malloc

malloc()

- En unix malloc() devuelve un bloque de size bytes alineado a 8-bytes (double word).
- No inicializa la memoria devuelta.
- Utiliza la system call sbrk o mmap.

free()

```
#include <stdlib.h>
void free(void *ptr);

Returns: nothing
```

Libera bloques reservados en el heap.

El ptr debe haber sido reservado previamente con malloc(), calloc() o realloc(). Si esto no sucede el comportamiento de free es INDEFINIDO

Ejemplo

```
#include <stdio.h>
#include <stdlib.h>
int main(void)
    int *p1 = malloc(4*sizeof(int)); // allocates enough for an array of 4 int
    int *p2 = malloc(sizeof(int[4])); // same, naming the type directly
    int *p3 = malloc(4*sizeof *p3); // same, without repeating the type name
    if(p1) {
        for(int n=0; n<4; ++n) // populate the array</pre>
            p1[n] = n*n;
        for(int n=0; n<4; ++n) // print it back out</pre>
            printf("p1[%d] == %d\n", n, p1[n]);
    free (p1);
    free (p2);
    free (p3);
```

$$ap = malloc(n)$$

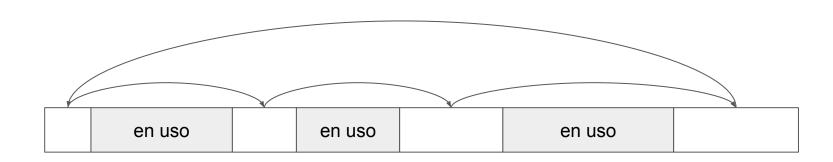
```
Header memoria alocada

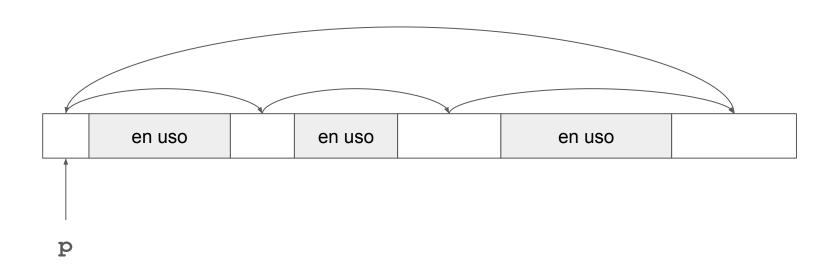
ap

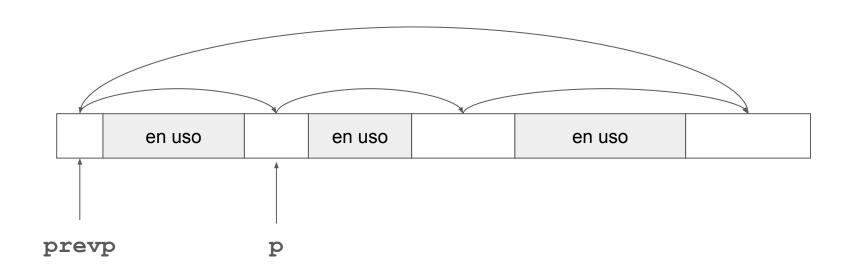
bp = (Header *) addr - 1
```

