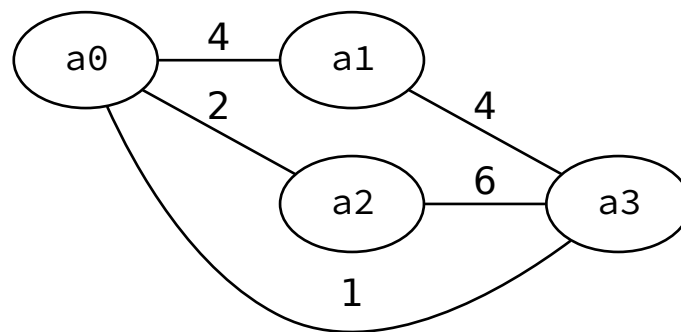


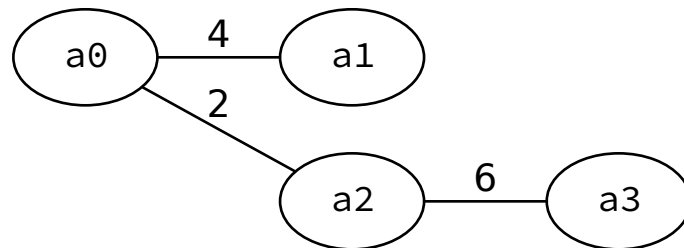
Weighted(Cost) Graph

>> Contains Cost or Weight while traversing from one node to another in the Graph.

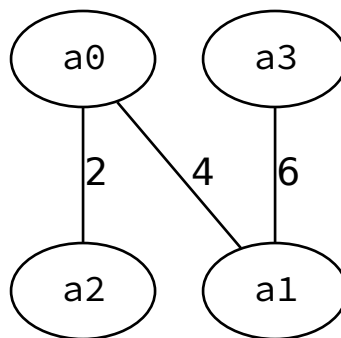


Spanning Tree

A tree representation, subset of a graph, that is constructed from minimum number of edges.



or



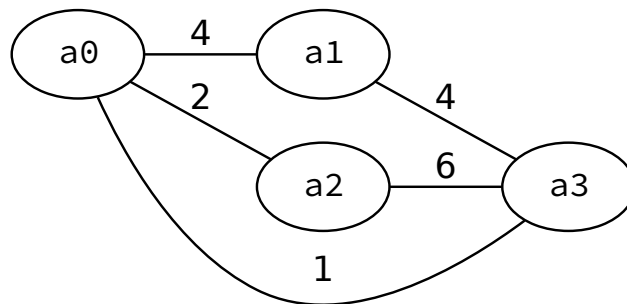
Total vertices = n

Total Edges = $n-1$

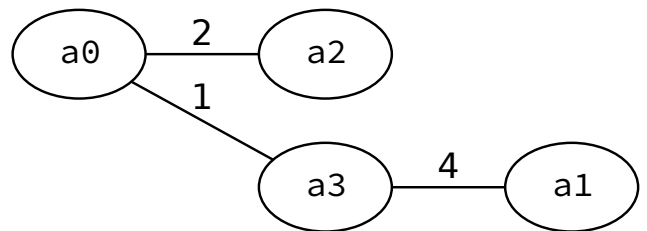
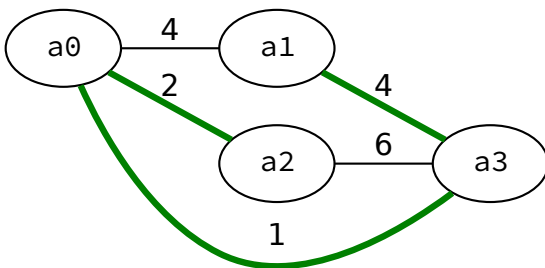
No. of Trees = n^{n-2}

Minimum Spanning Tree

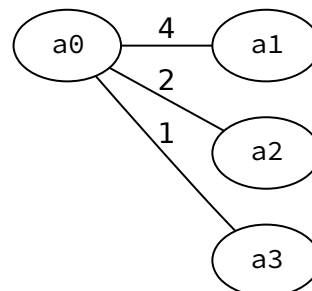
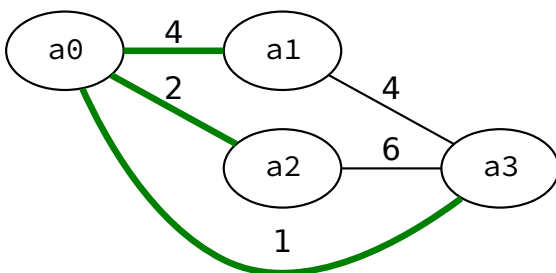
A spanning tree (does not have cycle), that is constructed from minimum cost edges.



