

② Build LPS using KMP

$S = c a c y c a c a$

$LPS[i] \rightarrow$

$LPS[i] \rightarrow LPS(s[0..i])$

$LPS(string) \rightarrow$ the length of longest prefix which is also a suffix.

NOTE: do not consider complete string!

$LPS(ababab) \rightarrow 4$

P	S
a	b x
ab	ab ✓
aba	bab x
abab	abab ✓
ababa	babab x
ababab	

$S = c \ a \ c \ g \ c \ a \ c \ a$
 $LPS[] \rightarrow 0 \ 0 \ 1 \ 0 \ 1 \ 2 \ 3 \ 2$

BF

$$TC = O(N^3)$$

$S: s_0 \ s_1 \ s_2 \ s_3 \ s_4 \ \dots \ s_{i-5} \ s_{i-4} \ s_{i-3} \ s_{i-2} \ s_{i-1} \ s_i \ \dots \ s_N$
 $LPS[]$

Assume: $LPS[i] = 5$

$$s_0 \ s_1 \ s_2 \ s_3 \ s_4 = s_{i-4} \ s_{i-3} \ s_{i-2} \ s_{i-1} \ s_i$$

$$LPS[i-1] \geq 4$$

$$\geq 5 - 1$$

$$LPS[i-1] \geq LPS[i] - 1$$

$$LPS[i] \leq LPS[i-1] + 1$$

$LPS[]$

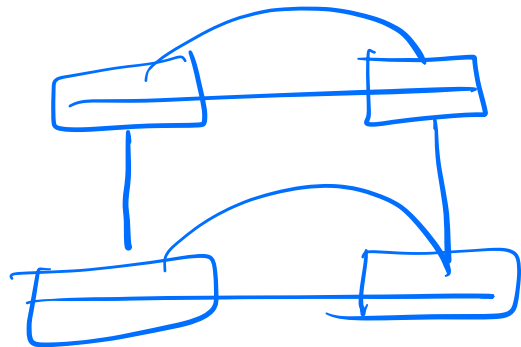
$$i-1 \quad i$$

$$6 \leq 7$$

$S =$ a b a y a b a ch
 LPS \rightarrow 0 0 1 0 1 2 3 ≤ 4
 if ($s[3] == ch$)
 LPS[i] $\rightarrow 4$

$S =$ b c a d c b c a d ch
 LPS[i] \rightarrow 0 0 0 0 0 1 2 3 4 ≤ 5
 if ($s[4] == ch$)
 LPS[i] $\rightarrow 5$

$S =$ c a c y c a c a b c a c y c a c y
 LPS[i]: 0 0 1 0 1 2 3 2 0 1 2 3 4 5 6 7 ≤ 8
 $x = \text{LPS}[i-1]$
 if ($s[i] == s[x]$)
 LPS $\rightarrow x+1$ X



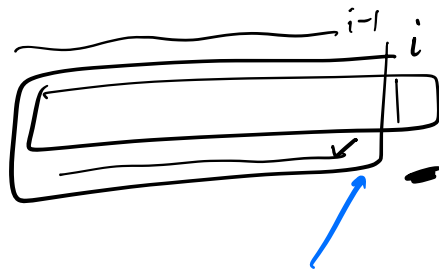
\rightarrow 10 11
 x k

CODE :

```

LPS[N];
LPS[0] = 0;
for (i = 1; i < N; i++) {
    n = LPS[i-1];
    while (s[n] != s[i]) {
        if (n == 0) { n = -1; break; }
        n = LPS[n-1];
    }
    LPS[i] = n + 1;
}
return LPS[];

```



TC :

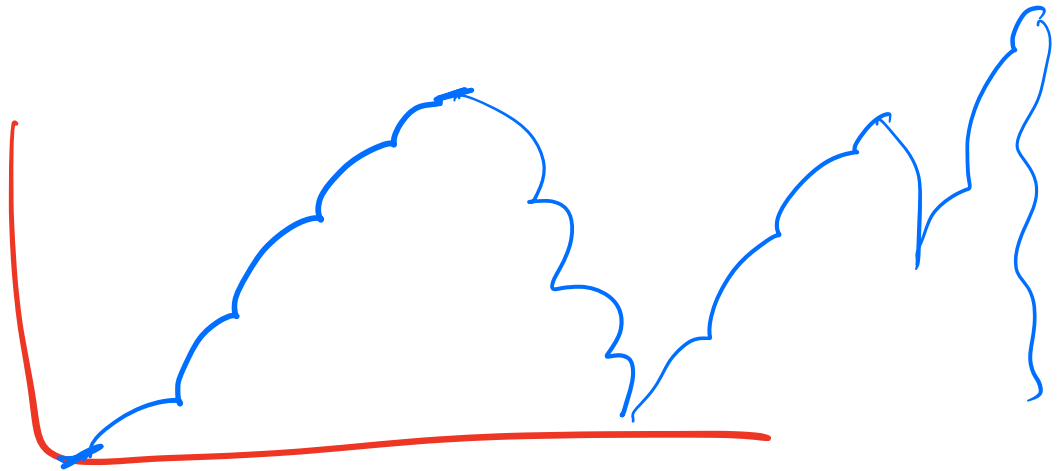
s : 0 1 2 3 4 0 1 2 3 4 - - 10 0
 LPS : 0 1 2 3 4 0 1 2 3 4 - - 10 0
 ops : 0 1 1 1 1 5 1 1 1 1 - - 11 - -

