Database Management System – 34 Database design – Closure and Equivalence set

Ajay James Asst. Prof in CSE Government Engineering College Thrissur

Outline

- Closure
- Closure of attributes algorithm
- Cover
- Equivalence set

Closure

- Closure of a set F of FDs is the set F⁺ of all FDs that can be inferred from F
- Attribute Closure
 - determine each set of attributes X that appears as a left-hand side of some FD in F
 - then determine the set of all attributes that are dependent on X

Closure of X

- For each such set of attributes X
- Determine the set X⁺ of attributes that are functionally determined by X based on F
- X⁺ is called the closure of X under F

Algorithm to calculate X⁺

 Input: A set F of FDs on a relation schema R, and a set of attributes X, which is a subset of R.

```
X^+:=X;

repeat

oldX^+:=X^+;

for each functional dependency Y \to Z in F do

if X^+ \supseteq Y then X^+:=X^+ \cup Z;

until (X^+=oldX^+);
```

$R(ABCDE) = \{AB->C, B->D, C->E, D->A\}$

```
    AB<sup>+</sup>

    X+:= AB

X^{+} := X;

    oldX<sup>+</sup>:= AB

repeat
                                               1. AB \rightarrow C, AB \supseteq AB,
   oldX^+:=X^+;
                                                    X^+:=ABC
   for each FD Y \rightarrow Z in F do
       if X^+ \supseteq Y then X^+ := X^+ \cup Z; 2. B \Rightarrow D, ABC \supseteq B,
                                                     X^+:=ABCD
   until (X^+ = oldX^+);
                                               3. C \rightarrow E, ABCD \supseteq C
                                                    X^+:=ABCDE
                                               4. D \rightarrowA, ABCDE \supseteq D
                                                    X^+:=ABCDE
```

Example

- CLASS (Classid, Course#, Instr_name, Credit_hrs, Text, Publisher, Classroom, Capacity)
- FD1: Classid → {Course#, Instr_name,
 Credit_hrs, Text, Publisher, Classroom, Capacity}
- FD2: Course# → Credit_hrs
- FD3: {Course#, Instr name} → {Text, Classroom}
- FD4: Text → Publisher
- FD5: Classroom → Capacity

Closure example

- { Classid } + = { Classid , Course#, Instr_name, Credit_hrs, Text, Publisher, Classroom, Capacity } = CLASS
- { Course#} + = { Course#, Credit_hrs}
- { Course#, Instr_name } + = { Course#, Instr_name, Credit_hrs, Text, Publisher, Classroom, Capacity }

Equivalence of Sets of Functional Dependencies

- COVER Definition
- A set of functional dependencies F is said to cover another set of functional dependencies G if every FD in G is also in F⁺
- that is, if every dependency in G can be inferred from F
- Alternatively, we say G is covered by F

Equivalence of Sets of Functional Dependencies

- Definition
- Two sets of functional dependencies G and F are equivalent if G⁺ = F⁺.
- Equivalence means that every FD in G can be inferred from F, and every FD in F can be inferred from G
- G is equivalent to F if both the conditions—G covers F and F covers G—hold

$$F = \{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$$

and $G = \{A \rightarrow CD, E \rightarrow AH\}$

- First Check G covers F or $(F \subseteq G^+)$
- Check each FD in F:
- F1: A → C
 - find A⁺ in **G**
 - A^+ in G = {ACD} which includes C
- F2:AC → D
 - get AC+ in G
 - AC⁺ in G = {ACD} which contains D
- F3: E → AD, E → H
 - get E+ in G
 - E^+ in G = {EACDH} which contains AD and H.

$$F = \{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$$

and $G = \{A \rightarrow CD, E \rightarrow AH\}$

- Second Check F covers G or $(G \subseteq F^+)$
- · Check each FD in G:
- G1: A → CD
 - find A+ in F
 - $-A^+$ in F = {ACD} which includes CD
- G2:E → AH
 - get E⁺ in F
 - E⁺ in F = {EACDH} which contains AH
- G covers F and F covers G, so, F and G are equivalent.

Summary

- Closure of X
- Cover
- Equivalent set

Reference

 Elmasri R. and S. Navathe, Database Systems: Models, Languages, Design and Application Programming, Pearson Education 6th edition and 7th edition Thank you