

CSE 431/531: Algorithm Analysis and Design (Fall 2022)

# Dynamic Programming

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*Department of Computer Science and Engineering  
University at Buffalo*

# Paradigms for Designing Algorithms

## Greedy algorithm

- Make a greedy choice
- Prove that the greedy choice is safe
- Reduce the problem to a sub-problem and solve it iteratively
- Usually for optimization problems

## Divide-and-conquer

- Break a problem into many **independent** sub-problems
- Solve each sub-problem separately
- Combine solutions for sub-problems to form a solution for the original one
- Usually used to design more efficient algorithms

# Paradigms for Designing Algorithms

## Dynamic Programming

- Break up a problem into many **overlapping** sub-problems
- Build solutions for larger and larger sub-problems
- Use a **table** to store solutions for sub-problems for reuse

## Recall: Computing the $n$ -th Fibonacci Number

- $F_0 = 0, F_1 = 1$
- $F_n = F_{n-1} + F_{n-2}, \forall n \geq 2$
- Fibonacci sequence: 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89,  $\dots$

### Fib( $n$ )

```
1:  $F[0] \leftarrow 0$   
2:  $F[1] \leftarrow 1$   
3: for  $i \leftarrow 2$  to  $n$  do  
4:    $F[i] \leftarrow F[i - 1] + F[i - 2]$   
5: return  $F[n]$ 
```

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- Store each  $F[i]$  for future use.

# Outline

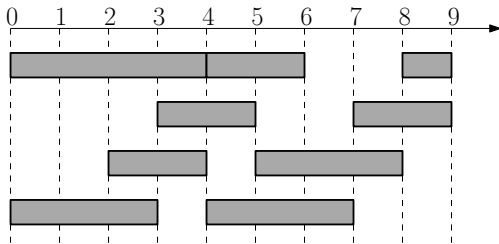
- 1 Weighted Interval Scheduling
- 2 Subset Sum Problem
- 3 Knapsack Problem
- 4 Longest Common Subsequence
  - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree
- 8 Summary

## Recall: Interval Scheduling

**Input:**  $n$  jobs, job  $i$  with start time  $s_i$  and finish time  $f_i$

$i$  and  $j$  are compatible if  $[s_i, f_i)$  and  $[s_j, f_j)$  are disjoint

**Output:** a maximum-size subset of mutually compatible jobs

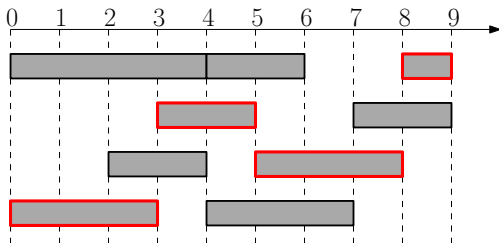


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## Weighted Interval Scheduling

**Input:**  $n$  jobs, job  $i$  with start time  $s_i$  and finish time  $f_i$

each job has a weight (or value)  $v_i > 0$

$i$  and  $j$  are compatible if  $[s_i, f_i)$  and  $[s_j, f_j)$  are disjoint

**Output:** a maximum-weight subset of mutually compatible jobs

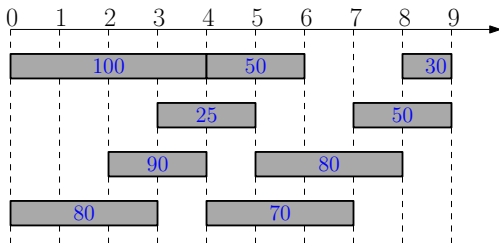
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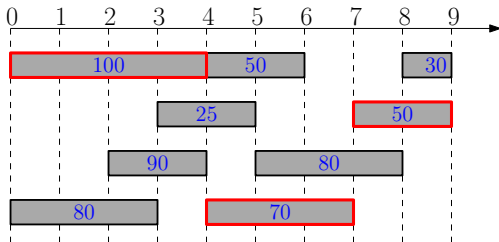
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Optimum value = 220

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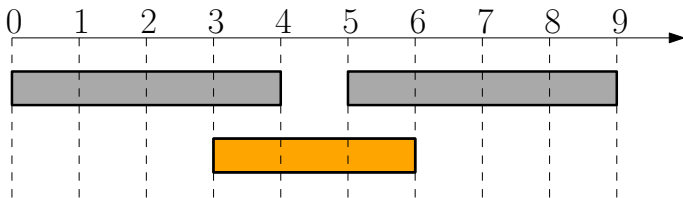
No, when weights are equal, this is the shortest job

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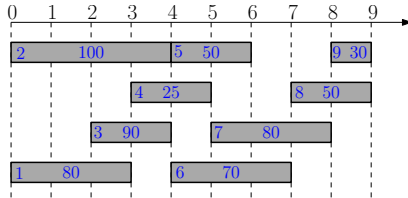
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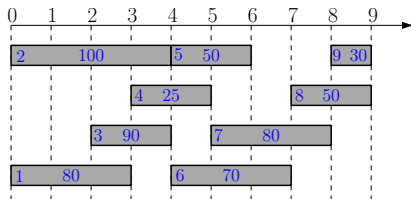
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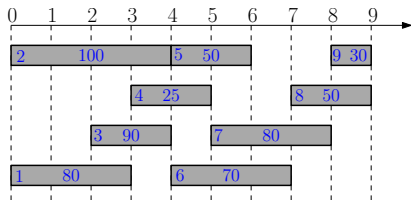
- Sort jobs according to non-decreasing order of finish times

# Designing a Dynamic Programming Algorithm



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- $opt[i]$ : optimal value for instance only containing jobs  $\{1, 2, \dots, i\}$

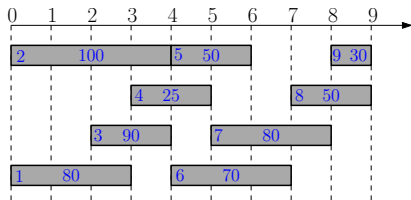
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$i$	$opt[i]$
0	
1	
2	
3	
4	
5	
6	
7	
8	
9	

# Designing a Dynamic Programming Algorithm

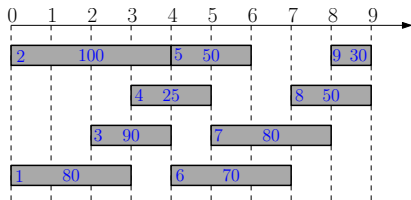


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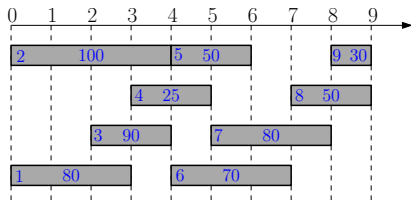
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- Sort jobs according to non-decreasing order of finish times
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$i$	$opt[i]$
0	0
1	80
2	
3	
4	
5	
6	
7	
8	
9	

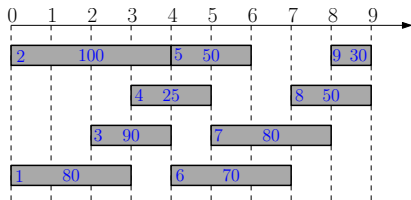
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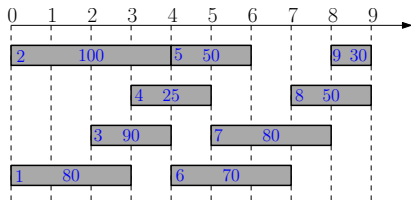
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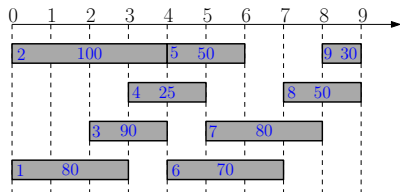
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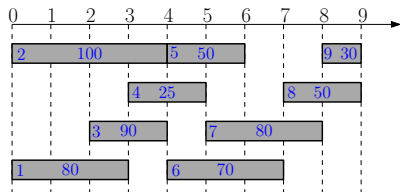
$i$	$opt[i]$
0	0
1	80
2	100
3	100
4	105
5	150
6	170
7	185
8	220
9	220

# Designing a Dynamic Programming Algorithm



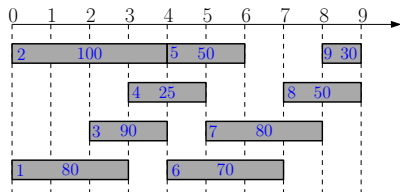
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# Designing a Dynamic Programming Algorithm



- Focus on instance  $\{1, 2, 3, \dots, i\}$ ,
- $opt[i]$ : optimal value for the instance
- assume we have computed  $opt[0], opt[1], \dots, opt[i - 1]$

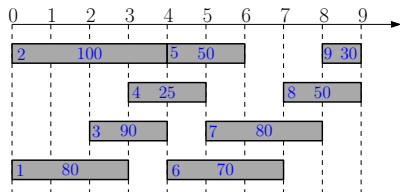
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**Q:** The value of optimal solution that **does not contain**  $i$ ?

