# COSC364 Internet Technologies and Engineering Assignment2 Report

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### **Formulations**

- F1 Minimize [x, c, d, r] r
- Subject -co
  F2 demand volume: \( \sum\_{\text{tal}} \times \text{tkj=hij, for \$i \in \{1, \dots, \dots\}, \j \in \{1, \dots, \dots\}} \)
- F3 (Source + transit)  $\sum_{j=1}^{Z} x_i k_j \leq C_i k$ , for  $i \in \{1, \dots, X\}$ ,  $k \in \{1, \dots, Y\}$
- F4 capacity 2:  $\sum_{j=1}^{\infty} x_i k_j \leq d_{k_j}$ , for  $k_{G_1}, \dots, Y_{j}^{\infty}$ ,  $j \in \{1, \dots, Z_{j}^{\infty}\}$
- F5 Transit load:  $\sum_{i=1}^{x} \sum_{j=1}^{z} x_{i} + \sum_{j=1}^{z} x_{$
- F6 Binary variable: \( \sum\_{K=1}^{Y} \text{Uikej} = \text{Nij}, \for i \in \left(1, \ldots, \text{X}\right), \( \frac{1}{6} \xi \ldots, \ldots, \ldots, \text{X}\right), \( \frac{1}{6} \xi \ldots, \ldots, \ldots, \ldots, \text{X}\right), \( \frac{1}{6} \xi \ldots, \ld
- F7 path flow:  $xikj = \frac{\text{Vikj} \cdot hij}{\text{Nij}}$ , for  $i \in \{1, ..., X\}$ ,  $k \in \{1, ..., X\}$ ,  $j \in \{1, ..., Z\}$ ,
- F8 Bounds , xxxx >0, for ie {1, ..., x}, ke {1, ..., Y}, ie {1, ..., Z}
- F9 Binaries: Use e[0,1], for ie[1,..., x], ke[1,..., Y], ie[1,..., Z]

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#### Formulation and explanation

variables we used:

i: index of source node

k: index of transit node

j: index of destination node

xikj: flow on path from source i using transit k to destination j

hij: demand volume from source i to destination j

cik: capacity on link from source i to transit k

dkj: capacity on link from transit k to destination j

uikj: binary variable on path from source i using transit k to destination j

nij: number of split paths

#### F1: Minimize objective function

We introduce an auxiliary r for the utilization of the load on each transit node, so we can make the load on transit nodes to balance by adjusting the value of r. Then introduce the new constraint for r which r should not be smaller than the total value of load on each transit node xikj. In this way, when r gets the minimized value, it means demand volume which allocates on each transit node is minimized, the utilization of transit nodes is optimized.

#### F2. Demand volume constraint

Each path flow xikj should add up to hij which is the given demand volume between node i and j.

#### F3 – F4: Capacity constraints

Each path flow xikj which involves the link ik or kj should not be greater than the capacity of this link. In this problem for the optimized situation, the capacity should be equal to the demand volume to make the cost of capacity of each link lowest.

#### F5: Load on each transit node constraint

The sum of flow on each transit node xikj should not be greater than the introduced variable r which means r is a upper bound for load on transit node. The eligible flow on each transit node is smaller, meaning the utilization of transit nodes is more balanced.

As the load on transit node influences path flow, we think it is more reasonable to put the constraints F3 – F5 before F6 and F7 which is a little different from the fomulations on booklet.

#### F6: Binary variable constraint

The binaray variable uikj is equal to 1 meaning the demand volume split into the path i-k-j is used and 0 means not into this path. As in this problem, each demand volume split over exactly three different paths, the amount of paths which is used is 3, the sum of binary variable uikj on each path between i and j is 3.

#### F7: Path flow equation

As this is an equal split load balancing problem, for each used path i - k - j, demand volume hij between node i and j should be equally splited into nij = 3 path flows xikj.

#### F8: Bounds

All the path flow xiki should be non-negative.

#### F9: Binaries

All the binary variables which are equal to one of used paths.

