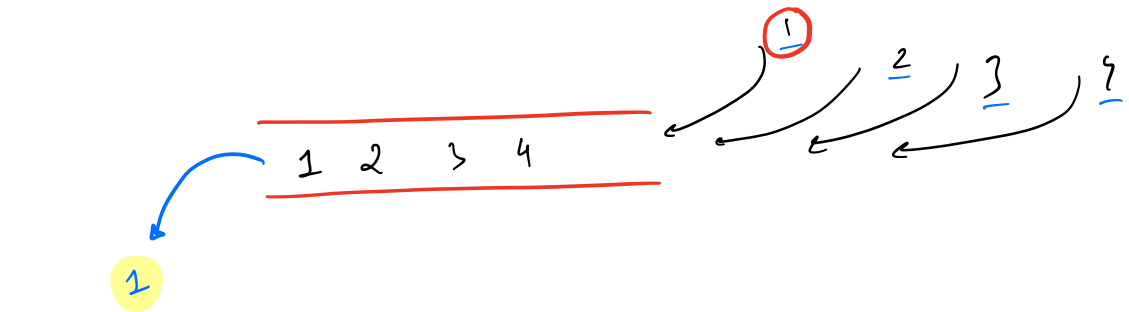
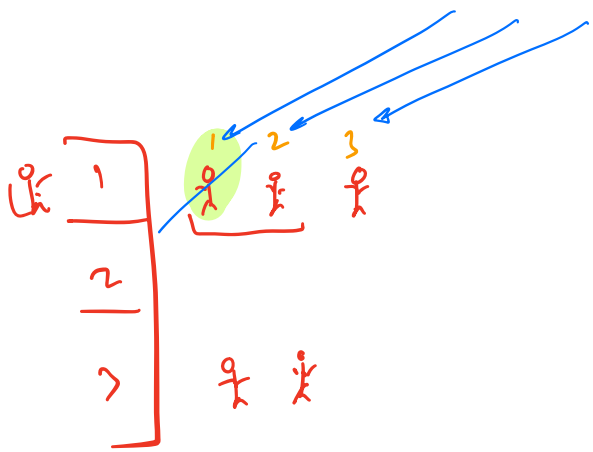
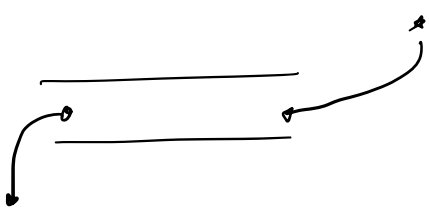
Queue → FIFO

First In First Out





⊛ Ops of Queue



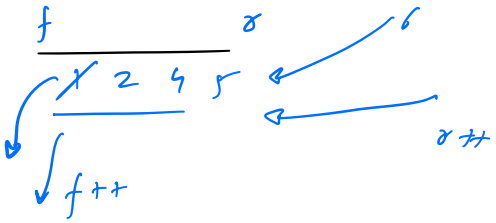
1) enqueue (x) : Insert x at the back of the Queue

2) dequeue () : Removing the frontmost.

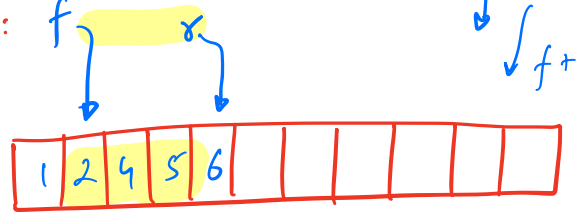
3) front () : get the frontmost.

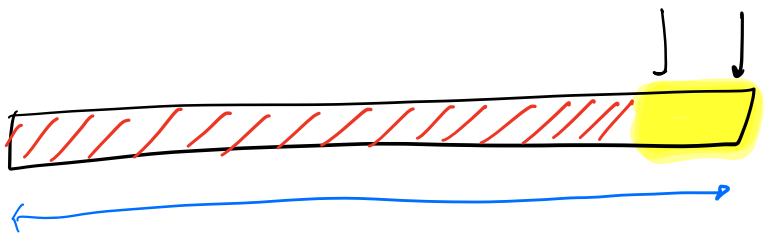
4) size () : get the # of elements..

