

I SANMACS

- 309 & 313 PRATAP CHAMBERS, GURUDWARA ROAD, KAROL BAGH, DELHI - 110005 PH.: 28757911, 28757630, 32917966
- ➤ SAI SANMACS PLAZA, PLOT NO. 6A, COMMUNITY CENTER, DDA, SECTOR 8, ROHINI, DELHI 110085 PH.: 32458354, 32522575

e-mail: info@sanmacs.com

website: www.sanmacs.com

Rewarding Career

HINTS AND SOLUTIONS

CTNC - III

- **1. (d)** 26 Capital letters +26 small letters + 10 digits +32 other characters.
- **2. (d)** Now 3 address instruction means the operation code followed by 3 address registers. 64 operation codes= 2^6 so 6 bits for op. code Now 2 MB = 2^{21} bytes. So no. of bits in MAR should be 21. So max. bits required for a 3 address be 21. So max. bits required for 3 Address instruction are 6 + 21 + 21 + 21 = 69.
- **3. (c)** It uses laser technology.
- **4. (d)** Output in NAND gate is 1 when at least one input is O. so, no. of combinations having at least 1 zero = total combinations -1 combination having all 1's. So, $2^{10} 1 = 1023$ because total combinations with 10 inputs is 2^{10}
- **5. (b)** Capacity = s * p * t * m

= 256 * 128 * 1000 * 10 = 327.68 MB.

10 Surfaces because 5 plates has 10 surfaces.

6. (c) A1. UV is in 32 base system. Now $32 = 2^5$ so we replace each digit by a 5 bit code to get. The binary equivalent of this no. So A1.UV = 0101000001.1111011111. Now to get the Octal no. take the combinations of 3 bits and then replace each combination by the corresponding octal digit 0101000 001.111 101 111 1 =

101000001.1111011111100 = (501.7574) in base8.

- 7. (d) $2MB = 2 * 1000 K = 2 * 2^{10} * 2^{10}$ bytes.
- $= 2^{21}$ bytes. So, 21 bits for MAR.

Each word of 1 byte and 1 byte = 8 bits.

So, 8 bits for MDR.

- **8. (b)** Minimum no. is = 0.100000000000E1111
- $=(2^{-1})*2^{-(2-1)}=2^{-1}*2^{-7}=2^{-8}$
- 9. (a) Average access time =avg. latency time+ avg. seek time. Avg. seek time = 10ms. Avg. latency time = Half the time taken by disk rotation. Time for 1 rotation = 1/10000 min. = 60/10000s = 6ms. So. Avg. latency time = 6/2 ms = 3 ms. So avg. access time = (10+3) ms = 13ms.
- **10. (c)** On adding the 2 exponent we get 11010110 in the exponent part which indicates over flow for 2 positive exponent.
- 11. (a) For unsigned integer 1110 = 14 in decimal in signed 2's complement representation. Now the sign bit is

1 so the no. is negative and we again have to take the 2's complement 2's compl. of 1110 is 0010. So the no. is -2.

12. (c)

13. (c) 16 million = 2^{24} . Nowany pixelcan have 1 out of 2^{24} colors. So, 24 bits will be there for each pixel in the memory .24 bits = 3 bytes.

Now total pixels = 1024 * 1280. So total buffer memory required = 1024*1280* 3bytes = 4MB.

- 14. (c) Cache memory is a buffer between main memory and the processor.
- 15. (c)

16. (d)

- 17. (b) Perl is a procedural scripting language.
- **18.** (d) All the statement are true.
- **19. (b)** It is a tab operator. So prints 1 to 5 in the same line but with fixed gap between them.

20. (d)

- **21. (b)** Multiple intialisation and incrementation expressions are allowed in FOR loop. However only 1 expression is allowed in test expression.
- **22. (c)** We can only have an int. constant or a char constant after case. So float constant will give a Syntax error.
- 23. (c) Range for integer constants is -32768 to 32767.

Now here i = 35,000 So, 35,000 - 32767 = 2233.

Now 32767 + 1 = -32768

So, 32767 + 2233 = -32768 + 2232 = -30536.

24. (b) Now 24 address lines means the maximum capacity is 2^{24} .

1 memory chip has capacity 64kb=2¹⁶ bytes.

So no. of chips = $2^{24}/2^{16} = 256$

25. (c) a = 12.5, b = 2.75. Using normalized floating Point representation with 4 bits for mantissa and exponent a = 1100.1 = 0.110E0100

b = 0010.11 = 0.101E0010

Now making the exponent of smaller no. same as that of larger. So b =0.001E0100now on adding the two terms mantissa becomes 0.110+0.001=0.111 and exponent is

0100. So the no. is 0.111E0100 which is 1110 in binary

and 14 in decimal.