

# Data Structures

## Important Points

- Big Theta notation is used to denote (tight) asymptotic bound (upper and lower bound)
- Big Omega notation is used to denote asymptotic lower bound.
- Big O notation is used to denote an asymptotic upper bound.

## Stack ADT operations

- Push - public
- Pop - public
- Peek - public
- isFull - public or private
- isEmpty - public or private
- Size - public or private

Internal data structure should be either array or linked list or both.



→ The main memory is organized for access and storage in row major order and column major order.

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Binary search	↔	Log base 2
Last in first out	↔	Recursion

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Modularity	↔	Testing of modules separate from rest of the system
Encapsulation	↔	Independence of module's behaviour
Separation of concerns	↔	Module responsible for one and only one feature
Abstraction	↔	Hiding details of the implementation

→ Stack ADT supports LIFO/FILO operation



→ Peek	↔	Return value of latest push operation
Top of stack	↔	Latest insertion in the stack
Pop	↔	Decrement top of stack
Push	↔	Increment top of stack

- Queue supports FIFO/LIFO model.
- Like circular queue, implementing a circular stack is possible.
- Having elements of a queue at contiguous locations does not reduce the time complexity of its operations.
- Stack is better suited than queue for reversing the order of elements that are sorted.



Null pointer	$\longleftrightarrow$	tail $\rightarrow$ next
Traversing a linked list	$\longleftrightarrow$	strictly linear order
Contiguous memory locations	$\longleftrightarrow$	Access via indexing
Element of type struct	$\longleftrightarrow$	both Linked List and Array

→ Space Complexity : Memory space required by an algorithm in its execution cycle. Space needed by an algorithm is equal to the sum of the following two components :

- Fixed part : Space required to store data and variables that are not dependent of the size of input.
- Variable part : Space required by variables, whose size is dependent on the size of the problem.
- Space complexity is usually denoted by Big O notation.



## → Postfix Examples

1.  $62+8/23+5/*$

Infix method

$$= (62+) 8/23+5/*$$

$$= (6+2) 8/23+5/*$$

$$= (88/) 23+5/*$$

$$= (8/8) 23+5/*$$

$$= 1(23+)5/*$$

$$= 1(2+3)5/*$$

$$= 1(55/) *$$

$$= 1(5/5) *$$

$$= 11* = 1 //$$

2.  $62/7*23-5+*$

Infix method

$$= (62/) 7*23-5+*$$

$$= (6/2) 7*23-5+*$$

$$= (37*) 23-5+*$$

$$= (3*7) (23-) 5+*$$

$$= 21(2-3)5+*$$

$$= 21(-15+)*$$

$$= 21(-1+5)*$$

$$= 214*$$

$$= 21*4$$

$$= 84 //$$

3.  $46+7+9+2-$

Infix method

$$= (46+) 7+9+2-$$

$$= (4+6) 7+9+2-$$

$$= (107+) 9+2-$$

$$= (10+7) 9+2-$$

$$= (179+) 2-$$

$$= (17+9) 2-$$

$$= (262-)$$

$$= 26-2 = 24 //$$



$$4. 32 * 955 + 2 // +$$

Infix method

$$= (32 *) 955 + 2 // +$$

$$= (3 * 2) 9 (55 +) 2 // +$$

$$= 6 9 (5 + 5) 2 // +$$

$$= 6 9 (10 2 /) / +$$

$$= 6 9 (10 / 2) / +$$

$$= \cancel{6 (9 / 5)}$$

$$= 6 (95 /) +$$

$$= 6 (9 / 5) +$$

$$= (6 1 \cdot 8) +$$

$$= (6 + 1 \cdot 8) = 7 \cdot 8 //$$

$d$  = current depth  
(range 0 to  $D-1$ )

Lookup operation

2D Array & 3D Array

$R$  = Total Rows

$C$  = Total Columns

$D$  = Total Depth

$a[x][y][z]$ ,

$$\text{index} = x * C + y$$

$a[x][y][z]$ ,

$$\text{index} = d * R * C + x * C + y$$

$x$  = Current row  
(range 0 to  $R-1$ )

$y$  = Current column  
(range 0 to  $C-1$ )

bottom right