

* Carry Forward

Count the pairs "ag"

Given a string, calculate the no. of pairs (i, j)

: $i < j$ & $s[i] == 'a'$ & $s[j] == 'g'$

0 1 2 3 4 5 6 7
s = b a a g d c a g

(1, 3), (1, 7)
(6, 7) (2, 3)
(2, 7) → 5

I) BF

```

cnt = 0;
for (i = 0; i < N; i++) {
    for (j = i + 1; j < N; j++) {
        if (s[i] == 'a' & s[j] == 'g') {
            cnt++;
        }
    }
}
return cnt;

```

TC: $O(N^2)$

SC: $O(1)$

```

    cnt = 0;
    for (i = 0; i < N; i++) {
        if (s[i] == 'a') {
            for (j = i + 1; j < N; j++) {
                if (s[j] == 'g') {
                    cnt++;
                }
            }
        }
    }
    return cnt;

```

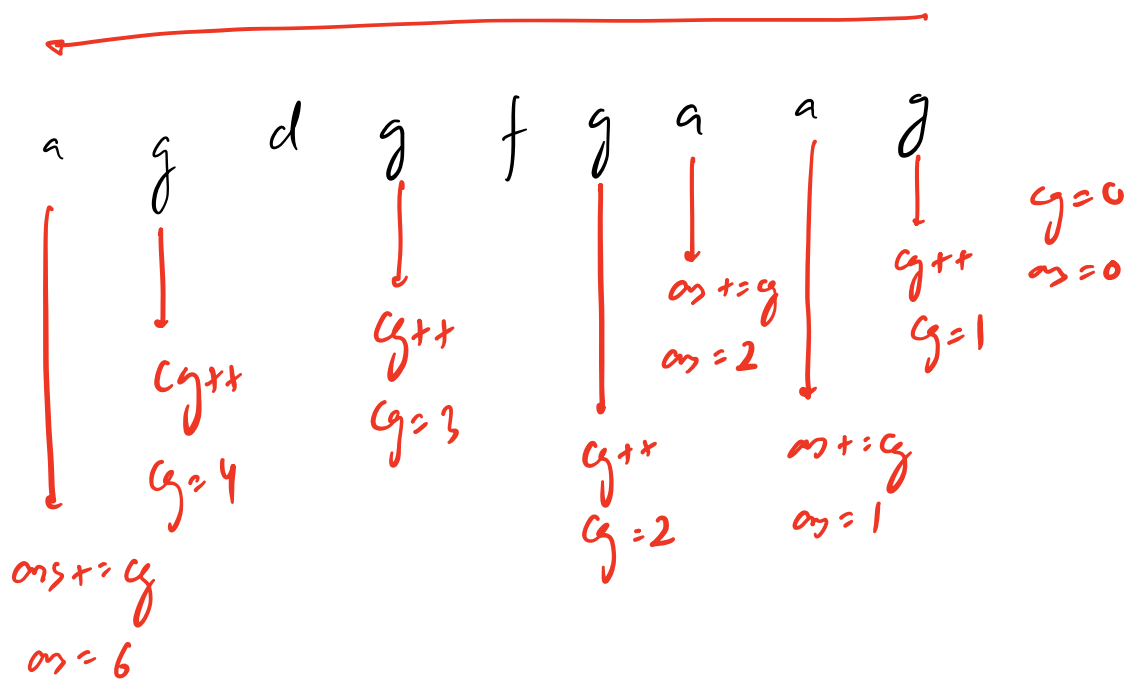
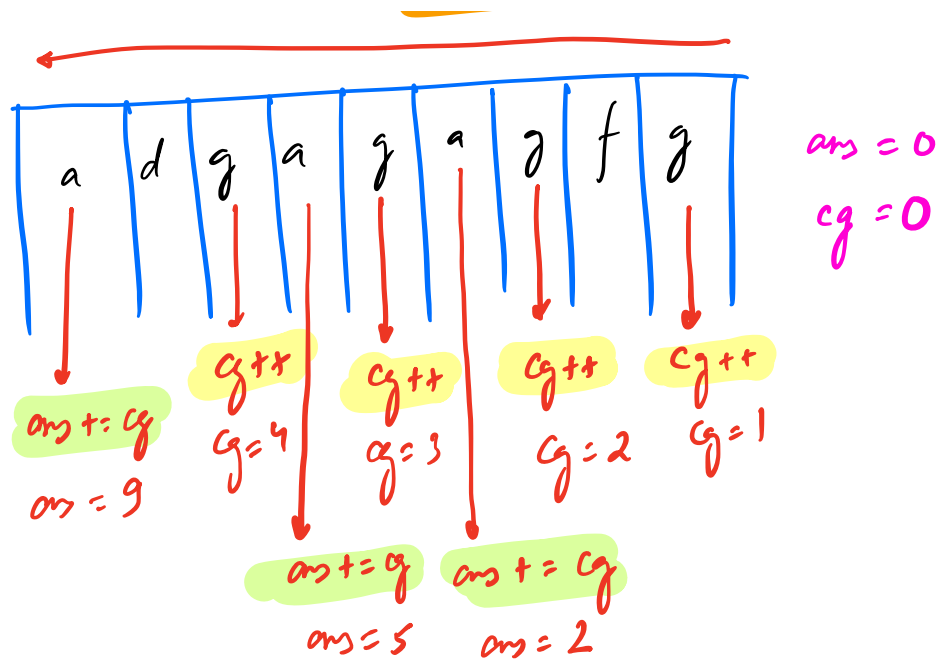
TC: $O(N^2)$