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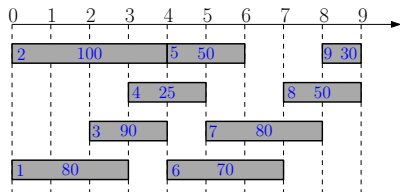
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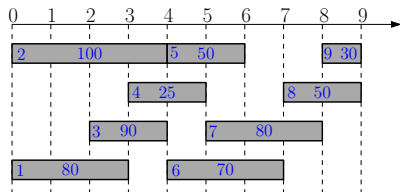
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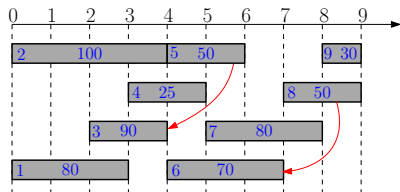
**A:**  $opt[i-1]$

**Q:** The value of optimal solution that **contains** job  $i$ ?

**A:**  $v_i + opt[p_i]$ ,

$p_i =$  the largest  $j$  such that  $f_j \leq s_i$

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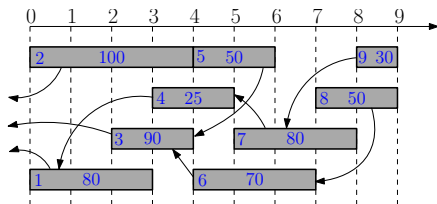
Recursion for  $opt[i]$ :

$$opt[i] = \max \{opt[i - 1], v_i + opt[p_i]\}$$

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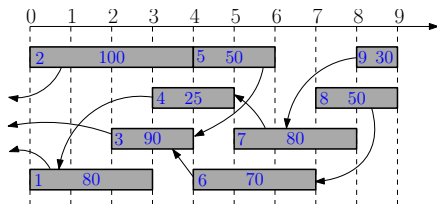


- $opt[0] = 0$
- $opt[1] = \max\{opt[0], 80 + opt[0]\} = 80$
- $opt[2] =$
- $opt[3] =$
- $opt[4] =$
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# Designing a Dynamic Programming Algorithm

Recursion for  $opt[i]$ :

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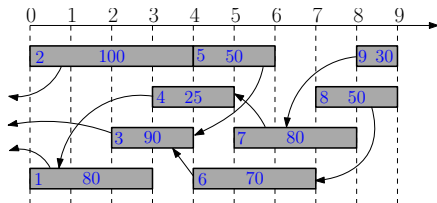


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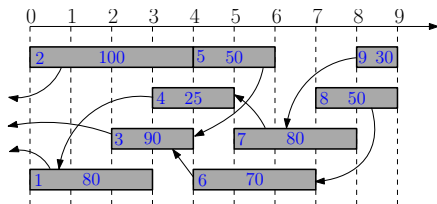
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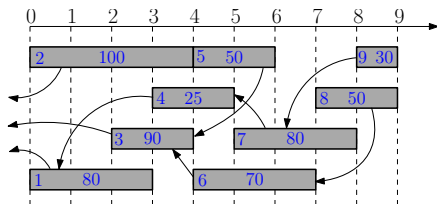


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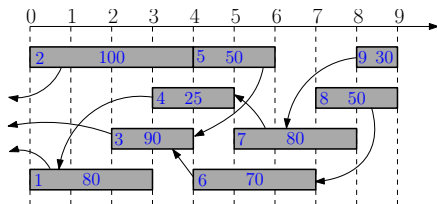


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- $opt[5] =$

# Designing a Dynamic Programming Algorithm

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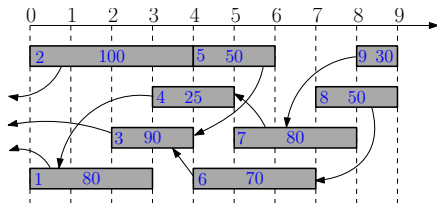


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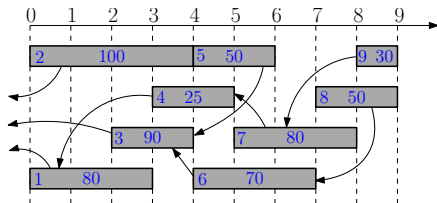


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- $opt[4] = \max\{opt[3], 25 + opt[1]\}$
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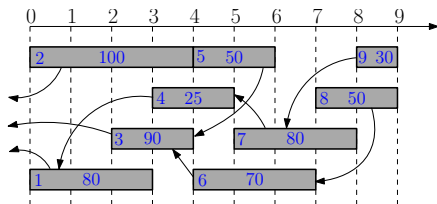


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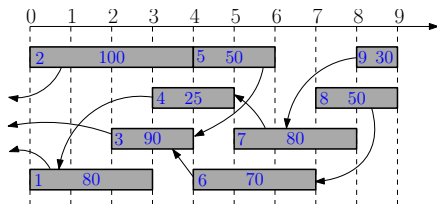


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- $opt[5] = \max\{opt[4], 50 + opt[3]\}$

# Designing a Dynamic Programming Algorithm

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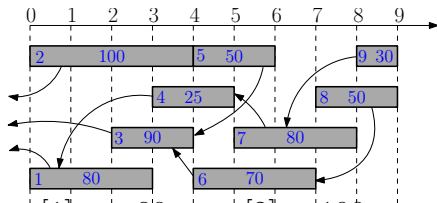


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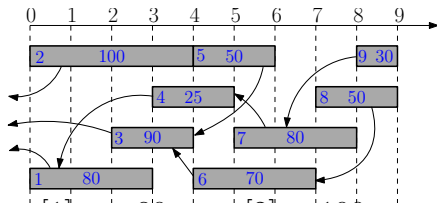
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- $opt[3] = 100, \quad opt[4] = 105, \quad opt[5] = 150$
- $opt[6] = \max\{opt[5], 70 + opt[3]\} = 170$
- $opt[7] = \max\{opt[6], 80 + opt[4]\} = 185$
- $opt[8] = \max\{opt[7], 50 + opt[6]\} = 220$
- $opt[9] = \max\{opt[8], 30 + opt[7]\} = 220$

# Dynamic Programming

- 1: sort jobs by non-decreasing order of finishing times
- 2: compute  $p_1, p_2, \dots, p_n$
- 3:  $opt[0] \leftarrow 0$
- 4: **for**  $i \leftarrow 1$  to  $n$  **do**
- 5:      $opt[i] \leftarrow \max\{opt[i - 1], v_i + opt[p_i]\}$

# Dynamic Programming

```
1: sort jobs by non-decreasing order of finishing times
2: compute  $p_1, p_2, \dots, p_n$ 
3:  $opt[0] \leftarrow 0$ 
4: for  $i \leftarrow 1$  to  $n$  do
5:    $opt[i] \leftarrow \max\{opt[i-1], v_i + opt[p_i]\}$ 
```

- Running time sorting:  $O(n \lg n)$
- Running time for computing  $p$ :  $O(n \lg n)$  via binary search
- Running time for computing  $opt[n]$ :  $O(n)$

# How Can We Recover the Optimum Schedule?

```
1: sort jobs by non-decreasing order of  
   finishing times  
2: compute  $p_1, p_2, \dots, p_n$   
3:  $opt[0] \leftarrow 0$   
4: for  $i \leftarrow 1$  to  $n$  do  
5:     if  $opt[i - 1] \geq v_i + opt[p_i]$  then  
6:        $opt[i] \leftarrow opt[i - 1]$   
7:  
8:     else  
9:        $opt[i] \leftarrow v_i + opt[p_i]$   
10:
```

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7:          $b[i] \leftarrow \text{N}$ 
8:     else
9:          $opt[i] \leftarrow v_i + opt[p_i]$ 
10:         $b[i] \leftarrow \text{Y}$ 
```

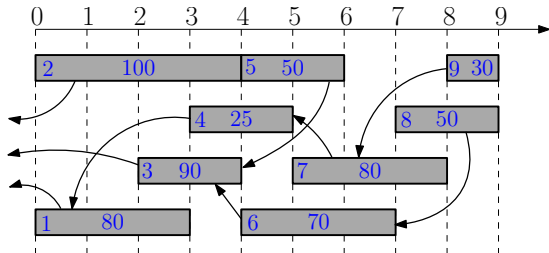
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8:     else
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10:         $b[i] \leftarrow Y$ 
```

```
1:  $i \leftarrow n, S \leftarrow \emptyset$ 
2: while  $i \neq 0$  do
3:     if  $b[i] = N$  then
4:          $i \leftarrow i - 1$ 
5:     else
6:          $S \leftarrow S \cup \{i\}$ 
7:          $i \leftarrow p_i$ 
8: return  $S$ 
```

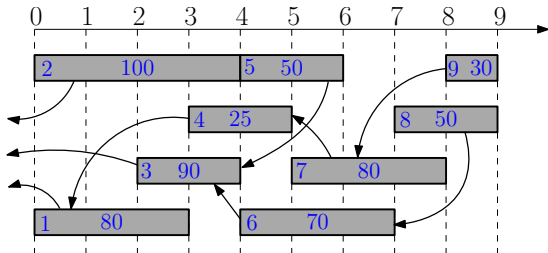
# Recovering Optimum Schedule: Example

$i$	$opt[i]$	$b[i]$
0	0	$\perp$
1	80	
2	100	
3	100	
4	105	
5	150	
6	170	
7	185	
8	220	
9	220	



