

Domino Tilings of the Chessboard

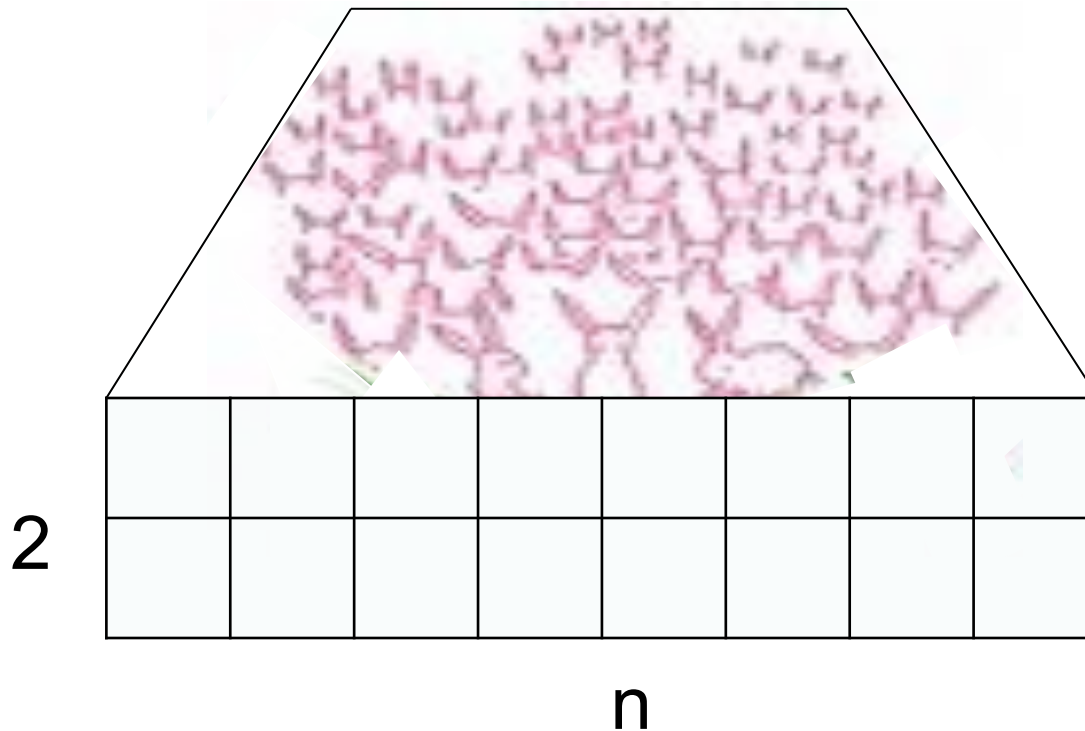
An Introduction to Sampling and Counting