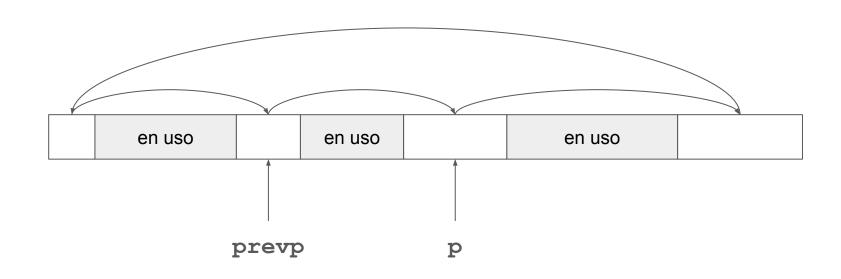
Estructura Header

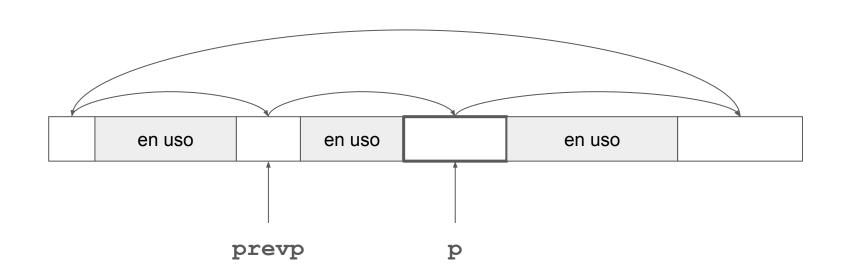
```
struct header {
  struct header *ptr;
 unsigned int size;
typedef struct header Header;
```



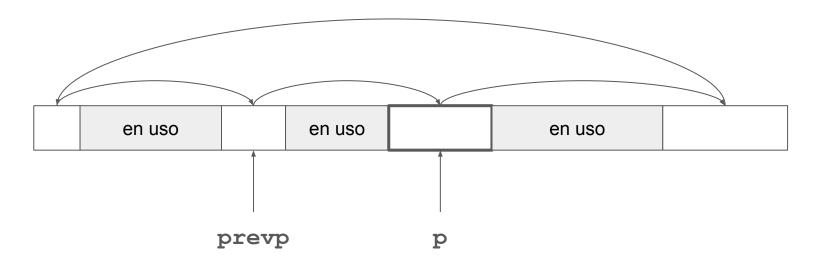






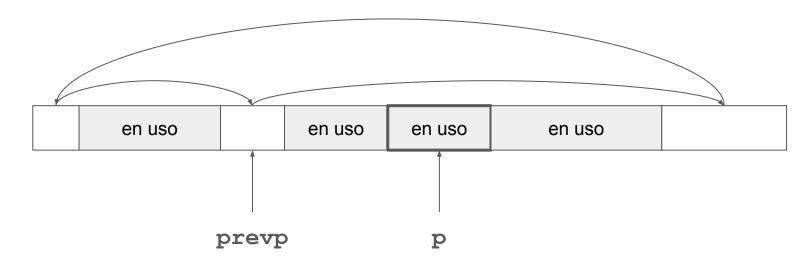


Opcion 1: Reservar el bloque completo



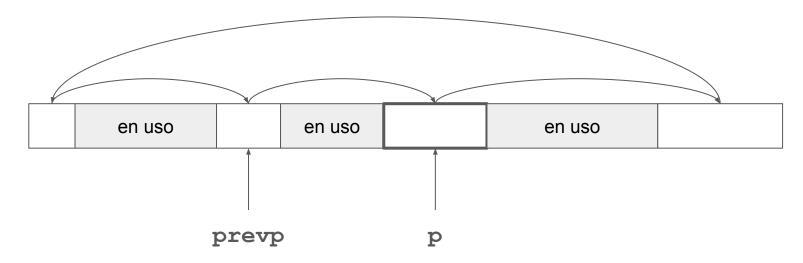
```
prevp->ptr = p->ptr;
return (p+1); // p es de tipo Header
```

Opcion 1: Reservar el bloque completo

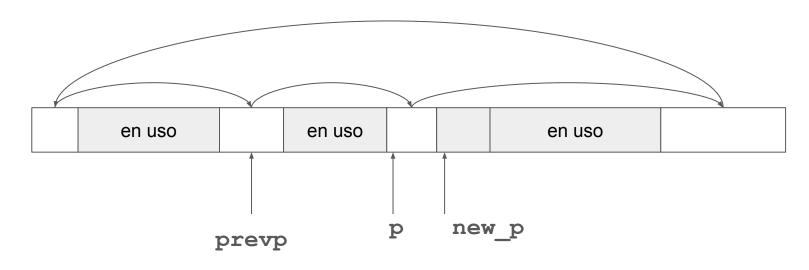


```
prevp->ptr = p->ptr;
return (p+1); // p es de tipo Header
```

