

# Robust Estimation

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Many examples of learning models from data:

Gaussians, Mixtures, ICA, topic models, dictionaries etc.

Efficient algorithms for them.

but what if not all the data is generated by the model? e.g.  $(1-\epsilon)$  is from Model  
 $\epsilon$  is arbitrary!

Even worse, malicious adversary replaces  $\epsilon$  fraction of data with points of their choice.  
Can we still learn/estimate the model parameters?

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Consider isotropic transformation or SVD.

- model has  $E(X) = \mu = 0$  (say)

if  $\epsilon N$  points can be replaced  $E(X)$  can be ...  
even if 1 point is replaced! ... anything

What about singular/eigenvectors?

— What about singular / eigenvectors?

Say model has  $\sigma_i = \ell_i = \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$

if we add a point

far enough away along  $\ell_2$ , then  $\tilde{\sigma}_i \rightarrow \ell_2$ !

So, mean, variance and low-degree sample moments are not robust estimates.

Can we estimate a k-GMM with noise?

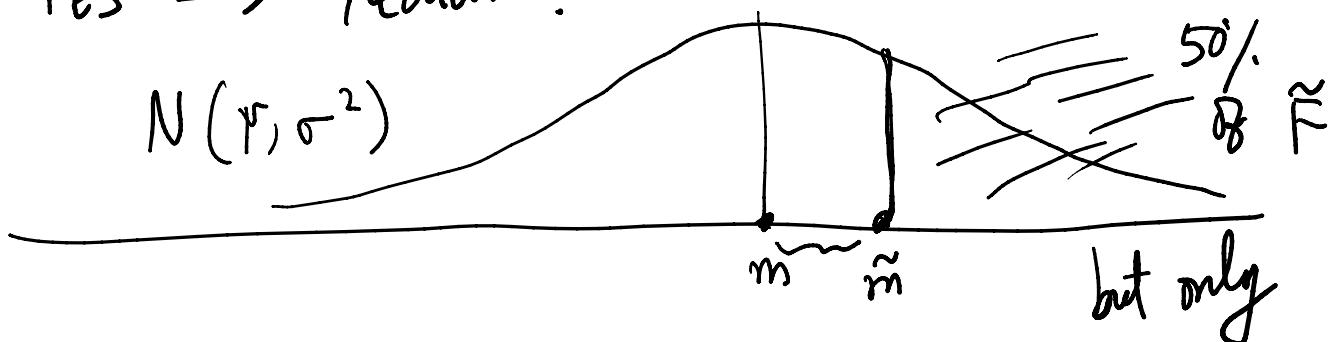
How about for  $k=1$ , i.e. a single Gaussian?

Let's even assume  $\Sigma = I$ , i.e.  $N(\mu, I)$ .

Can't do sample mean ....

How about in one dimension?

YES  $\rightarrow$  Median!



$$|m - \tilde{m}| \leq \varsigma$$

$$1 - c \cdot \underline{|m - \tilde{m}|} \text{ of } F.$$

$$\text{So } c \frac{|m - \bar{m}|}{\sigma} \leq \varepsilon \quad \frac{1}{2} - c \cdot \frac{|m - \bar{m}|}{\sigma} \text{ of } F.$$

$$|m - \bar{m}| = O(\varepsilon) \cdot \sigma .$$

This is best possible "agnostic" estimate in 1-d.

Higher dim? median along each coordinate?

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Robust Statistics Huber; Tukey. (~1960)

Tukey Ellipsoid : Smallest ellipsoid containing half the points.

best estimator, but hard to compute!



[2016] For a large class of distributions

including arbitrary Gaussians, we can robustly estimate mean, covariance up to information-theoretic limits.

## information-theoretic limits -

$$\left\| \sum_{j=1}^{\frac{1}{2}} (\hat{\mu}_j - \mu_j) \right\| \leq O(\varepsilon \sqrt{\log \frac{1}{\varepsilon}}).$$

Error does not grow with dimension!

Two simple algorithms:

## ① Iterative filtering

② recursive dimension halving.

Both based on following two ideas.

