Semiconductor Memory Organization & Addressing

Arquitectura de Computadoras 2022

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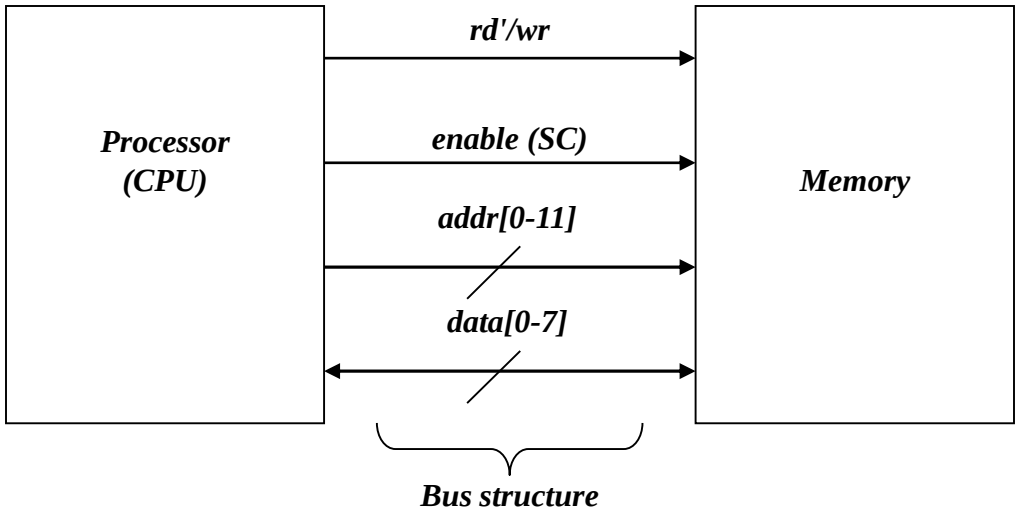
Introduction

- Computer systems functionality aspects
  - Processing
    - Transformation of data
    - Implemented using processors
  - Storage
    - Retention of data
    - Implemented using memory
  - Communication
    - Transfer of data between processors and memories
    - Implemented using buses
    - Called *interfacing*

A simple bus



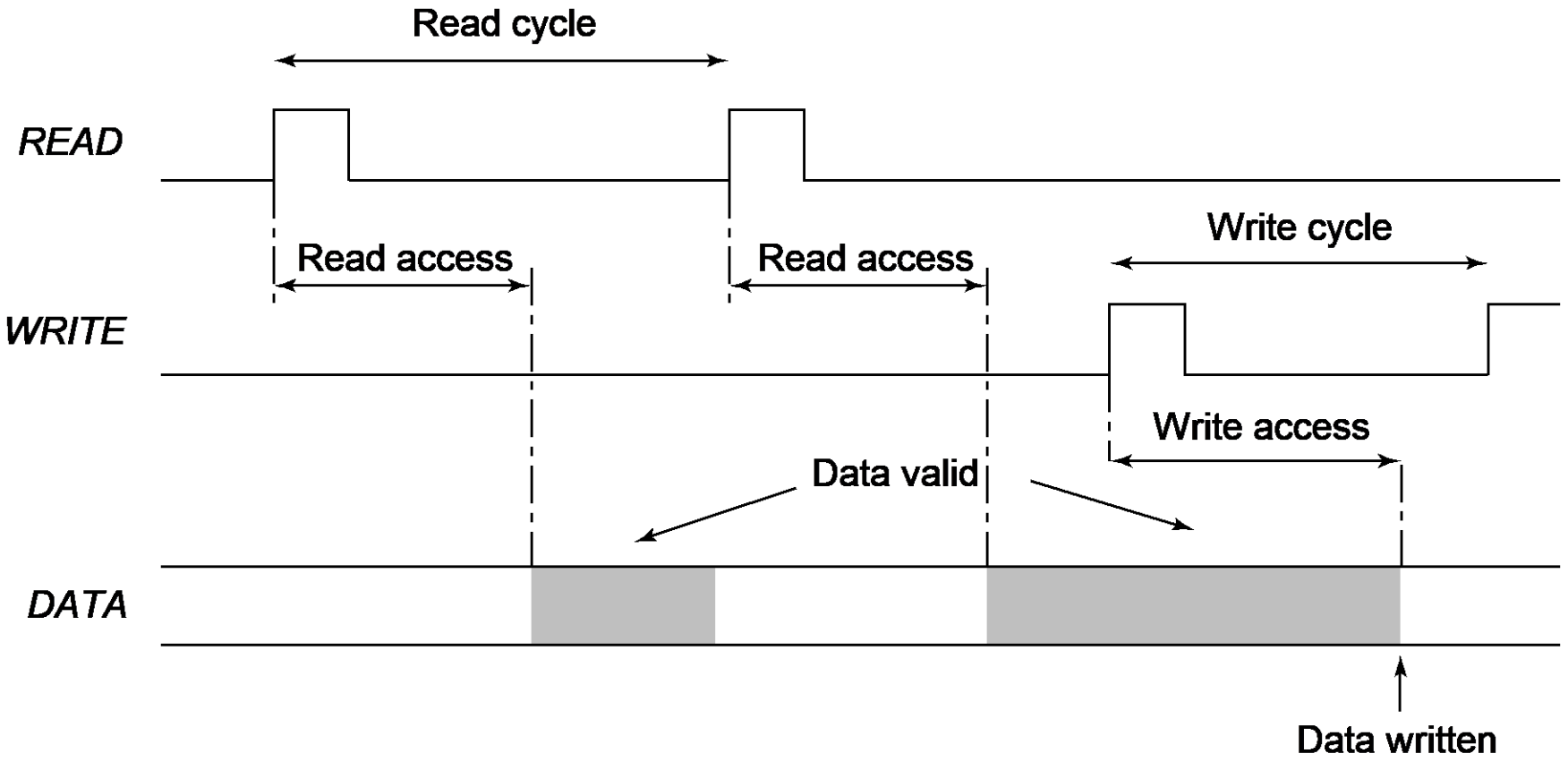
- Wires:
  - Uni-directional or bi-directional
  - One line may represent multiple wires
- Bus
  - Set of wires with a single function
    - Address bus, data bus
  - Or, entire collection of wires
    - Address, data and control
    - Associated protocol: rules for communication