Solution - 1

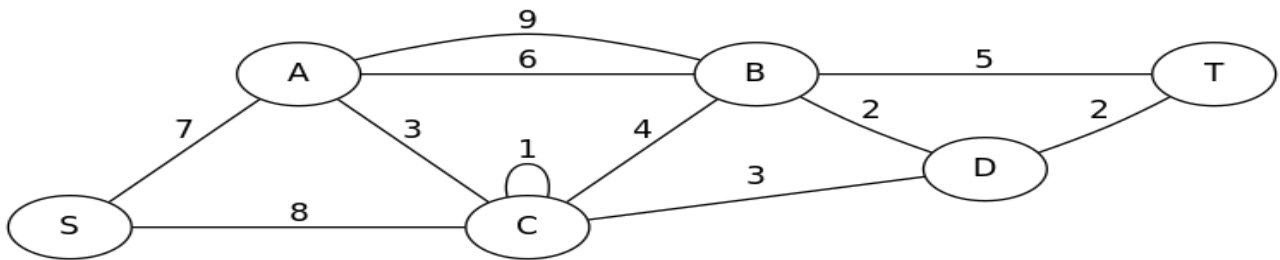


Solution - 2

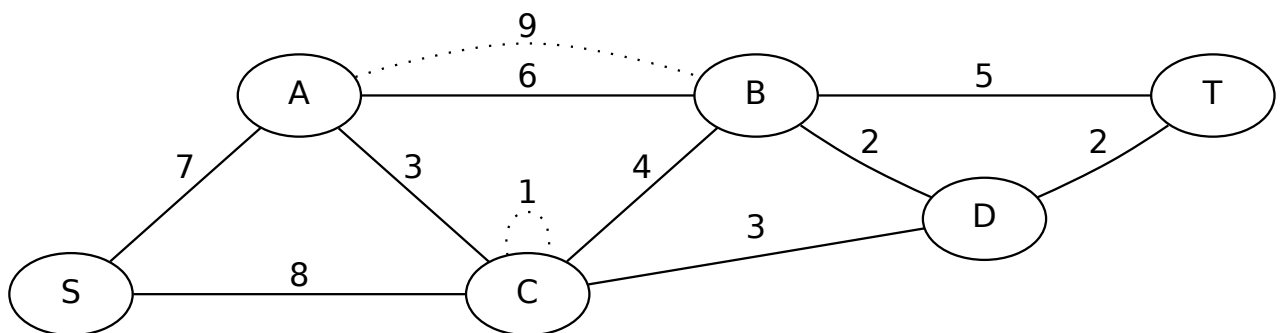


Kruskal's Algorithm

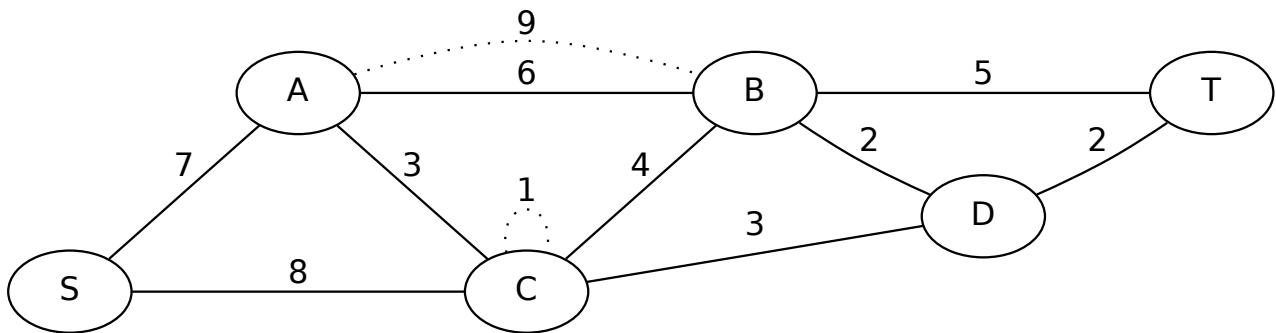
Example:



1. Remove self loop and parallel edges.

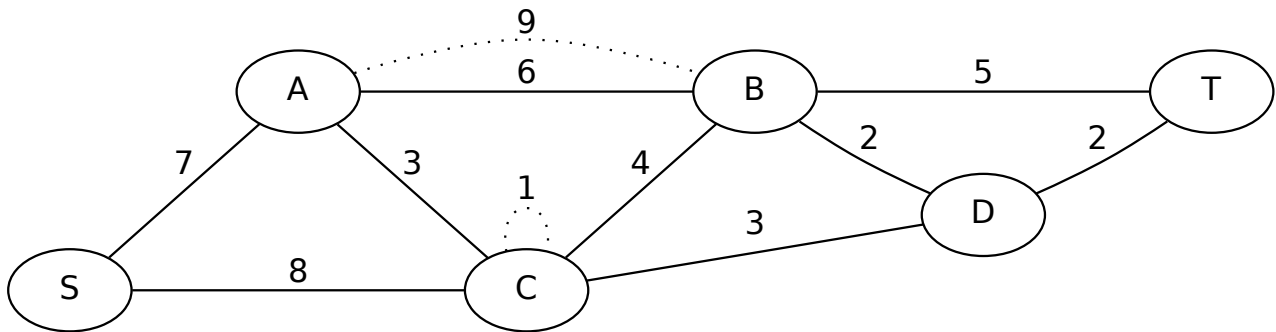


2. Arrange all edges with ascending order.

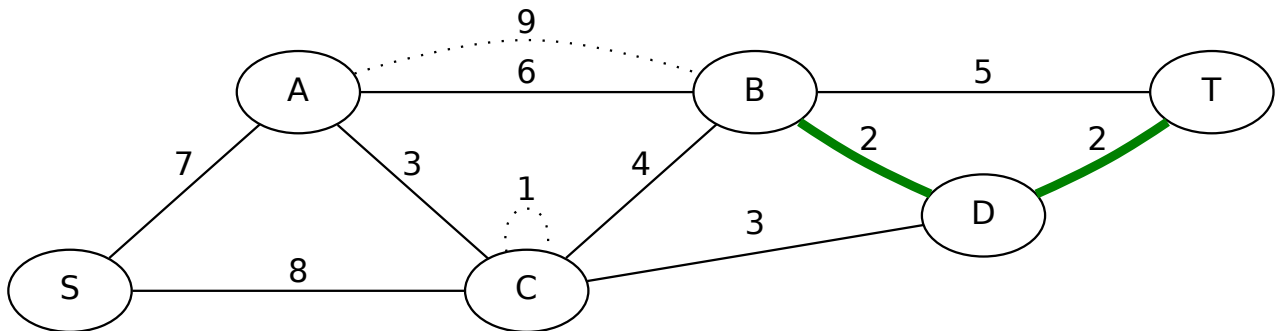


B→D	D→T	A→C	C→D	C→B	B→T	A→B	S→A	S→C
2	2	3	3	4	5	6	7	8

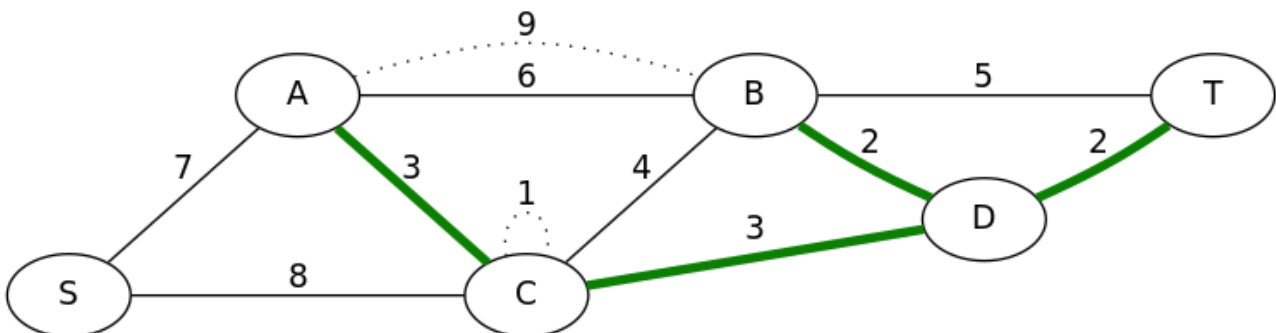
3. Add the edge with least weight.