SL = $O(1)$

Worst case: "a a a a — a a"

S = 0 1 2 3 4 5 6 7 8
 a d g a g a g f g
 ↑ ↑ ↑
 4 + 3 + 2 → 9 ANS

Obs: for every 'a', count the no of 'g' on the right side of it!



cg = 0, ans = 0

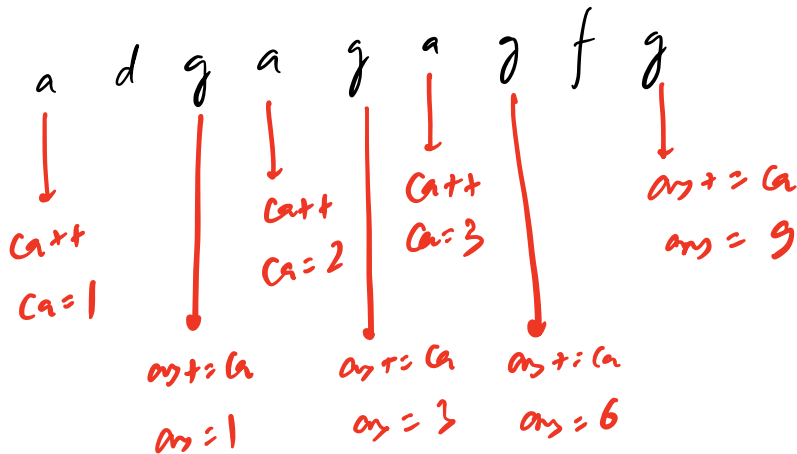
```
f(i = N-1; i >= 0; i--) {  
    if (s[i] == 'g') {  
        cg++;  
    }  
    else if (s[i] == 'a') {  
        ans += cg;  
    }  
}
```