Opcion 2: Split



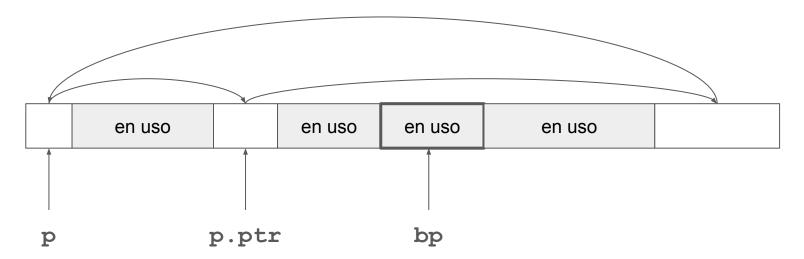
Opcion 2: Reservar parte del bloque



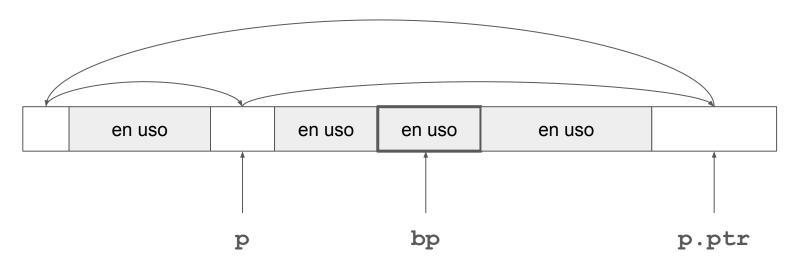
```
p->size =- nunits;
new_p = p + p->size;
new_p->size = nunits;
```

