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**Week 3 Quiz: Heuristic Search**

Quiz due Apr 11, 2017 05:00 IST

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## Week 3 Quiz: Heuristic Search

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### Optimality #1

10/10 points (graded)

Consider graph search algorithms for some search space. Suppose the branching factor  $b$  is finite, the shallowest goal is at finite depth  $d$ , and step costs are finite, greater than some small positive constant, but not necessarily all equal. Check all that apply:

☐ Depth-First Search is optimal

☐ Depth-Limited Search (limit  $> d$ ) is optimal

☐ Iterative-Deepening Search is optimal

☐ Breadth-First Search is optimal

☒ Uniform-Cost Search is optimal ✓



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You have used 4 of 4 attempts

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### Optimality #2

10.0/10.0 points (graded)

Consider graph search algorithms for some search space. Suppose the branching factor  $b$  is finite, the shallowest goal is at finite depth  $d$ , and step costs are finite, positive, and all identical. Check all that apply:

☐ Depth-First Search is optimal

☐ Depth-Limited Search (limit  $> d$ ) is optimal

- ▶ Week 5: Machine Learning 1

☒ Iterative-Deepening Search is optimal ✓

☒ Breadth-First Search is optimal ✓

- ▶ Week 6: Machine Learning 2

☒ Uniform-Cost Search is optimal ✓



- ▶ Week 7: Machine Learning 3

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- ▶ Week 8: CSP

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- ▶ Week 9: Reinforcement Learning

## Admissibility #1

10/10 points (graded)

Consider a finite search space. Suppose step costs are finite and greater than some small positive constant, but not necessarily all equal. Suppose  $h(n)$  is an admissible heuristic. Check all that apply:

- ▶ Week 10: Logical Agents

☒  $h(n)$  never overestimates the true cost  $h^*(n)$  from  $n$  to the goal ✓

☐  $f(n)$  never overestimates the true cost  $h^*(n)$  from  $n$  to the goal

☒  $f(n)$  never overestimates the true cost  $g(n) + h^*(n)$  from the root to the goal through  $n$  ✓

☐  $h(n)$  is "optimistic" because it acts as an upper bound on the true cost of reaching the goal

☐  $h(n)$  is "optimistic" because it always underestimates the true cost of reaching the goal

☐  $h(n)$  is "optimistic" because it is greedy



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## Admissibility #2

10.0/10.0 points (graded)

In lecture, we see an example of a heuristic for the map problem; that is, the straight line distances  $h\_SLD(n)$  from  $n$  to the goal. Check all that apply to this instance of the search problem (in particular, note that no edges are less than unit cost):

☒  $h\_SLD(n)$  is admissible ✓

☐  $h\_SLD(n)^2$  is admissible

☒  $\sqrt{h\_SLD(n)}$  is admissible ✓

☐  $h\_SLD(n)^2 - 99 * h\_SLD(n)$  is admissible

☒ 0 is admissible ✓



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## Completeness #1

10/10 points (graded)

Consider graph search algorithms for an infinite\* search space. Suppose the branching factor  $b$  is finite, the shallowest goal is at finite depth  $d$ , step costs are finite, greater than some small positive constant, but not necessarily all equal. Check all that apply:

\* As an example of an infinite search space, consider the 3-dimensional integer lattice  $\mathbb{Z}^3$ ; that is, the lattice in the Euclidean space  $\mathbb{R}^3$  whose lattice points are ordered triples of integers. A "state" may consist of a point in space (i.e. an ordered triple in  $\mathbb{Z}^3$ ), and a "transition" may consist of moving a certain number of units in the positive or negative directions parallel to one of the coordinate axes.

☐ Depth-First Search is complete

☒ Depth-Limited Search (limit  $> d$ ) is complete ✓

☒ Iterative-Deepening Search is complete ✓☒ Breadth-First Search is complete ✓☒ Uniform-Cost Search is complete ✓☒ A-Star Search is complete ✓

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## Completeness #2

10/10 points (graded)

Consider graph search algorithms for a finite search space. Suppose the branching factor  $b$  is finite, the shallowest goal is at finite depth  $d$ , step costs are finite, greater than some small positive constant, but not necessarily all equal. Check all that apply:

☒ Depth-First Search is complete ✓☒ Depth-Limited Search (limit  $> d$ ) is complete ✓☒ Iterative-Deepening Search is complete ✓☒ Breadth-First Search is complete ✓☒ Uniform-Cost Search is complete ✓☒ Greedy Best-First Search is complete ✓☒ A-Star Search is complete ✓

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You have used 3 of 4 attempts

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## Greedy Best-First Search #1

10.0/10.0 points (graded)

In lecture, we see an example of a heuristic for the map problem; that is, the straight line distances  $h\_SLD(n)$  from  $n$  to the goal. Consider Greedy Best-First Search applied to this instance of the search problem, using the straight-line distance heuristic (in particular, note that no edges are less than unit cost). Check all that apply:

☐ It always manages to reach the goal in the fewest number of steps

☐ It always manages to reach the goal through the least costly path

☒ At each step it tries to get as close to the goal as possible ✓

☐ At each step it always gets closer to the goal



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## Greedy Best-First Search #2

10.0/10.0 points (graded)

Consider Greedy Best-First Search on a search space where the branching factor  $b$  is finite, the shallowest goal is at finite depth  $d$ , step costs are finite, greater than some small positive constant, but not necessarily all equal. Check all that apply:

☒ Evaluates nodes by using just the heuristic function ✓

☐ Evaluates nodes by using the heuristic function plus path costs

☐ Is complete if the heuristic is admissible

☒ Is complete if the search space is finite ✓

☐ Is optimal if the heuristic is admissible☐ Is optimal if the search space is finite

You have used 3 of 4 attempts

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**i** Answers are displayed within the problem

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## A-Star Search #1

10.0/10.0 points (graded)

In lecture, we see an example of a heuristic for the map problem; that is, the straight line distances  $h\_SLD(n)$  from  $n$  to the goal. Consider A-Star Search applied to this instance of the search problem, using the straight-line distance heuristic (in particular, note that no edges are less than unit cost). Check all that apply:

☐ It always manages to reach the goal in the fewest number of steps☒ It always manages to reach the goal through the least costly path ☐ At each step it tries to get as close to the goal as possible☐ At each step it always gets closer to the goal

You have used 3 of 4 attempts

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## A-Star Search #2

10.0/10.0 points (graded)

Consider A-Star Search on a tree of finite depth, where the branching factor  $b$  is finite, the shallowest goal is at finite depth  $d$ , step costs are finite, greater than some small positive constant, but not necessarily all equal, and the tree contains no duplicate nodes. Check all that apply. It:

- ☐ Evaluates nodes by using just the heuristic function
- ☒ Evaluates nodes by using the heuristic function plus path costs ✓
- ☒ Is optimal if it has an admissible heuristic ✓
- ☒ Becomes Uniform-Cost Search if the heuristic  $h(n) = 0$  for all  $n$  ✓
- ☐ Becomes Uniform-Cost Search if the path cost  $g(n) = 0$  for all  $n$



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