Network Driver

OS Study Session #6

Labs
Part One: NIC

Problem Summary

- Filling the function e1000_recv() and e1000_transmit()
- Understanding overall process of sending / receiving data with NIC

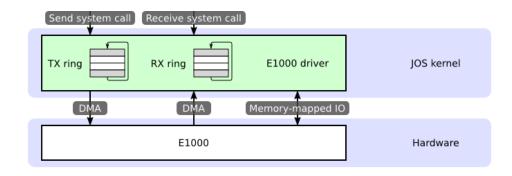
Role of Network Driver (E1000)

- Allocating memory for the transmit and receive queues
- Setting up DMA descriptors
- Configuring the E1000 with the location of these queues
- everything after that is asynchronous

DMA (Direct Memory Access): read and write packet data directly from memory without involving the CPU

Receive System Call

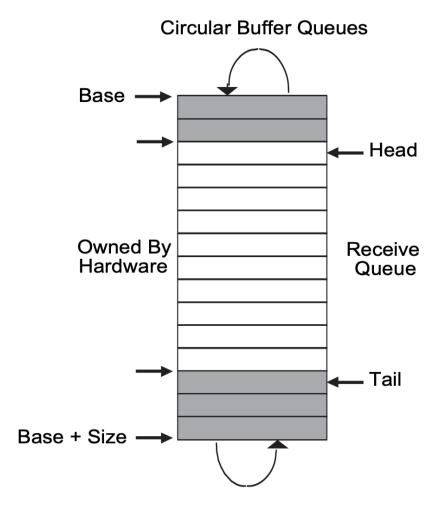
When the E1000 receives a packet, it copies it into the next DMA descriptor in the receive queue, which the driver can read from at its next opportunity.



https://pdos.csail.mit.edu/6.828/2018/labs/lab6/

- 1. NIC receives the new data packet in receiving buffer
- 2. An interrupt is triggered by NIC to notify
- Network driver calls e1000_recv() to handling this interrupt (e1000_recv() = interrupt handler)
- 4. e1000_recv() transfers the data to OS's network stack

Receive Descriptor Ring Structure



- The driver always add descriptors to the tail and moves the tail pointer
- Hardware always consumes descriptors from the head and moves the head pointer
- The contents of the queue are the descriptors between head and tail pointer
- Receive queue contents: Free descriptors that the card can receive packets into
- Transmit queue content : packets waiting to be sent

e1000_recv

```
#define E1000_RXD_STAT_DD 0x01 /* Descriptor Done */
#define E1000_RXD_STAT_EOP 0x02 /* End of Packet */
```

net_rx() : deliver a packet to the networking stack

- 1. Find an the next waiting received packet index
- 2. Check if a new packet is available
- 3. Deliver a packet
- 4. Allocate a new buffer (kalloc()) to replace the one just given to net_rx().
- 5. Update E1000_RDT register to the new index

```
static void
e1000 recv(void)
  while (1) {
   int idx = (regs[E1000 RDT] + 1) % RX RING SIZE;
    if ((rx ring[idx].status & E1000 RXD STAT DD) == 0) {
      return;
    if (rx ring[idx].status & E1000 RXD STAT E0P) {
      int len = rx ring[idx].length;
     net rx(rx bufs[idx], len);
      rx bufs[idx] = kalloc();
      if(rx bufs[idx] == 0){
        return;
      rx ring[idx].status = 0;
      rx ring[idx].addr = (uint64)rx bufs[idx];
    regs[E1000 RDT] = idx;
```

e1000_transmit

```
#define E1000_TXD_CMD_RS 0 \times 08000000 /* Report Status */ #define E1000_TXD_CMD_EOP 0 \times 01000000 /* End of Packet */
```

To transmit a packet, the driver copies it into the next DMA descriptor in the transmit queue and informs the E1000 that another packet is available;

https://pdos.csail.mit.edu/6.1810/2024/readings/8254x_GBe_SDM.pdf

- Finding the next packet index
 (by E1000_TDT control register)
- 2. Check overFlowed / not completed
- 3. Free the last buffer (kfree()) if there was
- 4. Fill the descriptor

```
e1000 transmit(char *buf, int len)
 // buf contains an ethernet frame; program it into
 // the TX descriptor ring so that the e1000 sends it. Stash
 // a pointer so that it can be freed after send completes.
 // Acquire lock
 acquire(&e1000 lock);
 int idx = regs[E1000 TDT];
 // check whether it is overflowed/not completed
 if ((tx ring[idx].status & E1000 TXD STAT DD) == 0) {
   release(&e1000 lock);
   return -1;
 // free the last buffer
 if (tx bufs[idx])
   kfree(tx bufs[idx]);
  tx bufs[idx] = buf;
 tx ring[idx].length = len;
 tx ring[idx].addr = (uint64) (buf);
 tx ring[idx].cmd = E1000 TXD CMD RS | E1000 TXD CMD EOP;
 tx ring[idx].status = 0;
 // update ring location
 regs[E1000 TDT] = (idx + 1) % TX RING SIZE;
 release(&e1000 lock);
  return 0;
```