① Implementation



① Array: f r



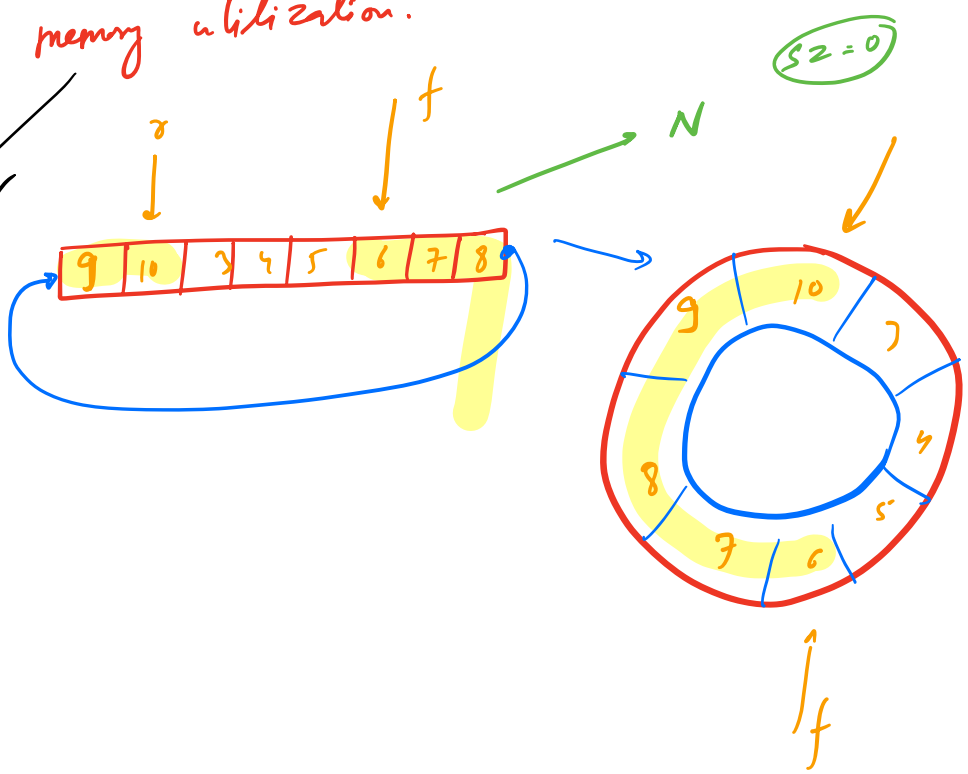


Problem:

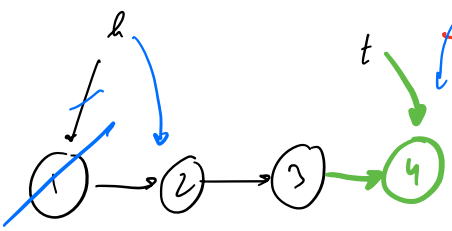
1) Need to know the # of enqueue() ! → Dynamic Array !

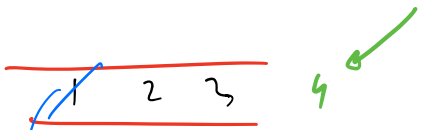
2) Proper memory utilization.

