

# Value Iteration with Guessing



Institute of  
Science and  
Technology  
Austria

K. Chatterjee, M. JafariRaviz, R. Saona, J. Svoboda  
Institute of Science and Technology Austria

## Objective

Compute the value of Markov Decision Processes,  
using only local updates,  
for reachability (and shortest stochastic path).

$$v \in \mathbb{R}^S \mapsto \text{Update}(v)(s) = \max_{a \in A} \left\{ \sum_{s' \in S} p(s' | s, a) v(s') \right\}$$

## Value Iteration

$$v_0(s) = 1[s \text{ is target}]$$

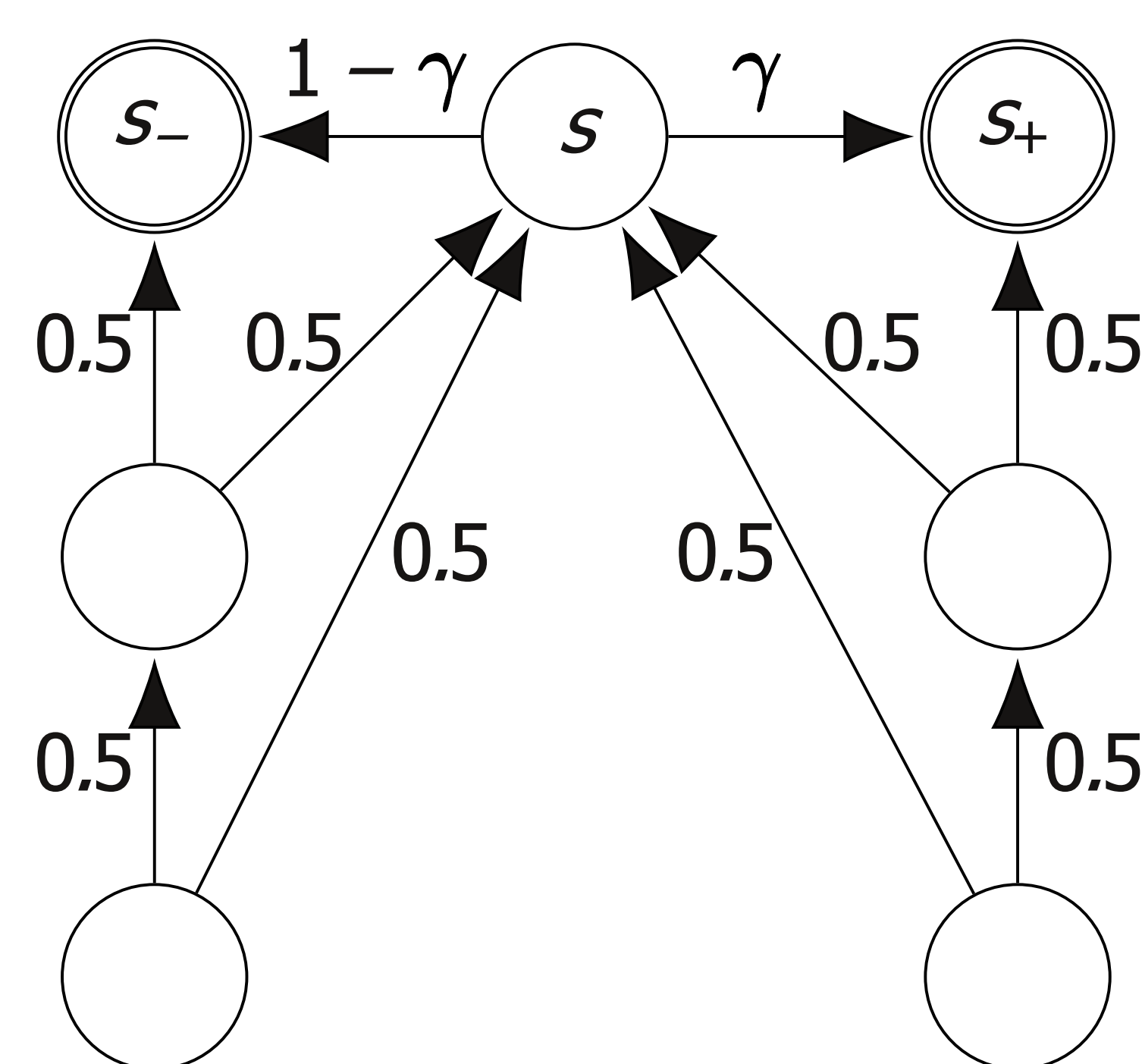
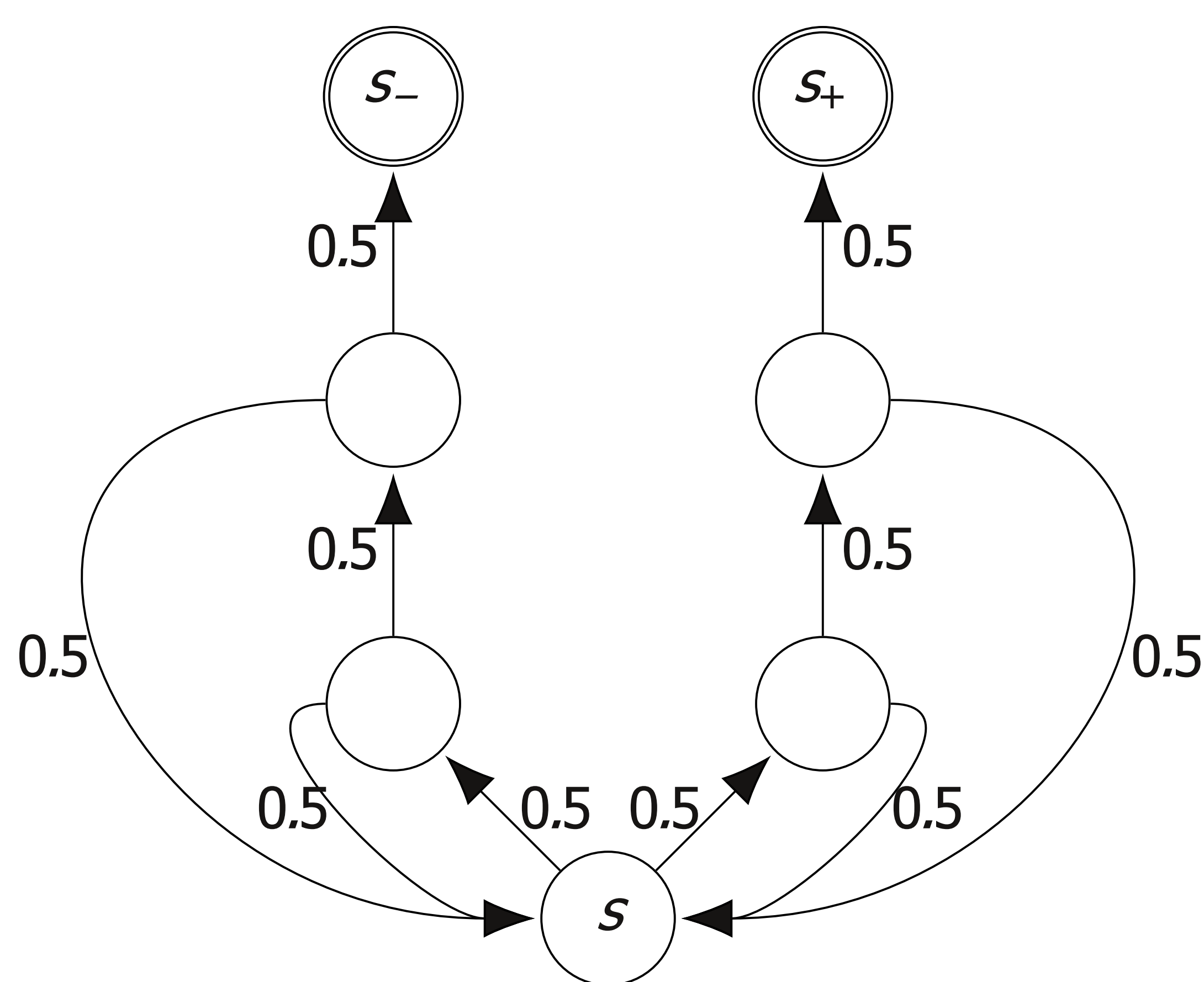
$$v_{i+1} = \text{Update}(v_i)$$