Semiconductor Memory Classification

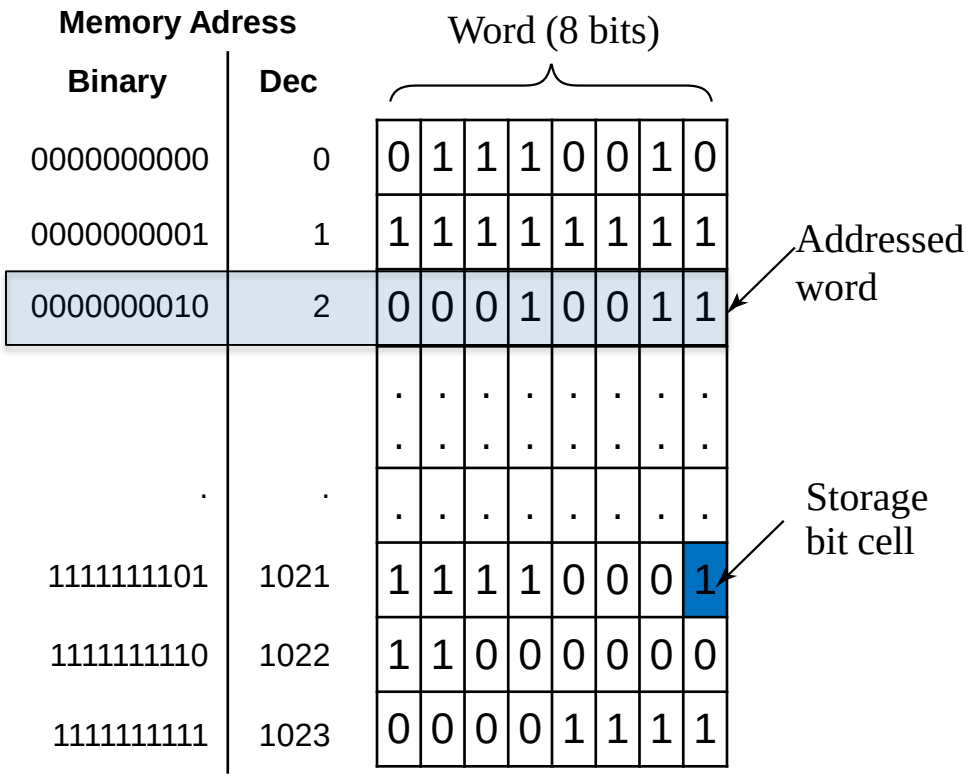
Read-Write Memory		Non-Volatile Read-Write Memory	Read-Only Memory
Random Access	**Non-Random Access**	EPROM E^2PROM FLASH	Mask-Programmed Programmable (PROM)
SRAM DRAM	FIFO LIFO Shift Register CACHE		

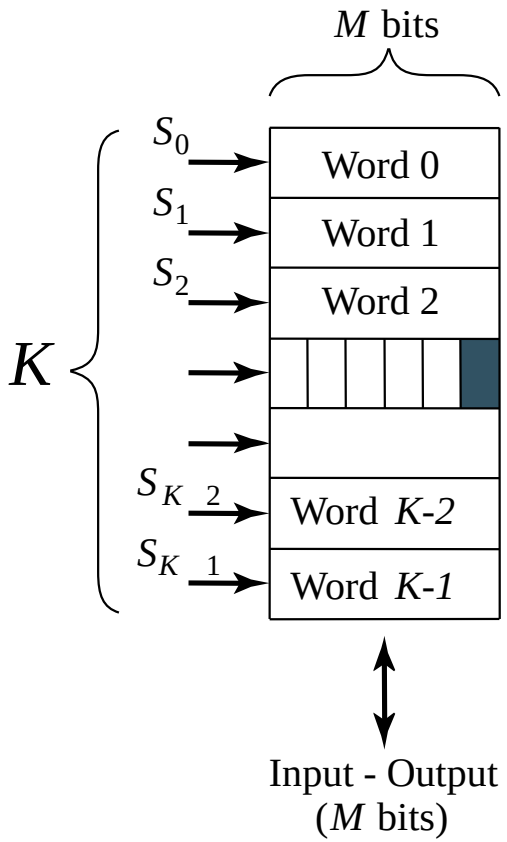
** Digital Integrated Circuits 2dn – Memories (SCU 2011, PhD Shoba Krishnan)*

