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December 9, 2021 • Arrays / Data Structure

Trapping Rainwater

Problem Statement: Given an array of non-negative integers representation elevation of ground. Your task is to find the water that can be trapped after rain.

Examples:

Example 1:

Input: height= [0,1,0,2,1,0,1,3,2,1,2,1]

Output: 6

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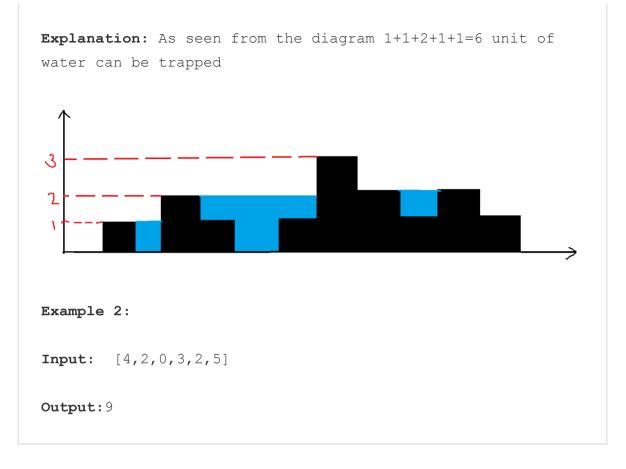
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Solution

Disclaimer: Don't jump directly to the solution, try it out yourself first.

Solution 1: Brute force

Approach: For each index, we have to find the amount of water that can be stored and we have to sum it up.If we observe carefully the amount the water

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Print N to 1 using Recursion

stored at a particular index is the minimum of maximum elevation to the left and right of the index minus the elevation at that index.

Code:

C++ Code

```
#include<bits/stdc++.h>
using namespace std;
int trap(vector < int > & arr) {
 int n = arr.size();
 int waterTrapped = 0;
 for (int i = 0; i < n; i++) {
   int j = i;
    int leftMax = 0, rightMax = 0;
    while (j \ge 0) {
      leftMax = max(leftMax, arr[j]);
      j--;
    j = i;
    while (j < n) {
      rightMax = max(rightMax, arr[j]);
      j++;
    waterTrapped += min(leftMax, rightMax) - arr[i];
 return waterTrapped;
```

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```
int main() {
  vector < int > arr;
  arr = {0,1,0,2,1,0,1,3,2,1,2,1};
  cout << "The water that can be trapped is " << trap(arr) << endl;
}</pre>
```

Time Complexity: O(N*N) as for each index we are calculating leftMax and rightMax so it is a nested loop.

Space Complexity: O(1).

Java Code

Time Complexity: O(N*N) as for each index we are calculating leftMax and rightMax so it is a nested loop.

Space Complexity: O(1).

Python Code

```
from typing import List

def trap(arr: List[int]) -> int:
    n = len(arr)
    waterTrapped = 0
```

```
for i in range(n):
        j = i
       leftMax = 0
        rightMax = 0
       while j \ge 0:
           leftMax = max(leftMax, arr[j])
           i -= 1
       j = i
       while j < n:
           rightMax = max(rightMax, arr[j])
           j += 1
       waterTrapped += min(leftMax, rightMax) - arr[i]
    return waterTrapped
if name == " main ":
    arr = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]
    print(f"The water that can be trapped is {trap(arr)}")
```

Time Complexity: O(N*N) as for each index we are calculating leftMax and rightMax so it is a nested loop.

Space Complexity: O(1).

Solution 2:Better solution

Intuition: We are taking O(N) for computing leftMax and rightMax at each index. The complexity can be boiled down to O(1) if we precompute the leftMax and rightMax at each index.

Approach: Take 2 array prefix and suffix array and precompute the leftMax and rightMax for each index beforehand. Then use the formula min(prefix[I], suffix[i])-arr[i] to compute water trapped at each index.

Code:

C++ Code

```
#include<bits/stdc++.h>

using namespace std;
int trap(vector < int > & arr) {
  int n = arr.size();
  int prefix[n], suffix[n];
  prefix[0] = arr[0];
  for (int i = 1; i < n; i++) {
    prefix[i] = max(prefix[i - 1], arr[i]);
  }
  suffix[n - 1] = arr[n - 1];
  for (int i = n - 2; i >= 0; i--) {
    suffix[i] = max(suffix[i + 1], arr[i]);
  }
  int waterTrapped = 0;
  for (int i = 0; i < n; i++) {
    waterTrapped += min(prefix[i], suffix[i]) - arr[i];
}</pre>
```

```
}
  return waterTrapped;
}
int main() {
  vector < int > arr;
  arr = {0,1,0,2,1,0,1,3,2,1,2,1};
  cout << "The water that can be trapped is " << trap(arr) << endl;
}</pre>
```

Time Complexity: O(3*N) as we are traversing through the array only once.

And O(2*N) for computing prefix and suffix array.

Space Complexity: O(N)+O(N) for prefix and suffix arrays.

Java Code

```
import java.util.*;
class TUF {
    static int trap(int[] arr) {
        int n = arr.length;
        int prefix[] = new int[n];
        int suffix[] = new int[n];
        prefix[0] = arr[0];
        for (int i = 1; i < n; i++) {
            prefix[i] = Math.max(prefix[i - 1], arr[i]);
        }
}</pre>
```

