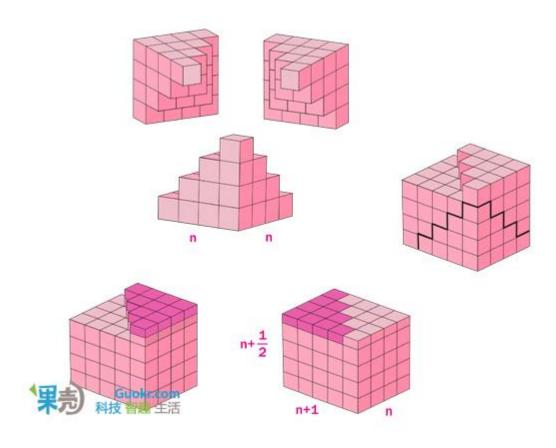
$$S_{0} = 0, \quad \therefore A = 0$$

$$S_{1} = 1, \quad B + C = 1$$

$$S_{2} = 3, \quad 2B + 4C = 3, \quad \therefore B = C = \frac{1}{2}$$
So
$$S_{n} = \frac{1}{2}n + \frac{1}{2}n^{2} = \frac{1}{2}n(n+1)$$
This proves
$$1 + 2 + 3 + \dots + n = \frac{1}{2}n(n+1)$$



$$1^{2}+2^{2}+3^{2}+\cdots+n^{2}=\frac{1}{3}n (n+\frac{1}{2}) (n+1)$$





Linear Homogeneous Recurrence Relation

Eg2: Calculate
$$S_n = \sum_{k=0}^n k^2$$

 $S_n = 1 + 2^2 + 3^2 + \dots + (n-1)^2 + n^2$ \therefore $S_n - S_{n-1} = n^2$
 $S_{n-1} = 1 + 2^2 + 3^2 + \dots + (n-1)^2$ Similarly $S_{n-1} - S_{n-2} = (n-1)^2$
Subtract, get $S_n - 2S_{n-1} + S_{n-2} = 2n - 1$
Similarly $S_{n-1} - 2S_{n-2} + S_{n-3} = 2(n-1) - 1$
Subtract, get $S_n - 3S_{n-1} + 3S_{n-2} - S_{n-3} = 2$
Similarly $S_{n-1} - 3S_{n-2} + 3S_{n-3} - S_{n-4} = 2$
 \therefore $S_n - 4S_{n-1} + 6S_{n-2} - 4S_{n-3} + S_{n-4} = 0$
 $S_0 = 0$, $S_1 = 1$, $S_2 = 5$, $S_3 = 14$



$$S_n - 4S_{n-1} + 6S_{n-2} - 4S_{n-3} + S_{n-4} = 0$$

$$S_0 = 0, \quad S_1 = 1, \quad S_2 = 5, \quad S_3 = 14$$

Correspondent characteristic equation is:

$$r^4 - 4r^3 + 6r^2 - 4r + 1 = (r - 1)^4 = 0$$

 $r = 1$ is a 4-multiple root
 $\therefore S_n = (A + Bn + Cn^2 + Dn^3)(1)^n$
As $S_0 = 0$, $S_1 = 1$, $S_2 = 5$, $S_3 = 14$ we have a equation group about A, B, C, D:
 $A = 0$
 $B + C + D = 1$
 $2B + 4C + 8D = 5$
 $3B + 9C + 27D = 14$

$$1^{2}+2^{2}+3^{2}+\cdots+n^{2}=\frac{1}{3}n (n+\frac{1}{2}) (n+1) \qquad D_{n}=\begin{vmatrix} 1 & 1 & \cdots & 1 \\ x_{1} & x_{2} & \cdots & x_{n} \\ x_{1}^{2} & x_{2}^{2} & \cdots & x_{n}^{2} \\ \vdots & \vdots & & \vdots \end{vmatrix} = \prod_{n\geq i>j\geq 1}(x_{i}-x_{j}).$$

$$\begin{cases} A = 0 \\ B + C + D = 1 \\ 2B + 4C + 8D = 5 \\ 3B + 9C + 27D = 14 \end{cases} \begin{vmatrix} 1 & 1 & 1 \\ 2 & 4 & 8 \\ 3 & 9 & 27 \end{vmatrix} = \frac{1}{6}$$

