

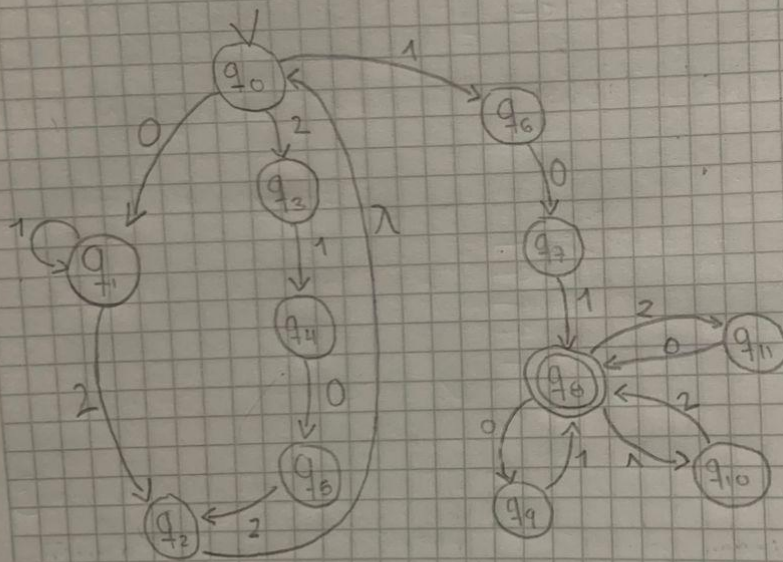
Computer Science III

Workshop 1

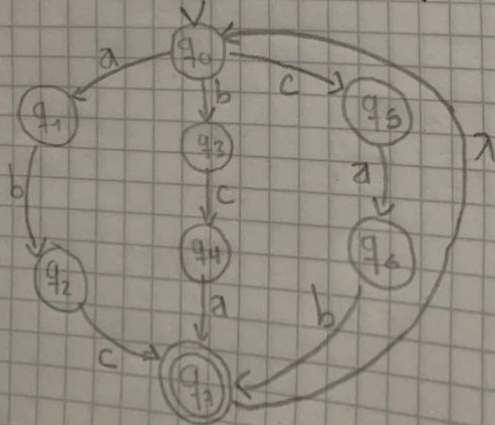
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1. For each of the following languages, define the corresponding finite-state machine:

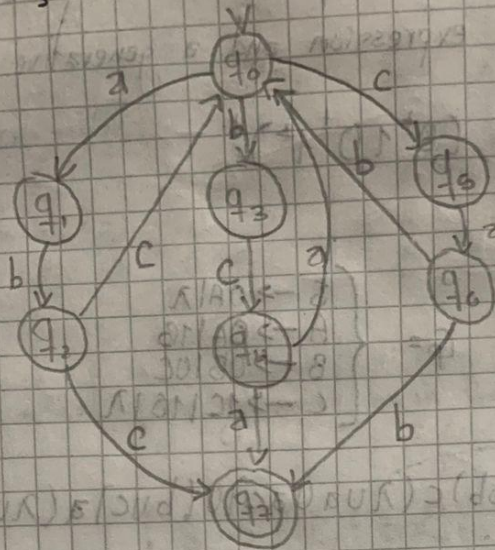
(i) $\Sigma = \{0, 1, 2\}$ $L = (01^*2 \cup 2102)^* 101(01 \cup 12 \cup 20)^*$