#### **Execution result**

Given X = Y = 7:

number of transit nodes(Y)

	3	4	5	6	7
Execution time(in seconds)	0.01338	0.08076	0.07471	0.07672	0.10540
NO. of links with non-zero capacities	42	56	70	83	95
Highest capacity	26	24	23	19	17
Links with highest capacity	c71, c72, c73, d17, d27, d37	d26	c75	c74, d67	c67, c75
Load on the transit nodes(excluding unused nodes)	1:130.667 2:130.667 3:130.667	1:98.000 2:97.999 3:98.000 4:98.0	1:78.333 2:78.667 3:78.667 4:77.667 5:78.667	1:65.333 2:65.333 3:65.333 4:65.333 5:65.333 6:65.333	1:55.999 2:56.000 3:55.999 4:56.000 5:56.000 6:56.0 7:56.000

#### Explaination of execution results:

The execution result is same as we expected. As the number of transit nodes increasing, the execution time becomes longer. This is because the number of variables has increased, which makes the problem more complex to solve and hence takes more time. The number of links with non-zero capacities has also increased because demand is distributed to more transit nodes in order to reduce the capacity on each link. When the demand volume is fixed, less capacity on a link means less cost. We can also clearly find that the load on the used transit nodes are almost equal thanks to load balancing. The load on each transit node is decreasing with the addition of transit nodes since load is distributed to more transit nodes. Thus each node is responsible for less load on average.