Memory Timing: Definitions



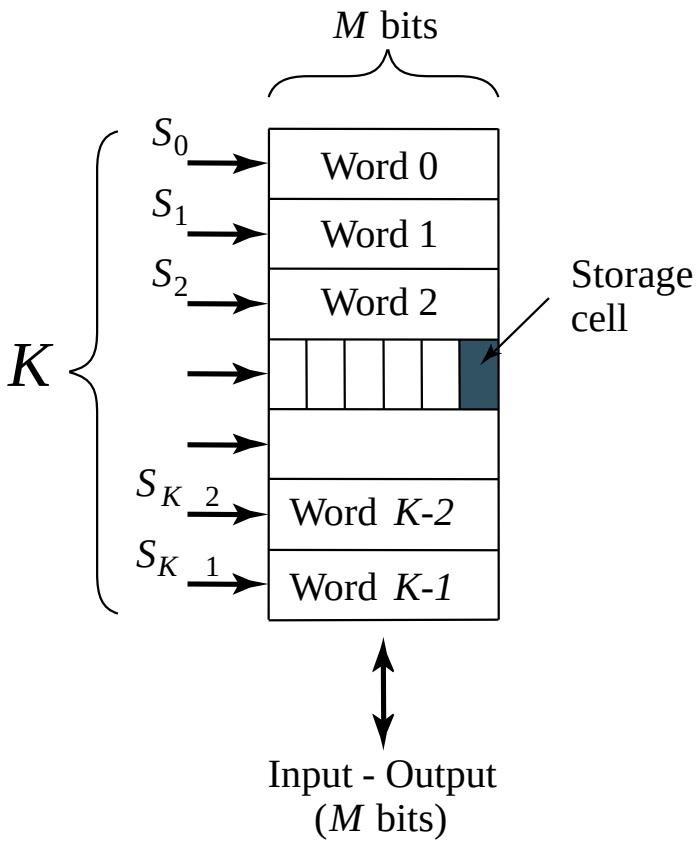
Memory Organization

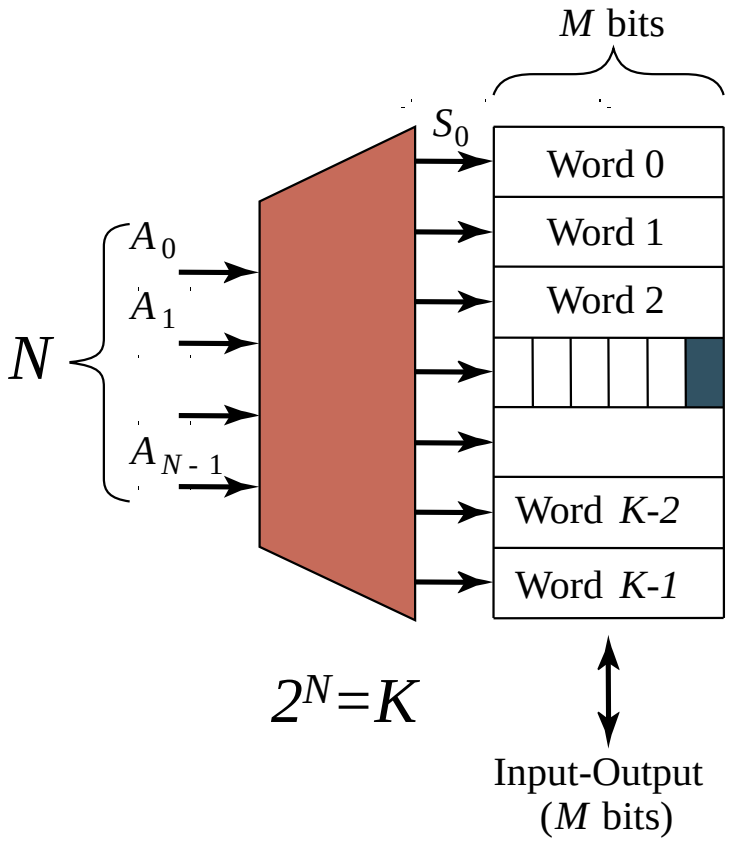




Example organization for
1Kword x 8 bits = 8K bits memory

Memory Architecture: Decoders





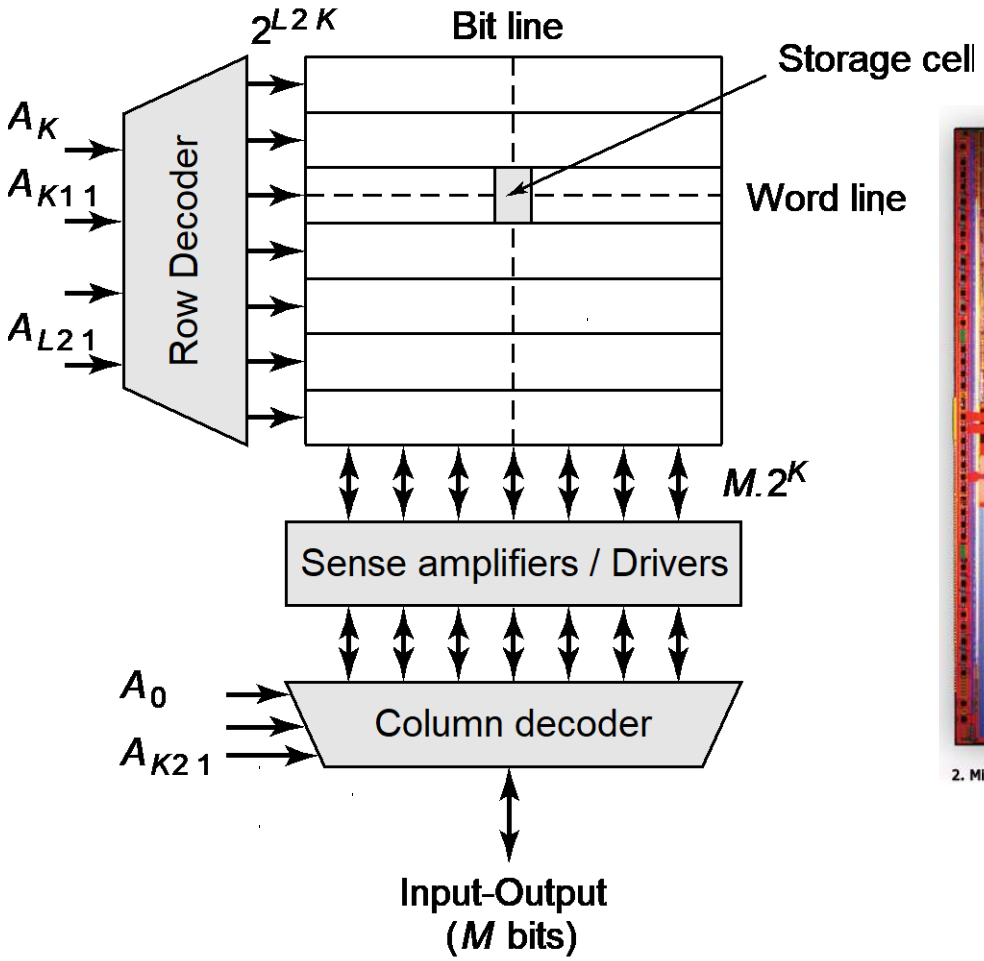
Intuitive architecture for K x M memory
Too many select signals:
K words == K select signals

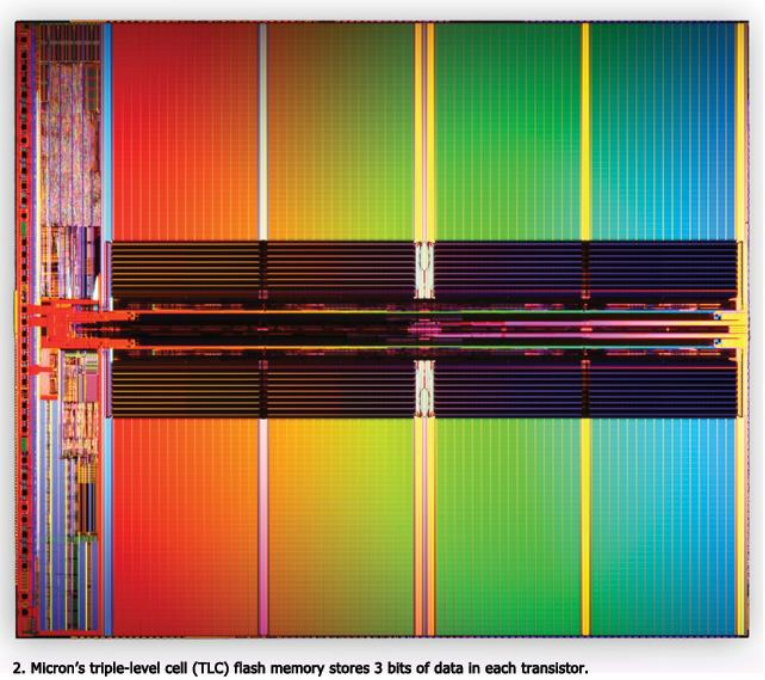
Decoder reduces the number of select signals