CIRCULAR ARRAY



⊙ LL





enqueue: O(1)

front : O(1)

dequeue : O(1)

size() : O(1)

s2 = ~~1 / 2 / 3~~ 3

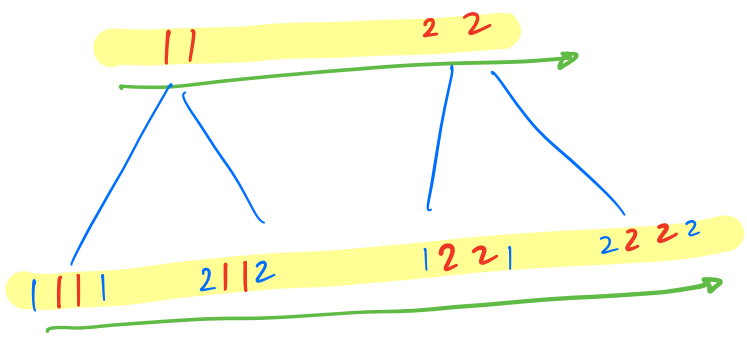
Q Given N. Generate first N perfect No's.

→ Even no. of digits

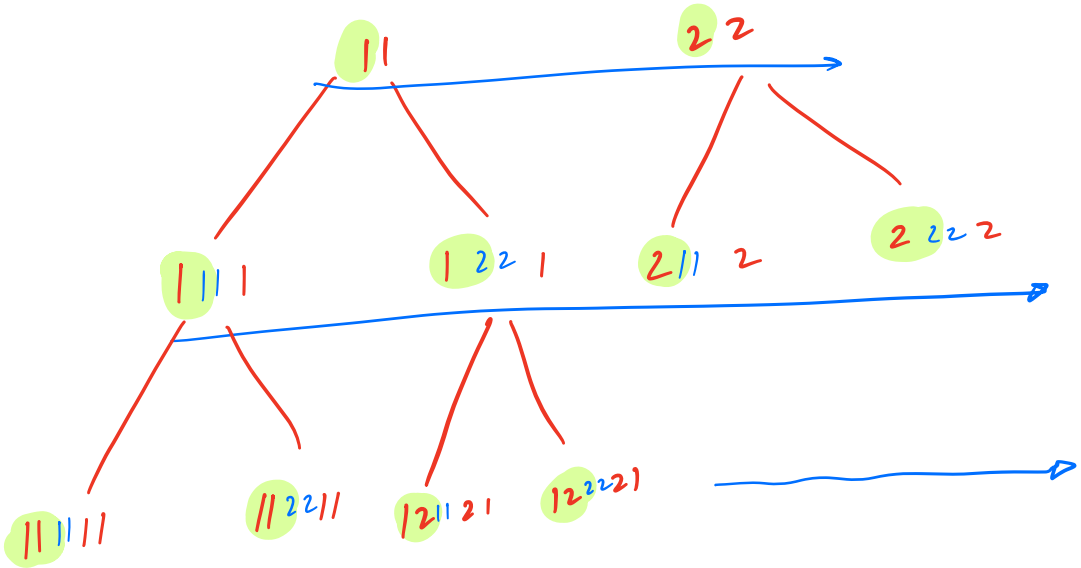
→ palindrome

→ digits ∈ (1, 2)

