## **CS 228 : Logic in Computer Science**

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#### **Rules for Natural Deduction**

The rule of double negation elimination  $\neg \neg e$ 

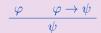
$$\frac{\neg\neg\varphi}{\varphi}$$

The rule of double negation introduction  $\neg \neg i$ 

$$\frac{\varphi}{\neg\neg\varphi}$$

#### **Rules for Natural Deduction**

#### The implies elimination rule or Modus Ponens MP



▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

1.  $p \rightarrow (q \rightarrow \neg \neg r)$  premise

2.

▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

- 1.  $p \rightarrow (q \rightarrow \neg \neg r)$  premise
- 2.  $p \rightarrow q$  premise
- 3.

▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

1.	p  o (q  o  eg  eg r)	premise
2	$n \rightarrow a$	nremise

3. *p* premise

▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

1.	$p  ightarrow (q  ightarrow \lnot \lnot r)$	premise
2.	$ extcolor{p}  ightarrow  extcolor{q}$	premise
3.	р	premise
4.	$q  ightarrow \lnot \lnot r$	MP 1,3
5		

▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

1.	p  o (q  o  eg  eg r)	premise
2.	extstyle p  o q	premise
3.	p	premise
4.	$q  ightarrow \neg \neg r$	MP 1,3

MP 2,3

5. 6.

▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

1.	p  o (q  o  eg  eg r)	premise
2.	${m  ho}  ightarrow {m q}$	premise
3.	р	premise
4.	$q  ightarrow \lnot \lnot r$	MP 1,3
5.	q	MP 2,3
6.	$\neg \neg r$	MP 4,5
7		

▶ Show that  $p, p \rightarrow q, p \rightarrow (q \rightarrow \neg \neg r) \vdash r$ 

1.	p  o (q  o  eg  eg r)	premise
2.	${m  ho}  ightarrow {m q}$	premise
3.	p	premise
4.	$q  ightarrow \lnot \lnot r$	MP 1,3
5.	q	MP 2,3
6.	$\neg \neg r$	MP 4,5
7.	r	¬¬ <i>e</i> 6

#### **Rules for Natural Deduction**

#### Another implies elimination rule or Modus Tollens MT



▶ Show that  $p \rightarrow \neg q, q \vdash \neg p$ 

1.  $p \rightarrow \neg q$  premise

2.

▶ Show that  $p \rightarrow \neg q, q \vdash \neg p$ 

- 1.  $p \rightarrow \neg q$  premise
- 2. *q* premise

3.

▶ Show that  $p \rightarrow \neg q, q \vdash \neg p$ 

1.	p  ightarrow  eg q	premise
2.	q	premise
3.	$\neg \neg q$	¬¬ <i>i</i> 2
1		

▶ Show that  $p \rightarrow \neg q, q \vdash \neg p$ 

1.	p  ightarrow  eg q	premise
2.	q	premise
3.	$\neg \neg q$	¬¬ <i>i</i> 2
4.	$\neg p$	MT 1,3

▶ Can we prove  $p \rightarrow q \vdash \neg q \rightarrow \neg p$ ?

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- ▶ So far, no proof rule that can do this.

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- ▶ Given  $p \rightarrow q$ , let us assume  $\neg q$ . Can we then prove  $\neg p$ ?

- ▶ Can we prove  $p \rightarrow q \vdash \neg q \rightarrow \neg p$ ?
- So far, no proof rule that can do this.
- ▶ Given  $p \rightarrow q$ , let us assume  $\neg q$ . Can we then prove  $\neg p$ ?
- ► Yes, using MT.

## The implies introduction rule $\rightarrow i$

1.	p  o q	premise
2.	$\neg a$	assumption

2. 
$$\neg q$$
 assumptio  
3.  $\neg p$  MT 1,2

$$\neg q 
ightarrow 
eg p 
ightarrow i$$
 2-3

- $\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ 
  - true

2.

premise

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

- 1. true premise 2.  $q \rightarrow r$  assumption
- g ,

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	$q \rightarrow r$	assumption
3.	eg q  o  eg p	assumption
4.		

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	q  o r	assumption
3.	eg q  o  eg p	assumption
4.	р	assumption
5.		

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	$q \rightarrow r$	assumption
3.	eg q  o  eg p	assumption
4.	p	assumption
5.	$  \   \   \ \neg \neg p$	¬¬ <i>i</i> 4
6.		

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	q  o r	assumption
3.	eg q  o  eg p	assumption
4.	p	assumption
5.	$  \   \   \ \neg \neg p$	¬¬ <i>i</i> 4
6.	$  \   \   \ \neg \neg q$	MT 3,5
7.		

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	$q \rightarrow r$	assumption
3.	eg q  ightarrow  eg p	assumption
4.	p	assumption
5.	$  \neg \neg \rho$	¬¬ <i>i</i> 4
6.		MT 3,5
7.	q	¬¬ <i>e</i> 6
8.		