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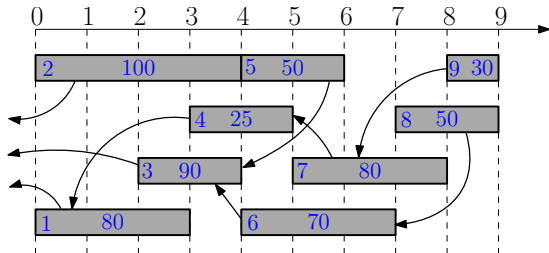
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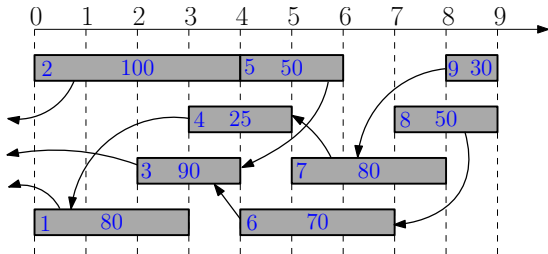
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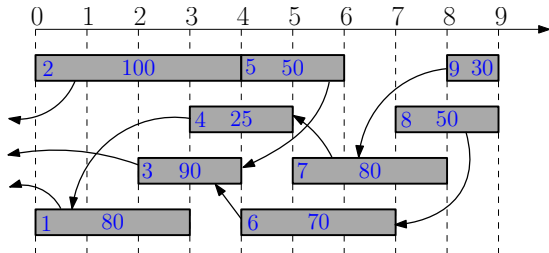
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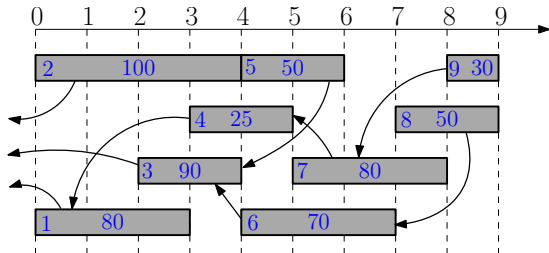
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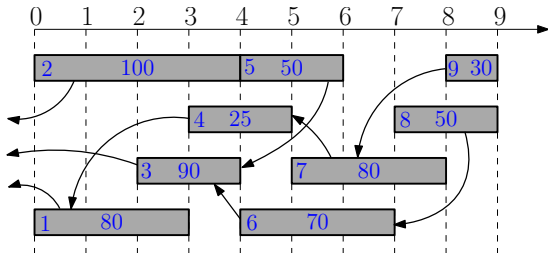
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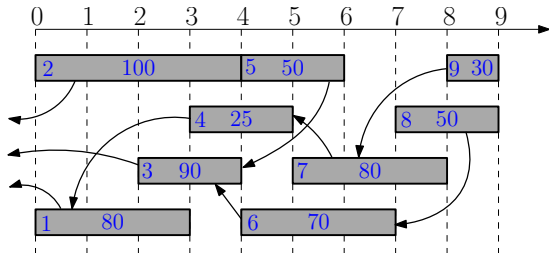
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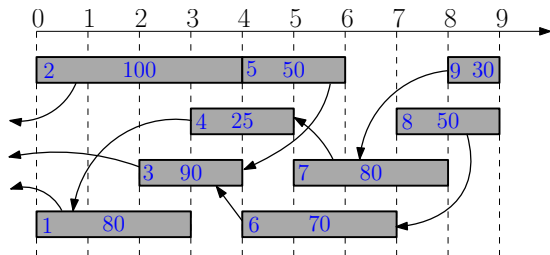
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7	185	Y
8	220	
9	220	



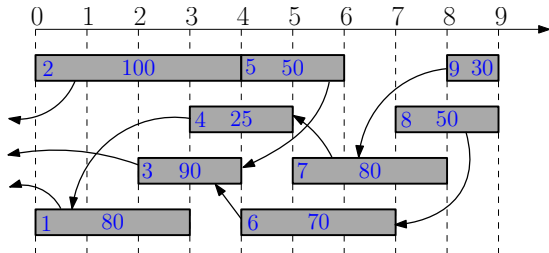
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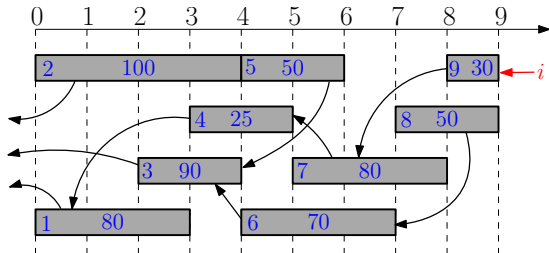
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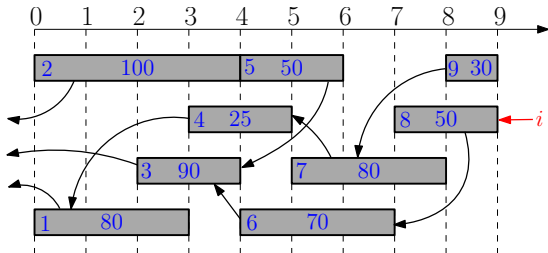
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6	170	Y
7	185	Y
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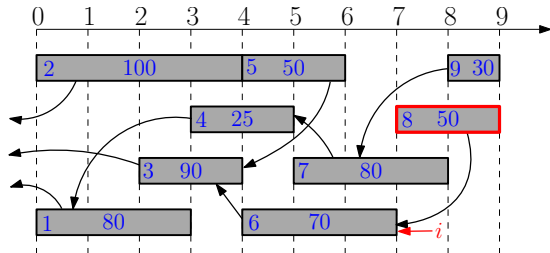
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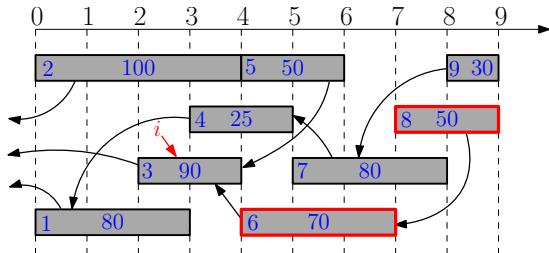
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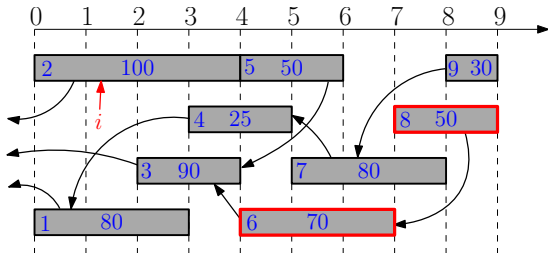
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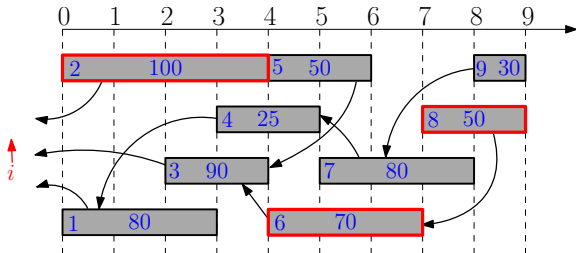
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8	220	Y
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# Dynamic Programming

- Break up a problem into many **overlapping** sub-problems
- Build solutions for larger and larger sub-problems
- Use a **table** to store solutions for sub-problems for reuse

# Outline

- 1 Weighted Interval Scheduling
- 2 Subset Sum Problem**
- 3 Knapsack Problem
- 4 Longest Common Subsequence
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## Subset Sum Problem

**Input:** an integer bound  $W > 0$

a set of  $n$  items, each with an integer weight  $w_i > 0$

**Output:** a subset  $S$  of items that

$$\text{maximizes } \sum_{i \in S} w_i \quad \text{s.t.} \quad \sum_{i \in S} w_i \leq W.$$

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### Example:

- $W = 35, n = 5, w = (14, 9, 17, 10, 13)$

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- Motivation: you have budget  $W$ , and want to buy a subset of items, so as to spend as much money as possible.

### Example:

- $W = 35, n = 5, w = (14, 9, 17, 10, 13)$
- Optimum:  $S = \{1, 2, 4\}$  and  $14 + 9 + 10 = 33$

# Greedy Algorithms for Subset Sum

## Candidate Algorithm:

- Sort according to non-increasing order of weights
- Select items in the order as long as the total weight remains below  $W$

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**A:** No.  $W = 100, n = 3, w = (51, 50, 50)$ .

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# Design a Dynamic Programming Algorithm

- Consider the instance:  $i, W', (w_1, w_2, \dots, w_i)$ ;
- $opt[i, W']$ : the optimum value of the instance

**Q:** The value of the optimum solution that **does not contain**  $i$ ?

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**Q:** The value of the optimum solution that **does not contain**  $i$ ?

**A:**  $opt[i - 1, W']$

**Q:** The value of the optimum solution that **contains**  $i$ ?

**A:**  $opt[i - 1, W' - w_i] + w_i$

# Dynamic Programming

- Consider the instance:  $i, W', (w_1, w_2, \dots, w_i)$ ;
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$$opt[i, W'] = \begin{cases} i = 0 \\ i > 0, w_i > W' \\ i > 0, w_i \leq W' \end{cases}$$

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# Dynamic Programming

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$$opt[i, W'] = \begin{cases} 0 & i = 0 \\ opt[i - 1, W'] & i > 0, w_i > W' \\ \max \left\{ \begin{array}{l} opt[i - 1, W'] \\ opt[i - 1, W' - w_i] + w_i \end{array} \right\} & i > 0, w_i \leq W' \end{cases}$$

# Dynamic Programming

```
1: for  $W' \leftarrow 0$  to  $W$  do  
2:    $opt[0, W'] \leftarrow 0$   
3: for  $i \leftarrow 1$  to  $n$  do  
4:   for  $W' \leftarrow 0$  to  $W$  do  
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$   
6:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$  then  
7:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$   
8: return  $opt[n, W]$ 
```

# Recover the Optimum Set

```
1: for  $W' \leftarrow 0$  to  $W$  do
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3: for  $i \leftarrow 1$  to  $n$  do
4:   for  $W' \leftarrow 0$  to  $W$  do
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$ 
6:      $b[i, W'] \leftarrow N$ 
7:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$ 
      then
8:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$ 
9:        $b[i, W'] \leftarrow Y$ 
10: return  $opt[n, W]$ 
```

# Recover the Optimum Set

```
1:  $i \leftarrow n, W' \leftarrow W, S \leftarrow \emptyset$   
2: while  $i > 0$  do  
3:   if  $b[i, W'] = Y$  then  
4:      $W' \leftarrow W' - w_i$   
5:      $S \leftarrow S \cup \{i\}$   
6:    $i \leftarrow i - 1$   
7: return  $S$ 
```

# Running Time of Algorithm

```
1: for  $W' \leftarrow 0$  to  $W$  do  
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8: return  $opt[n, W]$ 
```

- Running time is  $O(nW)$
- Running time is **pseudo-polynomial** because it depends on value of the input integers.