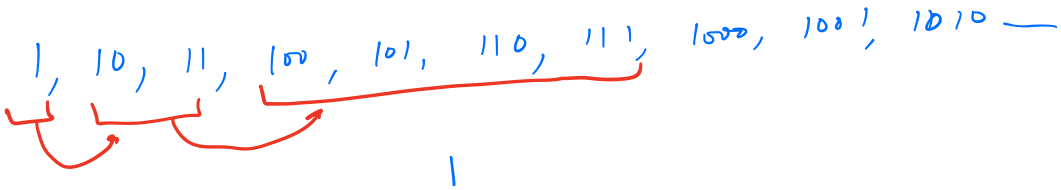
11, 22, 1111, 1221, 2112, 2222, 111111

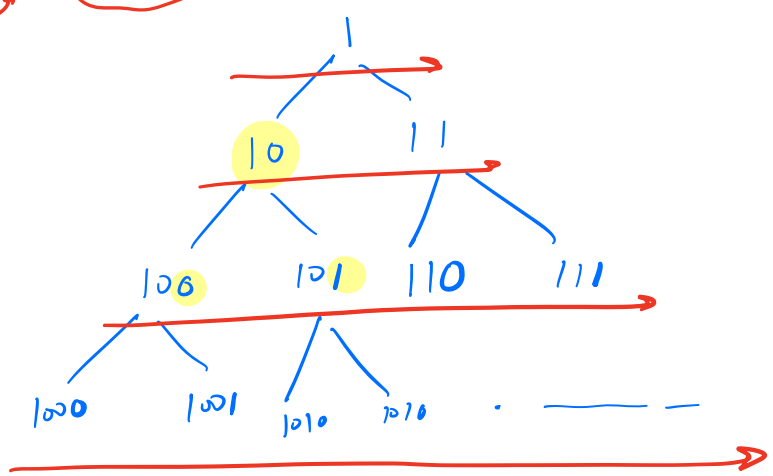


X



Q. Generate all binary no's from 1 — N

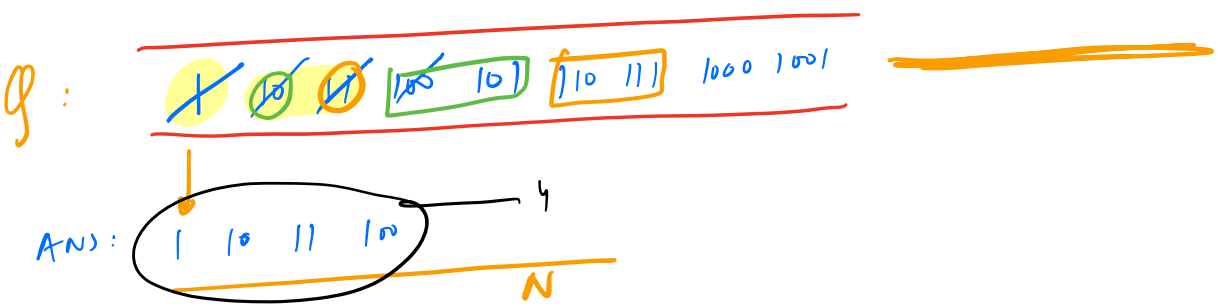




cnt → cnt the enqueue()

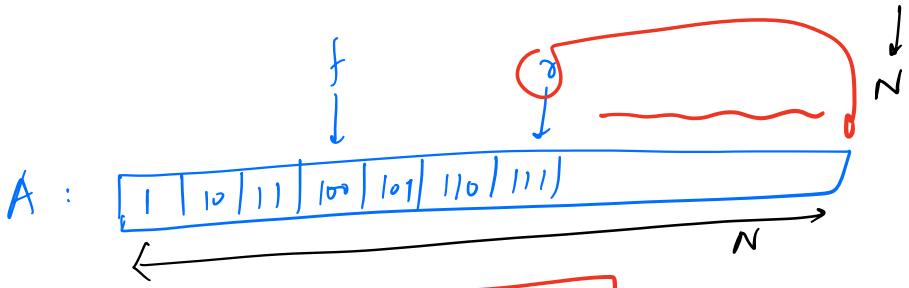
cnt == N

STOP!



ANS : 1 10 11 100 4

N



TC : O(N)

SC : O(1)

Q Given a stream of characters, after adding every character to a string, find the first non-repeating char so far.