return (new_p + 1); // new_p es de tipo Header

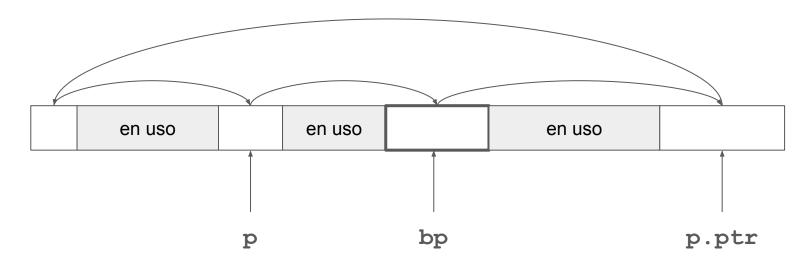
Liberar el bloque



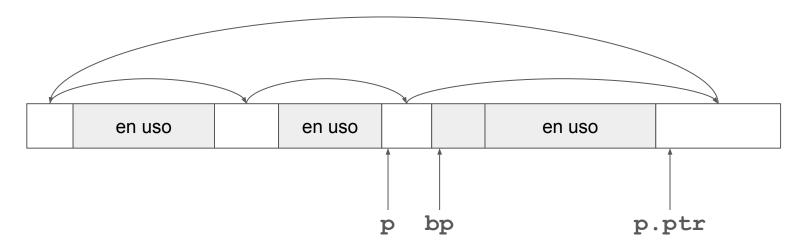
Liberar el bloque



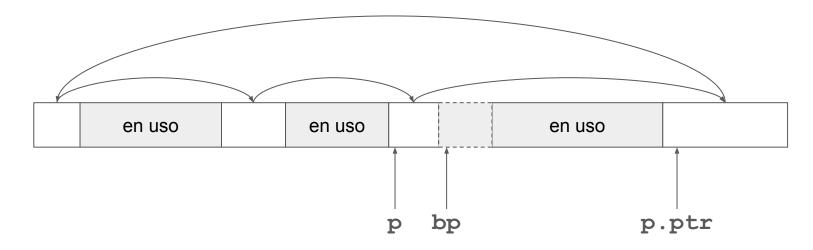
Liberar el bloque



Liberar bloque: Coalesce

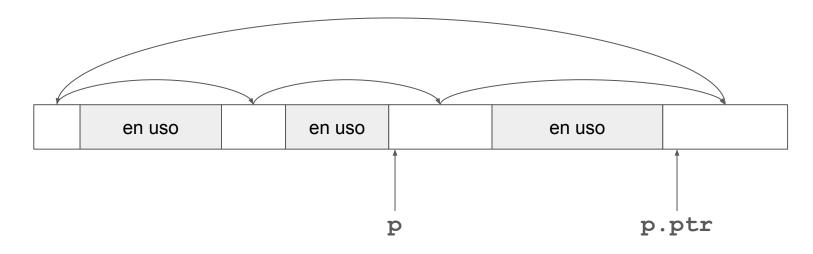


Liberar bloque: Coalesce



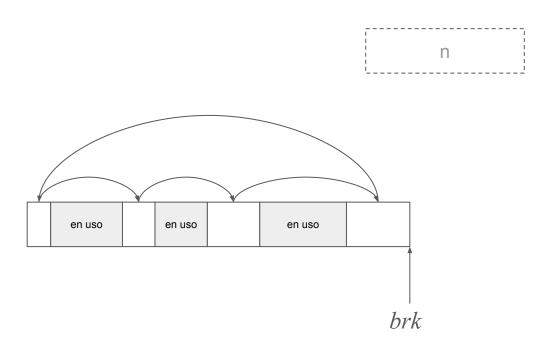
```
if (p + p->size == bp) {
    p->size += bp->size
}
```

Liberar bloque: Coalesce

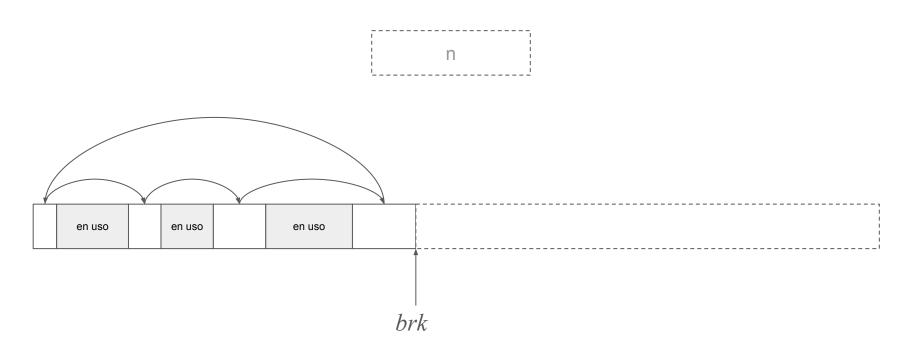


```
if (p + p->size == bp) {
    p->size += bp->size
}
```

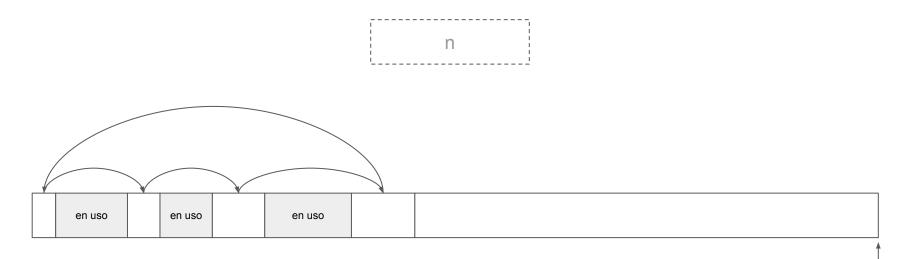
$$ap = malloc(n)$$



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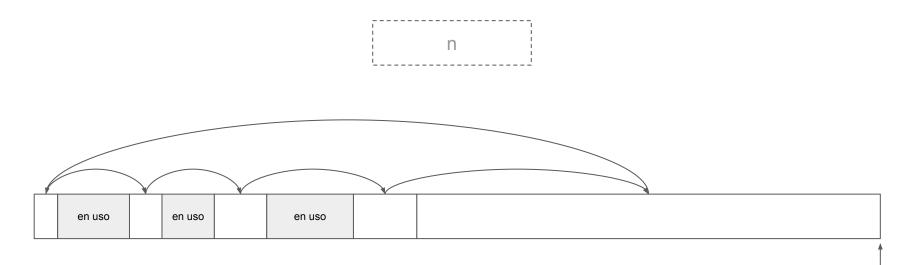


