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Aim: To use Backtracking to solve the N Queen problem

Theory:

Backtracking

Backtracking is an algorithmic technique for solving problems recursively by trying to build a solution incrementally, one piece at a time, removing those solutions that fail to satisfy the constraints of the problem at any point in time.

Backtracking can also be said as an improvement to the brute force approach. So basically, the idea behind the backtracking technique is that it searches for a solution to a problem among all the available options.

Initially, we start the backtracking from one possible option and if the problem is solved with that selected option then we return the solution else we backtrack and select another option from the remaining available options. There also might be a case where none of the options will give you the solution and hence we understand that backtracking won't give any solution to that particular problem

Algorithm: is

safe(chessboard,row, col, n)

- 1. for i = 0 to row 1
- 2. if chessboard[i][col] == 1
- 3. return false
- 4. for i = row, j = col; $i \ge 0$ and $j \ge 0$; i -, j -
- 5. if chessboard[i][j] == 1
- 6. return false
- 7. for i = row, j = col; $i \ge 0$ and j < n; i -, j + +
- 8. if chessboard[i][j] == 1



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9. return false 10. return true

```
solve_nqueen(chessboard, row, n)
1. if row == n
2.    print "Solution " ++ sol_count ++ ":"
3.    print_chessboard(chessboard, n)
4.    return
5. for i = 0 to n - 1
6.    if is_safe(chessboard, row, i, n)
7.         chessboard[row][i] = 1
8.         solve_nqueen(chessboard, row + 1, n)
9.         chessboard[row][i] = 0
```

PROGRAM:

```
#include <stdio.h>
#include <math.h>
#include <stdlib.h>

//board index --> row no , borad[index] --> col no at which queen is placed
int board[20], count;
int main()
{
    int n, i, j;
    void queen(int row, int n);

    printf(" - N Queens Problem Using Backtracking -");
    printf("\n\nEnter number of Queens:");
    scanf("%d", &n);
    queen(1, n);
    return 0;
}

// function for printing the solution
void print(int n)
{
    int i, j;
    printf("\n\nSolution %d:\n\n", ++count);
```

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```
for (i = 1; i <= n; ++i){
        printf("\t%d", i);
    for (i = 1; i \le n; ++i)
        printf("\n%d\t", i);
        for (j = 1; j \leftarrow n; ++j) // for nxn board
            if (board[i] == j)
                printf("Q\t");
            else
                printf(".\t"); // empty slot
        }
    printf("\n\n");
int place(int row, int column)
    int i;
    for (i = 1; i \le row - 1; ++i)
        if (board[i] == column)
            return 0;
        else if (abs(board[i] - column) == abs(i - row))
            return 0;
    return 1; // all cond met
void queen(int row, int n)
    int column;
    for (column = 1; column <= n; ++column)</pre>
```

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Output:

Solution 20:									
	1	2	3	4	5	6	7	8	
1			Q						
2						Q			
3		Q							
4							Q		
5					Q				
6	Q			•					
7			•	•	•	•	•	Q	
8	•			Q	•		•	•	
Solut	Solution 21:								
	1	2	3	4	5	6	7	8	
1		•	Q	•			•		
2				•	•	Q	•		
3	:	•	•	Q			•		
4	Q	•	•	•	•	•	•	•	
5	•		•	•	•	•	•	Q	
6	•	•	•	•	Q	•	•	•	
7 8	•	•	•	•		•	Q	•	
8	•	Q	•	•	•	•	•	•	



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Middle solution:

Whate Solution.										
Solution 55:										
	1	2	3	4	5	6	7	8		
1					Q					
2			Q							
3	Q									
4							Q			
5		Q								
6								Q		
7						Q				
8				Q						
Solut	Solution 56:									
	1	2	3	4	5	6	7	8		
1					Q					
2			Q		•					
3					•			Q		
4			•	Q	•					
5				•	•		Q	•		
6	Q		•		•			•		
7			•	•	•	Q	•	•		
8		Q	•	•	•					

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2 Intermediate solution:

Solution 92:									
	1	2	3	4	5	6	7	8	
1								Q	
2				Q					
3	Q								
4			Q				•		
5						Q			
6		Q		•					
7			•	•			Q		
8			•	•	Q		•		

Total solution count: 92

Observation:

- The n-queen problem has a significant computational complexity for large board sizes.
- The number of solutions increases rapidly with increasing board size.
- Backtracking algorithms like the one used in this implementation can efficiently find all solutions.

The time complexity is O(n!) as n queen is an combinatorics problem, we are just arranging n items.



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Conclusion:

In conclusion, the n-queen problem is a challenging and computationally complex problem. However, with backtracking algorithms, it is possible to efficiently find all solutions for small to moderate board sizes.