stream: a b c a d e d b e c

→ a a a b b b b c c #

a b c a d e d b e c

a a a b b b b c c #

CL

~~a~~ ~~b~~ ~~c~~ ~~d~~ ~~e~~

HM

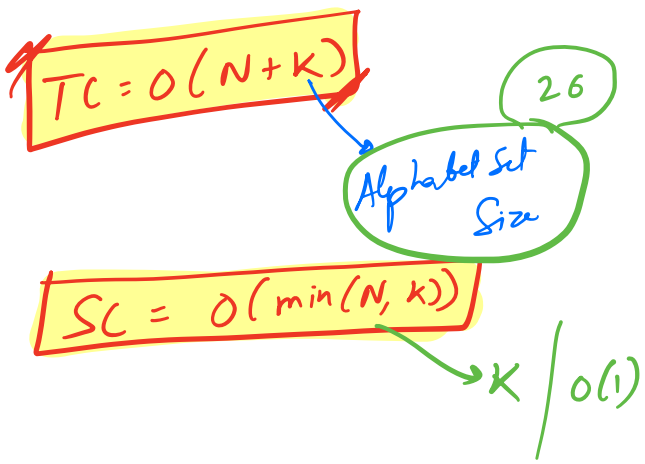
<a, 2>
<b, 2>
<c, 2>
<d, 2>
<e, 2>

```
CL : Queue
HashMap<char, int> hm;

for (i:0; i<N; i++) {
    ch = S[i];
    hm[ch]++;
    if (hm[ch] == 1) {
        CL.enqueue(ch);
    }
    while (!CL.isEmpty() && hm[CL.front()] > 1)
        CL.dequeue();

    if (!CL.isEmpty())
        ANS(i) = CL.front();
    else ANS(i) = '#';
}
```

N

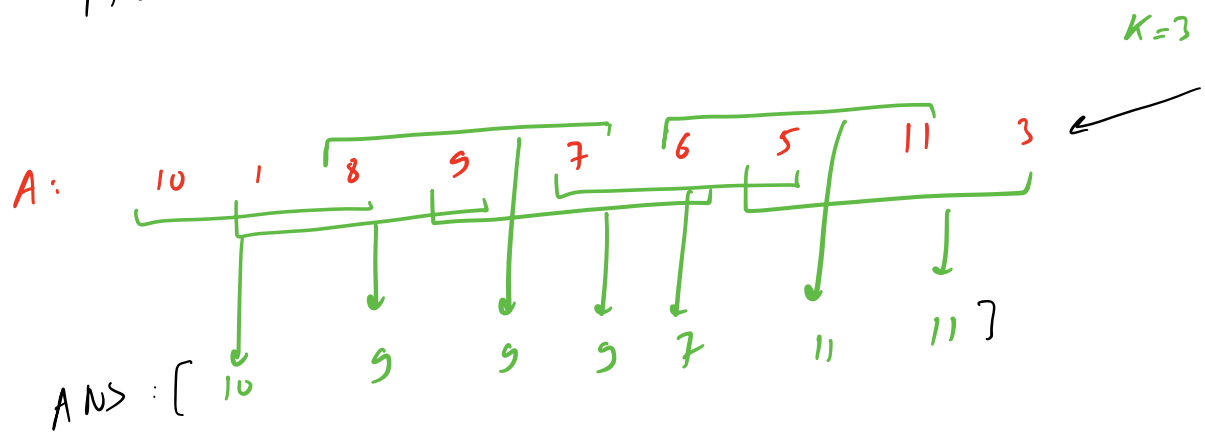
K : 26

K : 26



SLIDING WINDOW MAX

Q Given an Array A.
Find the max. element for all SAs of size K.



I) BF

∀ SA ⟶ N-K+1

find MAX ⟶ O(K)

(N-K)K

~ N² : K : N/2

II) Tree Map

TC = O(N log K)