# Avoiding Unnecessary Computation and Memory Using Memoized Algorithm and Hash Map

**compute-opt( $i, W'$ )**

```
1: if  $opt[i, W'] \neq \perp$  then return  $opt[i, W']$ 
2: if  $i = 0$  then  $r \leftarrow 0$ 
3: else
4:    $r \leftarrow \text{compute-opt}(i - 1, W')$ 
5:   if  $w_i \leq W'$  then
6:      $r' \leftarrow \text{compute-opt}(i - 1, W' - w_i) + w_i$ 
7:     if  $r' > r$  then  $r \leftarrow r'$ 
8:    $opt[i, W'] \leftarrow r$ 
9: return  $r$ 
```

- Use hash map for  $opt$



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## Knapsack Problem

**Input:** an integer bound  $W > 0$

a set of  $n$  items, each with an integer weight  $w_i > 0$

a value  $v_i > 0$  for each item  $i$

**Output:** a subset  $S$  of items that

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- Motivation: you have budget  $W$ , and want to buy a subset of items of maximum total value

# DP for Knapsack Problem

- $opt[i, W']$ : the optimum value when budget is  $W'$  and items are  $\{1, 2, 3, \dots, i\}$ .
- If  $i = 0$ ,  $opt[i, W'] = 0$  for every  $W' = 0, 1, 2, \dots, W$ .

$$opt[i, W'] = \begin{cases} i = 0 \\ i > 0, w_i > W' \\ i > 0, w_i \leq W' \end{cases}$$

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$$opt[i, W'] = \begin{cases} 0 & i = 0 \\ opt[i - 1, W'] & i > 0, w_i > W' \\ & i > 0, w_i \leq W' \end{cases}$$

# DP for Knapsack Problem

- $opt[i, W']$ : the optimum value when budget is  $W'$  and items are  $\{1, 2, 3, \dots, i\}$ .
- If  $i = 0$ ,  $opt[i, W'] = 0$  for every  $W' = 0, 1, 2, \dots, W$ .

$$opt[i, W'] = \begin{cases} 0 & i = 0 \\ opt[i - 1, W'] & i > 0, w_i > W' \\ \max \left\{ \begin{array}{l} opt[i - 1, W'] \\ opt[i - 1, W' - w_i] + v_i \end{array} \right\} & i > 0, w_i \leq W' \end{cases}$$

## Exercise: Items with 3 Parameters

**Input:** integer bounds  $W > 0$ ,  $Z > 0$ ,  
a set of  $n$  items, each with an integer weight  $w_i > 0$   
a size  $z_i > 0$  for each item  $i$   
a value  $v_i > 0$  for each item  $i$

**Output:** a subset  $S$  of items that

$$\begin{aligned} &\text{maximizes } \sum_{i \in S} v_i \quad \text{s.t.} \\ &\sum_{i \in S} w_i \leq W \text{ and } \sum_{i \in S} z_i \leq Z \end{aligned}$$



# Outline

- 1 Weighted Interval Scheduling
- 2 Subset Sum Problem
- 3 Knapsack Problem
- 4 Longest Common Subsequence
  - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree
- 8 Summary

# Subsequence

- $A = bacdca$
- $C = adca$

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- $A = bacdca$
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**Def.** Given two sequences  $A[1 .. n]$  and  $C[1 .. t]$  of letters,  $C$  is called a **subsequence** of  $A$  if there exists integers  $1 \leq i_1 < i_2 < i_3 < \dots < i_t \leq n$  such that  $A[i_j] = C[j]$  for every  $j = 1, 2, 3, \dots, t$ .

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- Exercise: how to check if sequence  $C$  is a subsequence of  $A$ ?

## Longest Common Subsequence

**Input:**  $A[1 \dots n]$  and  $B[1 \dots m]$

**Output:** the longest common subsequence of  $A$  and  $B$

### Example:

- $A = \text{'bacdca'}$
- $B = \text{'adbcda'}$

## Longest Common Subsequence

**Input:**  $A[1 \dots n]$  and  $B[1 \dots m]$

**Output:** the longest common subsequence of  $A$  and  $B$

### Example:

- $A = 'bacdca'$
- $B = 'adbcdca'$
- $LCS(A, B) = 'adca'$

## Longest Common Subsequence

**Input:**  $A[1 \dots n]$  and  $B[1 \dots m]$

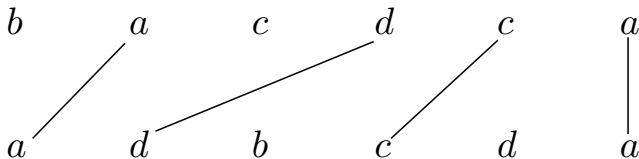
**Output:** the longest common subsequence of  $A$  and  $B$

### Example:

- $A = 'bacdca'$
  - $B = 'adbcdca'$
  - $LCS(A, B) = 'adca'$
- Applications: edit distance (diff), similarity of DNAs



# Matching View of LCS



- Goal of LCS: find a maximum-size non-crossing matching between letters in  $A$  and letters in  $B$ .

# Reduce to Subproblems

- $A = \text{'bacdca'}$
- $B = \text{'adbcdca'}$

# Reduce to Subproblems

- $A = \text{'bacdca'}$
- $B = \text{'adbcdca'}$

# Reduce to Subproblems

- $A = \text{'bacdc'}$
- $B = \text{'adbcd'}$

# Reduce to Subproblems

- $A = 'bacdc'$
- $B = 'adbcd'$
- either the last letter of  $A$  is not matched:
- or the last letter of  $B$  is not matched:

# Reduce to Subproblems

- $A = \text{'bacdc'}$
- $B = \text{'adbcd'}$
- either the last letter of  $A$  is not matched:
- need to compute  $\text{LCS}(\text{'bacd'}, \text{'adbcd'})$
- or the last letter of  $B$  is not matched:

# Reduce to Subproblems

- $A = \text{'bacdc'}$
- $B = \text{'adbcd'}$
- either the last letter of  $A$  is not matched:
  - need to compute  $\text{LCS}(\text{'bacd'}, \text{'adbcd'})$
- or the last letter of  $B$  is not matched:
  - need to compute  $\text{LCS}(\text{'bacdc'}, \text{'adbc'})$

# Dynamic Programming for LCS

- $opt[i, j], 0 \leq i \leq n, 0 \leq j \leq m$ : length of longest common sub-sequence of  $A[1 .. i]$  and  $B[1 .. j]$ .



# Dynamic Programming for LCS

- $opt[i, j], 0 \leq i \leq n, 0 \leq j \leq m$ : length of longest common sub-sequence of  $A[1 .. i]$  and  $B[1 .. j]$ .
- if  $i = 0$  or  $j = 0$ , then  $opt[i, j] = 0$ .

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- if  $i > 0, j > 0$ , then

$$opt[i, j] = \begin{cases} & \text{if } A[i] = B[j] \\ & \text{if } A[i] \neq B[j] \end{cases}$$

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- if  $i > 0, j > 0$ , then

$$opt[i, j] = \begin{cases} opt[i - 1, j - 1] + 1 & \text{if } A[i] = B[j] \\ & \text{if } A[i] \neq B[j] \end{cases}$$

# Dynamic Programming for LCS

- $opt[i, j], 0 \leq i \leq n, 0 \leq j \leq m$ : length of longest common sub-sequence of  $A[1 .. i]$  and  $B[1 .. j]$ .
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- if  $i > 0, j > 0$ , then

$$opt[i, j] = \begin{cases} opt[i - 1, j - 1] + 1 & \text{if } A[i] = B[j] \\ \max \begin{cases} opt[i - 1, j] \\ opt[i, j - 1] \end{cases} & \text{if } A[i] \neq B[j] \end{cases}$$

# Dynamic Programming for LCS

```
1: for  $j \leftarrow 0$  to  $m$  do  
2:    $opt[0, j] \leftarrow 0$   
3: for  $i \leftarrow 1$  to  $n$  do  
4:    $opt[i, 0] \leftarrow 0$   
5:   for  $j \leftarrow 1$  to  $m$  do  
6:     if  $A[i] = B[j]$  then  
7:        $opt[i, j] \leftarrow opt[i - 1, j - 1] + 1$   
8:     else if  $opt[i, j - 1] \geq opt[i - 1, j]$  then  
9:        $opt[i, j] \leftarrow opt[i, j - 1]$   
10:    else  
11:       $opt[i, j] \leftarrow opt[i - 1, j]$ 
```

# Dynamic Programming for LCS

```
1: for  $j \leftarrow 0$  to  $m$  do
2:    $opt[0, j] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:    $opt[i, 0] \leftarrow 0$ 
5:   for  $j \leftarrow 1$  to  $m$  do
6:     if  $A[i] = B[j]$  then
7:        $opt[i, j] \leftarrow opt[i - 1, j - 1] + 1$ ,  $\pi[i, j] \leftarrow \text{"↖"}$ 
8:     else if  $opt[i, j - 1] \geq opt[i - 1, j]$  then
9:        $opt[i, j] \leftarrow opt[i, j - 1]$ ,  $\pi[i, j] \leftarrow \text{"←"}$ 
10:    else
11:       $opt[i, j] \leftarrow opt[i - 1, j]$ ,  $\pi[i, j] \leftarrow \text{"↑"}$ 
```

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥						
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←					
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						



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	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$				
2	0 $\perp$						
3	0 $\perp$						
4	0 $\perp$						
5	0 $\perp$						
6	0 $\perp$						