Lemma: Suppose Gaussian has  $\Sigma = I$

$$\text{If } \|\tilde{\Sigma}\|_2 \leq 1 + \varepsilon \Rightarrow \|\tilde{\Gamma} - \Gamma\|_2 = O\left(\varepsilon \sqrt{\log \frac{1}{\varepsilon}}\right)$$

if noise is only additive, then  $\sigma = O(\epsilon)$ .

Proof:  $G \cup B = S$

Good /  
Bad /  
Noise .

Assume the  $\mu = 0$ .

$$\tilde{v} = (1-\varepsilon)v + \varepsilon v_B^T \text{ Noise.}$$

Assume true  $v=0$ .  
then  $\tilde{v} = \varepsilon v_B$ .

$$\tilde{\Sigma} = (1-\varepsilon)I + \varepsilon \sum_B + (\varepsilon - \varepsilon^2) v_B v_B^T$$

(in additive model)

$$\therefore \text{for } \Sigma_B \quad v = \frac{v_B}{\|v_B\|}$$

$$1 + \varrho \geq v^T \tilde{\Sigma} v \geq 1 - \varepsilon + (\varepsilon - \varepsilon^2) \|v_B\|^2$$

$$2\varepsilon \geq \frac{(\varepsilon - \varepsilon^2)}{\varepsilon^2} \|v\|^2$$

$$\Rightarrow \|v\| = O(\varepsilon).$$

With general noise  $\|v\| = O(\varepsilon \sqrt{\log \frac{1}{\varepsilon}})$ .

Idea 2: Remove points so that  $\|\Sigma\|_2$  is close to 1.

How?

Suppose  $\exists r$ :

$$r^T \Sigma r > 1 + C \varepsilon \sqrt{\log \frac{1}{\varepsilon}}.$$

Lemma.  $\Pr(X > t) \leq e^{-t^2/2}$

Lemma.  $\Pr[X \leq -\epsilon] \geq e^{-\frac{\epsilon^2}{2}}$

$$X \sim N(0, 1)$$

If something like this holds, then

$$v^\top E_B(x x^\top) v = O(\log \frac{1}{\epsilon})$$

and since  $\Sigma = I + \epsilon \Sigma_B + (\epsilon - \epsilon^2) F_B F_B^\top + O(\epsilon \log \frac{1}{\epsilon})$

$$\begin{aligned} \sqrt{\Sigma} v &\leq 1 + \epsilon \cdot v^\top E_B(x x^\top) v + O(\epsilon \log \frac{1}{\epsilon}) \\ &= 1 + O(\epsilon \log \frac{1}{\epsilon}). \end{aligned}$$

~~Suppose~~  $\exists t: P_S(X > t+2) > C \cdot e^{-t^2/2}$

remove all points outside  $t$ .

At least  $\frac{1}{2}$  the points removed are from  $B$ .

By the end  $\|\Sigma\|_2$  is small

and at most  $2\epsilon$  points removed.

Lemma.  $\lambda_{\min}(\Sigma) \geq 1 - \epsilon.$

$\Rightarrow r \mapsto 1 - \delta(1 + O(\epsilon))$

Remove points:  $|x| > C \Delta$ .

$$\text{Tr}(\Sigma) \leq d(1 + O(\epsilon))$$

$$\Rightarrow \lambda_{d/2}(\Sigma) \leq 1 + \epsilon.$$

Proof.  $\Sigma = (1-\epsilon)I + \epsilon \Sigma_B + (\epsilon - \epsilon^2)P_B P_B^T$

$$\lambda_{\min}(\Sigma) = \sqrt{\Sigma}V \geq (1-\epsilon),$$

$$\begin{aligned} \text{Tr}(\Sigma) &\leq (1-\epsilon)d + \epsilon \left( \frac{1}{d} \sum_{i=1}^d X_i X_i^T \right) \\ &\leq (1-\epsilon)d + O(\epsilon) \cdot d \\ &= (1+O(\epsilon))d. \end{aligned}$$

So in bottom  $d/2$  eigenspace, <sup>sample</sup> mean is a good approximation of true mean.

Recurse on top  $d/2$  eigenspace.

$$\rightarrow \text{remove } x: \|x\| > C\sqrt{\dim}$$

SVD.

$$\log d \text{ levels of recursion} \Rightarrow \text{cost: } O(\epsilon \sqrt{\log d})$$