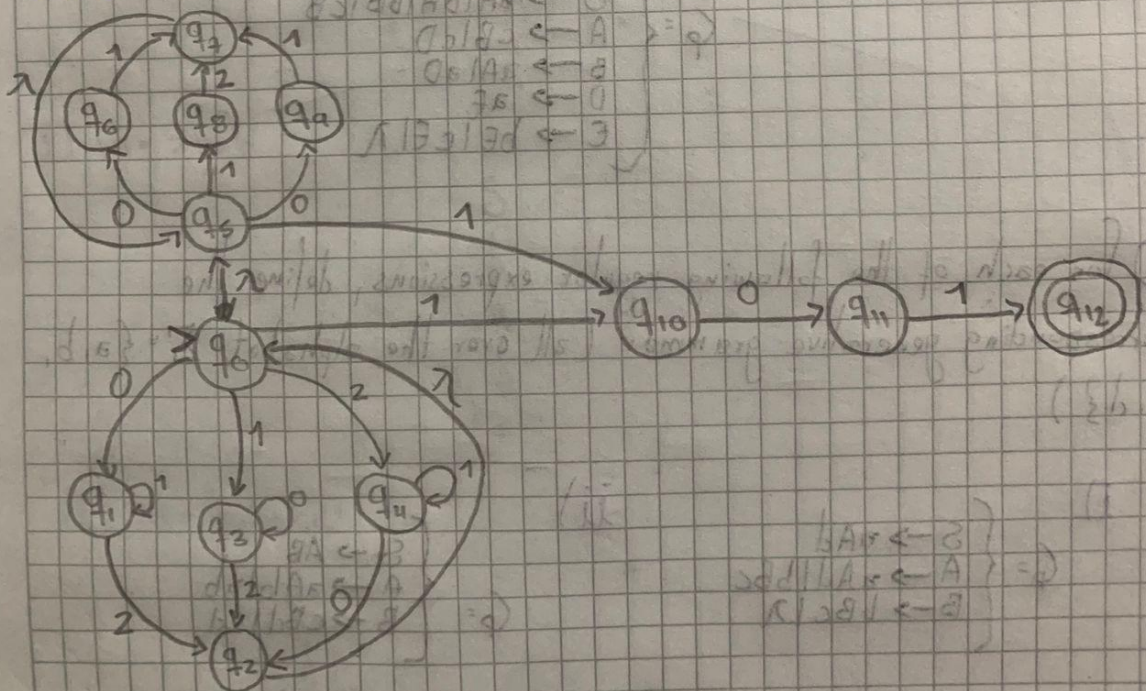
(ii) $\Sigma = \{a, b, c\}$ $L = (abc \cup bca \cup cab)(abc \cup bca \cup cab)^*$



iii) $\Sigma = \{a, b, c\}$ $L = (abc \cup bca \cup cab)^* (abc \cup bca \cup cab)$



iv) $\Sigma = \{0, 1, 2\}$ $L = (01^*2 \cup 10^*2 \cup 21^*0)^* (01 \cup 12 \cup 20)^* 101$



2. For each one of the following finite-state machines, define the corresponding regular expression and a generative grammar.

i) regular expression: $(0^+10^+1^+)^*$

generative grammar:

$$G = \begin{cases} S \rightarrow 0A1\lambda \\ A \rightarrow 0A1B \\ B \rightarrow 0B1C \\ C \rightarrow 1C1S1\lambda \end{cases}$$

ii) regular expression: $((a^+b^+)c(\lambda \cup a \cup ac)^+ \cup (b^+c^+)a(\lambda \cup c \cup ca)^+)((a^+b^+)^* \cup \lambda)$

generative grammar:

$$G = \begin{cases} S \rightarrow aA|bA|bB|cB \\ A \rightarrow cB|cD \\ B \rightarrow aA|aD \\ D \rightarrow aE \\ E \rightarrow bE|cE|\lambda \end{cases}$$

3) For each of the following regular expressions, define the corresponding generative grammar (all over the alphabet $\Sigma = \{a, b, c, d\}$)

i)

$$G = \begin{cases} S \rightarrow aAd \\ A \rightarrow aAd|bBc \\ B \rightarrow bBc|\lambda \end{cases}$$

ii)

$$G = \begin{cases} S \rightarrow AB \\ A \rightarrow aAb|ab \\ B \rightarrow cBd|cd \end{cases}$$

iii)

$$G = \begin{cases} S \rightarrow AID \\ A \rightarrow aBd \\ B \rightarrow aBd \mid bCc \\ C \rightarrow bCc \mid \lambda \\ D \rightarrow EF \\ E \rightarrow aEb \mid ab \\ f \rightarrow cfd \mid cd \end{cases}$$

