CFG $\leftrightarrow$ PDA

Tuesday, October 15, 2019 6:54 PM

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CFG

V: variables, $S \in V$

Σ: terminals

R: production rules $\subseteq V \to (V \cup \Sigma)^*$.

PDA

Q: states $q_0 \in Q$, $F \subseteq Q$.

Σ: input alphabet

Γ: stack alphabet

δ: transition relation

$$Q \times \Sigma \times \Gamma \longrightarrow 2^{Q \times \Gamma}$$

$$\delta(q, a, b) = \{ (q_1, b_1), (q_2, b_2) \dots \}$$

input top of stack.

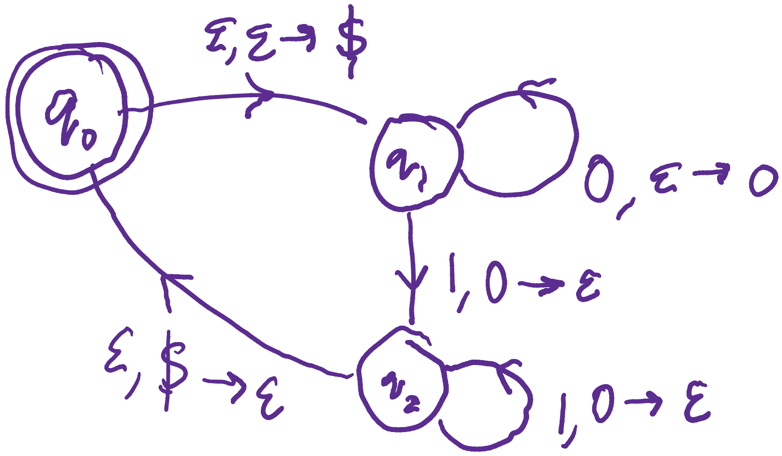
CFG for $(0^n 1^n)^*$

$S \to SS$

$S \to \varepsilon$

$S \to 0S1$

PDA:



Q. Is there always a PDA for a CFG and vice versa?

Given CFG, how to construct a PDA?

E.g.

$S \to SS$

$S \to (S)$

$S \to \varepsilon$

input: $(()(()))$

S

(S)

(SS)

$((S)S)$

What to keep in stack?

What to keep in stack?
everything in current string to the right of (and including) first variable.

((S)S)

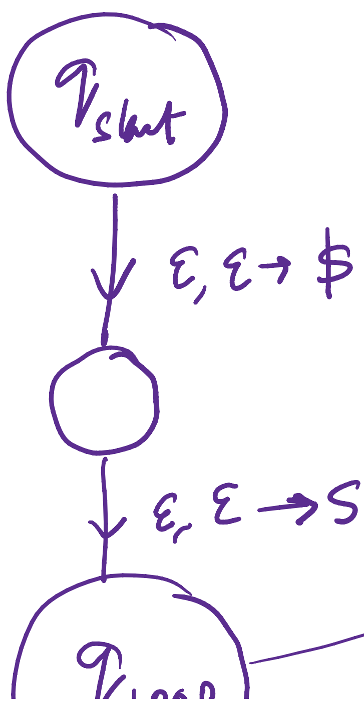
(()S)

(()(S))

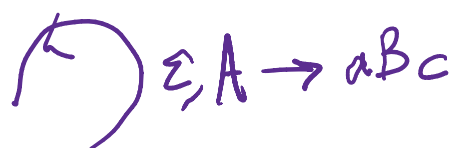
(()((S)))

(()(()))

- If top of stack is terminal, match with input.
- If ________ variable, apply rule