dec <= # inc

$\#inc \rightarrow N$
 $\#dn \rightarrow N$

$T = O(N)$

LP)

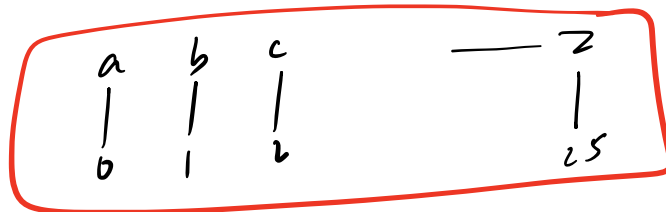


	$\nearrow n$						i
a	b	c	d	e	f	g	
0	0	0	0	0	0	0	0

a	a	a	a	a	a	a	
0	1	2	3	4	5	0	
	1	1	1	1	5		

$SC \rightarrow O(N)$

① Polynomial Rolling Hash
 → find the hash of a string



$abc : 0 + 1 + 2 \rightarrow 3$

$aaa : 0 + 0 + 0 \rightarrow 0$
 $aa : 0 + 0 \rightarrow 0$
 $a : 0 + 0 \rightarrow 0$



$aaa : 1 + 1 + 1 \rightarrow 3$
 $aa : 1 + 1 \rightarrow 2$
 $a : 1 \rightarrow 1$

$abc : 1 + 2 + 3 \rightarrow 6$
 $acb : \rightarrow 6$
 $bac : \rightarrow 6$
 $bca : \rightarrow 6$
 $cab : \rightarrow 6$
 $cba : \rightarrow 6$

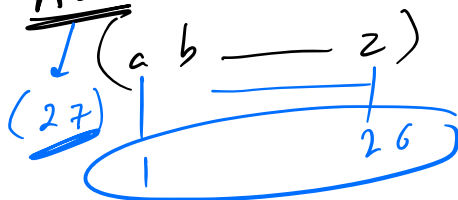
$aa \rightarrow 2$
 $b \rightarrow 2$

5416 :

DEC : $\rightarrow 5 \times 10^3 + 4 \times 10^2 + 1 \times 10^1 + 6 \times 10^0$

(10) 11

ALP:



~~0 — 9~~

0 — 9

3 2 1 0

a c d a

$$a \times 27^3 + c \times 27^2 + d \times 27^1 + a \times 27^0$$

$$1 \times 27^3 + 3 \times 27^2 + 4 \times 27^1 + 1 \times 27^0$$

integer

S: (s₀ s₁ s₂ — s_{N-1})

$$h(s) = s_0 + s_1 \times 27^1 + s_2 \times 27^2 + \dots + s_{N-1} \times 27^{N-1}$$

$$h(s) = (s_0 + s_1 \times 27^1 + s_2 \times 27^2 + \dots + s_{N-1} \times 27^{N-1}) \% (10^9 + 7)$$

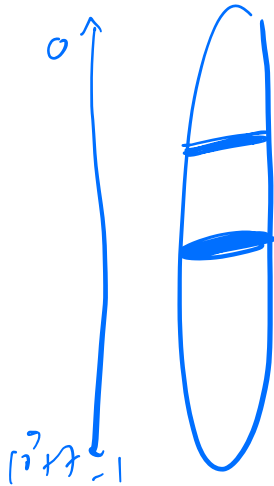
$$h(s) = (S_0 + S_1 \times M^1 + S_2 \times M^2 + \dots + S_{N-1} M^{N-1}) \% P$$