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$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

١.	true	premise
2.	$q \rightarrow r$	assumption
3.	eg q  ightarrow  eg p	assumption
4.	p	assumption
5.	$  \   \ \neg \neg p$	¬¬ <i>i</i> 4
6.	$  \   \   \ \neg \neg q$	MT 3,5
7.		¬¬ <i>e</i> 6
8.		MP 2.7

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	$q \rightarrow r$	assumption
3.	eg q  ightarrow  eg p	assumption
4.	p	assumption
5.		¬¬ <i>i</i> 4
6.		MT 3,5
7.	q	¬¬ <i>e</i> 6
8.	r	MP 2,7
9.	$p \rightarrow r$	<i>→ i</i> 4-8

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise
2.	q  o r	assumption
3.	eg q  o  eg p	assumption
4.	p	assumption
5.	$  \cdot   \cdot   \neg \neg p$	¬¬ <i>i</i> 4
6.	$  \cdot   \cdot \neg \neg q$	MT 3,5
7.		¬¬ <i>e</i> 6
8.	r	MP 2,7
9.	ho ightarrow r	→ <i>i</i> 4-8
10.	$(\neg q  ightarrow \neg p)  ightarrow (p  ightarrow r)$	→ <i>i</i> 3-9

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11.

$$\blacktriangleright \vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$$

1.	true	premise	
2.	q  ightarrow r	assumption	]
3.	eg q  ightarrow  eg p	assumption	
4.	p	assumption	
5.	$\neg \neg p$	¬¬ <i>i</i> 4	
6.		MT 3,5	
7.	q	¬¬ <i>e</i> 6	
8.	r	MP 2,7	
9.	$p \rightarrow r$	→ <i>i</i> 4-8	
10.	$(\neg q  ightarrow  eg  ho)  ightarrow ( ho  ightarrow r)$	→ <i>i</i> 3-9	
11.	$(a \rightarrow r) \rightarrow [(\neg a \rightarrow \neg p) \rightarrow (p \rightarrow r)]$	$\rightarrow$ i 2-10	

► Knowing the proof of  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ , can you prove

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- ► Knowing the proof of  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ , can you prove

  - $(q \to r), (\neg q \to \neg p) \vdash p \to r$

- ► Knowing the proof of  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ , can you prove

  - $(q \to r), (\neg q \to \neg p) \vdash p \to r$
  - $(\neg q \to \neg p) \vdash [(q \to r) \to (p \to r)]$

- ► Knowing the proof of  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ , can you prove

  - $(q \to r), (\neg q \to \neg p) \vdash p \to r$
  - $(\neg q \to \neg p) \vdash [(q \to r) \to (p \to r)]$
  - $(q \rightarrow r), (\neg q \rightarrow \neg p), p \vdash r$
- ► Knowing the proof of any of the above 4 sequents, can you prove  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ ?

## **Transforming Proofs**

► Knowing the proof of  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ , can you prove

```
▶ q \rightarrow r \vdash (\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)

▶ (q \rightarrow r), (\neg q \rightarrow \neg p) \vdash p \rightarrow r

▶ (\neg q \rightarrow \neg p) \vdash [(q \rightarrow r) \rightarrow (p \rightarrow r)]

▶ (q \rightarrow r), (\neg q \rightarrow \neg p), p \vdash r
```

- ► Knowing the proof of any of the above 4 sequents, can you prove  $\vdash (q \rightarrow r) \rightarrow [(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow r)]$ ?
- ► Transform any proof  $\varphi_1, \dots, \varphi_n \vdash \psi$  to  $\vdash \varphi_1 \to (\varphi_2 \to \dots (\varphi_n \to \psi) \dots)$  by adding n lines of the rule  $\to i$
- ► See an example : transform proof of  $p \to (q \to r) \vdash (p \land q) \to r$ into that of  $\vdash p \to (q \to r) \to [(p \land q) \to r]$

1.	p  o (q  o r)	premise
2.	$p \wedge q$	assumption
3.	p	<i>∧e</i> <sub>1</sub> 2
4.	q	<i>∧e</i> <sub>2</sub> 2
5.	q  ightarrow r	MP 1,3
6.	r	MP 4,5

7.

 $p \land q \rightarrow r \rightarrow i \ 2-6$ 

7.

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	ho  ightarrow (q  ightarrow r)	assumption
2		

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	p  o (q  o r)	assumption
3.	$p \wedge q$	assumption
4.		

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	p  o (q  o r)	assumption
3.	$p \wedge q$	assumption
4.	p	∧ <i>e</i> <sub>1</sub> 3
5.		

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	p  ightarrow (q  ightarrow r)	assumption
3.	$p \wedge q$	assumption
4.	p	∧ <i>e</i> <sub>1</sub> 3
5.	q	∧ <i>e</i> <sub>2</sub> 3
6.		

