

# Database Management System – 15 (Relational Model Constraints)

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## Outline

- Relational Integrity Constraints
- Key constraints
- Entity Integrity constraints
- Referential Integrity constraints

## Relational Integrity Constraints

- Constraints are *conditions* that must hold on *all* valid relation instances
  - *Inherent model-based constraints or implicit constraints*
  - *Schema-based constraints or explicit constraints*
  - *Application-based or semantic constraints or business rules*
- **Inherent model-based constraints**
  - Key constraints
  - Entity Integrity constraints
  - Referential Integrity constraints

## Key Constraints

- **Superkey of R:**
    - A set of attributes SK of R such that no two tuples *in any valid relation instance*  $r(R)$  will have the same value for SK
    - For any distinct tuples  $t_1$  and  $t_2$  in  $r(R)$ ,  $t_1[SK] \neq t_2[SK]$
  - **Key of R:**
    - A "minimal" superkey
    - A superkey K such that removal of any attribute from K results in a set of attributes that is not a superkey
- Example:** CAR(State, Reg#, SerialNo, Make, Model, Year)  
 has two keys Key1 = {State, Reg#}, Key2 = {SerialNo}, which are also superkeys  
 {SerialNo, Make} is a superkey but *not* a key
- If a relation has *several candidate keys*, one is chosen arbitrarily to be the **primary key**
  - Primary key attributes are *underlined*

## Example

Diagram illustrating the structure of a relation schema:

- Relation Name:** STUDENT
- Attributes:** Name, Ssn, Home\_phone, Address, Office\_phone, Age, Gpa
- Tuples:** The rows of data in the relation.

Name	Ssn	Home_phone	Address	Office_phone	Age	Gpa
Benjamin Bayer	305-61-2435	(817)373-1616	2918 Bluebonnet Lane	NULL	19	3.21
Chung-cha Kim	381-62-1245	(817)375-4409	125 Kirby Road	NULL	18	2.89
Dick Davidson	422-11-2320	NULL	3452 Elgin Road	(817)749-1253	25	3.53
Rohan Panchal	489-22-1100	(817)376-9821	265 Lark Lane	(817)749-6492	28	3.93
Barbara Benson	533-69-1238	(817)839-8461	7384 Fontana Lane	NULL	19	3.25

**CAR**

<u>License_number</u>	Engine_serial_number	Make	Model	Year
Texas ABC-739	A69352	Ford	Mustang	02
Florida TVP-347	B43696	Oldsmobile	Cutlass	05
New York MPO-22	X83554	Oldsmobile	Delta	01
California 432-TFY	C43742	Mercedes	190-D	99
California RSK-629	Y82935	Toyota	Camry	04
Texas RSK-629	U028365	Jaguar	XJS	04

## Entity Integrity

- Relational Database Schema:**

- A set  $S$  of relation schemas that belong to the same database.  $S$  is the *name* of the **database**

$$S = \{R_1, R_2, \dots, R_n\}$$

- Entity Integrity:**

- *Primary key attributes* PK of each relation schema  $R$  in  $S$  cannot have null values in any tuple of  $r$
- $t[PK] \neq \text{null}$  for any tuple  $t$  in  $r(R)$

- **Note:** Other attributes of  $R$  may be similarly constrained to disallow null values, even though they are not members of the primary key

# Example

## EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
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## DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
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## DEPT\_LOCATIONS

Dnumber	Dlocation
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## PROJECT

Pname	Pnumber	Plocation	Dnum
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## WORKS\_ON

Essn	Pno	Hours
------	-----	-------

## DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
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## EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1985-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

## DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

## DEPT\_LOCATIONS

Dnumber	Dlocation
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

WORKS_ON			PROJECT			
<u>Essn</u>	<u>Pno</u>	Hours	<u>Pname</u>	<u>Pnumber</u>	<u>Plocation</u>	<u>Dnum</u>
123456789	1	32.5	ProductX	1	Bellaire	5
123456789	2	7.5	ProductY	2	Sugarland	5
666884444	3	40.0	ProductZ	3	Houston	5
453453453	1	20.0	Computerization	10	Stafford	4
453453453	2	20.0	Reorganization	20	Houston	1
333445555	2	10.0	Newbenefits	30	Stafford	4
333445555	3	10.0				
333445555	10	10.0				
333445555	20	10.0				
999887777	30	30.0				
999887777	10	10.0				
987987987	10	35.0				
987987987	30	5.0				
987654321	30	20.0				
987654321	20	15.0				
888665555	20	NULL				

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
333445555	Alice	F	1986-04-05	Daughter
333445555	Theodore	M	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	M	1942-02-28	Spouse
123456789	Michael	M	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabeth	F	1967-05-05	Spouse