$$B = \frac{1}{12} \begin{vmatrix} 1 & 1 & 1 \\ 5 & 4 & 8 \\ 14 & 9 & 27 \end{vmatrix} = \frac{1}{6}$$

$$\begin{cases} A = 0 \\ B + C + D = 1 \\ 2B + 4C + 8D = 5 \end{cases} \begin{vmatrix} 1 & 1 & 1 \\ 2 & 4 & 8 \\ 3 & 9 & 27 \end{vmatrix} = \begin{vmatrix} x_1^n & x_2^n & \cdots & x_n^n & 1 \\ = (4-3)(4-2)(4-1)(3-2)(3-1)(2-1) \\ = 12 \\ = 12 \end{vmatrix}$$

$$\mathbf{B} = \frac{1}{12} \begin{vmatrix} 1 & 1 & 1 \\ 5 & 4 & 8 \\ 14 & 9 & 27 \end{vmatrix} = \frac{1}{6}$$

$$C = \frac{1}{12} \begin{vmatrix} 1 & 1 & 1 \\ 2 & 5 & 8 \\ 3 & 14 & 27 \end{vmatrix} = \frac{1}{2}$$

$$D = \frac{1}{12} \begin{vmatrix} 1 & 1 & 1 \\ 2 & 4 & 5 \\ 3 & 9 & 14 \end{vmatrix} = \frac{1}{3}$$

Applications of generating function and recurrence relation

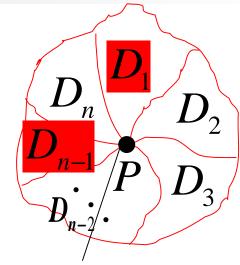
Eg: There's a point P on the plane. It's the cross of n fields $D_1, D_2, \dots D_n$. Color these n fields with k colors. We require the color of two adjacent areas to be different.

Calculate the number of arrangements.

Let a_n be the number of arrangement to color these areas. There are 2 situations:

Applications of generating function and recurrence relation

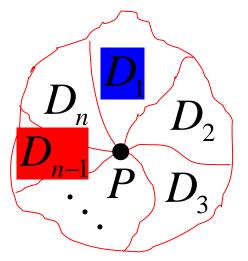
(1) D_1 and D_{n-1} have the same color; D_n has k-1 choices, which is all colors except the one used by D_1 and D_{n-1} ; the arrangements for D_{n-2} to D_1 的 are one-to-one correspondent to the arrangements for n-2 areas. $(k-1)a_{n-2}$



(2) D_1 and D_{n-1} have different colors. D_n has k-2 choices; the arrangements from D_1 to D_{n-1} are one-to-one correspondent to the arrangements for n-1 areas. $(k-2)a_{n-1}$

$$\therefore a_n = (k-2)a_{n-1} + (k-1)a_{n-2},$$

$$a_2 = k(k-1), \ a_3 = k(k-1)(k-2).$$



Applications of generating function and recurrence relation

 $\therefore a_n = (k-2)a_{n-1} + (k-1)a_{n-2},$

$$a_{2} = k(k-1), \ a_{3} = k(k-1)(k-2).$$

$$a_{1} = 0, a_{0} = k.$$

$$x^{2} - (k-2)x - (k-1) = 0,$$

$$x_{1} = k - 1, \quad x_{2} = -1.$$

$$a_{n} = A(k-1)^{n} + B(-1)^{n}$$

$$\begin{cases} A = 1, \\ B = k - 1. \end{cases}$$

$$\vdots \quad a_{n} = (k-1)^{n} + (k-1)(-1)^{n}, \ n \ge 2.$$

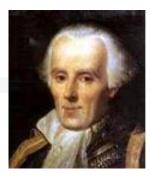
$$\vdots \quad a_{1} = k.$$



Def 2-1 For sequence $a_0, a_1, a_2...$, construct a function

$$G(x) = a_0 + a_1 x + a_2 x^2 + \dots,$$

Then G(x) is called the generating function of a_0, a_1, a_2



Laplace 1812 AD

Generating functions are a hanger to hang a serires of numbers.

