## DATA STRUCTURES (ITPC-203)

# Principles of Recursion



Mrs. Sanga G. Chaki

**Department of Information Technology** 

Dr. B. R. Ambedkar National Institute of Technology, Jalandhar

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### Recurrence Relations



- 1. A sequence is a discrete structure used to represent an ordered list.
- 2. A **recurrence relation** for the sequence  $\{a_n\}$  is an equation that expresses  $a_n$  in terms of one or more of the **initial terms** of the sequence, namely,  $a_0$ ,  $a_1$ , . . . . ,  $a_{n-1}$ , for all non-negative integers n
- 3. A sequence is called a **solution of a recurrence relation** if its terms satisfy the recurrence relation.
- 4. A recurrence relation is said to recursively define a sequence.
- 5. The **initial conditions** for a recursively defined sequence specify the terms that precede the first term where the recurrence relation takes effect.

## Recurrence Relations – Fibonacci Sequence



- 1. The Fibonacci sequence,  $f_0$ ,  $f_1$ ,  $f_2$ , . . . , is defined by the **initial conditions** 
  - $f_0 = 0$ ,  $f_1 = 1 \rightarrow$  The first two terms in the sequence
- 2. The recurrence relation

$$f_n = f_{n-1} + f_{n-2}$$
 for  $n = 2, 3, 4, ...$ 

- Each term is defined using the previous two terms of the sequence.
- So the sequence is defined in terms of itself.

#### Recurrence Relations



1. We say that we have solved the recurrence relation together with the initial conditions when we find an **explicit formula**, called a **closed formula**, for the terms of the sequence. – this formula does not use previous terms in the sequence to find each successive terms.

## Recurrence Relations



Closed

formula

Recurrent

formula

**1. Example**: Determine whether the sequence  $\{a_n\}$ , where  $a_n = 3n$  for every nonnegative integer n, is a solution of the recurrence relation  $\setminus$ 

$$a_n = 2a_{n-1} - a_{n-2}$$
 for  $n = 2, 3, 4...$ 

- Suppose that  $a_n = 3n$  for every nonnegative integer n.
- Then, for any  $n \ge 2$ , we see that
- $a_n = 3n$ ,
- $a_{n-1} = 3(n-1)$
- $a_{n-2} = 3(n-2)$  replace them in the recurrent relation

$$2a_{n-1} - a_{n-2} = 2(3(n-1)) - 3(n-2) = 3n = a_n$$

• Therefore, the sequence  $\{a_n\}$ , where  $a_n = 3n$ , is a solution of the given recurrence relation



## Principles of recursion



- In computer science, recursion is a method of solving a computational problem where the solution depends on solutions to smaller instances of the same problem.
- 2. Recursion solves such recursive problems by using functions that call themselves from within their own code.

- 1. A process by which a function calls itself repeatedly is called recursion
- 2. In the below example, the factorial of n is computed using the results of smaller instances of the same problem namely, the factorial of (n-1).
- 3. This again, would be computed using the results of fact(n-2) and so on.
  - Used for repetitive computations in which each action is stated in terms of a previous result.

```
fact(n) = n * fact(n-1)
```

- For a problem to be written in recursive form, two conditions are to be satisfied:
  - It should be possible to express the problem in recursive form.
  - The problem statement must include a stopping condition

```
fact(n) = 1, if n = 0
= n * fact(n-1), if n > 0
```

Some example problems which can be solved using recursion:

- 1. Factorial previous example.
- 2. Fibonacci Series (0, 1, 1, 2, 3, 5, 8, 13, ....) first example
- 3. GCD finding the greatest common divisor express this problem in recursive way.
- 4. Finding the length of a string express this problem in recursive way.
- 5. Find an example problem which cannot be solved using recursion.

## Recursion – Mechanism of Execution

```
long int fact (n)
int n;
{
    if (n == 0)
       return (1);
    else
      return (n * fact(n-1));
}
```

- Mechanism of execution
  - When a recursive program is executed, the recursive function calls are not executed immediately.
    - They are kept aside (on a stack) until the stopping condition is encountered.
    - The function calls are then executed in reverse order.

