## CS 228 Minor Autumn 2020 Assignment 1

- 1. Consider the 3 statements given.
  - (a) X : Y is the murderer
  - (b) Y: Z is the murderer
  - (c) Z : Neither X nor Y is the murderer

One of X, Y, or Z has committed the murder, and only one of the above statements is true. Using propositional logic satisfiability, find out the murderer.

- 2. Draw the parse tree for the formula  $\varphi = (r \to \neg s) \lor \neg (p \to (\neg q \lor (r \land (p \to (s \lor r)))))$ .
- 3. Write the formula  $\varphi$  given above in CNF. Call it  $\psi$ . Is  $\psi$  a Horn formula?
- 4. Using resolution, check for satisfiability of  $\varphi$  in question 2.
- 5. Consider the formula  $\varphi = (p \to (q \to p))$ . Show that
  - (a)  $p, \neg q \vdash \varphi$
  - (b)  $\neg p, \neg q \vdash \varphi$
  - (c)  $\neg p, q \vdash \varphi$
  - (d)  $p, q \vdash \varphi$

As seen in the proof of completeness, combine all 4 proofs to obtain a proof  $\vdash \varphi$ .

6. Which of the two formulae implies the other?

$$\varphi_1 = (p \leftrightarrow (q \leftrightarrow r))$$

$$\varphi_2 = ((p \land (q \land r)) \lor ((\neg p) \land ((\neg q) \land (\neg r))))$$

Is it the case that  $\varphi_1 \models \varphi_2$ , or  $\varphi_2 \models \varphi_1$ , both, or neither? Recall that  $\models$  stands for semantic entailment (hence you argue using assignments).

7. Consider a formula  $\varphi = C_1 \wedge C_2 \wedge \dots C_n$  where each clause  $C_i$  is of the form  $(\tau \to \alpha)$  or  $(\alpha_1 \wedge \dots \alpha_n \to \beta)$  or  $(\gamma \to \bot)$  where  $\alpha, \alpha_i, \beta, \gamma$  are literals. A logician wishes to apply HornSAT to this formula  $\varphi$  by renaming negative literals (if any) with fresh positive literals. Thus, if any  $\alpha, \alpha_i, \beta, \gamma$  was of the form  $\neg p$ , the logician will replace that  $\neg p$  with a fresh variable p'. The logician claims that he can check satisfiability of  $\varphi$  correctly by applying HornSAT on the new formula (call it  $\varphi'$ ) in the following way:  $\varphi$  is satisfiable iff HornSAT concludes that  $\varphi'$  is satisfiable, and  $\varphi$  is unsatisfiable iff HornSAT concludes that  $\varphi'$  is unsatisfiable.

- (a) Illustrate the logician's approach on an example (that is, take a  $\varphi$ , and the show how the corresponding  $\varphi'$  will look like).
- (b) Do you agree with the logician? If you do, clearly explain why. Your argument must work for all  $\varphi$  and the respective  $\varphi'$ . If you disagree with the logician, explain why, using your  $\varphi$  and  $\varphi'$  from part (a).