```
suffix[n - 1] = arr[n - 1];
for (int i = n - 2; i >= 0; i--) {
        suffix[i] = Math.max(suffix[i + 1], arr[i]);
}
int waterTrapped = 0;
for (int i = 0; i < n; i++) {
        waterTrapped += Math.min(prefix[i], suffix[i]) - arr[i];
}
return waterTrapped;
}

public static void main(String args[]) {
    int arr[] = {0,1,0,2,1,0,1,3,2,1,2,1};
    System.out.println("The duplicate element is " + trap(arr));
}</pre>
```

Time Complexity: O(3*N) as we are traversing through the array only once. And O(2*N) for computing prefix and suffix array.

Space Complexity: O(N)+O(N) for prefix and suffix arrays.

Python Code

from typing import List

```
def trap(arr: List[int]) -> int:
    n = len(arr)
    prefix = [0] * n
    suffix = [0] * n
    prefix[0] = arr[0]
    for i in range(1, n):
        prefix[i] = max(prefix[i - 1], arr[i])
    suffix[n - 1] = arr[n - 1]
    for i in range(n - 2, -1, -1):
        suffix[i] = max(suffix[i + 1], arr[i])
    waterTrapped = 0
    for i in range(n):
        waterTrapped += min(prefix[i], suffix[i]) - arr[i]
    return waterTrapped
if name _ == "__main__":
    arr = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]
    print(f"The water that can be trapped is {trap(arr)}")
```

Time Complexity: O(3*N) as we are traversing through the array only once.

And O(2*N) for computing prefix and suffix array.

Space Complexity: O(N)+O(N) for prefix and suffix arrays.

Solution 3:Optimal Solution(Two pointer approach)

Approach: Take 2 pointers I(left pointer) and r(right pointer) pointing to 0th and (n-1)th index respectively. Take two variables leftMax and rightMax and initialize them to 0. If height[I] is less than or equal to height[r] then if leftMax is less than height[I] update leftMax to height[I] else add leftMax-height[I] to your final answer and move the I pointer to the right i.e I++. If height[r] is less than height[I], then now we are dealing with the right block. If height[r] is greater than rightMax, then update rightMax to height[r] else add rightMax-height[r] to the final answer. Now move r to the left. Repeat these steps till I and r crosses each other.

Intuition: We need a minimum of leftMax and rightMax.So if we take the case when height[I]<=height[r] we increase I++, so we can surely say that there is a block with a height more than height[I] to the right of I. And for the same reason when height[r]<=height[I] we can surely say that there is a block to the left of r which is at least of height[r]. So by traversing these cases and using two pointers approach the time complexity can be decreased without using extra space.

Code:

C++ Code

```
#include<bits/stdc++.h>
using namespace std;
int trap(vector < int > & height) {
 int n = height.size();
  int left = 0, right = n - 1;
  int res = 0;
  int maxLeft = 0, maxRight = 0;
  while (left <= right) {</pre>
    if (height[left] <= height[right]) {</pre>
      if (height[left] >= maxLeft) {
        maxLeft = height[left];
      } else {
        res += maxLeft - height[left];
      left++;
    } else {
      if (height[right] >= maxRight) {
        maxRight = height[right];
      } else {
        res += maxRight - height[right];
      right--;
  return res;
int main() {
  vector < int > arr;
  arr = \{0,1,0,2,1,0,1,3,2,1,2,1\};
```

```
cout << "The water that can be trapped is " << trap(arr) << endl;
}</pre>
```

Time Complexity: O(N) because we are using 2 pointer approach.

Space Complexity: O(1) because we are not using anything extra.

Java Code

```
import java.util.*;
class TUF {
    static int trap(int[] height) {
        int n = height.length;
        int left = 0, right = n - 1;
        int res = 0;
        int maxLeft = 0, maxRight = 0;
        while (left <= right) {</pre>
            if (height[left] <= height[right]) {</pre>
                if (height[left] >= maxLeft) {
                    maxLeft = height[left];
                } else {
                    res += maxLeft - height[left];
                left++;
            } else {
                if (height[right] >= maxRight) {
                    maxRight = height[right];
```

Time Complexity: O(N) because we are using 2 pointer approach.

Space Complexity: O(1) because we are not using anything extra.

Python Code

```
from typing import List

def trap(height: List[int]) -> int:
```

```
n = len(height)
    left = 0
    right = n-1
    res = 0
    maxLeft = 0
    maxRight = 0
    while left <= right:
        if height[left] <= height[right]:</pre>
            if height[left] >= maxLeft:
                maxLeft = height[left]
            else:
                res += maxLeft - height[left]
            left += 1
        else:
            if height[right] >= maxRight:
                maxRight = height[right]
            else:
                res += maxRight - height[right]
            right -= 1
    return res
if name == " main ":
    arr = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]
    print(f"The water that can be trapped is {trap(arr)}")
```

Time Complexity: O(N) because we are using 2 pointer approach.

Space Complexity: O(1) because we are not using anything extra.

Special thanks to <u>Pranav Padawe</u> and <u>Sudip Ghosh</u> for contributing to this article on takeUforward. If you also wish to share your knowledge with the takeUforward fam, <u>please check out this article</u>

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