$N = log_2 K$

Array-Structured Memory Architecture

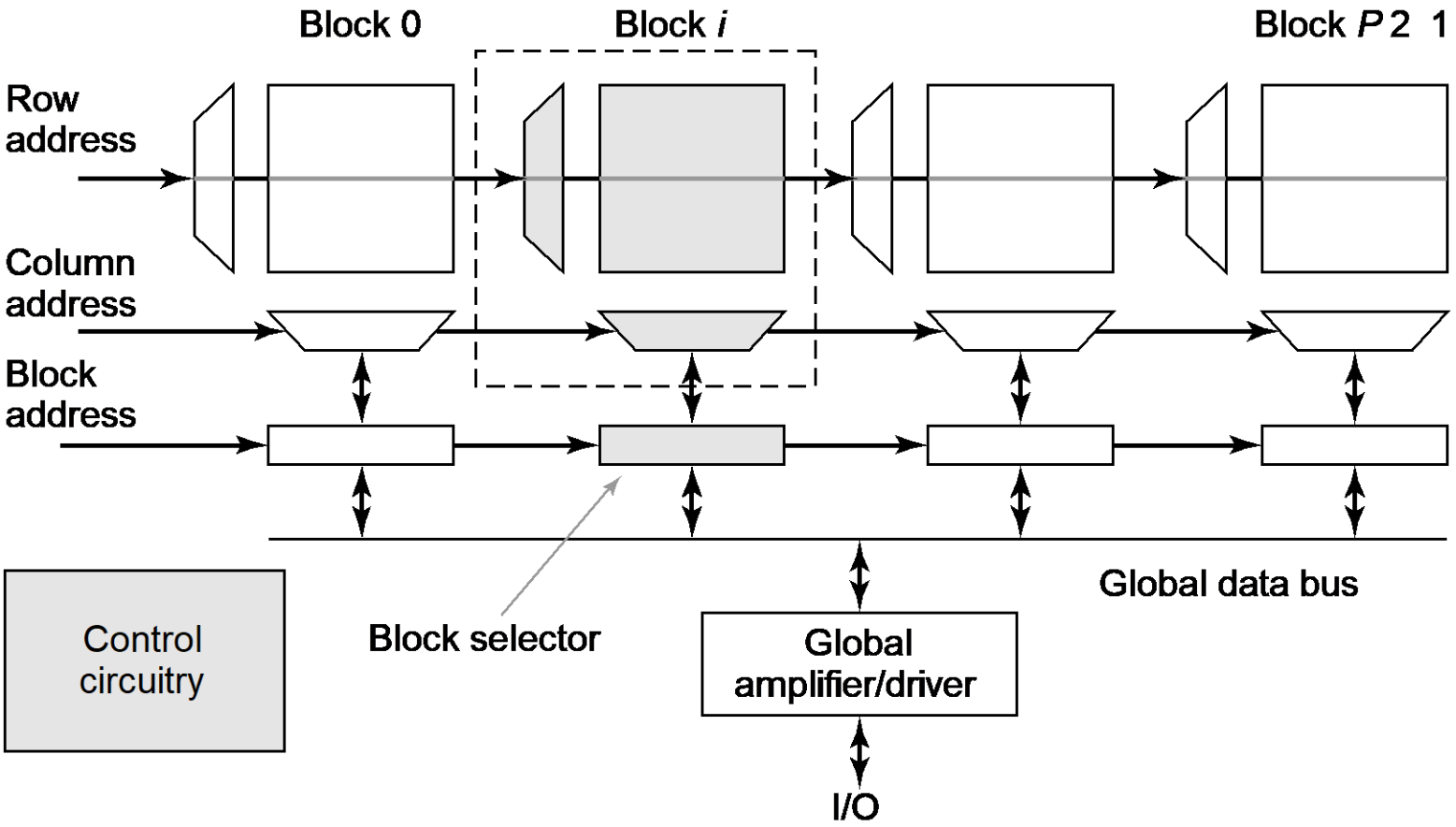
Problem: ASPECT RATIO or HEIGHT >> WIDTH





2. Micron's triple-level cell (TLC) flash memory stores 3 bits of data in each transistor.

Hierarchical Memory Architecture



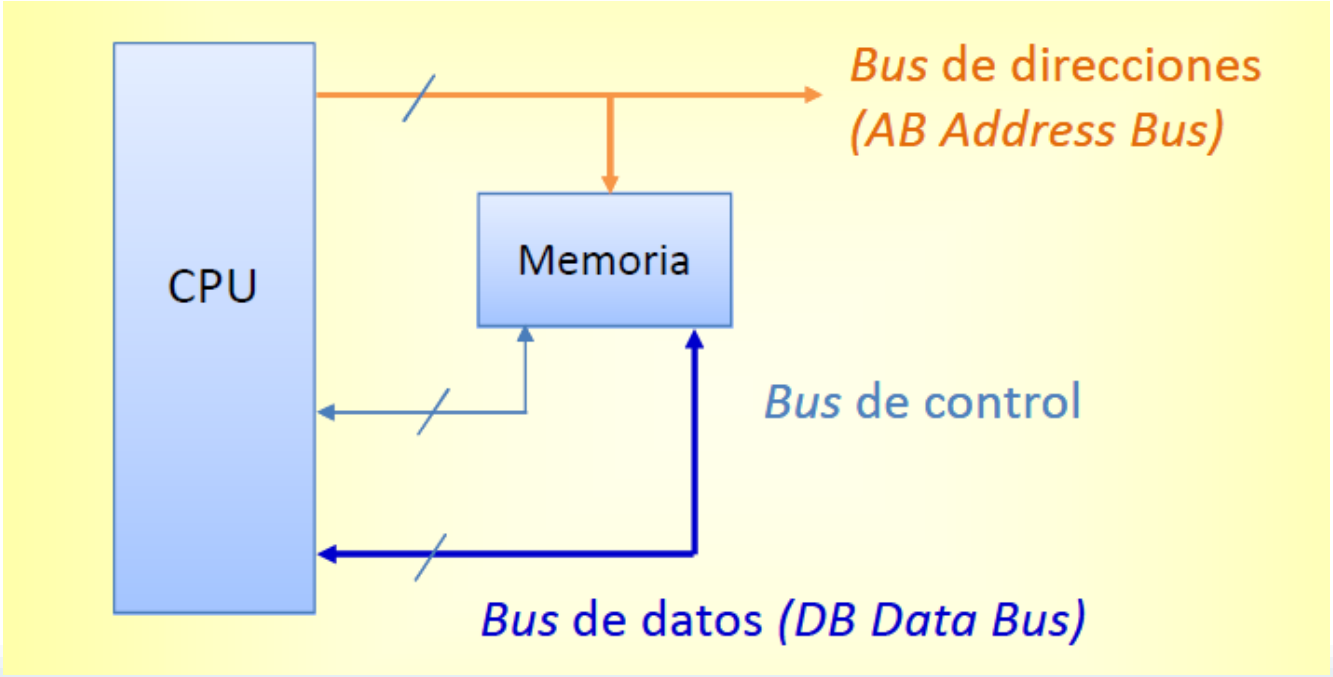
Advantages:

1. Shorter wires within blocks

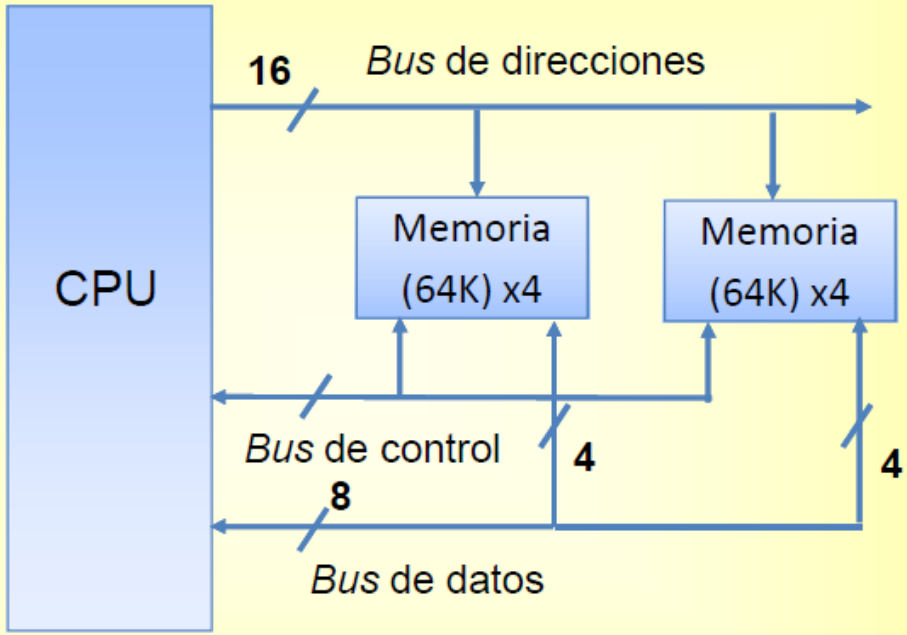
2. Block address activates only 1 block => power savings

Addressable Space

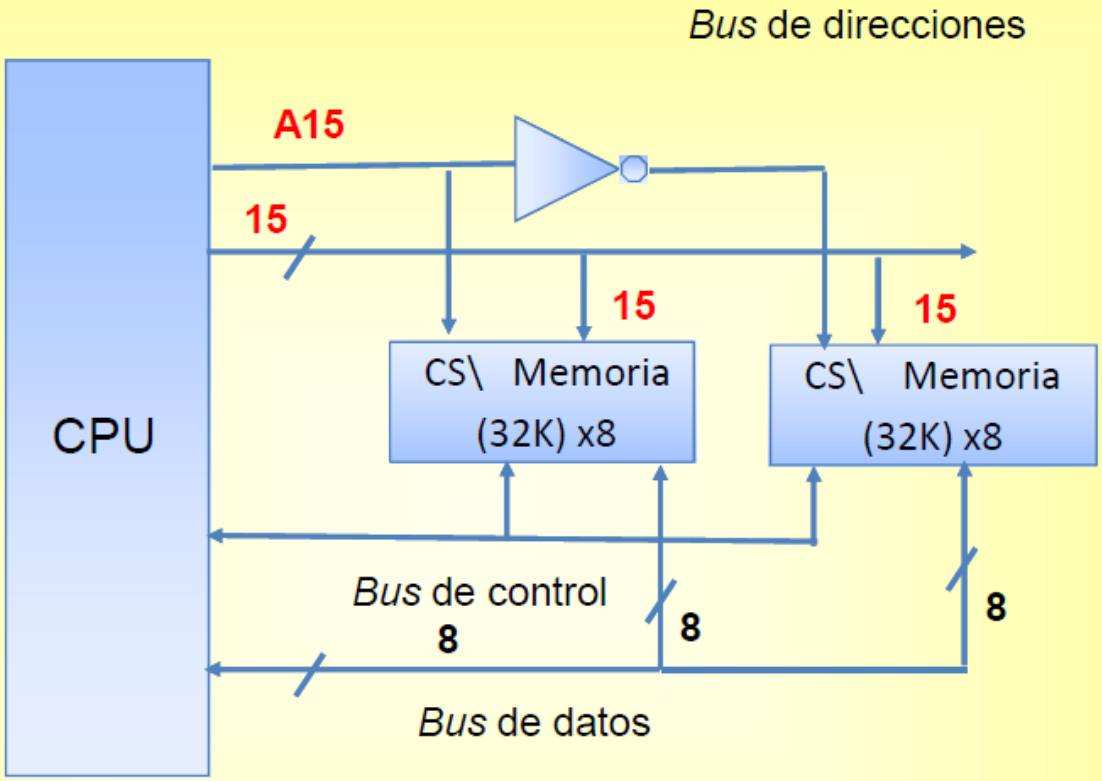
- It´s defined as the **total number** of addresses the CPU can access.
- It depends on the width (number of bits) of the address bits: n bits -> 2^n addresses



Memory Addressing (data bus)