## Referential Integrity

- Constraint involving *two* relations
- Used to specify a *relationship* among tuples in two relations: the **referencing relation** and the **referenced relation**
- Tuples in the *referencing relation*  $R_1$  have attributes FK (called **foreign key** attributes) that reference the primary key attributes PK of the *referenced relation*  $R_2$
- A tuple  $t_1$  in  $R_1$  is said to **reference** a tuple  $t_2$  in  $R_2$  if  $t_1[FK] = t_2[PK]$
- Displayed in a relational database schema as a directed arc from  $R_1.FK$  to  $R_2.PK$

## Example

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

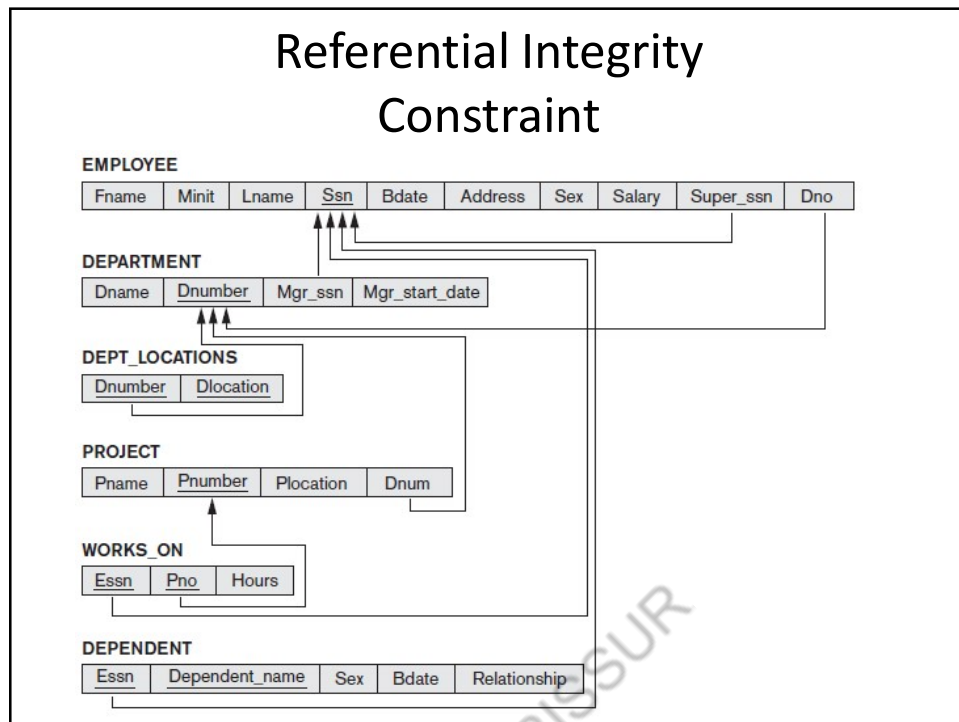
DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

## Referential Integrity Constraint

1. The attributes in FK have the same domain(s) as the primary key attributes PK of  $R_2$ ; the attributes FK are said to reference or refer to the relation  $R_2$
2. A value of FK in a tuple  $t_1$  of the current state  $r_1(R_1)$  either occurs as a value of PK for some tuple  $t_2$  in the current state  $r_2(R_2)$  or is NULL
  - In the former case, we have  $t_1[FK] = t_2[PK]$
  - In case (2), the FK in  $R_1$  should not be a part of its own primary key

## Referential Integrity Constraint



EMPLOYEE									
Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1985-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
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DEPARTMENT			
Dname	Dnumber	Mgr_ssn	Mgr_start_date
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DEPT_LOCATIONS	
Dnumber	Dlocation
1	Houston
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5	Bellaire
5	Sugarland
5	Houston

## Other Types of Constraints

- Semantic Integrity Constraints:
  - based on application semantics and cannot be expressed by the model
  - E.g., “the max. no. of hours per employee for all projects he or she works on is 56 hrs per week”
  - *A constraint specification language* may have to be used to express these

## Exercise - 1

1. Consider the following relations for a database that keeps track of student enrollment in courses and the books adopted for each course:

*STUDENT*(Ssn, Name, Major, Bdate)

*COURSE*(Course#, Cname, Dept)

*ENROLL*(Ssn, Course#, Quarter, Grade)

*BOOK\_ADOPTION*(Course#, Quarter, Book\_isbn)

*TEXT*(Book isbn, Book\_title, Publisher, Author)

Specify the foreign keys for this schema, stating any assumptions you make.



## Exercise - 1

2. Consider the following relations for a database that keeps track of business trips of salespersons in a sales office:

*SALESPERSON*(Ssn, Name, Start\_year, Dept\_no)

*TRIP*(Ssn, From\_city, To\_city, Departure\_date, Return\_date, Trip\_id)

*EXPENSE*(Trip\_id, Account#, Amount)

A trip can be charged to one or more accounts. Specify the foreign keys for this schema, stating any assumptions you make.

## Reference

- Elmasri R. and S. Navathe, Database Systems: Models, Languages, Design and Application Programming, Pearson Education 6<sup>th</sup> edition and 7<sup>th</sup> edition

Thank you

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