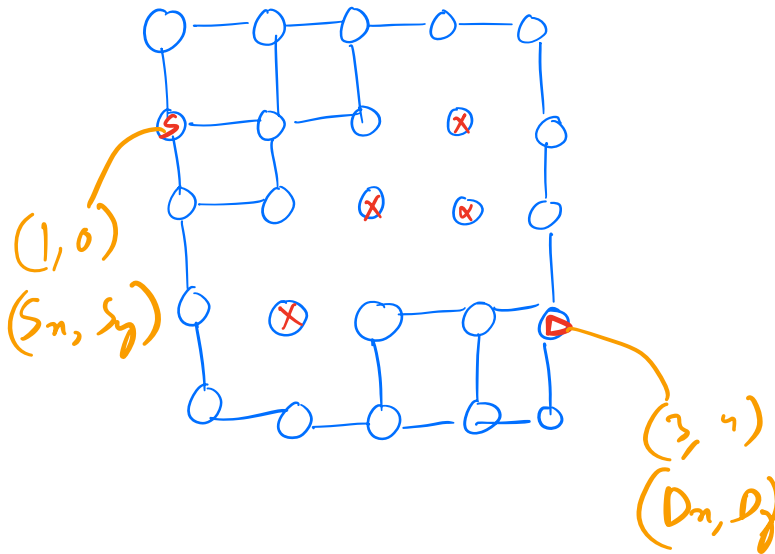
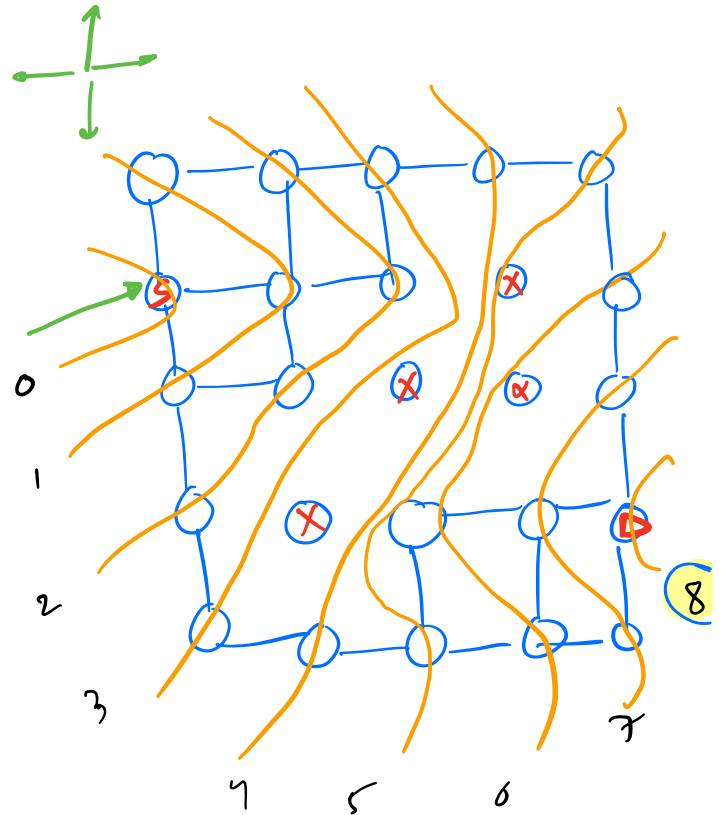
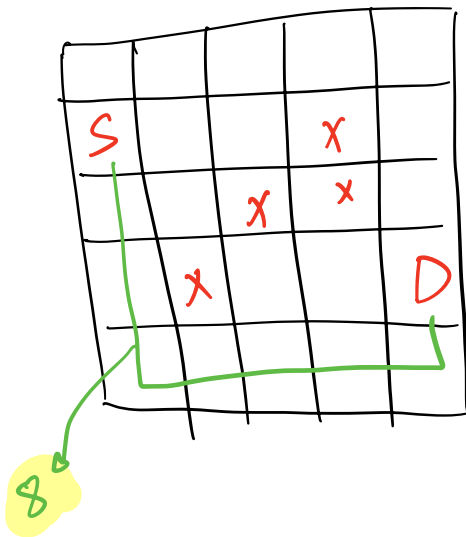


Q Given $N \times M$ matrix.
 Source & Destination.
 Find the shortest path b/w them

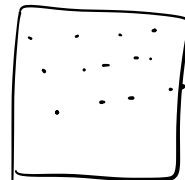
X \rightarrow Cannot visit.