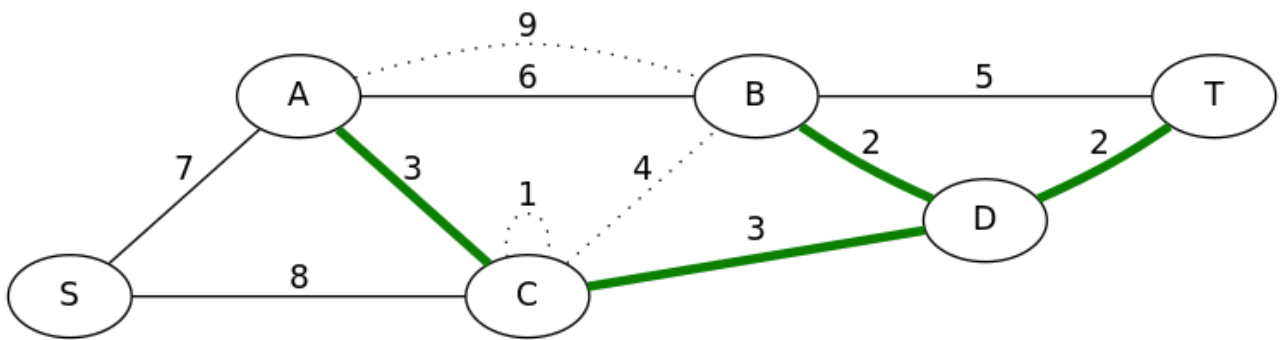
[least cost = 2]



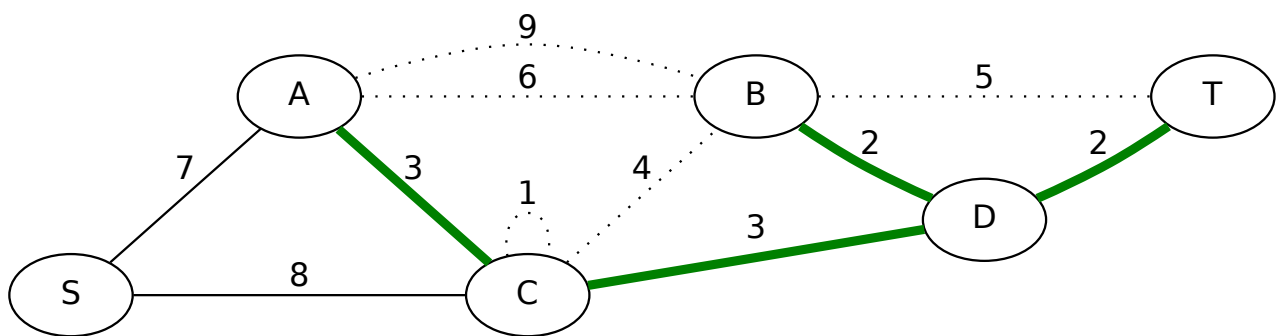
[least cost = 3]



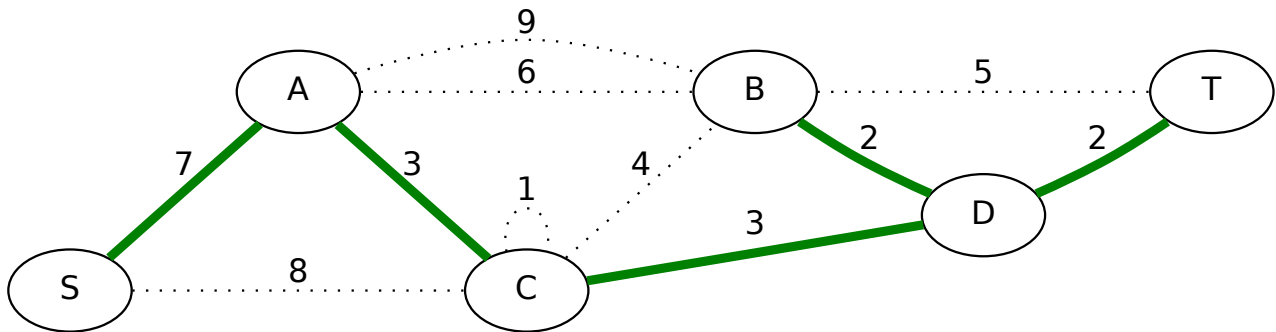
[least cost = 4, but makes cycle, so avoid it.]



