

$$1. A = \begin{bmatrix} 1 & 2 & 4 \\ 3 & 8 & 14 \\ 2 & 6 & 13 \end{bmatrix}$$

pivot entry of the lower triangle ( $L$ ) always has higher absolute values than its subcomponents.

$$\xrightarrow{\text{P}_1} \begin{bmatrix} 1 & 2 & 4 \\ 3 & 8 & 14 \\ 2 & 6 & 13 \end{bmatrix} \xrightarrow{\text{P}_1} \begin{bmatrix} 3 & 8 & 14 \\ 1 & 2 & 4 \\ 2 & 6 & 13 \end{bmatrix} \xrightarrow{\text{P}_1} \begin{bmatrix} 3 & 8 & 14 \\ \frac{1}{3} & \frac{2}{3} & \frac{4}{3} \\ \frac{2}{3} & \frac{4}{3} & \frac{13}{3} \end{bmatrix}$$

$$P_1 = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\xrightarrow{\text{P}_2} \begin{bmatrix} 3 & 8 & 14 \\ \frac{1}{3} & -\frac{2}{3} & \cancel{\frac{4}{3}} \\ \frac{2}{3} & -1 & \cancel{\frac{13}{3}} \end{bmatrix} \quad \frac{2}{3} \div (-\frac{2}{3}) = -1 \quad 4 - \frac{14}{3} = -\frac{2}{3} \quad \xrightarrow{\text{P}_2} \begin{bmatrix} 3 & 8 & 14 \\ \frac{1}{3} & -\frac{2}{3} & -\frac{2}{3} \\ \frac{2}{3} & -1 & 13 \end{bmatrix}$$

$$13 = \left[ \begin{array}{cc} \frac{2}{3} & -1 \end{array} \right] \left[ \begin{array}{c} 14 \\ -\frac{2}{3} \end{array} \right] = 3$$

no row pivoting necessary

$$\therefore P = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad L = \begin{bmatrix} 1 & & \\ \frac{1}{3} & 1 & \\ \frac{2}{3} & -1 & 1 \end{bmatrix}, \quad U = \begin{bmatrix} 3 & 8 & 14 \\ -\frac{2}{3} & -\frac{2}{3} & \\ & & 3 \end{bmatrix}$$

and we can check  $LU = PA$ .

Given  $A \in \mathbb{R}^{n \times n}$ , decompose  $A = UL$  where  $U \in \mathbb{R}^{n \times n}$  unit upper triangular  
 $L \in \mathbb{R}^{n \times n}$  lower triangular.

We partition the matrix A

$$A = \begin{bmatrix} A(1:k-1, 1:k) & \underline{A(1:k-1, k)} & A(1:k-1, k+1:n) \\ \underline{A(k, 1:k-1)} & A(k, k) & A(k, k+1:n) \\ A(k+1, 1:k-1) & A(k+1, k) & A(k+1, k+1:n) \end{bmatrix}$$

$$U = \begin{bmatrix} U(1:k-1, 1:k-1) & U(1:k-1, k) & U(1:k-1, k+1:n) \\ 0 & 1 & U(k, k+1:n) \\ 0 & 0 & U(k+1:n, k+1:n) \end{bmatrix}$$

$$L = \begin{bmatrix} L(1:k-1, 1:k-1) & 0 & 0 \\ L(k, 1:k-1) & L(k, k) & 0 \\ L(k+1:n, 1:k-1) & L(k+1:n, k) & L(k+1:n, k+1:n) \end{bmatrix}$$

$$A(k, 1:k-1) = L(k, 1:k-1) + U(k, k+1:n) \cdot L(k+1:n, 1:k-1)$$

$$A(k, k) = L(k, k) + U(k, k+1:n) \cdot L(k+1:n, k)$$

$$\Rightarrow A(k, 1:k) = L(k, 1:k) + U(k, k+1:n) \cdot L(k+1:n, 1:k)$$

Since  $a_{ij}$  can be overwritten by  $\begin{cases} u_{ij} & \text{if } i < j \\ l_{ij} & \text{if } i \geq j \end{cases}$

$$\Rightarrow L(k, 1:k) = A(k, 1:k) - U(k, k+1:n) \cdot L(k+1:n, 1:k)$$

$$\Rightarrow \underline{A(k, 1:k) = A(k, 1:k) - A(k, k+1:n) \cdot A(k+1:n, 1:k)}$$

For

$$A(1:k-1, k) = U(1:k-1, k) \cdot L(k, k) + U(1:k-1, k+1:n) \cdot L(k+1:n, k)$$

$$\Rightarrow U(1:k-1, k) = (A(1:k-1, k) - U(1:k-1, k+1:n) \cdot L(k+1:n, k)) / L(k, k)$$

$$\Rightarrow A(1:k-1, k) = (A(1:k-1, k) - A(1:k-1, k+1:n) \cdot A(k+1:n, k)) / A(k, k)$$

