

Median of BST

Find it in $O(1)$ space & $O(n)$ time complexity.

If I use a recursive function or auxiliary array extra space will be involved. So we use morris inorder traversal method.

Approach :

- 1) Count the number of nodes (n) in BST using morris traversal
- 2) Perform morris traversal again to find the median
 - If n is odd, return middle element $(\frac{n+1}{2})$
 - If n is even, return average of $\frac{n}{2}$ & $\frac{n}{2} + 1^{\text{th}}$ elements.

1) Count the total nodes

Counting is done using morris traversal.

Steps:

- 1) Start from root node
- 2) If the left child does not exist, increment the count & move to the right child.
- 3) If the left child exists, find its inorder predecessor (rightmost child node in left subtree)
 - If the predecessor does not have a link (right pointer to root), create a temporary link to the current node & move left.
 - If the predecessor does have a link, remove it, increment count, move right.
- 4) Continue until entire tree is traversed.

2) Find median

If n is odd \rightarrow median = $(\frac{n+1}{2})^{\text{th}}$ element.

If n is even \rightarrow median =
$$\frac{(\frac{n}{2})^{\text{th}} + (\frac{n+1}{2})^{\text{th}}}{2}$$

Morris traversal again

i) start from root

ii) If there is no left child, increase currCount.

iii) If there is a left child, find inorder predecessor.

- If no link exists, create one & move left

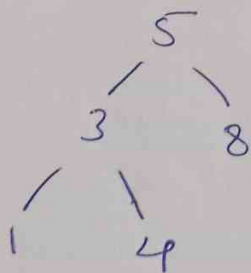
- If link exists, remove it, increase currCount, & move right.

iv) check for median condition

- If $\text{currCount} == \frac{n+1}{2}$ (for odd n) return value

- If $\text{currCount} == \frac{n}{2}$ or $\text{currCount} == \frac{n}{2} + 1$ (for even n) take average of the 2 values

e.g.



Morris inorder traversal

1) Start at root (5) (i.e., current = root)

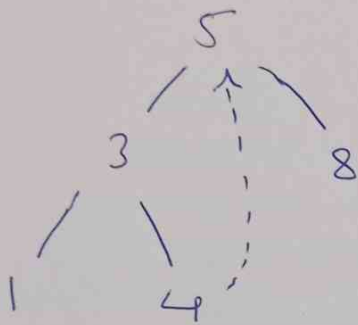
- Left subtree exists \rightarrow Find the inorder predecessor (right most node in left subtree)

Left subtree is

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graph TD; 3 --> 1; 3 --> 4;
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2) The inorder predecessor of 5 is 4 (rightmost node in 3's subtree)

- Create temporary link by pointing 4's right to 5.
 $\text{prev.right} = \text{current}$.



3) Move to left child (3) ($\text{current} = \text{current.left}$)

At node (3)

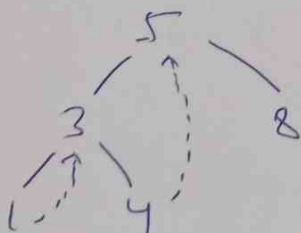
- Left subtree exists \rightarrow Find inorder predecessor (rightmost node in left subtree).

The left subtree of 3:

1

Inorder predecessor of 3 is 1. ($\text{prev} = 1$)

- Create temporary link by pointing 1's right to 3.
 $\text{prev.right} = \text{current}$.



4) Move to left child (1) ($\text{current} = \text{current.left}$) again
At node (1)

- No left subtree $\rightarrow \text{Count} = 1$

- move to the right (link exists because we created it in previous step) \rightarrow Go back to 3 (remove thread/link)

- Back to node(3)

 - \rightarrow Count = 2

 - move to right child (4)

- At node (4)

 - \rightarrow No left subtree \rightarrow Count = 3

 - move to the right (link is there pointing to 5) \rightarrow
Go back to 5 (remove link).

- Back to node (5)

 - \rightarrow Count = 4

 - move to the right child 8

- At Node (8)

 - No left subtree \rightarrow Count = 5.

 - Traversal ends

 - All nodes counted

This is how morris traversal works.

Do this traversal twice, one for counting nodes, another for finding median.

Then I will get median in BST.