TC: $O(N)$

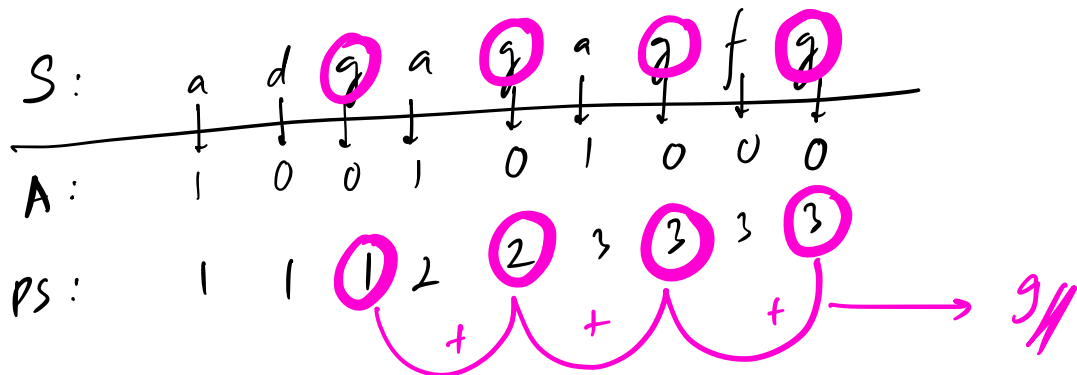
SC: $O(1)$

return ans;

cg = 0
ans = 0



III Prefix Sum



"a x g" (i, j, k) $i < j < k$
 $s(i) = 'a', s(j) = 'x', s(k) = 'g'$

$S:$	a	x	a	a	x	g	x	g	a	a	x
$PA:$	1	1	2	3	3	3	3	3	4	5	5
$SG:$	2	2	2	2	2	1	1	0	0	0	0

(Green vertical bars highlight the 'x' characters in S and the corresponding prefix sum values in PA and SG: 1, 3, 3, 3, 3, 5. Pink '+' signs are placed below the SG values: 2, +, 6, +, 3, +, 0. An arrow points from the final result 11 to the complexity analysis.)

TL: $O(N)$
 SC: $O(N)$
 $O(1)$

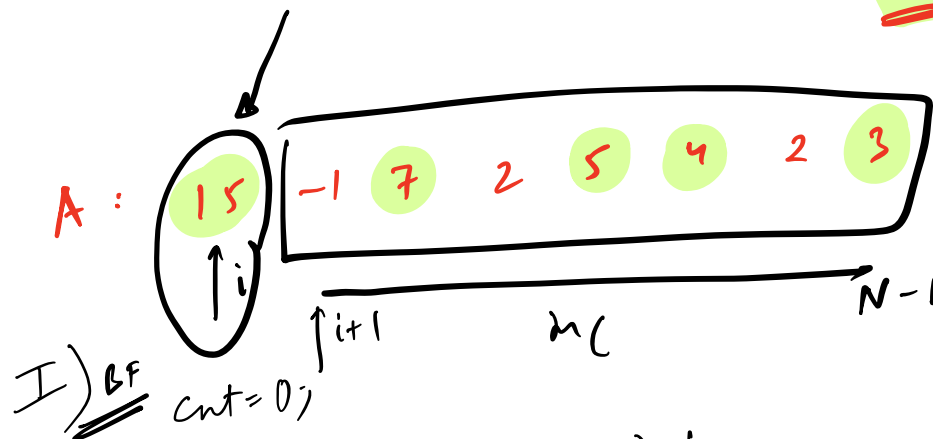
= 11

Q Given an array A.

Find the no. of leaders in the array!

An element is a leader if it is greater than all the elements on its right side!

NOTE: $A[N-1]$
is ALWAYS a leader!



