

## CS711 Assignment 3

Sankalp Gambhir  
180260032

October 5, 2021

### Question 1.

---

*Proof.* Given the set  $S_{d,n}$  as defined, suppose if possible that we cannot find an assignment  $\mathbf{b}$  in this set for some non-zero polynomial  $P$  of degree  $\leq d$  in  $n$  variables such that the evaluation is non-vanishing, i.e.,  $P(\mathbf{b}) = 0 \forall \mathbf{b} \in S_{d,n}$ .

Write the polynomial  $P$  defined by the  $\binom{n+d}{d}$  coefficients  $\{a_{\alpha_1 \alpha_2 \dots \alpha_n}\}$  as a linear functional acting on a vector of pre-evaluated monomials  $\{x_1^{\alpha_1} x_2^{\alpha_2} \dots x_n^{\alpha_n}\}$  with  $\sum \alpha_i \leq d$ ,

$$P(x_1, x_2, \dots, x_n) = (a_{00\dots 0} \ a_{10\dots 0} \ \dots \ a_{00\dots d})_{1 \times \binom{n+d}{d}} (1 \ x_1 \ \dots \ x_n^d)_{1 \times \binom{n+d}{d}}^\top$$

Writing the condition for  $\mathbf{b}$  in this notation, it resolves to the coefficient vector being acted on the right by the matrix  $M$  with columns as the x-vector pre-evaluated over each  $\mathbf{b}$ . The condition assumed above is equivalent to this matrix having a non-trivial kernel, i.e., there existing a coefficient vector which is non-zero, such that all its evaluations over this set are 0.

Consider any two distinct assignments  $A = (\alpha_1, \alpha_2, \dots, \alpha_n)$  and  $B = (\beta_1, \beta_2, \dots, \beta_n)$ . Suppose their pre-evaluated x-vectors (columns of  $M$ ) are linearly dependent, we have

$$\exists \lambda, \delta \in \mathbb{C} \ \lambda \vec{X}(A) + \delta \vec{X}(B) = \vec{0} .$$

Cherry-picking dimensions from this vector equation, first the degree 0 condition, and finally all the degree 1 conditions, we find

$$\begin{aligned} \lambda \cdot 1 + \delta \cdot 1 &= 0 , \\ \lambda &= -\delta , \\ \vec{X}(A) &= \vec{X}(B) , \text{ and} \\ \alpha_i &= \beta_i \forall i . \end{aligned}$$

However, we stipulated that these assignments were distinct, so this is a contradiction. Thus, columns of  $M$ , which represent distinct assignments, cannot be linearly dependent,  $M$  is of full-rank, and thus has a trivial kernel. This again, is a contradiction, so contrary to assumption, we must

be able to find  $\mathbf{b} \in S_{d,n}$  for every non-zero polynomial  $P$  of degree  $\leq d$  in  $n$  variables such that it evaluates to a non-zero quantity over  $\mathbf{b}$ .

□

### Question 2.

*Proof.* Given the set  $T_{d,n}$  as defined, suppose if possible, that we cannot find any vector in the set such that a given non-zero polynomial  $P$  of degree  $d$  in  $n$  variables evaluates to a non-zero value on it. I prove that this leads to a contradiction.

The evaluation condition implies  $P(\mathbf{b}) = 0$  for every  $\mathbf{b} \in T_{d,n}$ . Instead of writing it as a polynomial system, evaluate each possible variable in  $P(x_1, x_2, \dots, x_n)$  and write it as the following linear system

$$(a_{00\dots 0} \quad a_{10\dots 0} \quad \dots \quad a_{00\dots 0d}) (x_1^0 x_2^0 \dots x_n^0 \quad x_1^1 x_2^0 \dots x_n^0 \quad \dots \quad x_1^0 x_2^0 \dots x_n^d)^\top$$

This linear system has dimension  $\binom{n+d}{d}$ . Given that we have evaluations on  $\binom{n+d}{d}$  points as well, we can write this as the linear system

$$(a_{00\dots 0} \quad a_{10\dots 0} \quad \dots \quad a_{00\dots 0d}) \begin{pmatrix} p_1^{0\cdot 0} p_2^{0\cdot 0} \dots p_n^{0\cdot 0} & p_1^{0\cdot 1} p_2^{0\cdot 1} \dots p_n^{0\cdot 1} & \dots & p_1^{0\cdot t} p_2^{0\cdot t} \dots p_n^{0\cdot t} \\ p_1^{1\cdot 0} p_2^{0\cdot 0} \dots p_n^{0\cdot 0} & p_1^{1\cdot 1} p_2^{0\cdot 1} \dots p_n^{0\cdot 1} & \dots & p_1^{1\cdot t} p_2^{0\cdot t} \dots p_n^{0\cdot t} \\ \vdots & \ddots & \ddots & \vdots \\ p_1^{0\cdot 0} p_2^{0\cdot 0} \dots p_n^{d\cdot 0} & p_1^{0\cdot 1} p_2^{0\cdot 1} \dots p_n^{d\cdot 1} & \dots & p_1^{0\cdot t} p_2^{0\cdot t} \dots p_n^{d\cdot t} \end{pmatrix} = (0 \quad 0 \dots 0) ,$$

with  $t = \binom{n+d}{d}$ . This matrix must have a non-trivial kernel, since the coefficients were constrained to be non-zero. However, note that this is the Vandermonde matrix in the  $\binom{n+d}{d}$  integers  $\{p_1^{\alpha_1} p_2^{\alpha_2} \dots p_n^{\alpha_n}\}$ , with  $\sum \alpha_i \leq d$ . Since these are distinct primes, there are no duplicates in this set, and thus the determinant of the Vandermonde matrix must be non-zero. However, this would mean that it is of full rank and has a trivial kernel, this is a contradiction.