[least cost = 5 & 6, but they also makes cycle, so avoid it.]



[Now, the least cost = 7 from A \rightarrow S]



Hence, all the nodes are covered.

Total Nodes = 6.

Total Edges = 5[Green Colors]

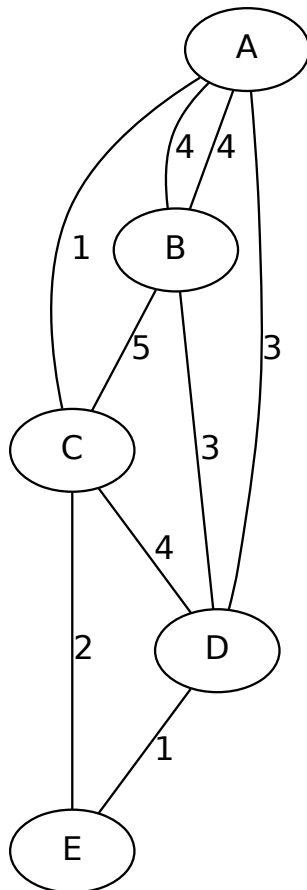
Total Cost = 2 + 2 + 3 + 3 + 7
= **17**

Complexity = $e \log e$, (e no. of edges.)

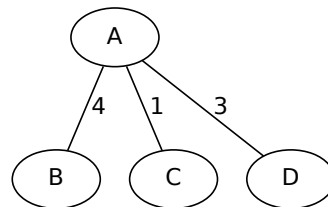
Round Robin Algorithm

(Tarjan And Cheriton's Algorithm For Finding MST. _pp.577 Of Book)

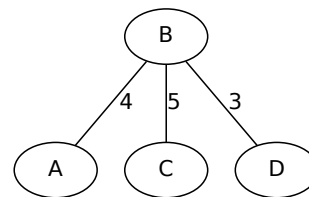
Generate all the priorities trees:



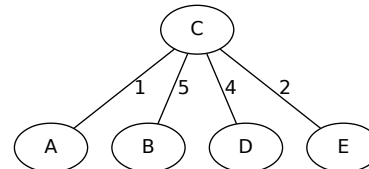
T1-Vertices(A) PQ: (A C 1), (A D 3), (A B 4)



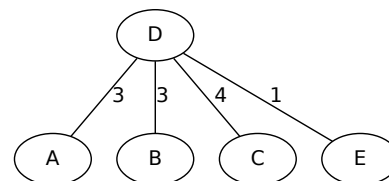
T2-Vertices(B) PQ: (B D 3), (B A 4), (B C 5)



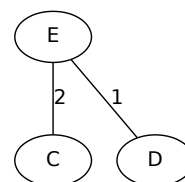
T3-Vertices(C) PQ: (C A 1), (C E 2), (C D 4), (C B 5)



T4-Vertices(D) PQ: (D E 1), (D A 3), (D B 3), (D C 4)



T5-Vertices(E) PQ: (E D 1), (E C 2)



Every partial tree has one vertex designated as its "root".

All Priorities Queues:

T1. Vertices: A PQ: (A C 1), (A D 3), (A B 4)
 T2. Vertices: B PQ: (B D 3), (B A 4), (B C 5)
 T3. Vertices: C PQ: (C A 1), (C E 2), (C D 4), (C B 5)
 T4. Vertices: D PQ: (D E 1), (D A 3), (D B 3), (D C 4)
 T5. Vertices: E PQ: (E D 1), (E C 2)

Now,

T1. Vertices: A P1: (A C 1), (A D 3), (A B 4)

Remove (A C 1), say α , v1 and v2 are two vertices of α .

α : (A C 1) v1: A v2: C

(A C 1) is a component of the MST

T3. Vertices: C PQ: (C A 1), (C E 2), (C D 4), (C B 5)

Combine AC, (CA and AC edges are same with 1 weight.)