Labs
Part Two: UDP Receive

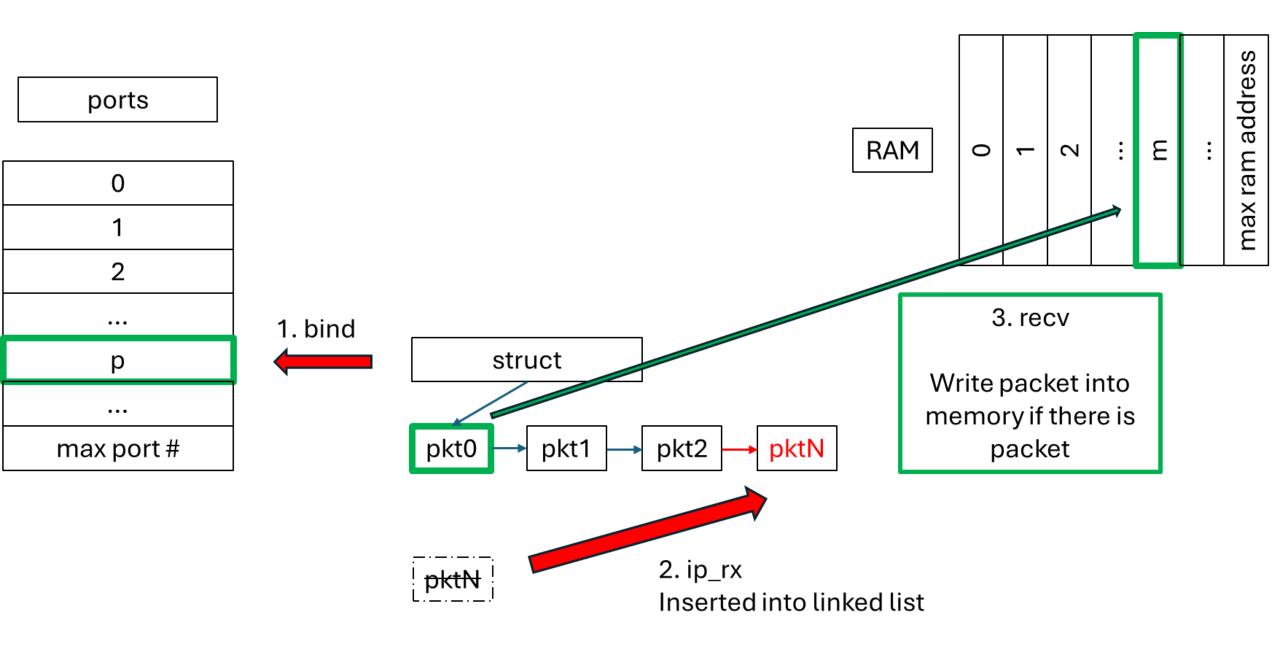
Problem Summary

- A system call for **sending UDP packets** has already been implemented.
- The objective is to create a system call for receiving UDP packets

Related system-call API and function

- send(sport, dst, dport, *buf, len) (→ Already implemented)
 - Sends a packet to the destination specified by (dst, dport).
 - o **Packet:** Starts at buf and has a length of len.
- recv(dport, *src, *sport, *buf, len) (→ To be implemented)
 - Receives packets arriving at dport.
 - Writes the received information into src and sprt, which are virtual addresses.
- bind (→ To be implemented)
 - A process must bind to a port before calling recv.
- ip_rx (→ To be implemented)
 - Restores the packet before calling recv.

ports Packet arrived to unbound port -> discarded pkt0 1. bind struct max port # pkt0 pkt1 pkt2 Packet arrived to bound port pktN



bind(port)

- Binding port before calling ip_rx and recv
- The maximum number of packets that can arrive is limited to 16
- If a packet arrives at an unbound port, the packet should be discarded

Data structure

- Create "struct bound_port"
- Each bound port must store waiting packets.
- A list of struct bound_port instances is stored in the udb_table.
- The function find_bound_port is called to check whether the given port is bound or not.

```
struct bound port
                  // Bound port number
 uint16 port;
 int pending count; // Number of pending packets
 struct spinlock lock;
 struct pending packet *head; // Packet queue head
 struct pending packet *tail; // Packet queue tail
#define MAX PORTS 64
struct
 struct spinlock lock;
 struct bound port ports[MAX PORTS];
 int port count;
 udp table;
netinit(void)
 initlock(&netlock, "netlock"); // Initialize lock
 udp table.port count = 0;
                                 // Initialize port counter
```

bind(port)