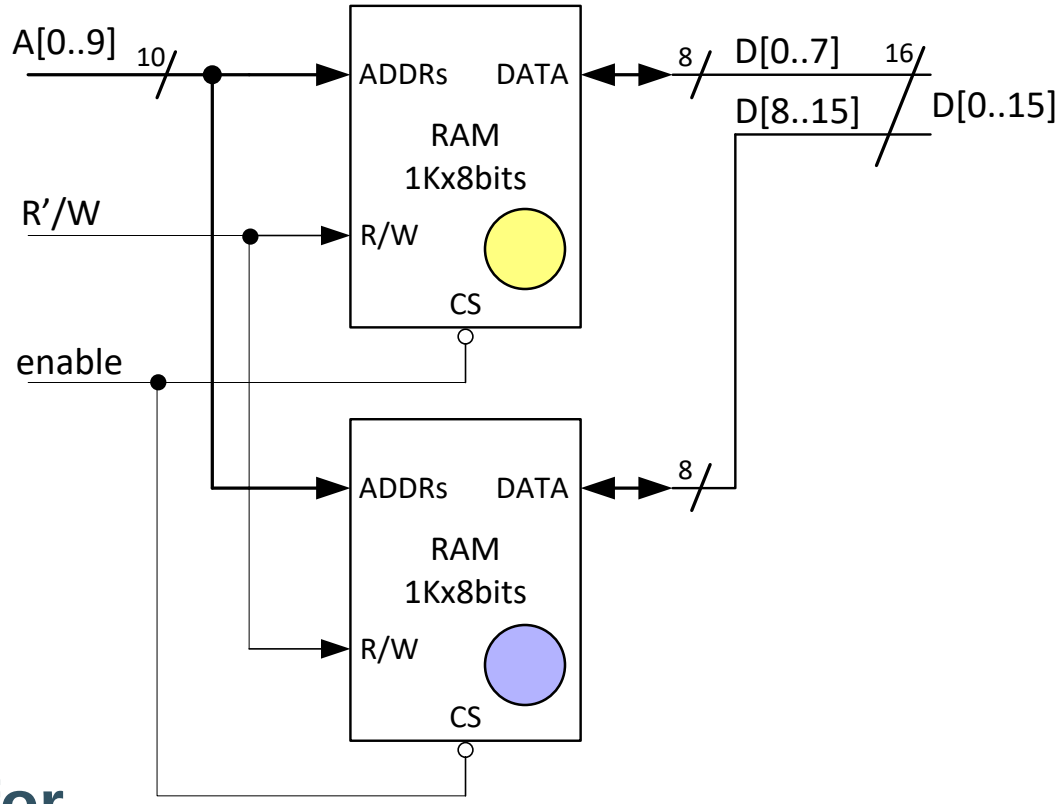


Memory Addressing (address bus)



Memory Addressing (parallel)

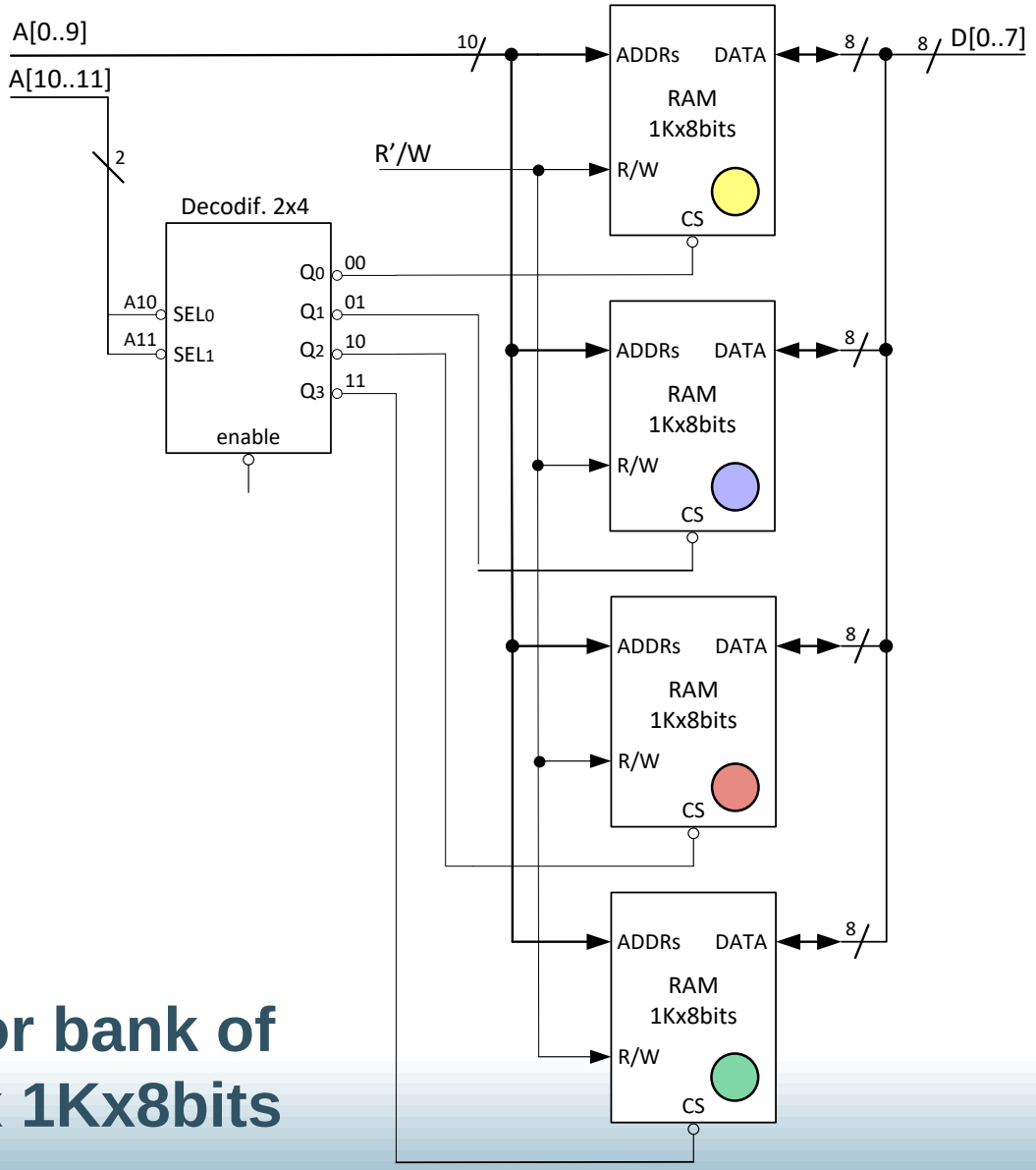
$A_9 \ldots A_0$	$D_{15} \ldots D_8$	$D_7 \ldots D_0$
0000000000	01.1	0.10
0000000001	11.1	1.11
0000000010	00.1	0.11
	.. ..	.. ..
	....	....
1111111101	11.1	0.01
1111111110	11.0	0.00
1111111111	00.0	1.11