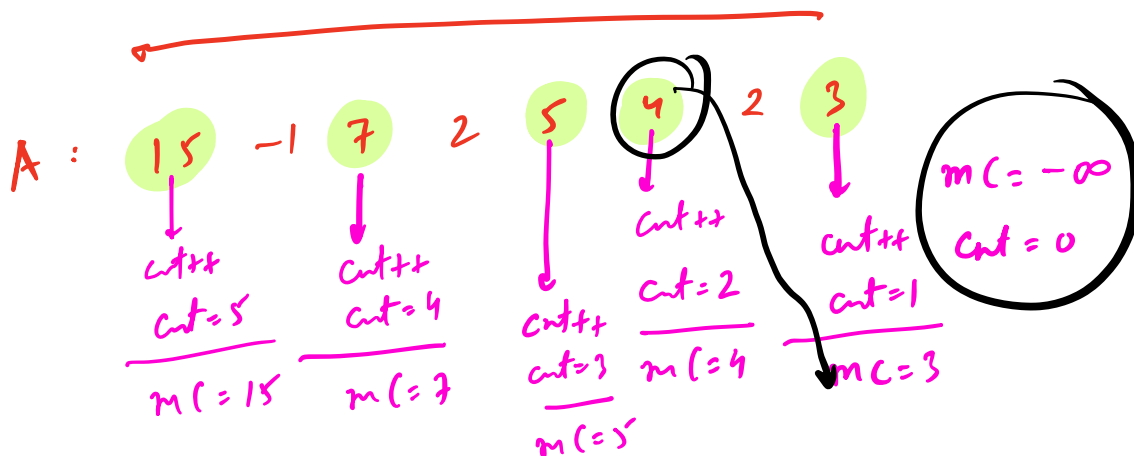
```
{ i = 0; i < N; i++ }  
  mC = -∞;  
  { j = i+1; j < N; j++ }  
    if ( A[j] > mC )  
      mC = A[j];  
  }  
  if ( A[i] > mC )  
    cnt++;  
}
```

```
return cnt;
```

TC = $O(N^2)$

SC = $O(1)$

II



```

mC = -∞, cnt = 0
for (i = N-1; i ≥ 0; i--) {
    if (A[i] > mC) {
        cnt++;
        mC = A[i];
    }
}
return cnt;

```

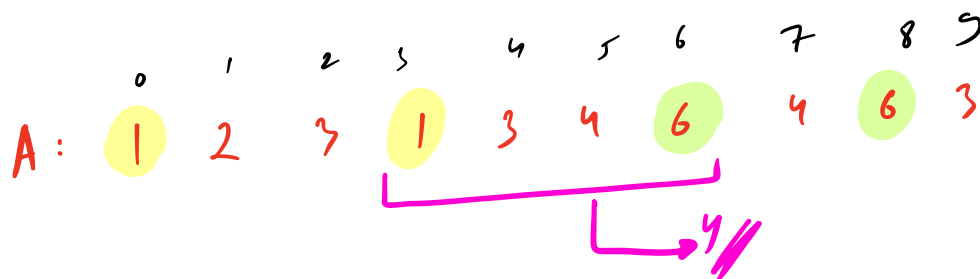
TC = $O(N)$

SC = $O(1)$

Q

Closest Min Max

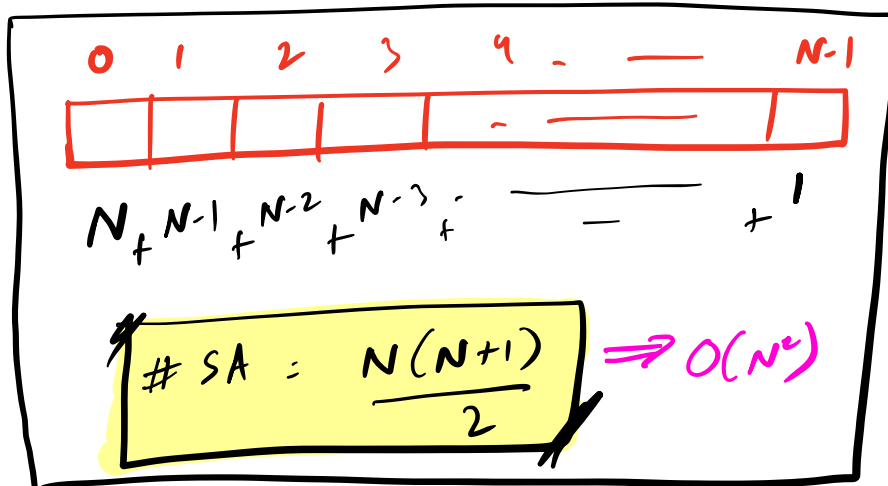
Given an array, find the length of the smallest subarray which contains both the MIN & MAX of the array!