## Recursion – Mechanism of Execution

## Example :: Calculating fact (4)

First, the function calls will be processed:

```
fact(4) = 4 * fact(3)
fact(3) = 3 * fact(2)
fact(2) = 2 * fact(1)
fact(1) = 1 * fact(0)
```

The actual values return in the reverse order:

```
fact(0) = 1

fact(1) = 1 * 1 = 1

fact(2) = 2 * 1 = 2

fact(3) = 3 * 2 = 6

fact(4) = 4 * 6 = 24
```

#### Recursion – Fibonacci Numbers

Fibonacci number f(n) can be defined as:

```
f(0) = 0

f(1) = 1

f(n) = f(n-1) + f(n-2), if n > 1
```

The successive Fibonacci numbers are:

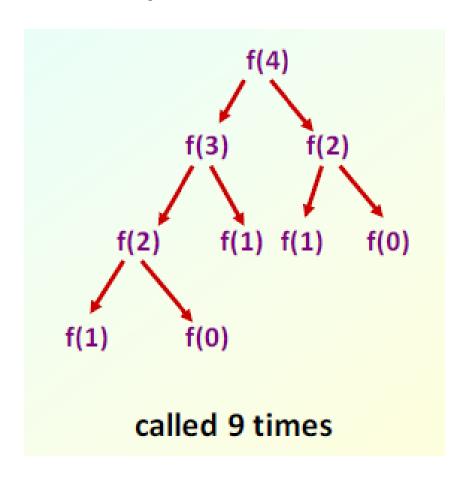
```
0, 1, 1, 2, 3, 5, 8, 13, 21, .....
```

Function definition:

```
int f (int n)
{
    if (n < 2) return (n);
    else return (f(n-1) + f(n-2));
}</pre>
```

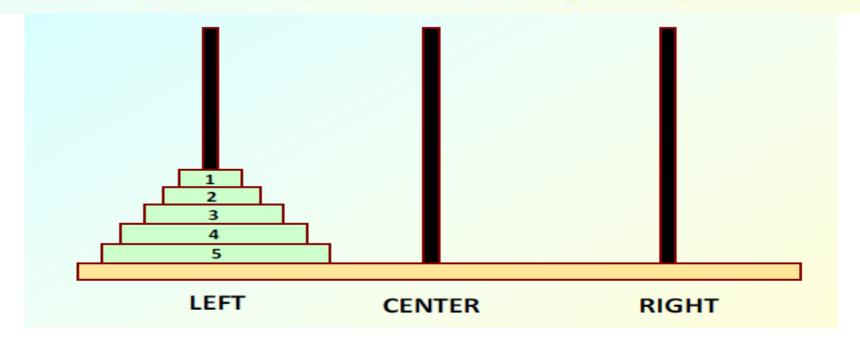
## Recursion – Fibonacci Numbers

How many times the function is called when evaluating f(4)?



- 1. Same thing is calculate multiple times inefficient.
- 2. Exponential "explosion" of calls
- 3. Takes more time and memory

- The problem statement:
  - Initially all the disks are stacked on the LEFT pole.
  - Required to transfer all the disks to the RIGHT pole.
    - Only one disk can be moved at a time.
    - A larger disk cannot be placed on a smaller disk.
  - CENTER pole is used for temporary storage of disks.



- Recursive statement of the general problem of n disks.
  - Step 1:
    - Move the top (n-1) disks from LEFT to CENTER.
  - Step 2:
    - Move the largest disk from LEFT to RIGHT.
  - Step 3:
    - Move the (n-1) disks from CENTER to RIGHT.

```
Move disk 1 from L to R
             Move disk 2 from L to C
Moved disk 3
 from L to R
             Move disk 1 from R to C
             Move disk 3 from L to R
Moved disk 2
             Move disk 1 from C to L
 from L to R
             Move disk 2 from C to R
Moved disk 1
             Move disk 1 from L to R
 from L to R
```

```
Move disk 1 from L to C
                       Move disk 2 from L to R
                       Move disk 1 from C to R
      Moved disk 4
                       Move disk 3 from L to C
        from L to R
                       Move disk 1 from R to L
                       Move disk 2 from R to C
                       Move disk 1 from L to C
                       Move disk 4 from L to R
                       Move disk 1 from C to R
       Moved disk 3
                       Move disk 2 from C to L
        from L to R
                       Move disk 1 from R to L
                       Move disk 3 from C to R
       Moved disk 2
                       Move disk 1 from L to C
         from L to R
                       Move disk 2 from L to R
                       Move disk 1 from C to R
Moved disk 1 from L to R
```

## Thanks!