## CS 6901 Capstone Exam Systems Fall 2013: Choose any 2 of the 3 problems.

- 1) Construct a combinational circuit that has 2 sets of 4-bit integer inputs  $(a_3a_2a_1a_0)$  and  $b_3b_2b_1b_0$ , a control input (x), and a 5-bit integer output  $(c_4c_3c_2c_1c_0)$ . If x=1, the output c is to be the difference of the integers a-b, while if x=0, the output c is to be the sum of the integers a+b. You may use full adders and up to 4 additional gates. Draw the circuit diagram.
- 2) Consider the semaphore solution to solve the producer/consumer problem with a buffer of n elements. Write the basic code for producers and consumers. Declare and initialize all semaphores.
- 3) Consider the following page replacement algorithms: FIFO (first in first out), LRU (least recently used), OPT (optimal replacement), and 2<sup>nd</sup> chance. Logical memory has 10 pages (pages 0 .. 9), while physical memory consists of 4 frames (frames 0 .. 3). The page reference string begins with 5, 3, 8, 4 to fill the four frames. Each part begins from this same initial point. On your solution page, show the 2 frame traces for each part. For 2<sup>nd</sup> chance, also show the reference bit values. Each reference bit value is indicated by a 1 or 0 in parentheses.
- a) Continue the page reference string with at most 4 additional terms where LRU will result in strictly fewer page faults than FIFO.

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b) Continue the page reference string with at most 3 additional terms where OPT will result in strictly fewer page faults than LRU.

c) Continue the page reference string with at most 5 additional terms where  $2^{nd}$  chance will result in strictly fewer page faults than FIFO.

ref. str.:	5	3	8	4	_	_	_	_	_
2 <sup>nd</sup> chance	5(1)	5(1) 3(1)	5(1) 3(1) 8(1)						
FIFO	5	5 3	5 3 8	5 3 8 4					

## CS 6901 Capstone Exam Data Structures Fall 2013: Choose any 2 problems.

- 1) Given a (possibly empty) binary search tree of integers, write an iterative function that inserts an integer x into the tree. Do not use recursion.
- 2) Consider the following quicksort code to sort an array of floats. The algorithm is begun by the call quicksort (a, 0, n-1).

```
void quicksort(float a[], int first, int last)
// Sort a[first]..a[last]; original call: quicksort(a, 0, n-1)
{
   if (first < last) {
     int splitpt = partition(a, first, last);
     quicksort(a, first, splitpt - 1);
     quicksort(a, splitpt + 1, last);
   }//endif
}</pre>
```

Write the function partition. The function returns splitpt and modifies a[first] .. a[last] in the following way: Let x be the float value in a[first] at the time partition is called. When partition has completed,

```
a[k] \le x for k = first, ..., splitpt
a[k] \ge x for k = splitpt, ..., last.
```

- 3) a) Write an algorithm that returns the median of 3 given integers a, b, and c. The average number of integer comparisons in your algorithm must be less than 3.
- b) Assuming the 3 inputs are random and distinct, determine the precise average number of comparisons in your algorithm. Justify your answer.

## Theory Exam

1. Give a state diagram for a *deterministic* finite automaton that recognizes the following language over  $\Sigma = \{0, 1\}$ :

{w: w has an even number of occurrences of the substring 01 and w has an odd length}

 Answer each of the following questions with <u>only</u> YES or NO to indicate whether or not the following languages are *decidable*. Do not guess if unsure, as wrong answers will lower your score!

Scoring: +2 points for correct answers; 0 points for no answers; -1 point for wrong answers

- a.  $\{D: D \text{ is a deterministic finite automaton and } L(D) = \emptyset\}$
- b.  $\{P: P \text{ is a pushdown automaton and } L(P) = \emptyset\}$
- c.  $\{M: M \text{ is a Turing machine and } L(M) = \emptyset\}$
- d.  $\{D_1, D_2: D_1 \text{ and } D_2 \text{ are deterministic finite automata and } L(D_1) = L(D_2)\}$
- e.  $\{P_1, P_2: P_1 \text{ and } P_2 \text{ are pushdown automata and } L(P_1) = L(P_2)\}$
- f.  $\{M_1, M_2: M_1 \text{ and } M_2 \text{ are Turing machines and } L(M_1) = L(M_2)\}$
- g.  $\{M: M \text{ is a Turing machine that has a state named } q_{27}\}$
- h.  $\{M: M \text{ is a Turing machine with a transition to } q_{27} \text{ in its delta function}\}$
- i.  $\{M, w: M \text{ is a Turing machine that enters its state } q_{27} \text{ when run on input string } w\}$
- j.  $\{M: M \text{ is a Turing machine that enters its state } q_{27} \text{ when run on any string} \}$
- 3. A vertex cover of a graph G = (V, E) is  $C \subseteq V$  such that every edge  $e \in E$  is adjacent to at least one  $c \in C$ . Let **VERTEX-COVER** =  $\{V, E, k: G = (V, E) \text{ is a graph that contains a vertex cover of size } k\}$ . Show that **VERTEX-COVER** is NP-Complete. You may assume the result of the Cook-Levin Theory.