$$ap = malloc(n)$$



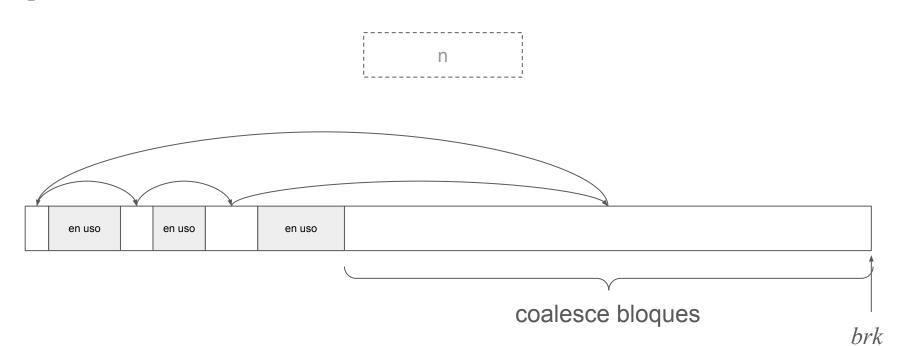
brk

$$ap = malloc(n)$$

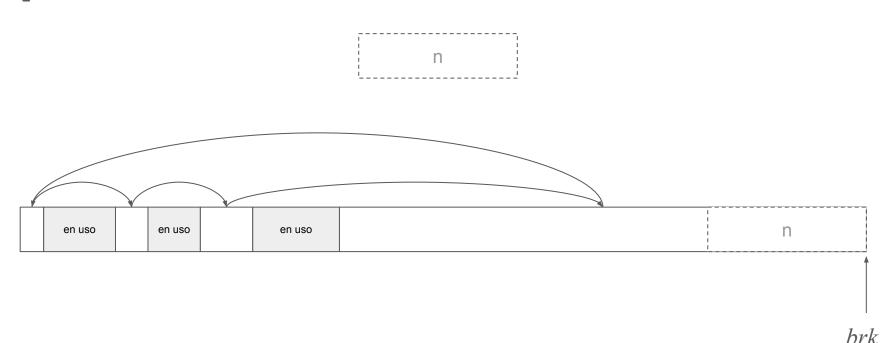


brk

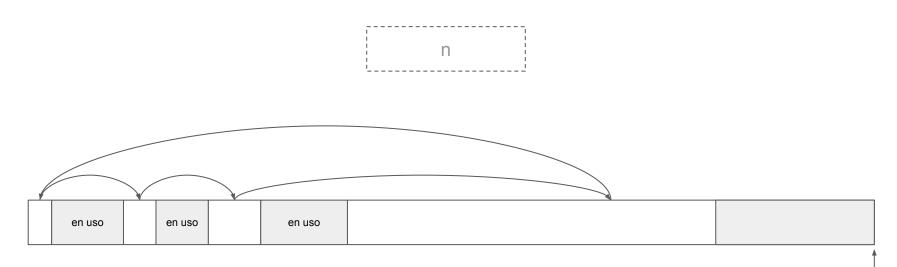
$$ap = malloc(n)$$



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brk

mmap

mmap

https://dashdash.io/2/mmap

Reservar un bloque con mmap

```
// mmap() returns a pointer to a chunk of free space
node_t *head = mmap(NULL, 4096, PROT_READ|PROT_WRITE,
MAP_ANONYMOUS|MAP_PRIVATE, -1, 0);
head->size = 4096 - sizeof(node_t);
head->next = NULL;
```

MAP_ANONYMOUS

The mapping is not backed by any file; its contents are initialized to zero. The *fd* argument is ignored [...] The *offset* argument should be zero.