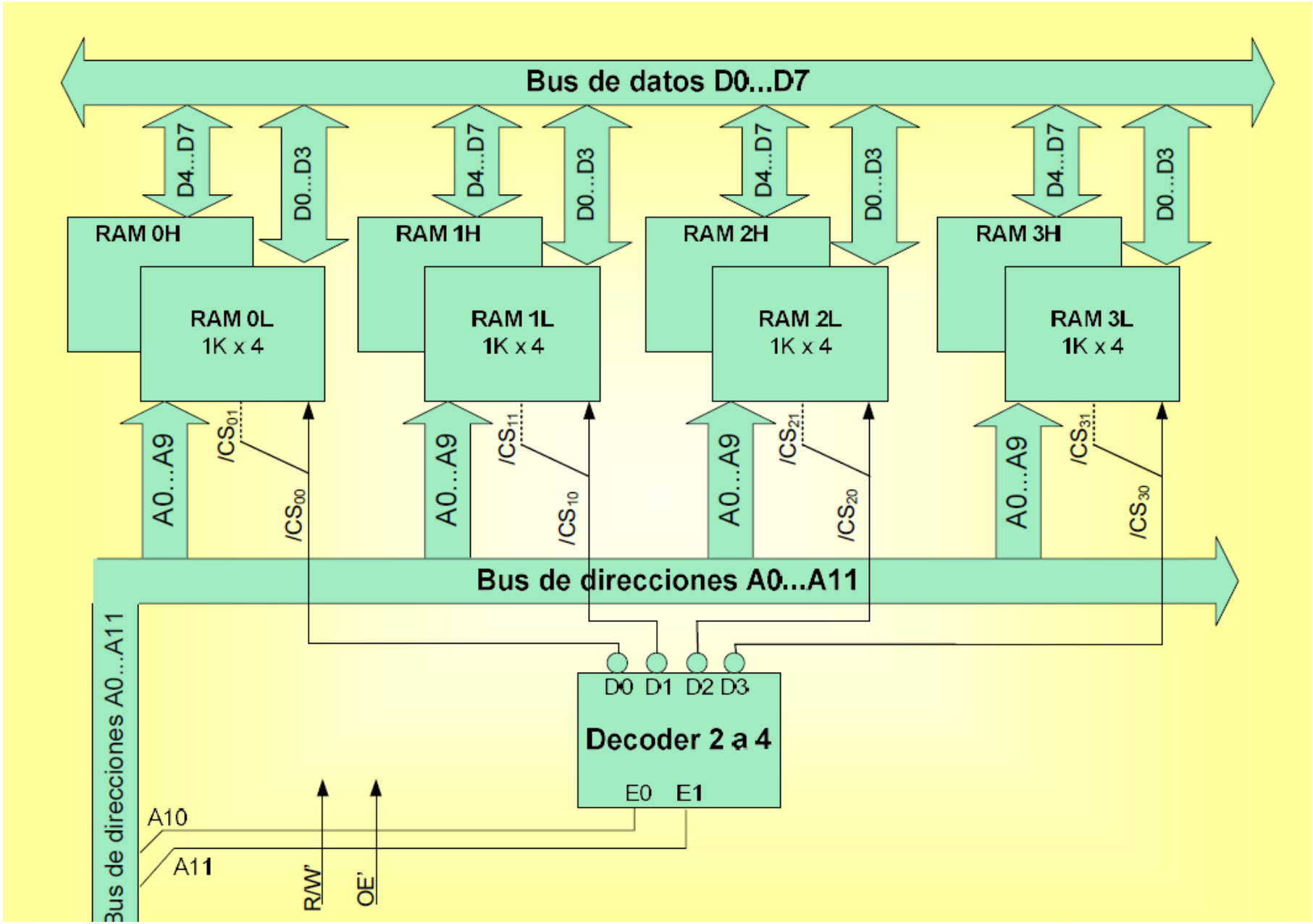
Example organization for bank of 1Kword x 16 bits from 2 x 1Kx8bits

Memory Addressing (serial)

