

COMP20003 Algorithms and Data Structures

Topological Sorting

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Directed

Acyclic

Useful for modelling many problems:

- Temporal dependencies
- Causalities
- Hierarchies
- Compiling modularized programs



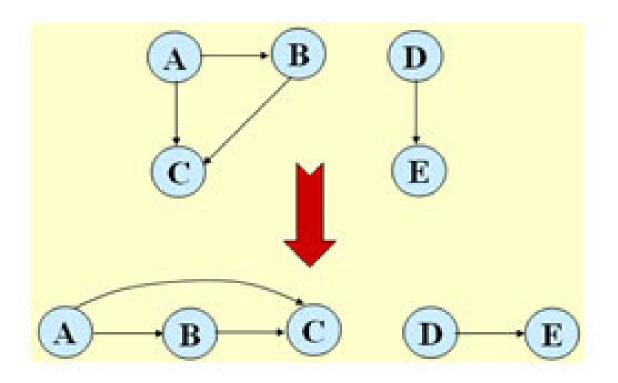
Topological sort: a partial ordering that fulfils certain constraints

All edges e(i,j) go in horizontal i→j direction

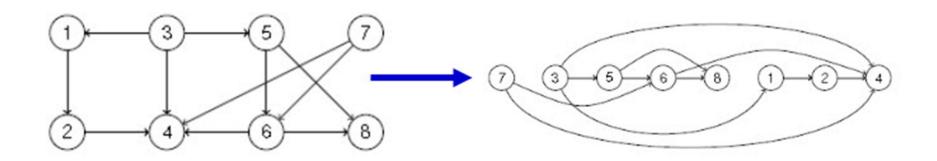
The output can be the schedule for:

- A builder
- A course plan
- etc.

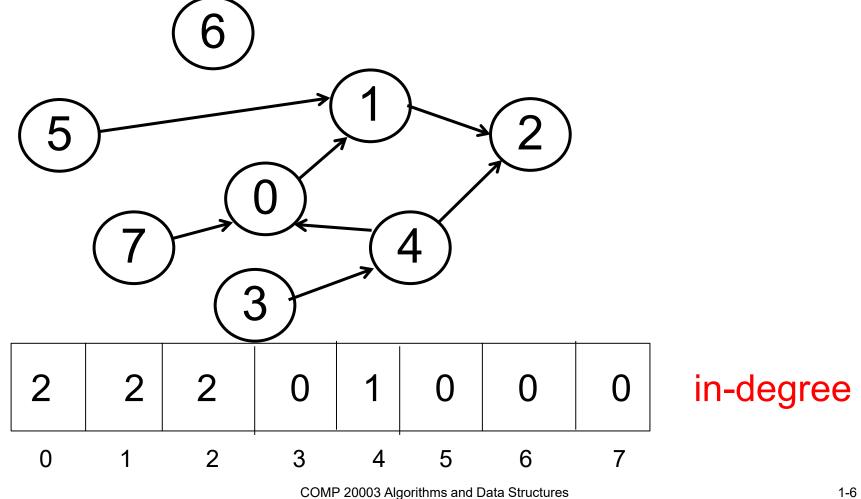




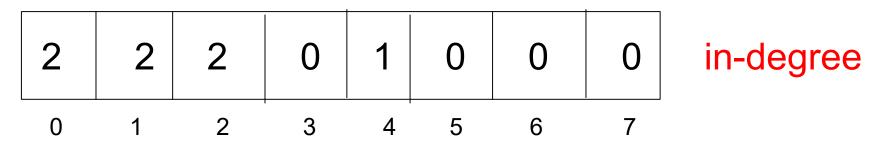










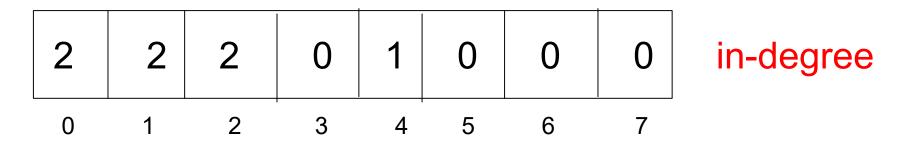


Identify a source (in-degree = 0)

- put that node in the topsort output
- remove that node from DAG
- update in-degree matrix

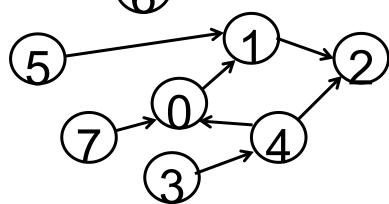
Indentify another source...





• 7-6-5-3-4(new src)-0(new src)-1(new src) - 2(new src)

- **6**-3-4-7-0-5-1-2
- 5-7-3-4-0-1-2-6
- etc.





There must be at least one source and one sink for this topological sorting algorithm to work.

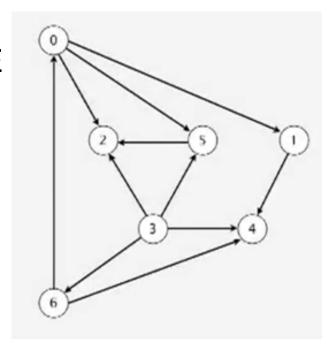
Is this a valid assumption?

Directed acyclic graph: must have at least one source and one sink



Complexity O()?

EXCERCISE





If a Hamiltonian path exists in the DAG, then the topological sort is unique.

Finding a Hamiltonian Path is NP-Hard

How hard is to prove Uniqueness?

Have to solve the decision problem: Given a DAG, Does a Hamiltonian Path exists?



The Hamiltonian Path problem has a property called polynomial verifiability

 Verifying the existence of a Hamiltonian path can be easier than determining its existence.

Given a topological sort, if two consecutive vertices are not connected, then you can swap them. Implies: Non Unique and No Ham. Path. Can be done in linear time.