$1 \leq N, M \leq 10^3$

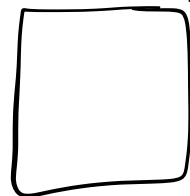


BFS

bool vis[][]



int dis[][]



int d[N][M] = {0};

bool vis[N][M];

// Sx, Sy, Dx, Dy.

Queue < pair < int, int > > q;

q.enqueue({ Sx, Sy });

d[Sx][Sy] = 0;

vis[Sx][Sy] = true;

while (! q.empty()) {

pair < int, int > p = q.front();

q.dequeue();

i = p.first, j = p.second;

for (k = 0; k < 4; k++) {

nI = i + dx[k];

nJ = j + dy[k];

if (check(nI, nJ) == true) {

q.enqueue({ nI, nJ });

d[nI][nJ] = 1 + d[i][j];

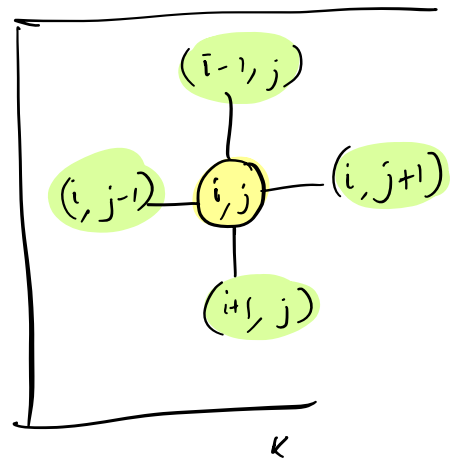
vis[nI][nJ] = true;

}

}

return d[dN][dJ];

pair < int, int > p;
↓ ↓
p.first p.second



dx[] = {-1, 0, 1, 0};
dy[] = {0, 1, 0, -1};

check (n, y)

1. inside matrix
2. not visited
3. not 'x'

```

bool check(n, y) {
    if (n >= 0 && n < N
        && y >= 0 && y < M
        && vis[n][y] == false
        && A[n][y] != 'X')
        return true;
    return false;
}

```

check (n, y)

- ✓ 1. Inside matrix
- ✓ 2. not visited
- ✓ 3. not 'X'

$G = \{V, E\}$

$V = N \times M$

$E = N \times M$

$TC = O(V+E)$

$NM + NM$

$TC = O(NM)$

$SC = O(V)$

$SC = O(NM)$

HW

Dfs on 2D Matrix !!!

Q Given a 2D matrix.


Residence (R)

Hospital (H)

for every (R), find the min. distance to a (H)

+

R	R	R	H
R	R	H	H
R	H	H	R



3	2	1	0
2	1	0	0
1	0	0	1

1) $\forall R$ do bfs to find the nearest (H)

\downarrow \downarrow

$N \times M$ \times $N \times M$

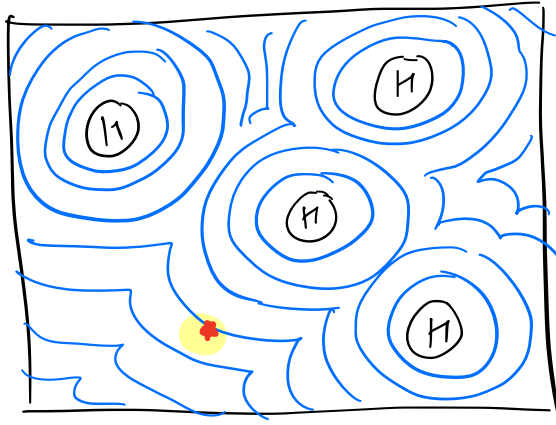
TC $\sim N^2 M^2$

2) $\forall H$ do bfs & update the shortest dist of every (R)