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	1	2	3	4	5	6
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖			
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nwarrow$	1 $\leftarrow$		
2	0 $\perp$						
3	0 $\perp$						
4	0 $\perp$						
5	0 $\perp$						
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	
2	0 $\perp$						
3	0 $\perp$						
4	0 $\perp$						
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0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nwarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$						
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nwarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$						
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0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nwarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nwarrow$					
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nwarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nwarrow$	1 $\leftarrow$				
3	0 $\perp$						
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<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$			
3	0 $\perp$						
4	0 $\perp$						
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nwarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nwarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$		
3	0 $\perp$						
4	0 $\perp$						
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	
3	0 $\perp$						
4	0 $\perp$						
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$						
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$						
4	0 $\perp$						
5	0 $\perp$						
6	0 $\perp$						

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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$	1 $\uparrow$					
4	0 $\perp$						
5	0 $\perp$						
6	0 $\perp$						

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	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$				
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$	2 $\leftarrow$	
4	0 $\perp$						
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0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$						
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# Example

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<i>A</i>	b	a	c	d	c	a
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	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$						
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<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑					
5	0 ⊥						
6	0 ⊥						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$				
5	0 $\perp$						
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$			
5	0 $\perp$						
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$		
5	0 $\perp$						
6	0 $\perp$						



# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	
5	0 ⊥						
6	0 ⊥						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$						
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$						
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$					
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑				
6	0 ⊥						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$			
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$		
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	
6	0 $\perp$						



# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$						

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$	1 $\swarrow$					

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$	1 $\swarrow$	2 $\uparrow$				

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$	1 $\swarrow$	2 $\uparrow$	2 $\leftarrow$			

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$	1 $\swarrow$	2 $\uparrow$	2 $\leftarrow$	3 $\uparrow$		

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\swarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\swarrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\swarrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$	1 $\swarrow$	2 $\uparrow$	2 $\leftarrow$	3 $\uparrow$	3 $\leftarrow$	

# Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$	0 $\perp$
1	0 $\perp$	0 $\leftarrow$	0 $\leftarrow$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$
2	0 $\perp$	1 $\nearrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$
3	0 $\perp$	1 $\uparrow$	1 $\leftarrow$	1 $\leftarrow$	2 $\nearrow$	2 $\leftarrow$	2 $\leftarrow$
4	0 $\perp$	1 $\uparrow$	2 $\nearrow$	2 $\leftarrow$	2 $\leftarrow$	3 $\nearrow$	3 $\leftarrow$
5	0 $\perp$	1 $\uparrow$	2 $\uparrow$	2 $\leftarrow$	3 $\nearrow$	3 $\leftarrow$	3 $\leftarrow$
6	0 $\perp$	1 $\nearrow$	2 $\uparrow$	2 $\leftarrow$	3 $\uparrow$	3 $\leftarrow$	4 $\nearrow$



## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

# Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖



## Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

# Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
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6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

# Find Common Subsequence

```
1:  $i \leftarrow n, j \leftarrow m, S \leftarrow ()$ 
2: while  $i > 0$  and  $j > 0$  do
3:   if  $\pi[i, j] = "\nwarrow"$  then
4:     add  $A[i]$  to beginning of  $S, i \leftarrow i - 1, j \leftarrow j - 1$ 
5:   else if  $\pi[i, j] = "\uparrow"$  then
6:      $i \leftarrow i - 1$ 
7:   else
8:      $j \leftarrow j - 1$ 
9: return  $S$ 
```

# Variants of Problem

## Edit Distance with Insertions and Deletions

**Input:** a string  $A$

each time we can delete a letter from  $A$  or insert a letter to  $A$

**Output:** minimum number of operations (insertions or deletions) we need to change  $A$  to  $B$ ?

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**Obs.**  $\#OPs = \text{length}(A) + \text{length}(B) - 2 \cdot \text{length}(\text{LCS}(A, B))$

# Variants of Problem

## Edit Distance with Insertions, Deletions and Replacing

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### Example:

- $A = \text{ocurrance}$ ,  $B = \text{occurrence}$ .
- 2 operations: insert 'c', change 'a' to 'e'
- Not related to LCS any more

# Edit Distance (with Replacing)

- $opt[i, j], 0 \leq i \leq n, 0 \leq j \leq m$ : edit distance between  $A[1 \dots i]$  and  $B[1 \dots j]$ .

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## Exercise: Longest Palindrome

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**Input:** a sequence  $A$

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- Input: acbcedeacab

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## Example:

- Input: **a**cb**ced**e**a**cb
- Output: acedeca

# Outline

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# Computing the Length of LCS

```
1: for  $j \leftarrow 0$  to  $m$  do
2:    $opt[0, j] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:    $opt[i, 0] \leftarrow 0$ 
5:   for  $j \leftarrow 1$  to  $m$  do
6:     if  $A[i] = B[j]$  then
7:        $opt[i, j] \leftarrow opt[i - 1, j - 1] + 1$ 
8:     else if  $opt[i, j - 1] \geq opt[i - 1, j]$  then
9:        $opt[i, j] \leftarrow opt[i, j - 1]$ 
10:    else
11:       $opt[i, j] \leftarrow opt[i - 1, j]$ 
```

**Obs.** The  $i$ -th row of table only depends on  $(i - 1)$ -th row.

# Reducing Space to $O(n + m)$

**Obs.** The  $i$ -th row of table only depends on  $(i - 1)$ -th row.

**Q:** How to use this observation to reduce space?

## Reducing Space to $O(n + m)$

**Obs.** The  $i$ -th row of table only depends on  $(i - 1)$ -th row.

**Q:** How to use this observation to reduce space?

**A:** We only keep two rows: the  $(i - 1)$ -th row and the  $i$ -th row.

# Linear Space Algorithm to Compute Length of LCS

```
1: for  $j \leftarrow 0$  to  $m$  do
2:    $opt[0, j] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:    $opt[i \bmod 2, 0] \leftarrow 0$ 
5:   for  $j \leftarrow 1$  to  $m$  do
6:     if  $A[i] = B[j]$  then
7:        $opt[i \bmod 2, j] \leftarrow opt[i - 1 \bmod 2, j - 1] + 1$ 
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12: return  $opt[n \bmod 2, m]$ 
```



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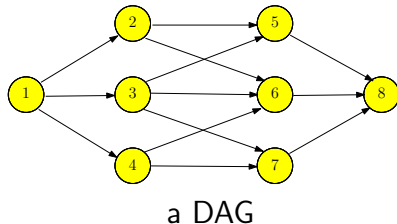
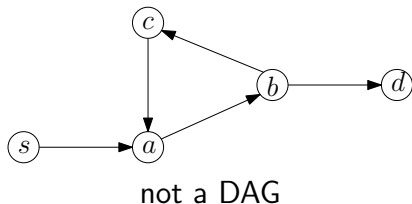
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  - Time:  $O(nm)$

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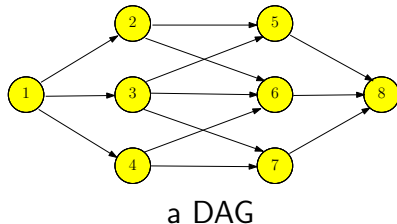
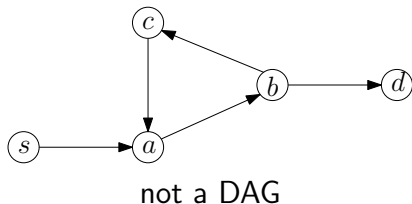
# Directed Acyclic Graphs

**Def.** A directed acyclic graph (DAG) is a directed graph without (directed) cycles.



# Directed Acyclic Graphs

**Def.** A directed acyclic graph (DAG) is a directed graph without (directed) cycles.



**Lemma** A directed graph is a DAG if and only if its vertices can be topologically sorted.

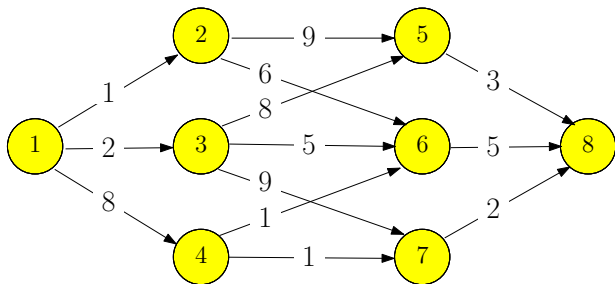


## Shortest Paths in DAG

**Input:** directed acyclic graph  $G = (V, E)$  and  $w : E \rightarrow \mathbb{R}$ .

Assume  $V = \{1, 2, 3, \dots, n\}$  is topologically sorted: if  $(i, j) \in E$ , then  $i < j$

**Output:** the shortest path from 1 to  $i$ , for every  $i \in V$

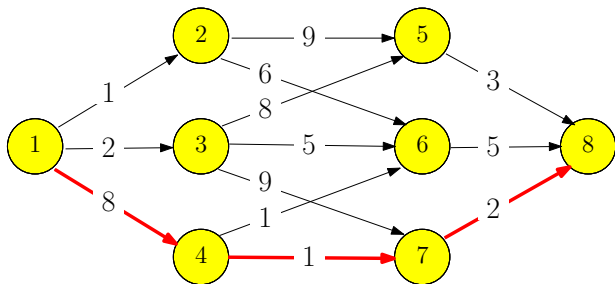


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# Shortest Paths in DAG

- $f[i]$ : length of the shortest path from 1 to  $i$

$$f[i] = \begin{cases} & i = 1 \\ & i = 2, 3, \dots, n \end{cases}$$

# Shortest Paths in DAG

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# Shortest Paths in DAG

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$$f[i] = \begin{cases} 0 & i = 1 \\ \min_{j:(j,i) \in E} \{f(j) + w(j,i)\} & i = 2, 3, \dots, n \end{cases}$$