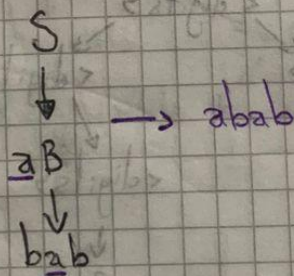
iv)

$$G = \begin{cases} S \rightarrow AB \\ A \rightarrow aAc \mid \lambda \\ B \rightarrow bBc \mid bc \end{cases}$$

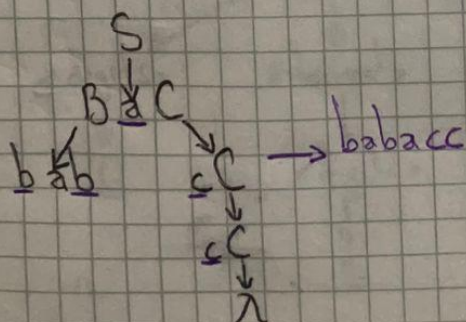
4. Let G a context-free grammar with the following productions:

$$G = \begin{cases} S \rightarrow ABC \mid BAC \mid aB \\ A \rightarrow Aa \mid a \\ B \rightarrow BAB \mid bab \\ C \rightarrow cC \mid \lambda \end{cases}$$

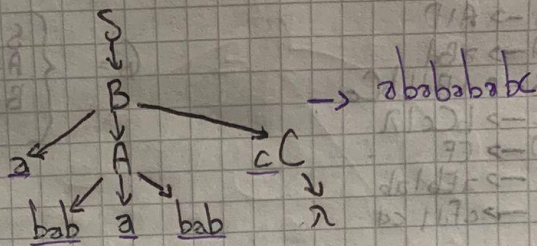
i) $w_1 = abab$



ii) $w_2 = babacc$

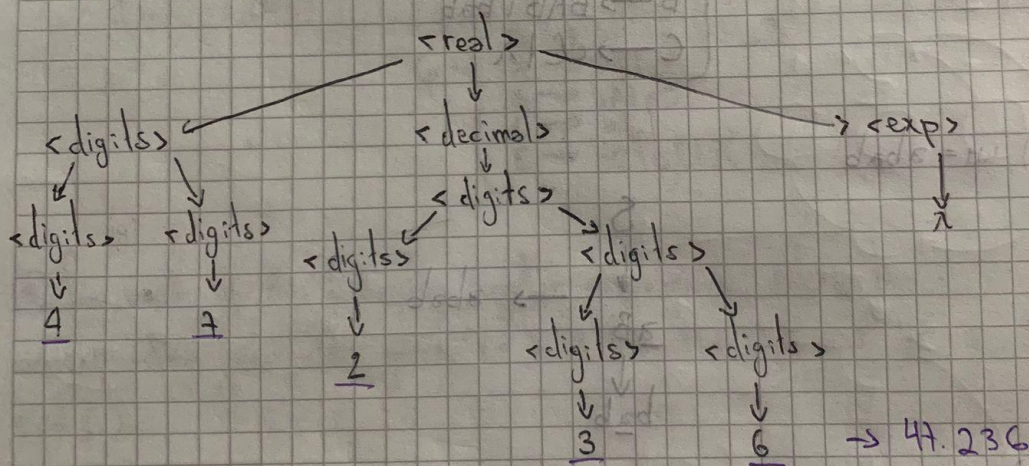