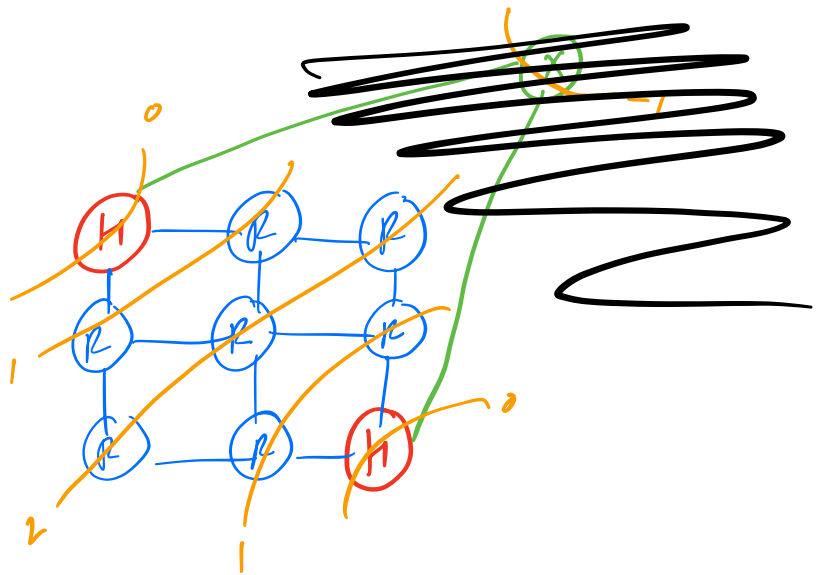
\downarrow \downarrow

$N \times M$ $N \times M$

TC $\sim N^2 M^2$



H	R	R
R	R	R
R	R	H



Initial :

q. enqueue (All the **H** coordinates)

→ $dis[i][j] = 0$ (\forall **H**)

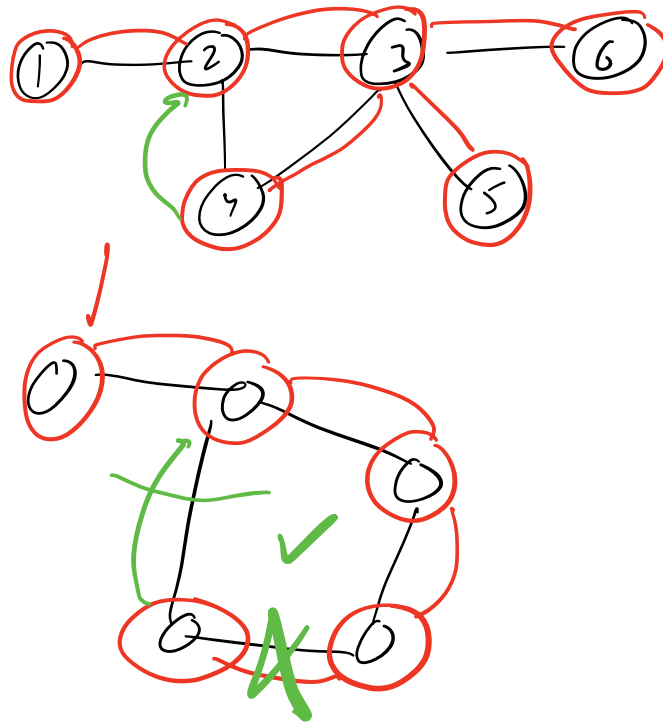
→ $vis[i][j] = true$ (\forall **H**)

start while (! q. isEmpty()) {

≡

}

Q Given an undirected graph $[1-N]$.
Check if a cycle is present or not!



Idea: If a node has any neighbour: other than the node
which is already visited
 \Rightarrow cycle!

bool isCycle = false;

```
void dfs(int v, int p) {
```

```
    vis[v] = true;
```

```
    for (u: adj[v]) {
```

```
        if (u != p) {
```

```
            if (vis[u] == false) {
```

```
                dfs(u, v);
```

```
            }
```

```
        } else {
```

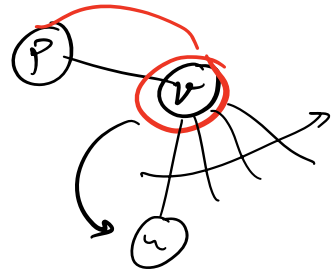
```
            isCycle = true;
```

```
        }
```

```
    }
```

```
}
```

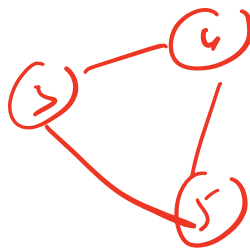
```
}
```



$T = O(V + E)$

$SC = O(V)$

$dfs(1, 0);$

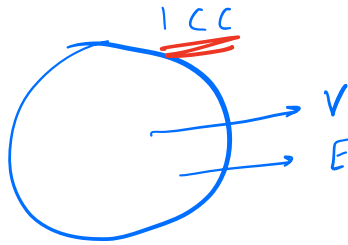


```
f(i2 | -> n)
  if (!vis[i3])
    dfs(i, 0)
```

NOTE:

REMOVE BIAS
GRAPH MIGHT NOT
ALWAYS BE
CONNECTED!