M, P : prime No

M : 29, 31

P : $10^9 + 7$, $10^9 + 9$

$$N \leq 10^5$$



$$\begin{aligned} S_1 &\rightarrow h_1 \rightarrow \frac{P}{0} \\ S_2 &\rightarrow h_2 \rightarrow \frac{1}{10^9+7} \\ S_3 &\rightarrow h_3 \rightarrow \frac{2}{10^9+7} \\ &\vdots \\ S_N &\rightarrow h_N \rightarrow \frac{N-1}{10^9+7} \end{aligned}$$

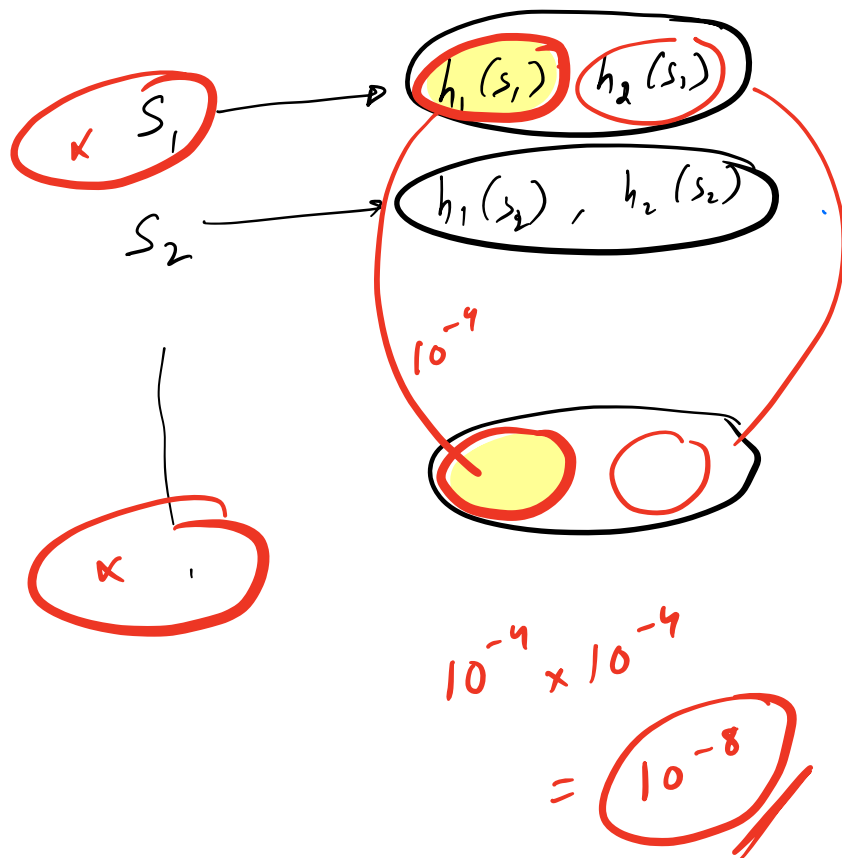
$$\sim \frac{10^5}{10^9} \sim 10^{-4}$$

$$h(s) = (s_0 + s_1 \times m^1 + s_2 \times m^2 + \dots + s_{n-1} m^{n-1}) \% p$$

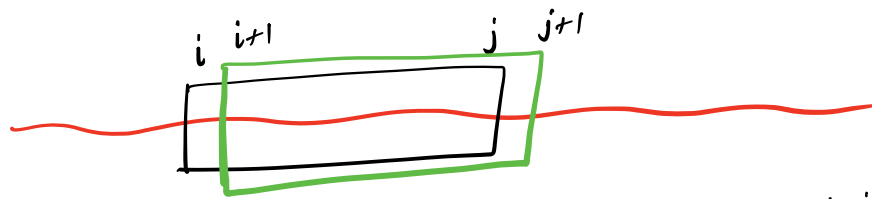
$$h_1(s) \longrightarrow M: 29, \quad P = 10^9 + 7$$

$$\begin{array}{l} h_1(s) \xrightarrow{\quad} \\ h_2(s) \xrightarrow{\quad} \end{array} \quad M: 31, \quad p = 10^3 + 9$$