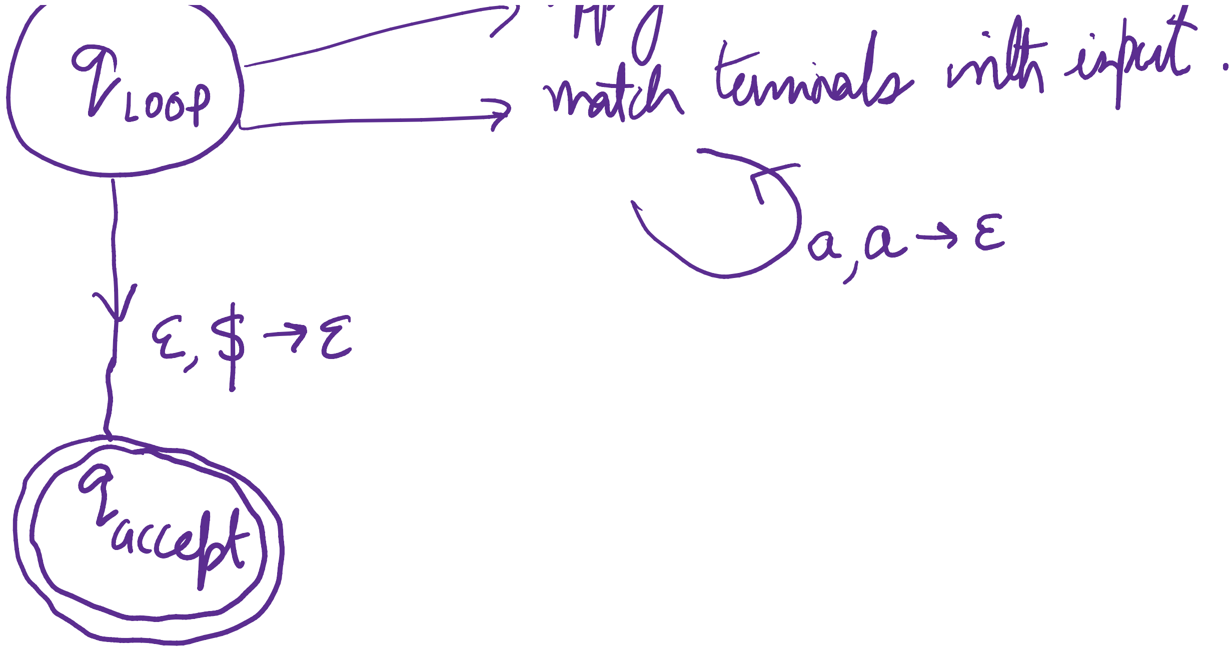


$A \to aBc$

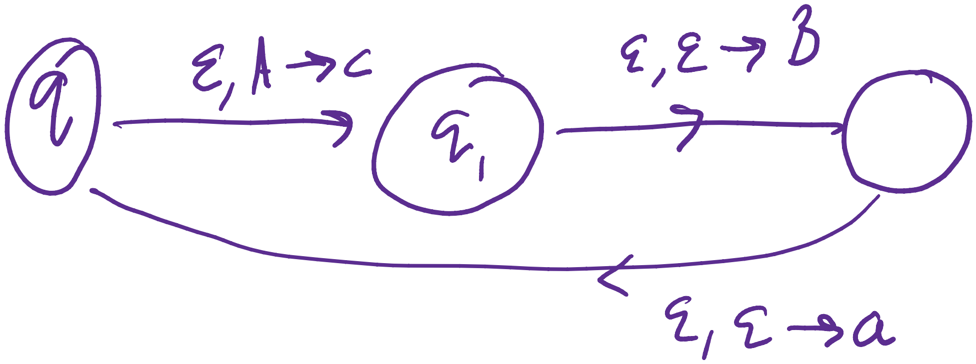


→ apply rules

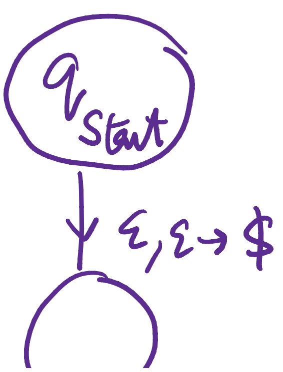
- match terminals with input.