ii) $w_3 = ababababc$

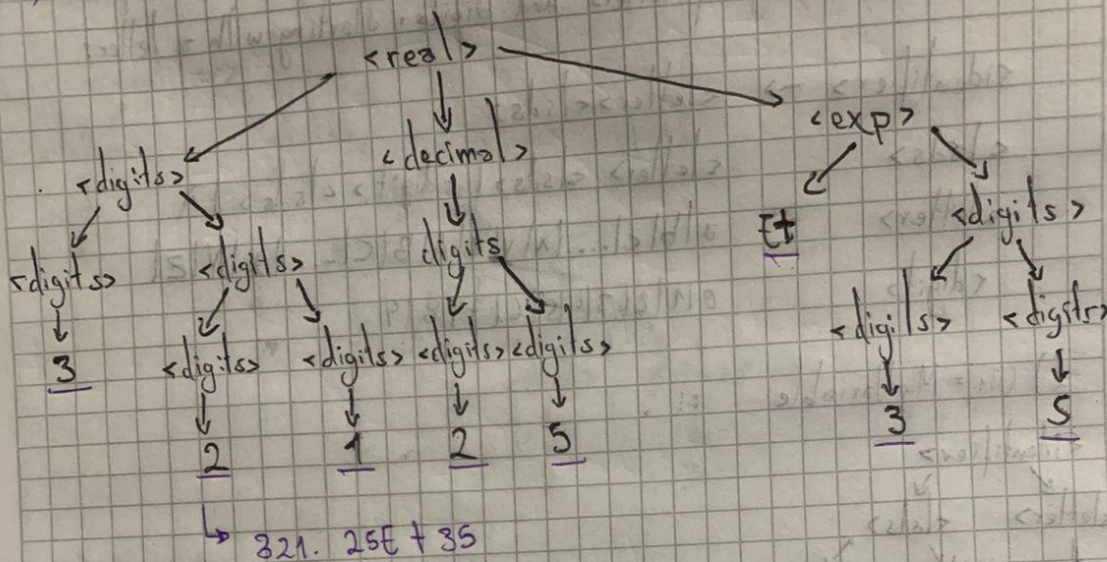


5. As follows there is a context-free grammar to generate real numbers without sign, the alphabet is $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \dots, \cdot, -\}$, E :

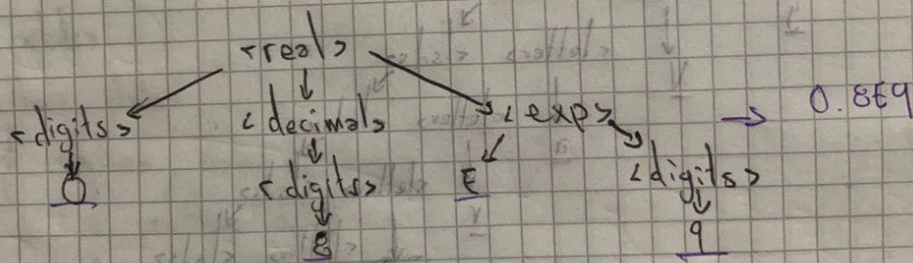
i) $w_1 = 47.236$



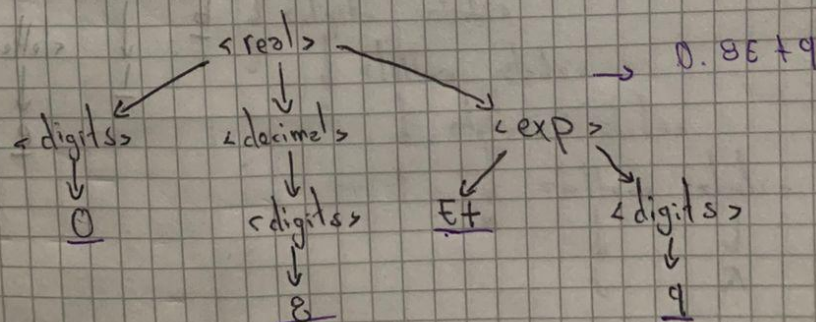
ii) $W_2 = 321.25E+35$



iii) $W_3 = 0.8E9$



iv) $W_4 = 0.8E+9$



6. The following is a context-free grammar to generate identifiers, identifiers are strings of letters and digits, starting with a letter:

