Verus: Verifying Rust Programs Using Linear Ghost Types

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Others: &a' T, &mut a' T, *const T, *mut T, Box<T>, Arc<T>, Mutex<T>

Pop Quiz

Which of the following Rust programs should compile?

```
1 fn main() {
2    let mut x = 5;
3    let r1 = &mut x;
4    let r2 = &mut x;
5    println!("{} {}", r1, r2);
6 }
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3    &s
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```
3 fn main() {
     let mut x = 5;
     {
     let r1 = \&mut x;
        *r1 += 1;
     let r2 = \&mut x;
     *r2 += 1;
10
  println!("{}", x);
11
12 }
```

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On the other hand...

```
1 struct Boxed(i32);
3 impl Boxed {
  fn inc(&mut self) -> () {
self.0 += 1;
6 }
   fn consume(self) -> () {
   ()
8
9
    fn inc_immutable(self) -> Self {
10
   Boxed(self.0 + 1)
12 }
13 }
14
15 fn main() {
let x = Boxed(5); // immutable struct created
   let mut x = x.inc_immutable(); // consumed and bound to a new
     mutable variable
    x.consume(); // moved and consumed
18
    x.inc(); // error: use after move
19
20 }
```

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Verus builds on the existing type and borrow checking done by the Rust compiler. A program is safe iff it passes the Rust compiler *and* Verus.

Writing a simple program with Verus $+\ \mathsf{SMT}\ \mathsf{encoding}$

	1		
	spec	proof	exec
Call spec	1	1	✓
Call proof	X	✓	1
Call exec	X	X	✓

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Compiled to machine code	X	X	✓

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Compiled to machine code	X	X	✓
Mutation	X	✓	✓

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Borrow-checking	X	✓	✓

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Call exec	X	X	✓
Compiled to machine code	X	X	✓
Mutation	X	1	✓
"SMT effects"	X	✓	✓
Borrow-checking	X	1	✓
SMT Types	1	1	X
Pre/post-conditions	X	✓	✓

Mutexes in Rust

```
1 const N: usize = 10;
2 let data = Arc::new(Mutex::new(0));
3
4 let (tx, rx) = channel();
5 for _ in 0..N {
      let (data, tx) = (Arc::clone(&data), tx.clone());
      thread::spawn(move | | {
           let mut data = (* data).lock().unwrap();
8
          *data += 1;
          if *data == N {
10
               tx.send(()).unwrap();
11
          }
12
      });
13
14 }
15
16 rx.recv().unwrap();
```

Interior Mutability and the Myth of Safe Rust

```
impl<T> Mutex<T> {
    // ...

pub fn lock(&self) -> LockResult<MutexGuard<'_, T>> {
    unsafe {
        self.inner.lock();
        MutexGuard::new(self)
    }
}
```

Unsafe code in Verus

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These objects can be verified as destroyed by consuming the permission object (see PPtr).

On (Local) Invariants

It is called an invariant because it additionally ensures that the value stored in the cell satisfies the type invariant:

```
impl < K, V, Pred: InvariantPredicate < K, V >> LocalInvariant < K, V, Pred >
  pub proof fn new(k: K, tracked v: V, ns: int) -> tracked i:
    LocalInvariant < K, V, Pred >
    requires
    Pred::inv(k, v),
    ensures
    i.constant() == k,
    i.namespace() == ns,
```

The invariant checks are in proof mode and are erased. InvariantPredicate is a typeclass just providing the inv function.

Thanks! Questions?

If we have more time we'll discuss and look at how certain things are implemented in the (Rust or Verus) standard library or play with the verifier.