Administración de memoria en xv6

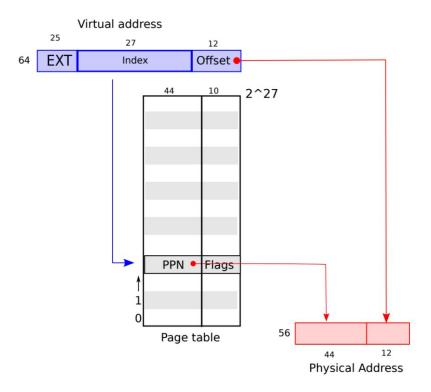


Figure 3.1: RISC-V virtual and physical addresses, with a simplified logical page table.

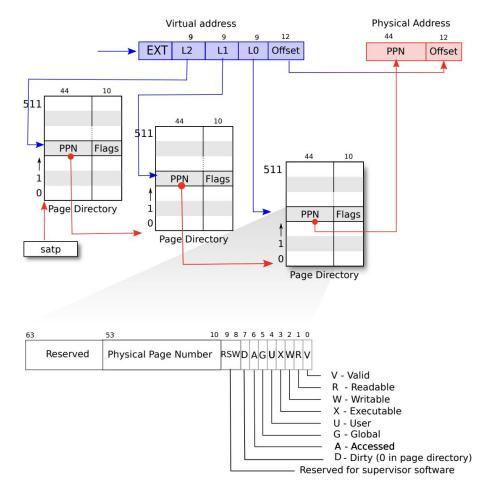


Figure 3.2: RISC-V address translation details.

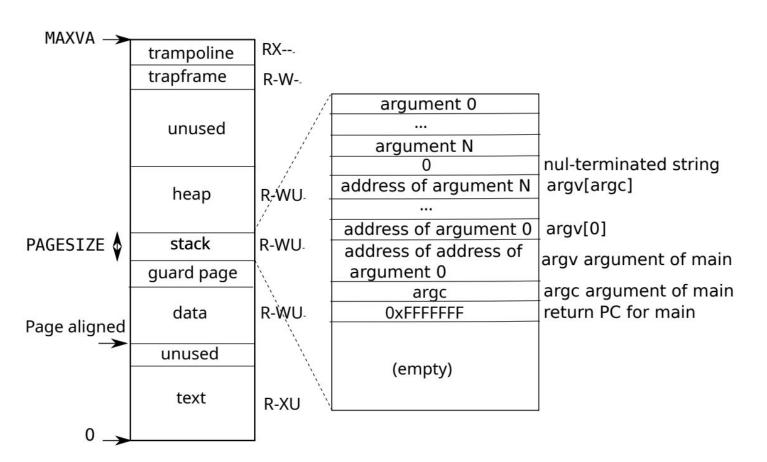


Figure 3.4: A process's user address space, with its initial stack.

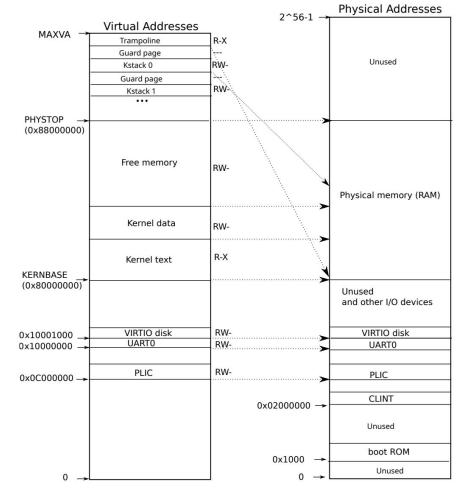


Figure 3.3: On the left, xv6's kernel address space. RWX refer to PTE read, write, and execute permissions. On the right, the RISC-V physical address space that xv6 expects to see.