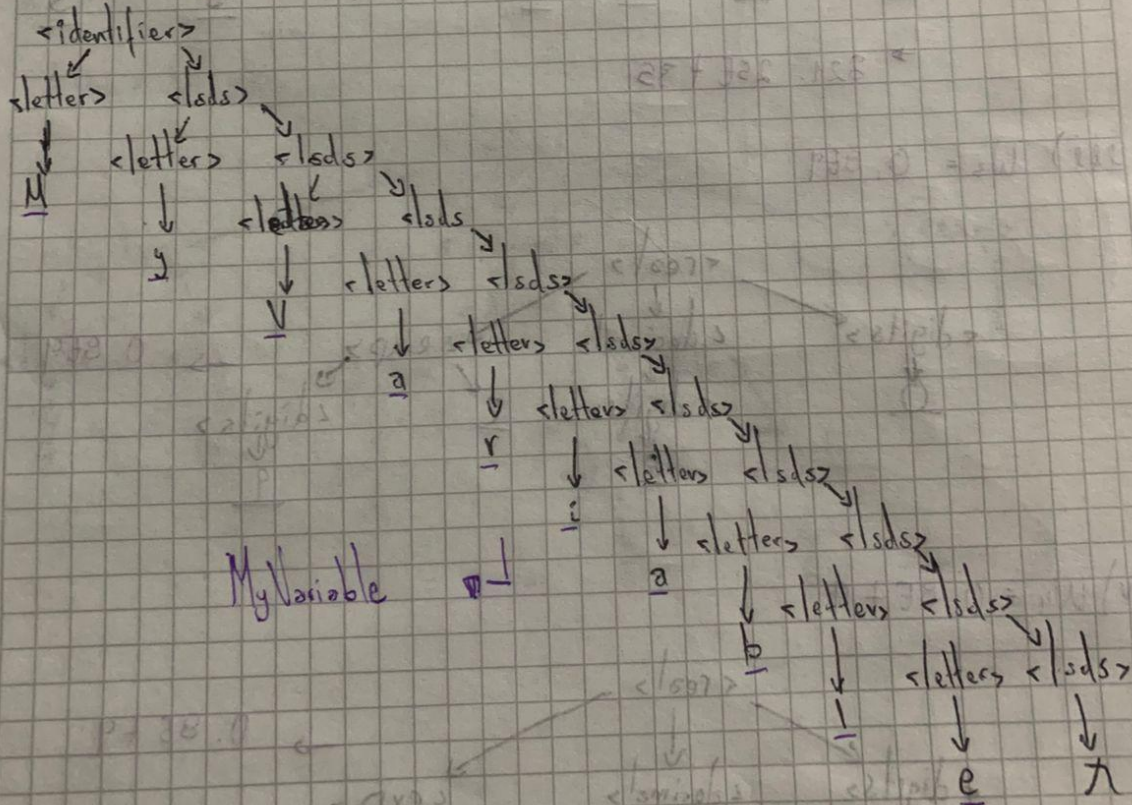
$\langle \text{identifier} \rangle \rightarrow \langle \text{letter} \rangle \langle \text{sds} \rangle$

$\langle \text{sds} \rangle \rightarrow \langle \text{letter} \rangle \langle \text{sds} \rangle \mid \langle \text{digit} \rangle \langle \text{sds} \rangle \mid \lambda$

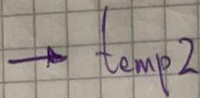
$\langle \text{letter} \rangle \rightarrow a|b|c|\dots|x|y|z|A|B|C|\dots|X|Y|Z|$

$\langle \text{digit} \rangle \rightarrow 0|1|2|3|4|5|6|7|8|9$

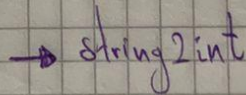
Q) $W_1 = \text{MyVariable}$



$\frac{1}{2} \times \frac{1}{2} = \frac{1}{4}$



sidentifizier



iv) $W_4 = 2\text{NotAVariable} \rightarrow$ The generative grammar given, can't generate this name, since it starts with a number