I) BF

1. Find MIN $\rightarrow O(N)$
2. — MAX $\rightarrow O(N)$
3. Consider every S.A. $[L, R]$ $\rightarrow N^2$
 - \rightarrow check if it contains MIN $\rightarrow O(N)$
 - \rightarrow — MAX
 - \rightarrow YES: $ANS = \min(ANS, R-L+1)$
 - \rightarrow NO: X

TC: $O(N^3)$



A: 8 8 8 8 \rightarrow ANS = 1

if (MIN == MAX)
ANS = 1

A: 2 2 6 4 5 1 5 2 6 4 1
MAX
MIN
MAX
MIN

Obs

1.

- MIN - MAX - MIN - X

- MAX - MAX - MIN - X

the ANS S.A. would contain
EXACTLY 1 occ of both MIN & MAX!

2.

- MIN - MAX - X

MIN - MAX

the ANS S.A. would contain
MIN & MAX at boundaries!

↳ Gen I :

MIN MAX

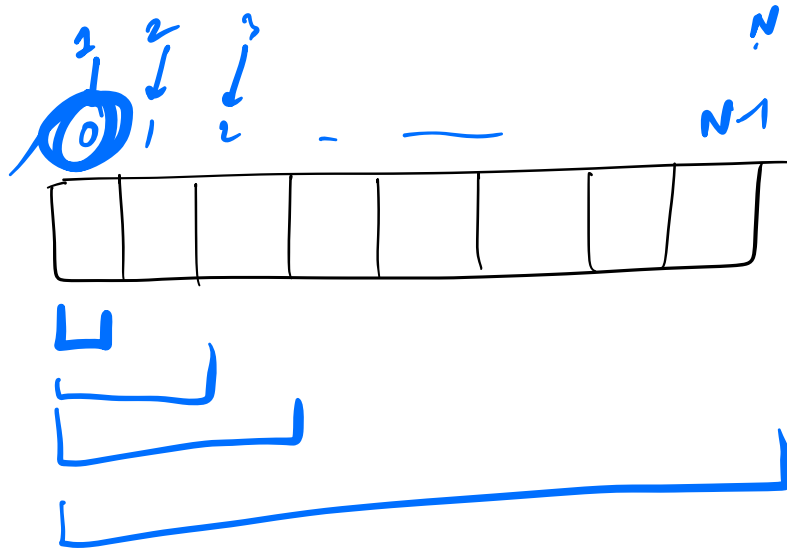
↳ Gen II :

MAX MIN

A: 0 1 2 3 4 5 6 7 8
 1 3 6 1 3 1 6 4 6

MIN_i MAX_i

ANS = 3/2



1. Find MIN $\rightarrow N$

2. Find MAX $\rightarrow N$ \rightarrow if (MIN == MAX) ret 1;
 i, j

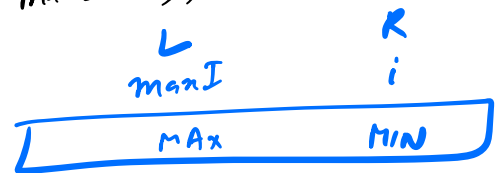
3. minI = -1, maxI = -1, ans = N

$f(i=0; i < N; i++) \{ \rightarrow N \times O(1)$

if (A[i] == MIN) {
 if (maxI != -1) {
 ans = min(ans, i - maxI + 1);
 }
 minI = i;
}

else if (A[i] == MAX) {
 if (minI != -1) {
 ans = min(ans, i - minI + 1);
 }
 maxI = i;
}

}
 ret ans;



$$N + N + N = 3N$$

$$TC: O(N)$$

$$SC: O(1)$$