# Shortest Paths in DAG

- Use an adjacency list for incoming edges of each vertex  $i$

## Shortest Paths in DAG

```
1:  $f[1] \leftarrow 0$ 
2: for  $i \leftarrow 2$  to  $n$  do
3:    $f[i] \leftarrow \infty$ 
4:   for each incoming edge  $(j, i)$  of  $i$  do
5:     if  $f[j] + w(j, i) < f[i]$  then
6:        $f[i] \leftarrow f[j] + w(j, i)$ 
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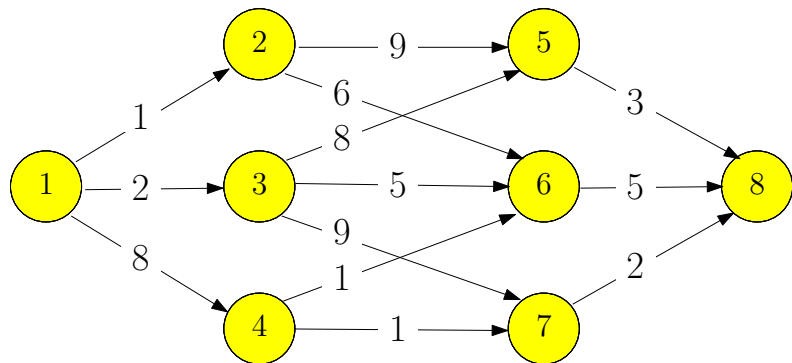
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## print-path( $t$ )

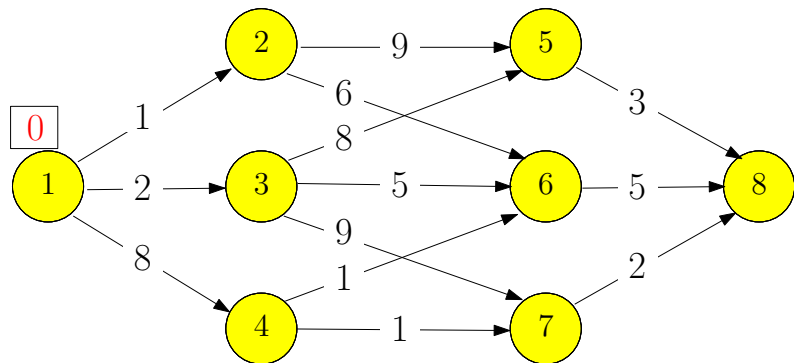
```
1: if  $t = 1$  then
2:   print(1)
3:   return
4: print-path( $\pi(t)$ )
5: print(", ",  $t$ )
```



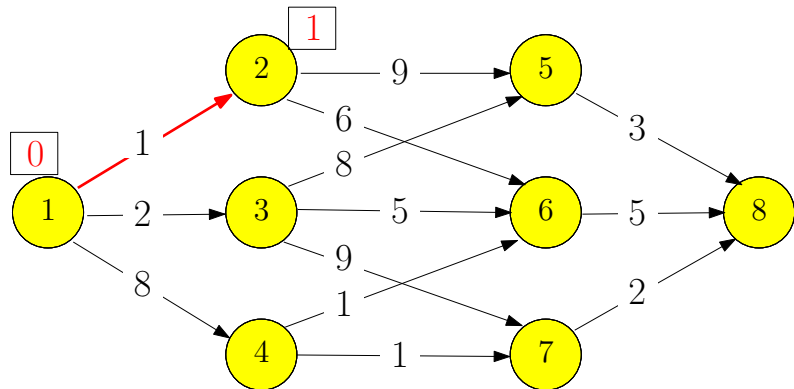
# Example



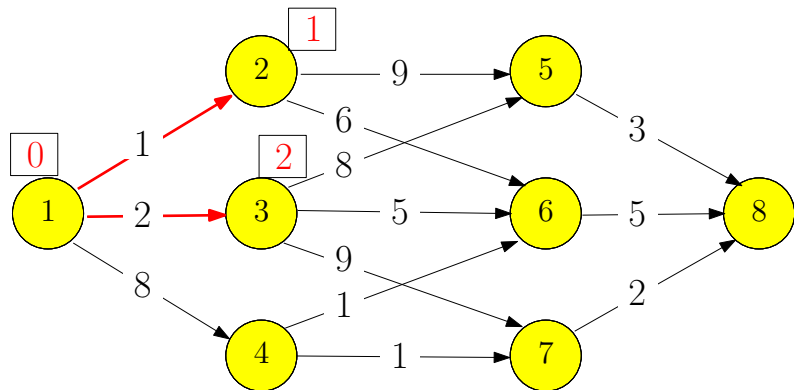
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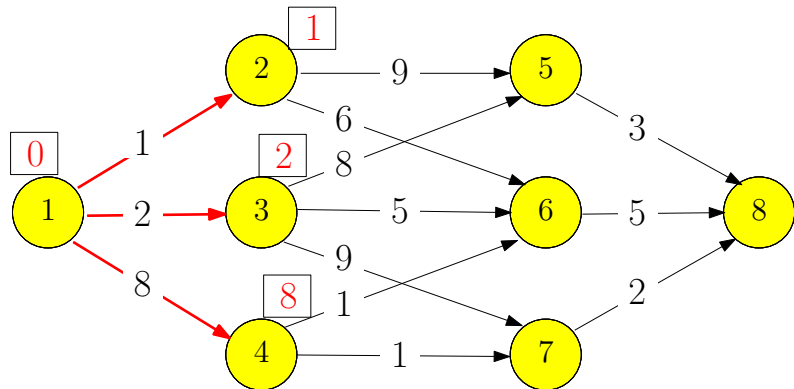
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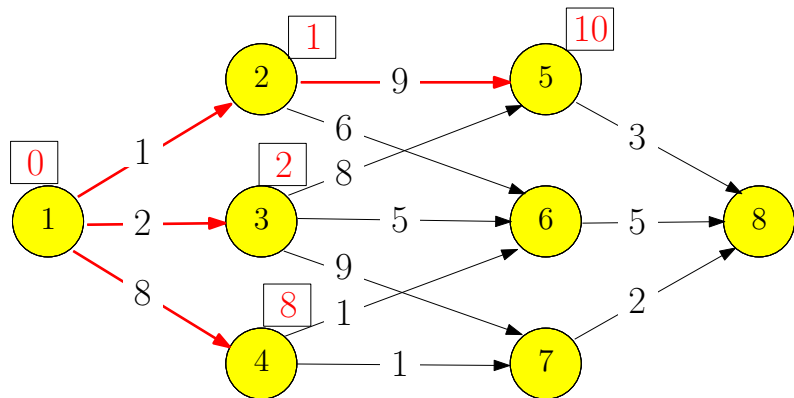
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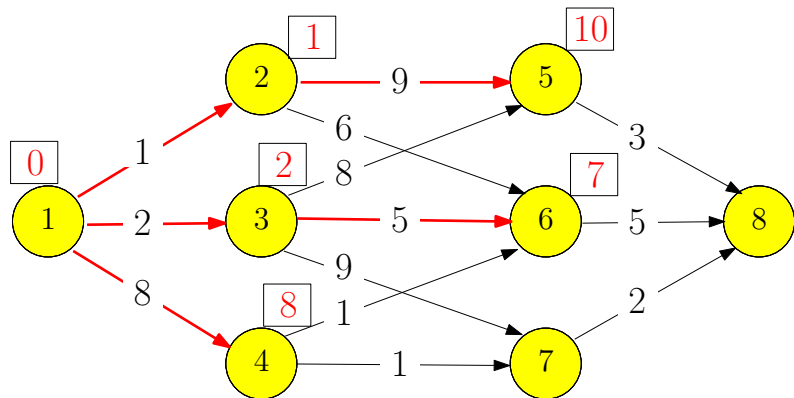
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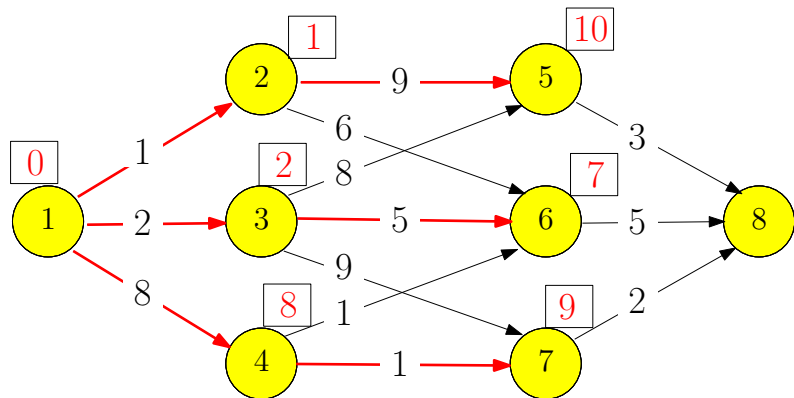
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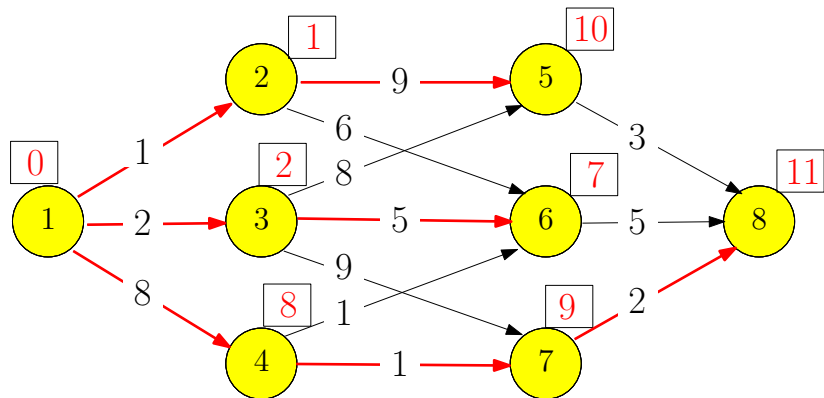


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# Variant: Heaviest Path in a Directed Acyclic Graph

## Heaviest Path in a Directed Acyclic Graph

**Input:** directed acyclic graph  $G = (V, E)$  and  $w : E \rightarrow \mathbb{R}$ .

Assume  $V = \{1, 2, 3, \dots, n\}$  is topologically sorted: if  $(i, j) \in E$ , then  $i < j$

**Output:** the path with the **largest** weight (the **heaviest** path) from 1 to  $n$ .

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$$f[i] = \begin{cases} & i = 1 \\ & i = 2, 3, \dots, n \end{cases}$$

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# Matrix Chain Multiplication

## Matrix Chain Multiplication

**Input:**  $n$  matrices  $A_1, A_2, \dots, A_n$  of sizes

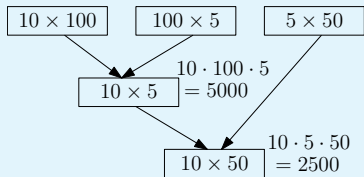
$r_1 \times c_1, r_2 \times c_2, \dots, r_n \times c_n$ , such that  $c_i = r_{i+1}$  for every  $i = 1, 2, \dots, n - 1$ .

**Output:** the order of computing  $A_1 A_2 \dots A_n$  with the minimum number of multiplications

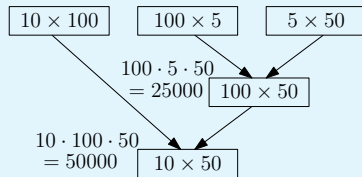
**Fact** Multiplying two matrices of size  $r \times k$  and  $k \times c$  takes  $r \times k \times c$  multiplications.

## Example:

- $A_1 : 10 \times 100$ ,  $A_2 : 100 \times 5$ ,  $A_3 : 5 \times 50$



$$\text{cost} = 5000 + 2500 = 7500$$

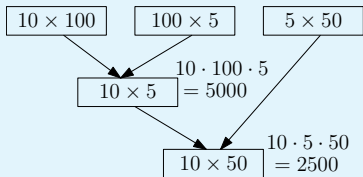


$$\text{cost} = 25000 + 50000 = 75000$$

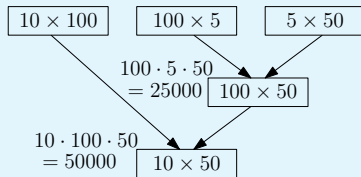
- $(A_1 A_2) A_3: 10 \times 100 \times 5 + 10 \times 5 \times 50 = 7500$
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- Cost of last step:  $r_1 \times c_i \times c_n$
- Optimality for sub-instances: we need to compute  $A_1 A_2 \cdots A_i$  and  $A_{i+1} A_{i+2} \cdots A_n$  optimally

# Matrix Chain Multiplication: Design DP

- Assume the last step is  $(A_1 A_2 \cdots A_i)(A_{i+1} A_{i+2} \cdots A_n)$
- Cost of last step:  $r_1 \times c_i \times c_n$
- Optimality for sub-instances: we need to compute  $A_1 A_2 \cdots A_i$  and  $A_{i+1} A_{i+2} \cdots A_n$  optimally
- $opt[i, j]$  : the minimum cost of computing  $A_i A_{i+1} \cdots A_j$

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- $opt[i, j]$  : the minimum cost of computing  $A_i A_{i+1} \cdots A_j$

$$opt[i, j] = \begin{cases} 0 & i = j \\ \min_{k: i \leq k < j} (opt[i, k] + opt[k + 1, j] + r_i c_k c_j) & i < j \end{cases}$$

# Matrix Chain Multiplication: Design DP

**matrix-chain-multiplication( $n, r[1..n], c[1..n]$ )**