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	p  o (q  o r)	assumption
3.	$p \wedge q$	assumption
4.	p	∧ <i>e</i> <sub>1</sub> 3
5.		∧ <i>e</i> <sub>2</sub> 3
6.	$  \mid q \rightarrow r$	MP 2,4
7.		

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	p  o (q  o r)	assumption
3.	$p \wedge q$	assumption
4.	p	∧ <i>e</i> <sub>1</sub> 3
5.	q	∧ <i>e</i> <sub>2</sub> 3
6.	$  q \rightarrow r$	MP 2,4
7.	r	MP 5,6
8.		

$$\blacktriangleright \vdash [p \to (q \to r)] \to [(p \land q) \to r]$$

1.	true	premise
2.	ho  ightarrow (q  ightarrow r)	assumption
3.	$p \wedge q$	assumption
4.	p	∧ <i>e</i> <sub>1</sub> 3
5.	q	<i>∧e</i> <sub>2</sub> 3
6.	$  q \rightarrow r$	MP 2,4
7.	r	MP 5,6
8.	$p \wedge q  ightarrow r$	→ <i>i</i> 3-7
9.	$[p \rightarrow (q \rightarrow r)] \rightarrow [(p \land q) \rightarrow r]$	$\rightarrow$ i 2-8

### **More Rules**

#### The or introduction rule $\vee i_1$

$$\frac{\varphi}{\varphi\vee\psi}$$

#### The or introduction rule $\vee i_2$

$$\frac{\psi}{\varphi\vee\psi}$$

### **More Rules**

#### The or elimination rule ∨*e*

$$\begin{array}{ccc} \varphi \lor \psi & \varphi \vdash \chi & \psi \vdash \chi \\ \hline \chi & \end{array}$$

$$P q \to r \vdash (p \lor q) \to (p \lor r)$$

- 1.  $q \rightarrow r$
- 2

premise

$$P q \rightarrow r \vdash (p \lor q) \rightarrow (p \lor r)$$

1.	$q \rightarrow r$	premise
2.	$p \lor q$	assumpti
3.		

$$P q \to r \vdash (p \lor q) \to (p \lor r)$$

1.	q  o r	premise
2.	$p \lor q$	assumption
3.	p	assumption
4.		

$$P q \to r \vdash (p \lor q) \to (p \lor r)$$

۱.	q  o r	premise
2.	$p \lor q$	assumption
3.	р	assumption
4.	$p \lor r$	∨ <i>i</i> <sub>1</sub> 3
5.		

$$P q \to r \vdash (p \lor q) \to (p \lor r)$$

١.	q  o r	premise
2.	$p \lor q$	assumption
3.	p	assumption
1.	p∨r	∨ <i>i</i> <sub>1</sub> 3
5.	q	assumption
3.		

$$P q \to r \vdash (p \lor q) \to (p \lor r)$$

1.	q  o r	premise
2.	$p \lor q$	assumption
3.	р	assumption
4.	<i>p</i> ∨ <i>r</i>	∨ <i>i</i> <sub>1</sub> 3
5.	q	assumption
6.	r	MP 1,5
7.		

$$P q \rightarrow r \vdash (p \lor q) \rightarrow (p \lor r)$$

1.	$q \rightarrow r$	premise
2.	$p \lor q$	assumption
3.	р	assumption
4.	$p \lor r$	∨ <i>i</i> <sub>1</sub> 3
5.	q	assumption
6.	r	MP 1,5
7.	p∨r	∨ <i>i</i> <sub>2</sub> 6

$$P q \rightarrow r \vdash (p \lor q) \rightarrow (p \lor r)$$

1.	q  o r	premise
2.	$p \lor q$	assumption
3.	р	assumption
4.	$p \lor r$	∨ <i>i</i> ₁ 3
5.	q	assumption
6.	r	MP 1,5
7.	$p \lor r$	∨ <i>i</i> <sub>2</sub> 6
8.	p∨r	∨ <i>e</i> 2, 3-4, 5-7

$$P q \rightarrow r \vdash (p \lor q) \rightarrow (p \lor r)$$

1.	q  o r	premise
2.	$p \lor q$	assumption
3.	p	assumption
4.	$p \lor r$	√ <i>i</i> <sub>1</sub> 3
5.	q	assumption
6.	r	MP 1,5
7.	$p \lor r$	∨ <i>i</i> <sub>2</sub> 6
8.	$p \vee r$	∨ <i>e</i> 2, 3-4, 5-7
_ `		

$$P q \rightarrow r \vdash (p \lor q) \rightarrow (p \lor r)$$

1.	$q \rightarrow r$	premise
2.	$p \lor q$	assumption
3.	р	assumption
4.	p∨r	∨ <i>i</i> <sub>1</sub> 3
5.	q	assumption
6.	r	MP 1,5
7.	p∨r	∨ <i>i</i> <sub>2</sub> 6
8.	p∨r	∨ <i>e</i> 2, 3-4, 5-7
9.	$(p \lor q) \to (p \lor r)$	$\rightarrow$ i 2-8