- A lock is necessary since other processes may try to bind the same port.
- If the port is already bound or the number of bound ports has already exce eded MAX_PORTS, the binding fails.

```
uint64
sys bind(void)
  int port; argint(0, &port);
  acquire(&udp table.lock);
    // Check if the port is already bound
    if (find bound port(port) != 0 || udp table.port count >= MAX PORTS)
      release(&udp table.lock);
      return -1;
    // Initialize a new bound port structure
    struct bound port *bp = &udp table.ports[udp table.port count++];
    bp->port = port;
    bp->pending count = 0;
    bp->head = bp->tail = 0;
    initlock(&bp->lock, "bound port");
  release(&udp table.lock);
  return 0;
```

ip_rx(buf, len)

- This function must receive the packet and store it in the queue within bound_port.
- Starting from the buf address, the addresses of the Ethernet packet, IP packet, and UDP packet are stored.
- Packets are represented as structures in net.h.

```
... buf
struct * eth
struct * ip
struct * udp
...
```

```
// a UDP packet header (comes after an IP header).
struct udp {
  uint16 sport; // source port
  uint16 dport; // destination port
  uint16 ulen; // length, including udp header, not including IP header
  uint16 sum; // checksum
};
```

Step1 of ip_rx(buf, len)

- Get the destination port from the UDP header.
- The function ntohs rearranges the byte order.
 - Packets use Big-endian
 - o CPU use Little-endian

```
ip rx(char *buf, int len)
 // don't delete this printf; make grade depends on it
 static int seen ip = 0;
 if(seen ip == 0)
   printf("ip rx: received an IP packet\n");
 seen ip = 1;
 struct eth *ethhdr = (struct eth *)buf:
 struct ip *iphdr = (struct ip *)(ethhdr + 1);
 struct udp *udphdr = (struct udp *)(iphdr + 1);
 uint16 dport = ntohs(udphdr->dport);
 acquire(&udp table.lock);
 struct bound port *bp = find bound port(dport);
 if (bp == 0)
   kfree(buf):
   release(&udp table.lock);
   return;
 acquire(&bp->lock);
  release(&udp table.lock);
```

```
struct udp {
    uint16 sport; // source port
    uint16 dport; // destination port
    uint16 ulen; // length, including udp header, not including IP header
    uint16 sum; // checksum
};
```

Step2 of ip_rx(buf, len)

 If the number of waiting packets is equal to MAX_PENDING_PACKETS, return from the function.

If not, initialize the packet.

```
if (bp->pending count >= MAX PENDING PACKETS)
  kfree(buf);
  release(&bp->lock);
  return;
int payload len = ntohs(udphdr->ulen) - sizeof(struct udp);
char *data = kalloc();
struct pending packet *pp = kalloc();
memmove(data, (char *)(udphdr + 1), payload len);
memset(pp, 0, sizeof(*pp));
pp->data = data;
pp->len = payload len;
pp->src ip = ntohl(iphdr->ip src);
pp->src port = ntohs(udphdr->sport);
pp->next = 0;
```

Step3 of ip_rx(buf, len)

- Insert the prepared packet into the queue.
- It has to call the wakeup function.
 - I will explain later.

```
Update the queue
  (bp->tail == 0)
  bp->head = bp->tail = pp;
else
  bp->tail->next = pp;
  bp->tail = pp;
bp->pending count++;
// Wake up the waiting process
wakeup(bp);
kfree(buf);
release(&bp->lock);
```

recv(dport, *src, *sport, *buf, len)

- Extract a packet stored in the queue.
- If no packet has arrived, wait until a new packet arrives.
- Store the IP address of the packet in src and store the port of the packet in sport.

Step1 of recv(dport, *src, *sport, *buf, len)

- Wait until the packet is received.
- If no packet has arrived, wait (by using sleep).
- It can be awakened by the previously mentioned wakeup.
 - The wakeup function is called if a packet arrives.
 - Exit the while loop and proceed to the next step.

```
acquire(&udp table.lock);
if ((bp = find bound port(dport)) == 0)
  release(&udp table.lock);
  return -1;
acquire(&bp->lock);
release(&udp table.lock);
while (bp->head == 0)
  if (p->killed)
    release(&bp->lock);
    return -1;
  sleep(bp, &bp->lock);
```

Step2 of recv(dport, *src, *sport, *buf, len)

- Src and sport are virtual addresses, which will store the source IP and source port, respectively.
- Extract the stored packet from the queue.
- Store the source IP and source port in the given virtual addresses.
 - Using copyout api

```
pp = bp->head;
bp->head = pp->next;
if (bp->head == 0)
{
    bp->tail = 0;
}
bp->pending_count--;
int copy_len = pp->len;
if (copy_len > maxlen)
    copy_len = maxlen;

if (copyout(p->pagetable, src_arg, (char *)&pp->src_ip, sizeof(pp->src_ip)) < 0 ||
    | copyout(p->pagetable, sport_arg, (char *)&pp->src_port, sizeof(pp->src_port)) < 0)
{
    goto bad;
}

if (copyout(p->pagetable, buf, pp->data, copy_len) < 0)
{
    goto bad;
}</pre>
```