```
1: let  $opt[i, i] \leftarrow 0$  for every  $i = 1, 2, \dots, n$ 
2: for  $\ell \leftarrow 2$  to  $n$  do
3:   for  $i \leftarrow 1$  to  $n - \ell + 1$  do
4:      $j \leftarrow i + \ell - 1$ 
5:      $opt[i, j] \leftarrow \infty$ 
6:     for  $k \leftarrow i$  to  $j - 1$  do
7:       if  $opt[i, k] + opt[k + 1, j] + r_i c_k c_j < opt[i, j]$  then
8:          $opt[i, j] \leftarrow opt[i, k] + opt[k + 1, j] + r_i c_k c_j$ 
9: return  $opt[1, n]$ 
```

# Recover the Optimum Way of Multiplication

**matrix-chain-multiplication( $n, r[1..n], c[1..n]$ )**

```
1: let  $opt[i, i] \leftarrow 0$  for every  $i = 1, 2, \dots, n$ 
2: for  $\ell \leftarrow 2$  to  $n$  do
3:   for  $i \leftarrow 1$  to  $n - \ell + 1$  do
4:      $j \leftarrow i + \ell - 1$ 
5:      $opt[i, j] \leftarrow \infty$ 
6:     for  $k \leftarrow i$  to  $j - 1$  do
7:       if  $opt[i, k] + opt[k + 1, j] + r_i c_k c_j < opt[i, j]$  then
8:          $opt[i, j] \leftarrow opt[i, k] + opt[k + 1, j] + r_i c_k c_j$ 
9:          $\pi[i, j] \leftarrow k$ 
10: return  $opt[1, n]$ 
```



# Constructing Optimal Solution

## Print-Optimal-Order( $i, j$ )