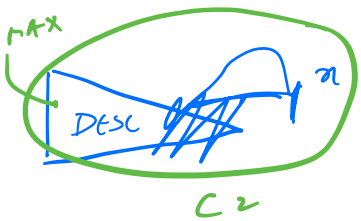
III

K = 3

10 1 8 9 7 6 5 11 3

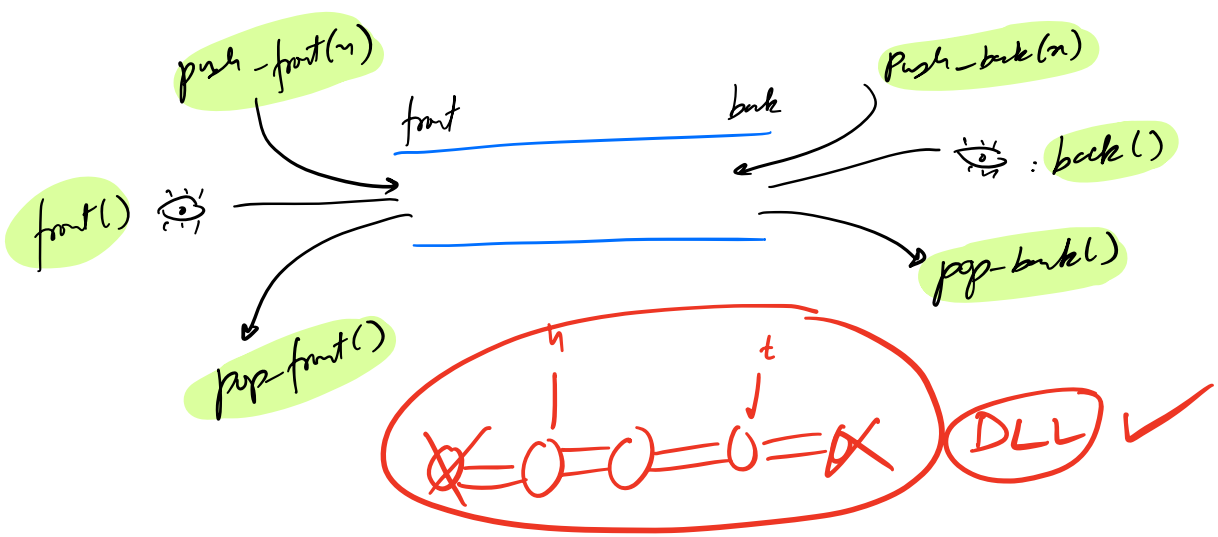
ANS: 10 9 5 9 7 11 11

CL: ~~10~~ X ~~8~~ ~~5~~ ~~7~~ ~~6~~ ~~8~~ 11 3

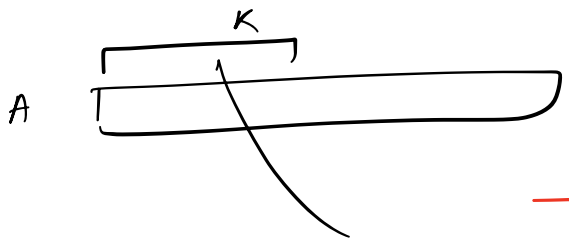


Requirement of CL ?

DEQUE : Doubly Ended Queue



```
// A[N], K
deque<int> CL;
vector<int> ans;

for(i:0; i<K; i++) {
    while( !(CL.isEmpty() && A[CL.back()] <= A[i]) {
        CL.pop_back();
    }
    CL.push_back(i)
}
ans.push_back(A[CL.front()]);

L=1, R=K;
while(R<N) {
    if( CL.front() == L-1) {
        CL.pop_front();
    }
    while( !(CL.isEmpty() && A[CL.back()] <= A[R]) {
        CL.pop_back();
    }
    CL.push_back(R);
    ans.push_back(A[CL.front()]);
    L++; R++;
}
```

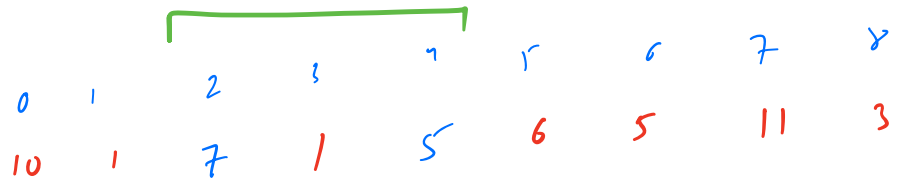


O(K)

N-K

TC = O(N)

SC = O(K)



ANS: 10 7 7

CL: 0 1 2 3 4

(INDEX)