$$\|v_{|S|i} - \text{val}\|_{\infty} \leq \left(1 - p_{\min}^{|S|}\right)^i \|v_0 - \text{val}\|_{\infty}$$

Number of updates required for approximation:

$$(-\log(\varepsilon)|S|/p_{\min})^{\Omega(|S|)}$$

## Guessing to simplify structure



## Inequalities

Guessed values can be verified with the recursive solution.

$$\text{Update}(\gamma, \text{val}[s = \gamma])(s) > \gamma \Leftrightarrow \text{val}(s) > \gamma$$

Approximate verification is possible.

$$\text{Update}(\gamma, \text{val}[s = \gamma])(s) > \gamma - \varepsilon p_{\min}^{|S|} \Rightarrow \text{val}(s) > \gamma - \varepsilon$$

## Markov Chains

We can precompute which updates to ask for using linear space and almost linear time  
to obtain an approximation of the value using few updates.

Precomputation

$$\mathcal{O}((|S| + |E|) \log |S|)$$

Numer of updates

$$(-\log(\varepsilon)|S|/p_{\min})^{\mathcal{O}(\sqrt{|S|})}$$

## Markov Decision Processes

Group 1: 170 instances

All algorithms take less than 0.1 secs

Group 2: 135 instance

Fastest and slowest algorithms are  
only at most 1.10 times of each other

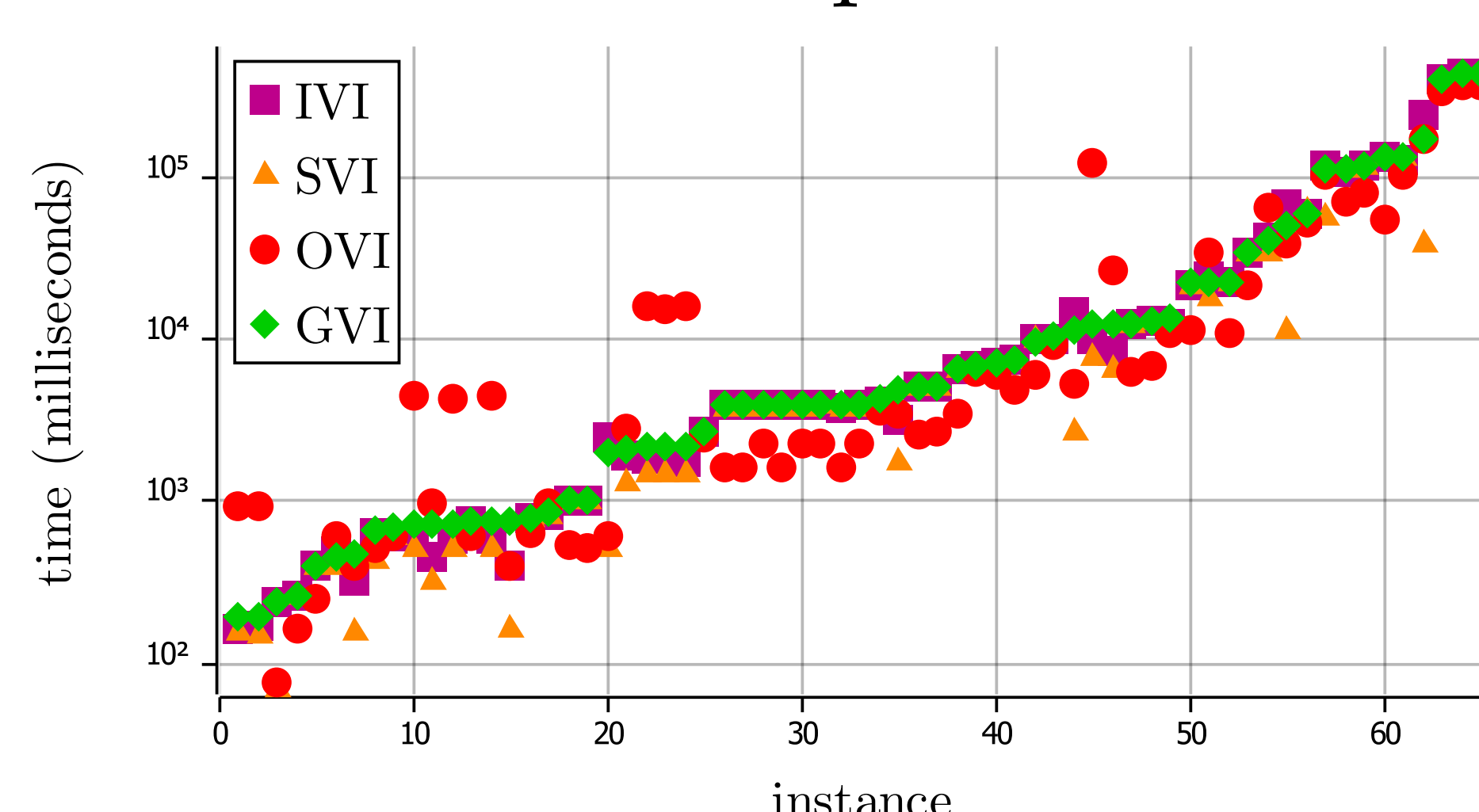
Group 3: 139 instances

Guessing VI is not the fastest

Group 4: 30 instances

The rest, our improvement

Group 3



Group 4