$$N \leq 10^5$$



Polynomial Rolling Hash



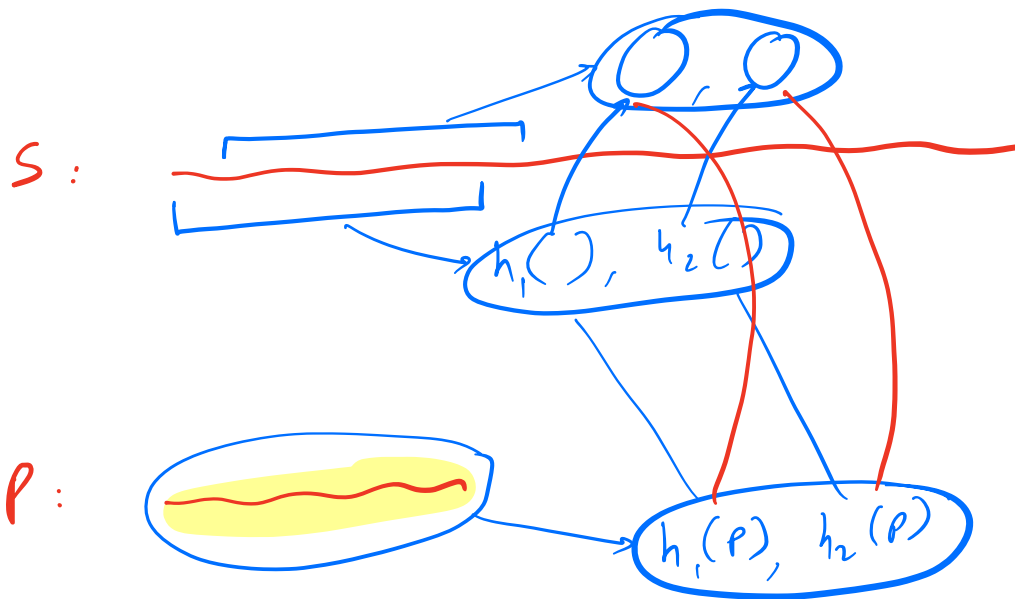
$$h(s[i-j]) = (s[i] + s[i+1]M + \dots + s[j] \cdot M^{j-i}) \% P$$

$$h(s[i+1-j+1]) = (s[i+1] + s[i+2]M + \dots + s[j+1]M^{j-i}) \% P$$

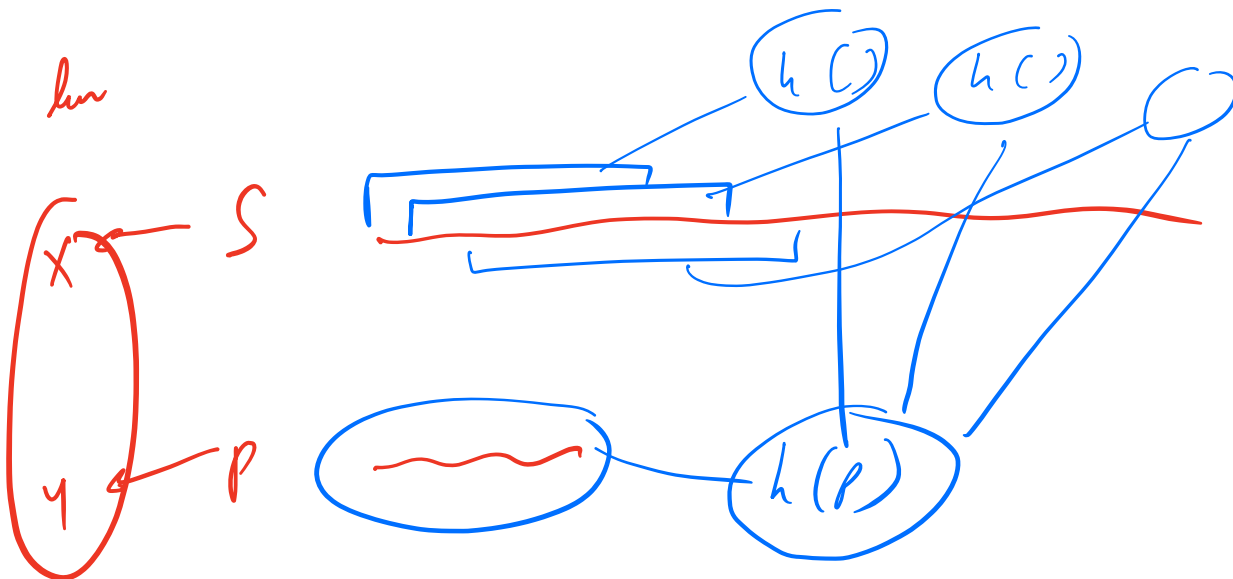
$$h(s[i+1-j+1]) = \left[\frac{(h(s[i-j]) - s[i])}{M} + s[j+1]M^{j-i} \right] \% P$$

Inverse Modulo $\rightarrow \log_2(P) / O(1)$

Q Given a text (T) & pattern (P).
Find if P is in T as a substring!



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$$TC = y + y + (x - y + 1) \times O(1) + h(P)$$

$$\boxed{TC = O(x + y + h(P))}$$