[(A C 1) is used to combine the partial Tree.]

T2. Vertices: B PQ: (B D 3), (B A 4), (B C 5)
 T4. Vertices: D PQ: (D E 1), (D A 3), (D B 3), (D C 4)
 T5. Vertices: E PQ: (E D 1), (E C 2)
 T13. Vertices: AC PQ: (C A 1), (C E 2), (A D 3), (A B 4), (C D 4),
 (C B 5)

Merging BD,

[(B D 3) is used to combine the partial Tree.]

T5. Vertices: E PQ: (E D 1), (E C 2)
 T13. Vertices: AC
 PQ: (C A 1), (C E 2), (A D 3), (A B 4), (C D 4), (C B 5)
 T24. Vertices: BD
 PQ: (D E 1), (D A 3), (D B 3), (B A 4), (D C 4), (B C 5)

Merging EBD,

[(E, D, 1) is used to combine the partial Tree.]

T13. Vertices: AC

PQ: (C A 1), **(C E 2)**, (A D 3), (A B 4), (C D 4), (C B 5)

T524. Vertices: EBD

PQ: (D E 1), (E C 2), (D A 3), (D B 3), (B A 4), (D C 4), (B C 5)

Repeating,

[(C, E, 2) is used to combine the partial Tree.]

T13524-Vertices: ACEBD

PQ: (D E 1), (E C 2), (A D 3), (D A 3), (D B 3),
(A B 4), (C D 4), (B A 4), (D C 4), (C B 5), (B C 5)

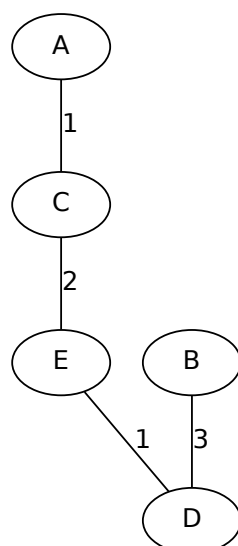
Result Tree i.e. MST:

Edges: {

(A C 1),
(B D 3),
(E D 1),
(C E 2)

}

And the MST is: $O(e \cdot \log \log n)$



Round Robin Algorithm

Algorithm 4.4.4: CHERITONTARJAN(G)

```

initialize  $Tree$  to contain no edges
for each  $u \in V(G)$ 
    { initialize  $T_u$  to contain only  $u$ 
      create  $PQ_u$ 
    do { for each  $v \rightarrow u$ 
        do add  $uv$  to  $PQ_u$ 
        comment: each edge will appear in two priority queue
      }
    comment: the forest currently has  $|G|$  nodes and 0 edges
 $t \leftarrow 0$ 
while  $t < |G| - 1$ 
    { select a component  $T_u$ 
      repeat
        select the minimum edge  $xy \in PQ_u$  and determine
        which components  $x$  and  $y$  are in, say  $x \in T_u$  and  $y \in T_v$ 
      until  $T_u \neq T_v$ 
    do { comment:  $xy$  connects two different components
        merge  $T_u$  and  $T_v$ 
        merge  $PQ_u$  and  $PQ_v$ 
        add  $xy$  to  $Tree$ 
         $t \leftarrow t + 1$ 
    }
  
```

Source: http://books.google.co.uk/books?id=zxSmHAoMiRUC&pg=PA78&lpg=PA78&dq=cheriton-tarjan%20algorithm&source=bl&ots=LQWYqzK7bq&sig=mmKjYZ7pevW3vkPD1Xmg7GC5B0I&hl=en&sa=X&ei=0AX4Usyk0_Cp7AbWiICADg&ved=0CC0Q6AEwAA#v=onepage&q=cheriton-tarjan%20algorithm&f=false

Greedy Algorithm

- uses local optimum solution hoping to get global optimum solution.
- this happens by chance only, not favorable for all the problems.

Currency Problem Using Greedy Algo:

>> Accumulate Rs 18 from the available currencies [Rs1, Rs2, Rs5, Rs10.] with minimum numbers.

Solution: Rs[10 + 5 + 2 + 1], **Worked.**

>> Accumulate Rs 15 from the available currencies [Rs1, Rs7, Rs10].

Solution: Rs[10 + 1 + 1 + 1 + 1 + 1]

Not Worked.

Could be: Rs[7 + 7 + 1]