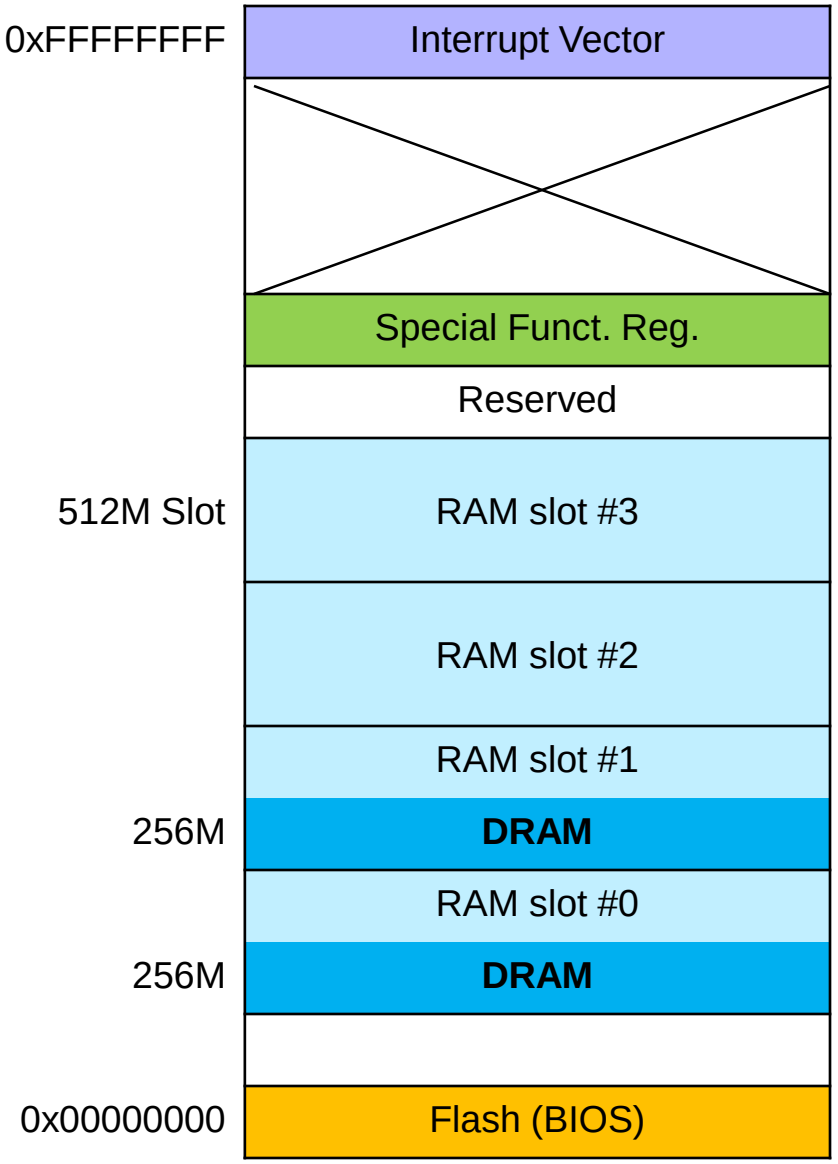
A_{11} A_{10}	A_9 . . . A_0	D_7 . . . D_0
0 0	0000000000	0 1 1 1 0 0 1 0
		. . . .
0 0	1111111111	0 0 0 1 0 1 1 1
0 1	0000000000	0 1 1 1 0 0 1 0
		. . . .
0 1	1111111111	0 0 0 1 0 1 1 1
1 0	0000000000	0 1 1 1 0 0 1 0
		. . . .
1 0	1111111111	0 0 0 1 0 1 1 1
1 1	0000000000	0 1 1 1 0 0 1 0
		. . . .
1 1	1111111111	0 0 0 1 0 1 1 1



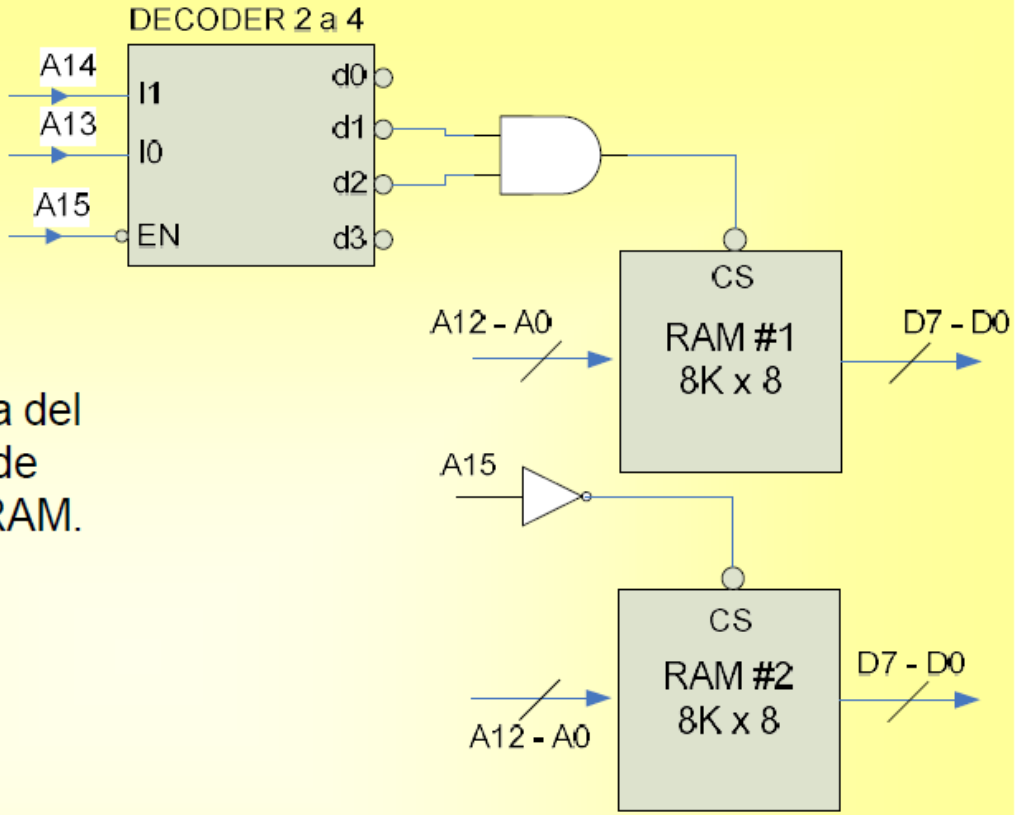
Example organization for bank of 4Kword x 8 bits from 4 x 1Kx8bits


Bus de datos D0...D7
D4...D7
D0...D3
RAM 0H
RAM 0L
1K x 4
RAM 1H
RAM 1L
1K x 4
RAM 2H
RAM 2L
1K x 4
RAM 3H
RAM 3L
1K x 4
A0...A9
/CS01
/CS00
/CS11
/CS10
/CS21
/CS20
/CS31
/CS30
Bus de direcciones A0...A11
Bus de direcciones A0...A11
D0 D1 D2 D3
Decoder 2 a 4
E0 E1
A10
A11
R/W'
OE'

Example: Conventional PC Memory map



‘Espejos’ en el mapa de memoria



Caso: Hallar el mapa de memoria del circuito indicando las posiciones de memoria ocupadas por cada CI RAM.

A15	A14	A13	Decoder	Mem	Posición en el mapa	Bytes
0	0	0	d0	-	H0000 – H1FFF	8K
	0	1	d1	RAM1	H2000 – H3FFF	8K
	1	0	d2	RAM1	H4000 – H5FFF	8K
	1	1	d3	-	H6000 – H7FFF	8K
1	X	X		RAM2	H8000 - HFFFF	32K

Por ejemplo, a la primer posición de memoria de RAM #2 se accede indistintamente con las direcciones:

	A_{15} ...A_{12}	A_{11} ..... A_0
8000	1 0 0 **0**	**0000 0000 0000**
A000	1 0 1 **0**	**0000 0000 0000**
C000	1 1 0 **0**	**0000 0000 0000**
E000	1 1 1 **0**	**0000 0000 0000**