# Appendix 1

```
Source code:
createLPfile.py (used to generate lp file )
#generate the nodes of source, transit, destination
X = int(input('The amount of source nodes:'))
Y = int(input('The amount of transit nodes:'))
Z = int(input('The amount of destination nodes:'))
N = 3
print('Source nodes: {}\ntransit nodes: {}\nDestination node: {}'.format(X, Y, Z))
print('----')
def DV contraint():
  "return the demand volume constraint:
  Xikj = hij
  which Xiki means the sum of load between source i to destination j'''
  DV = []
  demand equation = []
  for i in range(1, X + 1):
     for j in range(1, Z + 1):
       dv = []
       for k in range(1, Y + 1):
          dv.append("x{}{}){}".format(i, k, j))
       DV = ' + '.join(dv) + ' = {}'.format(i + j)
       demand equation.append(DV)
  demand constraint = '\n'.join(demand equation)
  return demand constraint
def ST capp constraint():
  "return the cappacity constraint from source to transit node"
  ST = []
  capp1 equation = []
  for i in range (1, X + 1):
     for k in range (1, Y + 1):
       st = []
       for j in range(1, Z+1):
          st.append('x{)}{}'.format(i, k, j))
       ST = ' + '.join(st) + ' - c\{\}\{\} \le 0'.format(i, k)
       capp1 equation.append(ST)
  capp1 constraint = '\n'.join(capp1 equation)
  return capp1 constraint
def TD capp constraint():
  "return the cappacity constraint from transit node to dest node"
```

```
TD = []
  capp2 equation = []
  for k in range(1, Y+1):
     for j in range(1, Z+1):
       td = []
       for i in range(1, X+1):
          td.append('x{}{}{}).format(i, k, j))
       TD = ' + '.join(td) + ' - d\{\}\{\} \le 0'.format(k, j)
       capp2 equation.append(TD)
  capp2 constraint = '\n'.join(capp2 equation)
  return capp2 constraint
def TN constraint():
  "return the constraint of load of transit nodes which should be minimize"
  TN = []
  tn equation = []
  for k in range(1, Y + 1):
     tn = []
     for j in range(1, Z + 1):
       for i in range(1, X + 1):
          tn.append('x{}{}{})  '.format(i, k, j))
     TN = ' + '.join(tn) + ' - r \le 0'.format(k, j)
     tn equation.append(TN)
  tn constraint = \n'.join(tn equation)
  return tn constraint
def BV constriant():
  "return the binary variables constraint:
  Uikj = Nk
  which Uiki means the sum of used paths and Nk = N in this problem'"
  BV = []
  binary equation = []
  for i in range(1, X + 1):
     for j in range(1, Z + 1):
       bv = []
       for k in range(1, Y + 1):
          bv.append("u{}{}\".format(i, k, j))
       BV = ' + '.join(bv) + ' = {}'.format(N)
       binary equation.append(BV)
  binary constraint = '\n'.join(binary equation)
  return binary constraint
def DF constraint():
  "return the demand flow constraint:
  Nk * Xikj = hij * Uikj'''
  DF = []
  flow equation = []
```

```
for i in range(1, X + 1):
     for j in range(1, Z + 1):
       for k in range(1, Y + 1):
          DF = \{ \{ \} \} \} \} - \{ \} \{ \} \} \} = 0'.format(N, i, k, j, i + j, i, k, j)
          flow equation.append(DF)
  demand flow constraint = '\n'.join(flow equation)
  return demand flow constraint
def Bounds variable():
  "return the bounds of demand variable: Xikj of this problem"
  bound x = []
  bound unequation x = []
  for i in range(1, X + 1):
     for j in range(1, Z + 1):
       for k in range(1, Y + 1):
          bound x = 0 \le x\{\}\{\}\}'.format(i, k, j)
          bound unequation x.append(bound x)
  Bounds x = \frac{\ln(1 - x)}{\ln(1 - x)}
  return Bounds x
def binary constraint():
  "return binary constraints"
  bc = "
  for i in range(1, X+1):
     for k in range(1, Y+1):
       for j in range(1, Z+1):
          bc += 'u\{\}\{\}\{\} \setminus n'.format(i,k,j)
  return bc
def createLP():
  "create a LP file for this problem"
  f = open(filename, 'w')
  content = \
  "Minimize
Subject to
demand volume: \n{}
srouce to tranfer node capp1: \n{}
transit to destination node capp2: \n{}
transit nodes: \n{}
binary variables: \n{}
demand flow: n
Bounds
{}
0 <= r
Binaries
{}
```

```
End".format(demand volume, ST capacity,
        TD capacity, transit nodes, binary variables, demand flow, bounds x, binaries)
  f.write(content)
  f.close
def set filename():
  "return the filename: XYZ.lp which Y belongs to {3, 4, 5, 6, 7}"
  filename = '{}{}.lp'.format(X, Y, Z)
  return filename
demand volume = DV contraint()
ST capacity = ST capp constraint()
TD capacity = TD capp constraint()
transit nodes = TN constraint()
binary variables = BV constriant()
demand flow = DF constraint()
bounds x = Bounds variable()
binaries = binary constraint()
filename = set filename()
createLP()
analyser.py (used to run CPLEX with given lp files and extract the useful information)
import subprocess
import time
import ison
def CPLEX(filename):
  "run CPLEX to solve the problem of given lp file"
  #args1 = ['/Users/mac/Desktop/364/a2/cplex', '-c', 'read /Users/mac/Desktop/364/a2/' +
filename, 'optimize',
       #'display solution variable -']
  args1 = ['/Users/zelta/Desktop/cplex', '-c', 'read /Users/zelta/Desktop/' + filename,
'optimize',
       'display solution variable -']
  time1 = time.time()
  process1 = subprocess.Popen(args1, stdout = subprocess.PIPE)
  output, error = process1.communicate()
  time2 = time.time()
  execution time = time2 - time1
 # print('Execution time: '+str("%.5f" % (execution time))+'seconds')
  result = output.decode("utf-8").split()
```

```
start = result.index("Incumbent")
  load = \{\}
  links = []
  link count = 0
  cappacities = {}
  cappacities = \{\text{'max':}[0, \lceil\rceil]\}
  max cappacity = 0.0
  for i in range(1,8):
     load[i] = float(0)
  for n in result[start:]:
     if n.startswith('x'):
       transit node = int(n[2])
       x index = result.index(n)
       load[transit node] += float(result[x index+1])
     if (n.startswith('c') or n.startswith('d')) and len(n) == 3:
       capp index = result.index(n)
       cappacity = float(result[capp index+1])
       if cappacity > 0:
          link count += 1
          links.append(n)
       if cappacity == max cappacity:
          cappacities['max'][1].append(n)
       if cappacity > max cappacity:
          max cappacity = cappacity
          cappacities['max'][0] = max cappacity
          cappacities['max'][1]=[n]
  result = 'Execution_time: '+str("%.5f" % (execution time))+'seconds\n'+'Load on transit
nodes: '+json.dumps(\overline{load})+'\n'+'Maximum cappacity: '+json.dumps(cappacities)+'\nNon-zero
capacity link count: '+str(link count)+'\nNon-zero capacity links: '+' '.join(links)
  return result
def main():
  for n in range(3, 8):
     filename = '7{}7.lp'.format(n)
     print(filename)
     result = CPLEX(filename)
     f = open(filename + '_analyser.txt','w')
     f.write(result)
     f.close()
main()
```