```
1: if  $i = j$  then  
2:   print( "A" $i$  )  
3: else  
4:   print( "(" )  
5:   Print-Optimal-Order( $i, \pi[i, j]$  )  
6:   Print-Optimal-Order( $\pi[i, j] + 1, j$  )  
7:   print( ")" )
```

<b>matrix</b>	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
<b>size</b>	$3 \times 5$	$5 \times 2$	$2 \times 6$	$6 \times 9$	$9 \times 4$

$$opt[1, 2] = opt[1, 1] + opt[2, 2] + 3 \times 5 \times 2 = 30, \quad \pi[1, 2] = 1$$

$$opt[2, 3] = opt[2, 2] + opt[3, 3] + 5 \times 2 \times 6 = 60, \quad \pi[2, 3] = 2$$

$$opt[3, 4] = opt[3, 3] + opt[4, 4] + 2 \times 6 \times 9 = 108, \quad \pi[3, 4] = 3$$

$$opt[4, 5] = opt[4, 4] + opt[5, 5] + 6 \times 9 \times 4 = 216, \quad \pi[4, 5] = 4$$

$$\begin{aligned} opt[1, 3] &= \min\{opt[1, 1] + opt[2, 3] + 3 \times 5 \times 6, \\ &\quad opt[1, 2] + opt[3, 3] + 3 \times 2 \times 6\} \\ &= \min\{0 + 60 + 90, 30 + 0 + 36\} = 66, \quad \pi[1, 3] = 2 \end{aligned}$$

$$\begin{aligned} opt[2, 4] &= \min\{opt[2, 2] + opt[3, 4] + 5 \times 2 \times 9, \\ &\quad opt[2, 3] + opt[4, 4] + 5 \times 6 \times 9\} \\ &= \min\{0 + 108 + 90, 60 + 0 + 270\} = 198, \quad \pi[2, 4] = 2, \end{aligned}$$

<b>matrix</b>	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
<b>size</b>	$3 \times 5$	$5 \times 2$	$2 \times 6$	$6 \times 9$	$9 \times 4$

$$\begin{aligned}
 \text{opt}[3, 5] &= \min\{\text{opt}[3, 3] + \text{opt}[4, 5] + 2 \times 6 \times 4, \\
 &\quad \text{opt}[3, 4] + \text{opt}[5, 5] + 2 \times 9 \times 4\} \\
 &= \min\{0 + 216 + 48, 108 + 0 + 72\} = 180,
 \end{aligned}$$

$$\pi[3, 5] = 4,$$

$$\begin{aligned}
 \text{opt}[1, 4] &= \min\{\text{opt}[1, 1] + \text{opt}[2, 4] + 3 \times 5 \times 9, \\
 &\quad \text{opt}[1, 2] + \text{opt}[3, 4] + 3 \times 2 \times 9, \\
 &\quad \text{opt}[1, 3] + \text{opt}[4, 4] + 3 \times 6 \times 9\} \\
 &= \min\{0 + 198 + 135, 30 + 108 + 54, 66 + 0 + 162\} = 192,
 \end{aligned}$$

$$\pi[1, 4] = 2,$$

<b>matrix</b>	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
<b>size</b>	$3 \times 5$	$5 \times 2$	$2 \times 6$	$6 \times 9$	$9 \times 4$

$$\begin{aligned}
 \text{opt}[2, 5] &= \min\{\text{opt}[2, 2] + \text{opt}[3, 5] + 5 \times 2 \times 4, \\
 &\quad \text{opt}[2, 3] + \text{opt}[4, 5] + 5 \times 6 \times 4, \\
 &\quad \text{opt}[2, 4] + \text{opt}[5, 5] + 5 \times 9 \times 4\} \\
 &= \min\{0 + 180 + 40, 60 + 216 + 120, 198 + 0 + 180\} = 220,
 \end{aligned}$$

$$\begin{aligned}
 \text{opt}[1, 5] &= \min\{\text{opt}[1, 1] + \text{opt}[2, 5] + 3 \times 5 \times 4, \\
 &\quad \text{opt}[1, 2] + \text{opt}[3, 5] + 3 \times 2 \times 4, \\
 &\quad \text{opt}[1, 3] + \text{opt}[4, 5] + 3 \times 6 \times 4, \\
 &\quad \text{opt}[1, 4] + \text{opt}[5, 5] + 3 \times 9 \times 4\} \\
 &= \min\{0 + 220 + 60, 30 + 180 + 24, \\
 &\quad 66 + 216 + 72, 192 + 0 + 108\} \\
 &= 234,
 \end{aligned}$$

$$\pi[1, 5] = 2.$$

<b>matrix</b>	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
<b>size</b>	$3 \times 5$	$5 \times 2$	$2 \times 6$	$6 \times 9$	$9 \times 4$

$opt, \pi$	$j = 1$	$j = 2$	$j = 3$	$j = 4$	$j = 5$
$i = 1$	0, /	30, 1	66, 2	192, 2	234, 2
$i = 2$		0, /	60, 2	198, 2	220, 2
$i = 3$			0, /	108, 3	180, 4
$i = 4$				0, /	216, 4
$i = 5$					0, /

$opt, \pi$	$j = 1$	$j = 2$	$j = 3$	$j = 4$	$j = 5$
$i = 1$	0, /	30, 1	66, 2	192, 2	234, 2
$i = 2$		0, /	60, 2	198, 2	220, 2
$i = 3$			0, /	108, 3	180, 4
$i = 4$				0, /	216, 4
$i = 5$					0, /

Print-Optimal-Order(1,5)

    Print-Optimal-Order(1, 2)

        Print-Optimal-Order(1, 1)

        Print-Optimal-Order(2, 2)

    Print-Optimal-Order(3, 5)

        Print-Optimal-Order(3, 4)

            Print-Optimal-Order(3, 3)

            Print-Optimal-Order(4, 4)

        Print-Optimal-Order(5, 5)

Optimum way for multiplication:  $((A_1 A_2)((A_3 A_4) A_5))$

# Outline

- 1 Weighted Interval Scheduling
- 2 Subset Sum Problem
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- 4 Longest Common Subsequence
  - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree**
- 8 Summary

# Optimum Binary Search Tree

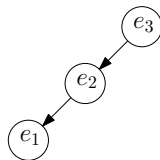
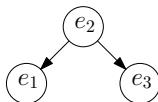
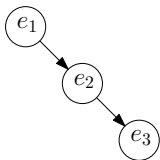
- $n$  elements  $e_1 < e_2 < e_3 < \dots < e_n$
- $e_i$  has frequency  $f_i$
- goal: build a binary search tree for  $\{e_1, e_2, \dots, e_n\}$  with the minimum accessing cost:

$$\sum_{i=1}^n f_i \times (\text{depth of } e_i \text{ in the tree})$$



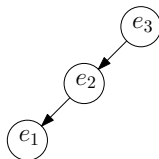
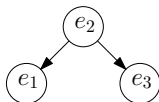
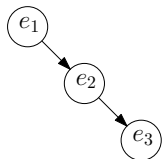
# Optimum Binary Search Tree

- Example:  $f_1 = 10, f_2 = 5, f_3 = 3$



# Optimum Binary Search Tree

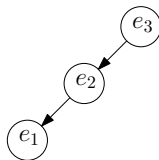
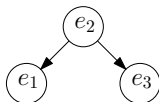
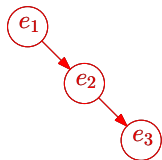
- Example:  $f_1 = 10, f_2 = 5, f_3 = 3$



- $10 \times 1 + 5 \times 2 + 3 \times 3 = 29$
- $10 \times 2 + 5 \times 1 + 3 \times 2 = 31$
- $10 \times 3 + 5 \times 2 + 3 \times 1 = 43$

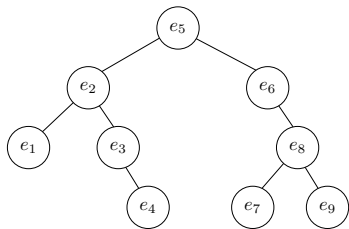
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- Example:  $f_1 = 10, f_2 = 5, f_3 = 3$

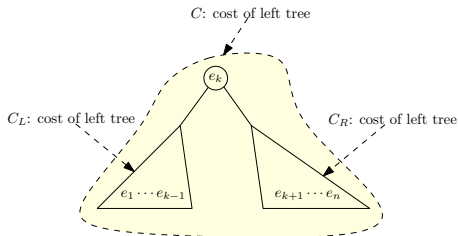


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- suppose we decided to let  $e_k$  be the root
- $e_1, e_2, \dots, e_{k-1}$  are on left sub-tree
- $e_{k+1}, e_{k+2}, \dots, e_n$  are on right sub-tree
- $d_j$ : depth of  $e_j$  in our tree
- $C, C_L, C_R$ : cost of tree, left sub-tree and right sub-tree



- $d_1 = 3, d_2 = 2, d_3 = 3, d_4 = 4, d_5 = 1,$
- $d_6 = 2, d_7 = 4, d_8 = 3, d_9 = 4,$
- $C = 3f_1 + 2f_2 + 3f_3 + 4f_4 + f_5 + 2f_6 + 4f_7 + 3f_8 + 4f_9$
- $C_L = 2f_1 + f_2 + 2f_3 + 3f_4$
- $C_R = f_6 + 3f_7 + 2f_8 + 3f_9$
- $C = C_L + C_R + \sum_{j=1}^9 f_j Z$



$$\begin{aligned}
 C &= \sum_{\ell=1}^n f_{\ell} d_{\ell} = \sum_{\ell=1}^n f_{\ell} (d_{\ell} - 1) + \sum_{\ell=1}^n f_{\ell} \\
 &= \sum_{\ell=1}^{k-1} f_{\ell} (d_{\ell} - 1) + \sum_{\ell=k+1}^n f_{\ell} (d_{\ell} - 1) + \sum_{\ell=1}^n f_{\ell} \\
 &= C_L + C_R + \sum_{\ell=1}^n f_{\ell}
 \end{aligned}$$

$$C = C_L + C_R + \sum_{\ell=1}^n f_{\ell}$$

- In order to minimize  $C$ , need to minimize  $C_L$  and  $C_R$  respectively

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- In order to minimize  $C$ , need to minimize  $C_L$  and  $C_R$  respectively
- $opt[i, j]$ : the optimum cost for the instance  $(f_i, f_{i+1}, \dots, f_j)$

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$$opt[1, n] = (opt[1, k-1] + opt[k+1, n]) + \sum_{\ell=1}^n f_{\ell}$$



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$$opt[1, n] = \min_{k: 1 \leq k \leq n} (opt[1, k-1] + opt[k+1, n]) + \sum_{\ell=1}^n f_{\ell}$$

- In general,  $opt[i, j] =$

$$\begin{cases} 0 & \text{if } i = j + 1 \\ \min_{k: i \leq k \leq j} (opt[i, k-1] + opt[k+1, j]) + \sum_{\ell=i}^j f_{\ell} & \text{if } i \leq j \end{cases}$$

## Optimum Binary Search Tree

```
1:  $fsum[0] \leftarrow 0$ 
2: for  $i \leftarrow 1$  to  $n$  do  $fsum[i] \leftarrow fsum[i - 1] + f_i$ 
    $\triangleright fsum[i] = \sum_{j=1}^i f_j$ 
3: for  $i \leftarrow 0$  to  $n$  do  $opt[i + 1, i] \leftarrow 0$ 
4: for  $\ell \leftarrow 1$  to  $n$  do
5:   for  $i \leftarrow 1$  to  $n - \ell + 1$  do
6:      $j \leftarrow i + \ell - 1, opt[i, j] \leftarrow \infty$ 
7:     for  $k \leftarrow i$  to  $j$  do
8:       if  $opt[i, k - 1] + opt[k + 1, j] < opt[i, j]$  then
9:          $opt[i, j] \leftarrow opt[i, k - 1] + opt[k + 1, j]$ 
10:         $\pi[i, j] \leftarrow k$ 
11:     $opt[i, j] \leftarrow opt[i, j] + fsum[j] - fsum[i - 1]$ 
```

# Printing the Tree

## Print-Tree( $i, j$ )

```
1: if  $i > j$  then  
2:   return  
3: else  
4:   print('(')  
5:   Print-Tree( $i, \pi[i, j] - 1$ )  
6:   print( $\pi[i, j]$ )  
7:   Print-Tree( $\pi[i, j] + 1, j$ )  
8:   print('')
```

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## Dynamic Programming

- Break up a problem into many **overlapping** sub-problems
- Build solutions for larger and larger sub-problems
- Use a **table** to store solutions for sub-problems for reuse

## Comparison with greedy algorithms

- Greedy algorithm: each step is making a small progress towards constructing the solution
- Dynamic programming: the whole solution is constructed in the last step

## Comparison with divide and conquer

- Divide and conquer: an instance is broken into many **independent** sub-instances, which are solved separately.
- Dynamic programming: the sub-instances we constructed are overlapping.

# Definition of Cells for Problems We Learnt

- Weighted interval scheduling:  $opt[i] =$  value of instance defined by jobs  $\{1, 2, \dots, i\}$
- Subset sum, knapsack:  $opt[i, W'] =$  value of instance with items  $\{1, 2, \dots, i\}$  and budget  $W'$
- Longest common subsequence:  $opt[i, j] =$  value of instance defined by  $A[1..i]$  and  $B[1..j]$
- Shortest paths in DAG:  $f[v] =$  length of shortest path from  $s$  to  $v$
- Matrix chain multiplication, optimum binary search tree:  
 $opt[i, j] =$  value of instances defined by matrices  $i$  to  $j$