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Algorithm for  $A = UL$  decomposition

for  $k = n:-1:1$

$$A(k, 1:k) = A(k, 1:k) - A(k, k+1:n) \cdot A(k+1:n, 1:k);$$

$$A(1:k-1, k) = [A(1:k-1, k) - A(1:k-1, k+1:n) \cdot A(k+1:n, k)] / A(k, k);$$

end

Example :  $A = \begin{pmatrix} 5 & 3 & 1 \\ 3 & 3 & 1 \\ 1 & 1 & 1 \end{pmatrix}$

$$\left[ \begin{array}{ccc} 5 & 3 & 1 \\ 3 & 3 & 1 \\ 1 & 1 & 1 \end{array} \right] \rightarrow \left[ \begin{array}{ccc} 5 & \frac{2}{3} & 1 \\ 0 & 2 & 1 \\ 1 & 1 & 1 \end{array} \right] \rightarrow \left[ \begin{array}{ccc} 5 & 1 & 1 \\ 0 & 2 & 1 \\ 1 & 1 & 1 \end{array} \right] \quad 2 \div 2 = 1$$

$$5 - [2 \ 1] \begin{bmatrix} 1 \\ 1 \end{bmatrix} = 5 - 3 = 2$$

$$\rightarrow \left[ \begin{array}{ccc} 2 & 1 & 1 \\ 0 & 2 & 1 \\ 1 & 1 & 1 \end{array} \right] \quad \therefore U = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 2 & 1 \\ 1 & 1 & 1 \end{bmatrix}, \quad L = \begin{bmatrix} 2 & 0 & 0 \\ 0 & 2 & 0 \\ 1 & 1 & 1 \end{bmatrix}$$

and the algorithm coded gives the same result.

3.  $A \in \mathbb{R}^{n \times n}$  a tridiagonal matrix (i.e.  $a_{ij} = 0$  if  $|i-j| > 1$ )  
and SPD.

Since  $A$  is a SPD tridiagonal matrix, for all  $A_{ij}$  such that  
 $i \geq j+2$ , their values should equal to 0.

Also, due to SPD, by Cholesky decomposition,  $A$  can be decomposed  
into  $A = LL^T$  where  $L$  is a lower triangular matrix with  
positive diagonals.

Assume  $L_{ij}$  for  $i \geq j+2$  is nonzero.

$$A_{ij} = L_{i\text{th row}} \cdot L_{j\text{th column}}^T \quad (\because A = LL^T)$$

Then, since  $L_{ij} \neq 0$  and  $L_{jj}^T$  is always nonzero,

therefore  $A_{ij} (i \geq j+2) \neq 0$  which contradicts that  
 $A$  is a tridiagonal matrix.

$L$  is bi-diagonal.

b) Since in a) we proved that  $L$  is a bi-diagonal matrix for  
the Cholesky decomposition of tridiagonal, SPD matrix  $A$ ,  
the substep calculations which were originally vector vector multiplication  
has reduced to scalar scalar multiplication.

for  $k = 1:n$

$$A(k,k) = \left( A(k,k) - A(k,k-1)^2 \right)^{\frac{1}{2}}$$

$$A(k+1,k) = A(k+1,k) / A(k,k)$$

end

pseudo code  
concept

for  $k = 1:n$

$$\text{if } k == 1 \\ A(k,k) = (A(k,k))^{\frac{1}{2}}$$

$$A(k+1,k) = A(k+1,k) / A(k,k)$$

else if  $k == n$

$$A(k,k) = (A(k,k) - A(k,k-1)^2)^{\frac{1}{2}}$$

Complexity

$$1+1=2$$

$$1+1+1=3$$

else

$$A(k,k) = (A(k,k) - A(k,k-1)^2)^{\frac{1}{2}}$$

$$A(k+1,k) = A(k+1,k) / A(k,k)$$

$$4(n-2)$$

$$= 4n-3$$

end

end

c)  $Ax=b$

Continuing from 3. b),

% forward substitution

for  $i = 1:n$

if  $i == 1$

$$b(i) = b(i) / A(i,i)$$

|

else

$$b(i) = (b(i) - A(i,i-1) \cdot b(i-1)) / A(i,i)$$

$$3 \cdot (n-1)$$

$$= 3n-2$$

end

end

% backward substitution

for  $i=n:-1:1$

if  $i == n$

$$b(i) = b(i) / A(i,i)$$

|

else

$$b(i) = (b(i) - b(i+1) \cdot A(i,i+1)) / A(i,i)$$

$$\frac{3(n-1)}{= 3n-2}$$

end

end

$$\therefore \text{Total} : 4n-3 + 3n-2 + 3n-2 = 10n-7 \approx 10n + O(1)$$

#### 4. Discrete 1-D Laplacian Equation $Ax = b$

$$A = \begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & \ddots & \ddots & -1 \\ & & & \ddots & -1 \\ & & & & 2 \end{bmatrix} \in \mathbb{R}^{n \times n}, \quad b = \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix} \in \mathbb{R}^n$$

a) It can directly be shown that  $A$  is a symmetric matrix (i.e.  $A = A^T$ )

$$\begin{aligned} x^T A x &= [x_1 \ x_2 \ \dots \ x_n] \begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & \ddots & \ddots & -1 \\ & & & \ddots & -1 \\ & & & & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \\ &= [2x_1 - x_2, -x_1 + 2x_2 - x_3, -x_2 + 2x_3 - x_4, \dots, -x_{n-1} + 2x_n] \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \\ &= \underbrace{2x_1^2 - x_1x_2 - x_1x_2 + 2x_2^2}_{x_1^2 + x_2^2 - 2x_1x_2 + x_2^2 + x_2^2} - x_2x_3 - x_2x_3 + 2x_3^2 \dots - x_{n-1}x_n + 2x_n^2 \\ &= x_1^2 + (x_1 - x_2)^2 + (x_2 - x_3)^2 + \dots + (x_{n-1} - x_n)^2 + x_n^2 \end{aligned}$$

Since all of them are square terms,  $x^T A x > 0$  for any  $x \in \mathbb{R}^n$  and  $x \neq 0$ .

∴ Since  $A = A^T$  (symmetric) and  $x^T A x > 0$  (positive definite),

$A$  is a symmetric positive definite (SPD) matrix.