# Appendix 2

```
324.lp (lp file for situation where X = 3, Y = 2, Z = 4)
```

```
Minimize
Subject to
demand volume:
x111 + x121 = 2
x112 + x122 = 3
x113 + x123 = 4
x114 + x124 = 5
x211 + x221 = 3
x212 + x222 = 4
x213 + x223 = 5
x214 + x224 = 6
x311 + x321 = 4
x312 + x322 = 5
x313 + x323 = 6
x314 + x324 = 7
source to transit node capp1:
x111 + x112 + x113 + x114 - c11 \le 0
x121 + x122 + x123 + x124 - c12 \le 0
x211 + x212 + x213 + x214 - c21 \le 0
x221 + x222 + x223 + x224 - c22 \le 0
x311 + x312 + x313 + x314 - c31 \le 0
x321 + x322 + x323 + x324 - c32 \le 0
transit to destination node capp2:
x111 + x211 + x311 - d11 \le 0
x112 + x212 + x312 - d12 \le 0
x113 + x213 + x313 - d13 \le 0
x114 + x214 + x314 - d14 \le 0
x121 + x221 + x321 - d21 \le 0
x122 + x222 + x322 - d22 \le 0
x123 + x223 + x323 - d23 \le 0
x124 + x224 + x324 - d24 \le 0
transit nodes:
x111 + x211 + x311 + x112 + x212 + x312 + x113 + x213 + x313 + x114 + x214 + x314 - r
x121 + x221 + x321 + x122 + x222 + x322 + x123 + x223 + x323 + x124 + x224 + x324 - r
<=0
binary variables:
u111 + u121 = 3
u112 + u122 = 3
u113 + u123 = 3
u114 + u124 = 3
u211 + u221 = 3
u212 + u222 = 3
u213 + u223 = 3
u214 + u224 = 3
```

- u311 + u321 = 3
- u312 + u322 = 3
- u313 + u323 = 3
- u314 + u324 = 3

#### demand flow:

- $3 \times 111 2 \times 111 = 0$
- $3 \times 121 2 \times 121 = 0$
- $3 \times 112 3 \times 112 = 0$
- $3 \times 122 3 \times 122 = 0$
- $3 \times 113 4 \times 113 = 0$
- $3 \times 123 4 \times 123 = 0$
- $3 \times 114 5 \times 114 = 0$
- $3 \times 124 5 \times 124 = 0$
- $3 \times 211 3 \times 211 = 0$
- $3 \times 221 3 \times 221 = 0$
- $3 \times 212 4 \times 212 = 0$
- $3 \times 222 4 \times 222 = 0$
- $3 \times 213 5 \times 213 = 0$
- $3 \times 223 5 \times 223 = 0$
- $3 \times 214 6 \times 214 = 0$
- $3 \times 214 0 \times 214 0$
- $3 \times 224 6 \times 224 = 0$
- $3 \times 311 4 \times 311 = 0$
- $3 \times 321 4 \times 321 = 0$
- $3 \times 312 5 \times 312 = 0$
- $3 \times 322 5 \times 322 = 0$
- $3 \times 313 6 \times 313 = 0$
- $3 \times 323 6 \times 323 = 0$
- $3 \times 314 7 \times 314 = 0$
- $3 \times 314 7 \times 1314 = 0$  $3 \times 324 - 7 \times 1324 = 0$
- Bounds
- $0 \le x111$
- $0 \le x121$
- $0 \le x112$
- $0 \le x122$
- $0 \le x113$
- $0 \le x123$
- $0 \le x114$
- $0 \le x124$
- $0 \le x211$
- $0 \le x221$
- $0 \le x212$
- 0 < x2120 <= x222
- $0 \le x213$
- $0 \le x223$
- $0 \le x214$
- $0 \le x224$
- $0 \le x311$
- $0 \le x321$
- $0 \le x312$
- $0 \le x322$

- $0 \le x313$
- $0 \le x323$
- $0 \le x314$
- $0 \le x324$
- 0 <= r

#### Binaries

- u111
- u112
- u113
- u114
- u121
- u121
- u122
- u123
- u124
- u211
- u212
- u213
- u214
- u217
- u221
- u222
- u223
- u224
- u311
- u312
- u313
- u314
- u321
- u322
- u323
- u324

End