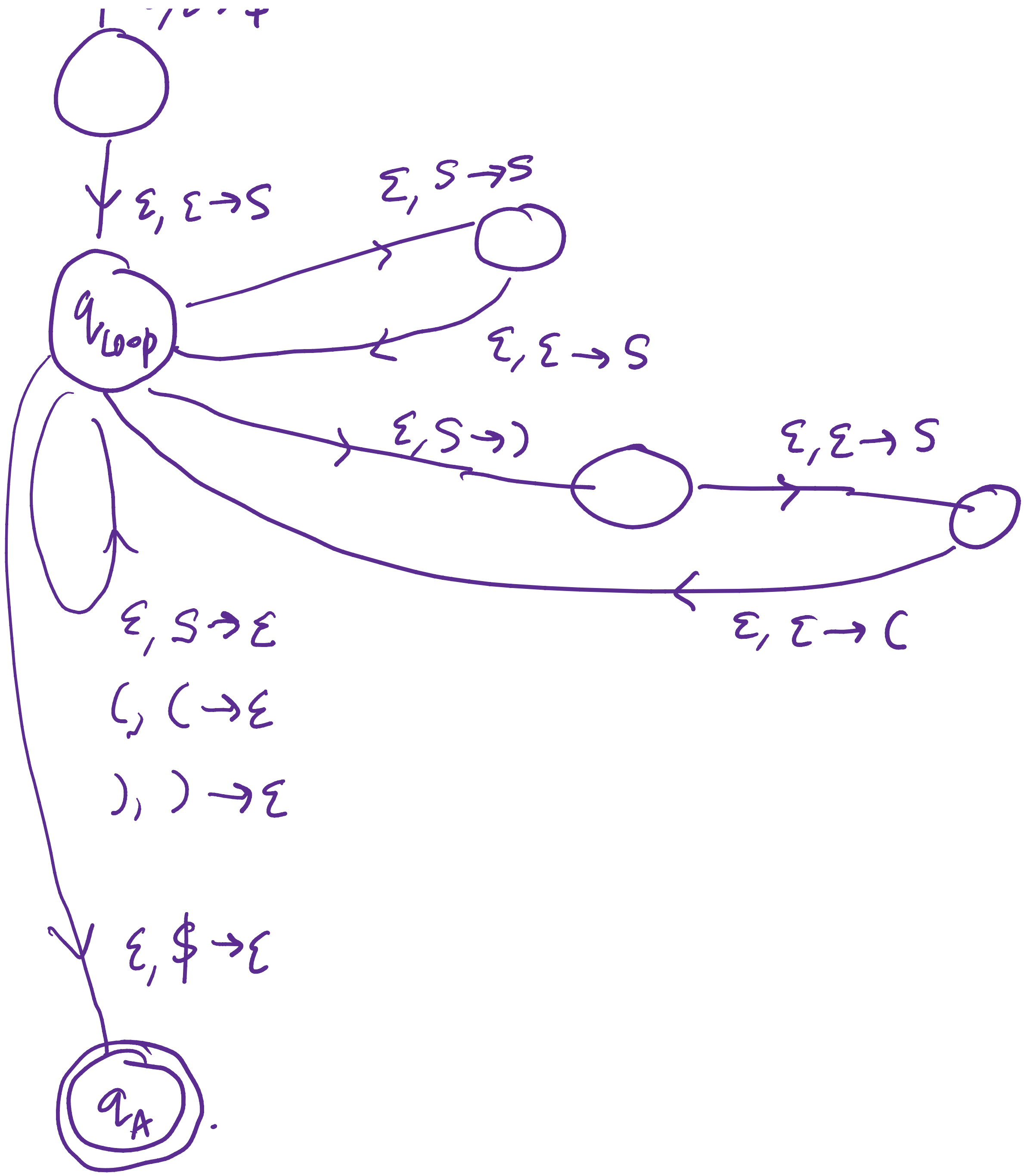


$A \to aBc$



So, for "balanced parenthesis",




ε, ε → S
ε, S → S
q_{Loop}
ε, ε → S
ε, S →)
ε, ε → S
ε, ε → (
ε, S → ε
(, (→ ε
),) → ε
ε, $ → ε
q_A