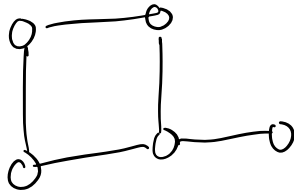
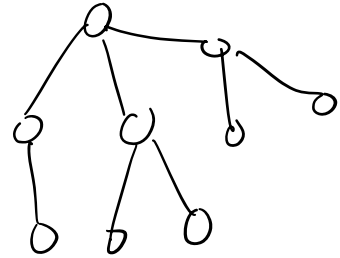
II



$E > V - 1$
→ Cycle!

tree

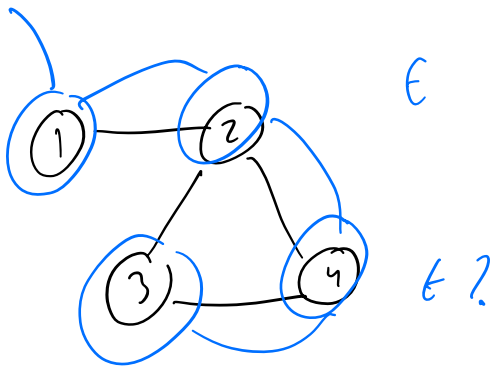
V, E



$V = 5$
 $E = 5$

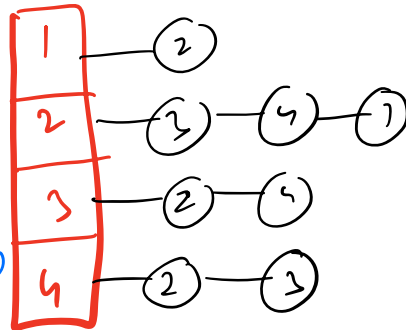
$E = V - 1$ No cycle

$E > V - 1$: Cycle!



$V \rightarrow \text{dfs} : \text{cnt}++$

$E \rightarrow \text{dfs} : \frac{\text{sz} += \text{adj}[v].\text{size}()}{2}$

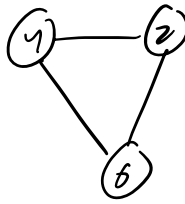
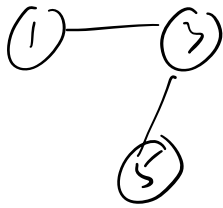


```
dfs(v) {
    vis[v] = true;
    cnt++;
    sz += adj[v].size();
    for (u: adj[v]) {
        if (vis[u] == false) {
            dfs(u);
        }
    }
}
```

$V = \text{cnt};$
 $E = \text{sz}/2;$

if ($E > V - 1$)
→ Cycle!

else
→ No cycle!



```

f ( i = 1 → N ) {
  if ( !vis[i] ) {
    cnt = 0; sz = 0;
    dp(i);
  }
}
  
```

$V = cnt;$

$E = sz/2;$

if ($E > V - 1$)
 → cycle!

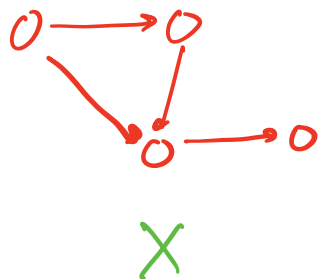
}

→ NO cycle!

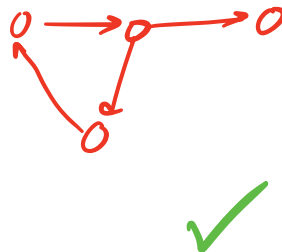
$T.C = O(V + E)$

$S.C = O(V)$

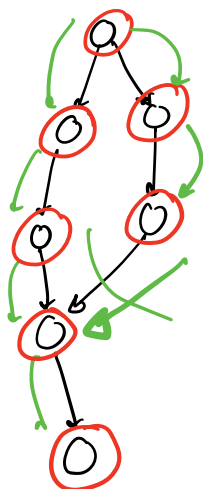
I Given a directed graph. Detect cycle!



X

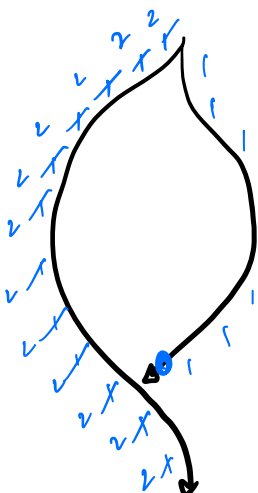


✓



VIS

0 : NOT VIS
1 : VIS & IN STACK
2 : VIS & OUT



X



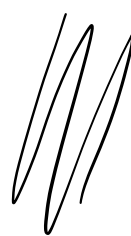
✓

```

void dfs(v) {
    vis[v] = 1;
    f(u: adj[v]) {
        if (vis[u] == 0) {
            dfs(u);
        }
        else if (vis[u] == 1) {
            isCyclic = true;
        }
    }
    vis[v] = 2;
}

```

$TC = O(V + E)$
 $SC = O(V)$

$f(i: 1 \rightarrow v)$
 CC