

# CS372

## Lab2 Report

### Quick sort runtime analysis

The quick sort algorithm has a worst-case runtime of  $O(n^2)$ . When an already sorted array is put through the quick sort algorithm, using the lowest index as the pivot, then the runtime will reach  $O(n^2)$ . Quick sort has an average case and best case of  $O(n \log n)$ . For the running time complexities, it depends on how even the partitions are. In a worst case the partitions are most unbalanced as the array is recursively called for  $n-1$  each time. In a best case, the size of the partitions is as evenly distributed as they can be, each subproblem is cut in half each time. As for the average case, there is a 50/50 chance that the partition will be the worst case or 3-to-1 split or better.

### Merge sort runtime analysis

The worst-case runtime for merge sort is  $O(n \log n)$ . Similarly, the runtimes for the best and average cases are also  $O(n \log n)$ . Unlike quick sort, merge sort will always divide the subproblems evenly in halves. Similar to a best-case quick sort runtime, each input can only be divided  $\log n$  amount of times and because it runs through the entire array it is multiplied by  $n$  giving it a complexity of  $O(n \log n)$ .

### Sorting 2gb of data

I would use quick sort to sort the 2gb of data. In the average case quick sort would give a runtime of  $O(n \log n)$ . Quick sort is an in-place sorting algorithm, meaning that it will save memory when executing a large set of data. Merge sort would not be as practical since it creates more arrays to work on its sub problems giving it higher space complexity. Lastly, I would not use bubble sort because its average time complexity is  $O(n^2)$ .

### Observation

Bubble sort shows exponential growth and quick sort and merge sort show a much smaller running time even as time progresses just as it was analyzed previously

Running times of 3 different sorting algorithms