Thus, our assumption of being unable to produce a non-zero evaluation over the given set was incorrect, and there exists a vector in  $T_{d,n}$  for every polynomial of maximum total degree  $d$  and  $n$  variables such that it evaluates to a non-zero complex quantity over the given vector provided that it is not identically zero.

□

**Question 3.**

*Proof.* Consider a multilinear polynomial  $P \in \mathbb{F}[x_1, x_2, \dots, x_n]$ . Since it is multilinear, and thus linear in  $x_1$ , we can express it as:

$$P(x_1, x_2, \dots, x_n) = x_1^0 P_1(x_2, x_3, \dots, x_n) + x_1^1 P'_1(x_2, x_3, \dots, x_n) .$$

As this decomposes into the  $(n-1)$  problem, it is easy to see that if either of the polynomials  $P_1$  or  $P'_1$  have non-zero assignments in  $\{0, 1\}^{n-1}$ , then we can generate a non-zero assignment including  $x_1 \in \{0, 1\}$ . The proof proceeds by induction.

*Base case:*  $n = 1$ .

$P_1, P'_1$  in this case are constants, with atleast one of them non-zero. If  $P_1$  is 0, choose  $x_1 = 1$ , else choose  $x_1 = 0$ .  $P \neq 0$  guarantees this assignment leads to a non-vanishing result.

*Induction:* Given that  $\forall m \in \mathbb{N}, m < n$ , every non-zero multilinear polynomial in  $m$  variables has a non-zero satisfying assignment, the same holds for non-zero multilinear polynomials in  $n$  variables.

As before, write

$$P(x_1, x_2, \dots, x_n) = x_1^0 P_1(x_2, x_3, \dots, x_n) + x_1^1 P'_1(x_2, x_3, \dots, x_n) ,$$

where  $P_1, P'_1$  depend on atmost  $n-1$  variables. If  $P_1 \neq 0$ , by the induction hypothesis, it has an assignment such that it is non-vanishing. Append to such an assignment,  $x_1 = 0$ . Else,  $P_1 \equiv 0$ , so we must have  $P'_1 \neq 0$  since  $P$  is given to be non-zero. Append  $x_1 = 1$  to a non-vanishing assignment of  $P'_1$ .

This produces a satisfying assignment for any non-zero multilinear polynomial of degree  $n$ .

□

**Question 4.**

Given the matrix  $M$  as defined and the functions  $\{f_i(\mathbf{y})\}$ , note that the non-vanishing of the  $\det(M)$  implies the linear independence of  $\{(f_i(\mathbf{x}_j))\}$  as vectors of functions. I prove instead the converse, that their linear *dependence* correspond to each other.

Forward Implication  $\Rightarrow$ : Linear dependence of  $\{f_i(\mathbf{y})\}$  implies the vanishing of  $\det(M)$ .

Given that the functions with the  $y$ 's as variables are linearly dependent, we can find a set of coefficients  $\{\alpha_i\}$  such that

$$\sum \alpha_i f_i(\mathbf{y}) = 0 .$$

Substituting one-by-one  $\mathbf{y}$  by each  $\mathbf{x}_i$ , we obtain  $n$  linear dependence equations, which combined imply

$$\sum \alpha_i \begin{pmatrix} f_1(\mathbf{x}_i) & f_2(\mathbf{x}_i) & \dots & f_n(\mathbf{x}_i) \end{pmatrix}^\top = 0 .$$

This is precisely the condition for linear dependence of the columns of  $M$ , and thus its determinant vanishes.

Backward Implication  $\Leftarrow$ : Vanishing of  $\det(M)$  implies the linear dependence of  $\{f_i(\mathbf{y})\}$ .

Given that the determinant vanishes, we know that the *rows* must be linearly dependent, i.e., we can find a set  $\{\alpha_i\}$  such that

$$\sum \alpha_i \begin{pmatrix} f_i(\mathbf{x}_1) & f_i(\mathbf{x}_2) & \dots & f_i(\mathbf{x}_n) \end{pmatrix}^\top = 0 .$$

Pick an arbitrary coordinate and substitute for the  $\mathbf{x}_j$  with  $y$ , and we recover the dependence criterion for the  $f_i(\mathbf{y})$ .

This proves that the linear dependence of the  $\{f_i(\mathbf{y})\}$  is equivalent to the vanishing of  $\det(M)$ , and by extension, the linear independence is equivalent to its non-vanishing.