Dana Randall
Schools of
Computer Science
and Mathematics
Georgia Tech

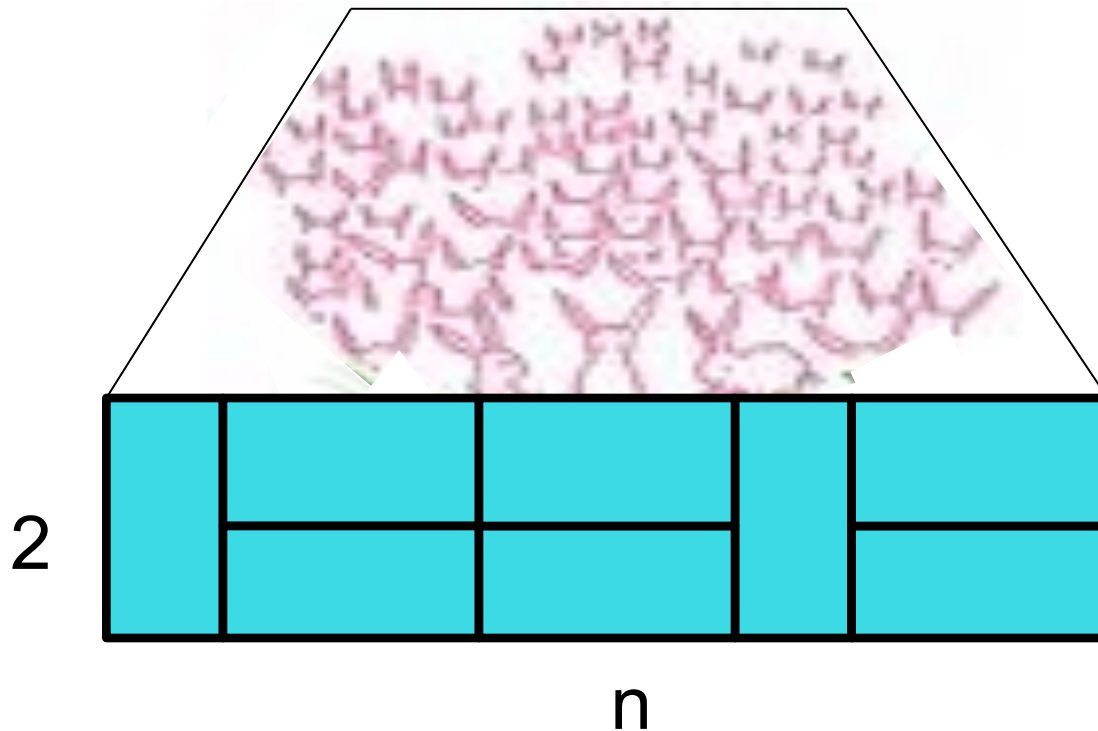
Building short walls

How many ways are there to
build a $2 \times n$ wall with 1×2 bricks?



Building short walls

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build a $2 \times n$ wall with 1×2 bricks?



Building short walls

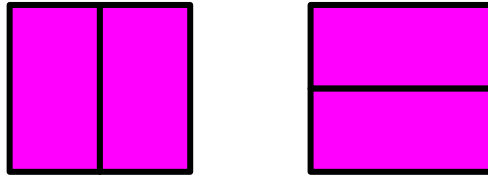
$n=0$:



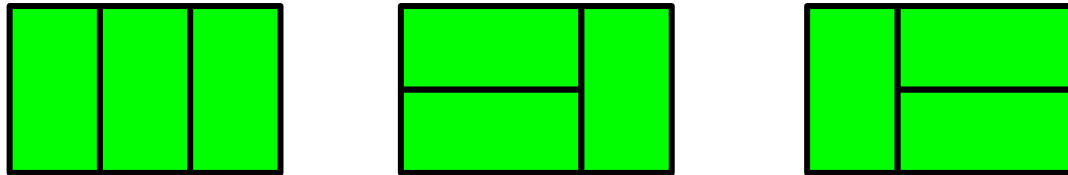
$n=1$:



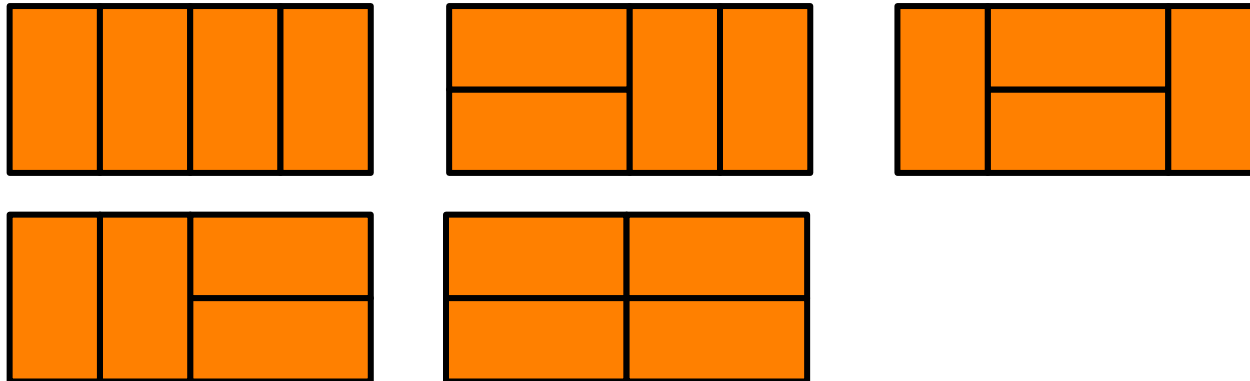
$n=2$:



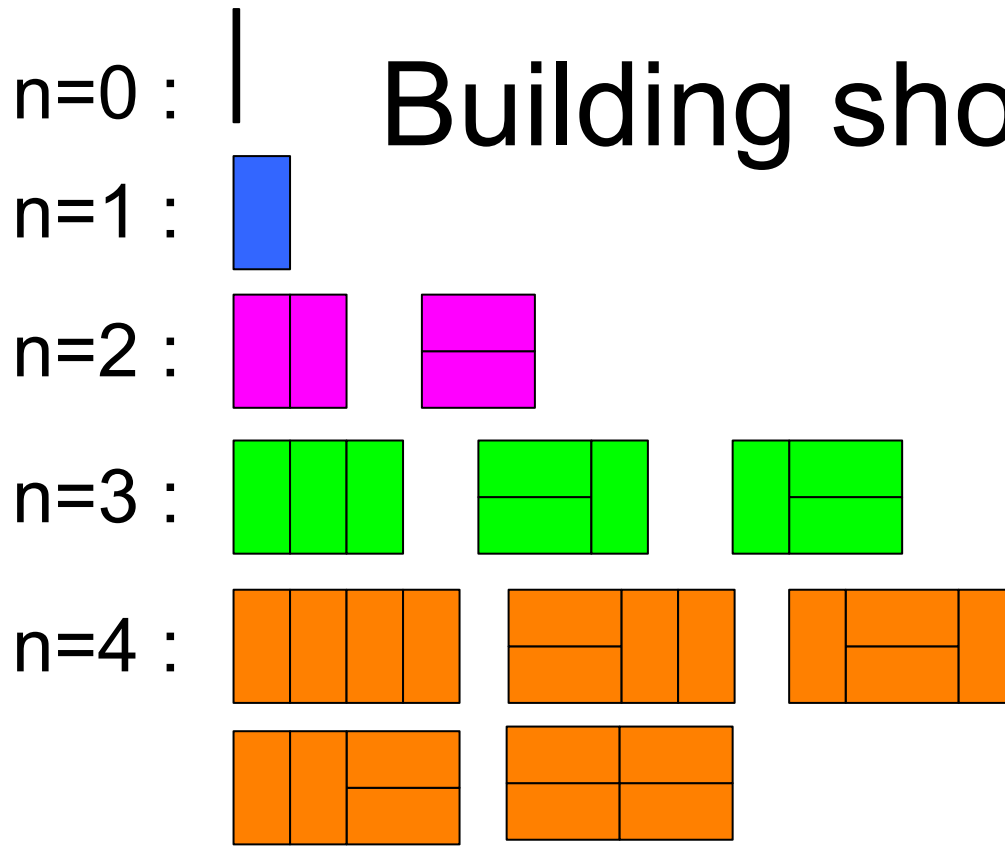
$n=3$:



$n=4$:



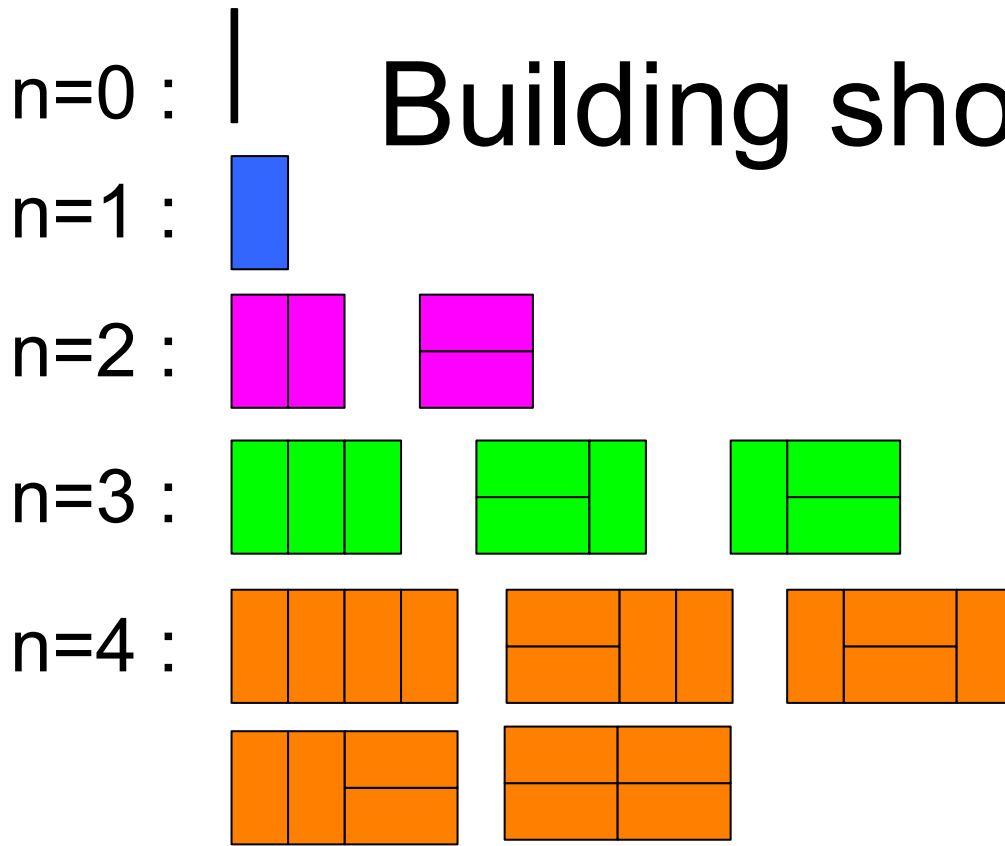
Building short walls



The number of walls equal:

$$f_n = 1, 1, 2, 3, 5$$

Building short walls




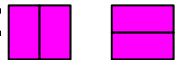
The number of walls equal:

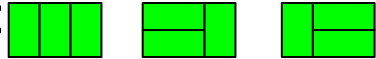
$$f_n = 1, 1, 2, 3, 5, 8, 13, 21, \dots$$

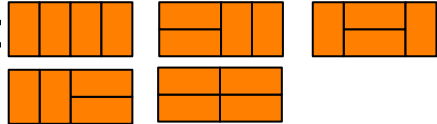


n=0 :

n=1 : 

n=2 : 

n=3 : 

n=4 : 

Building short walls



The number of walls equals:

$$f_n = 1, 1, 2, 3, 5, 8, 13, 21, \dots$$

n

$$f_n = \left\{ \begin{array}{l} \text{[Diagram: A long light blue rectangle with a blue square at the left end containing a yellow circle. The blue square is 2 units high.] } \\ \text{[Diagram: A long light blue rectangle with a blue square at the left end containing a yellow circle. The blue square is 2 units wide.] } \end{array} \right.$$

n=0 :

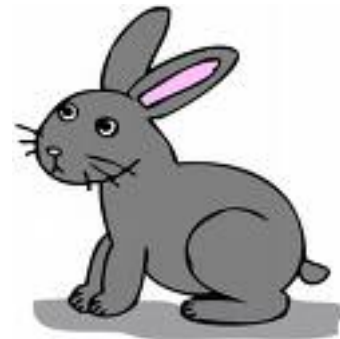
n=1 :

n=2 :

n=3 :

n=4 :

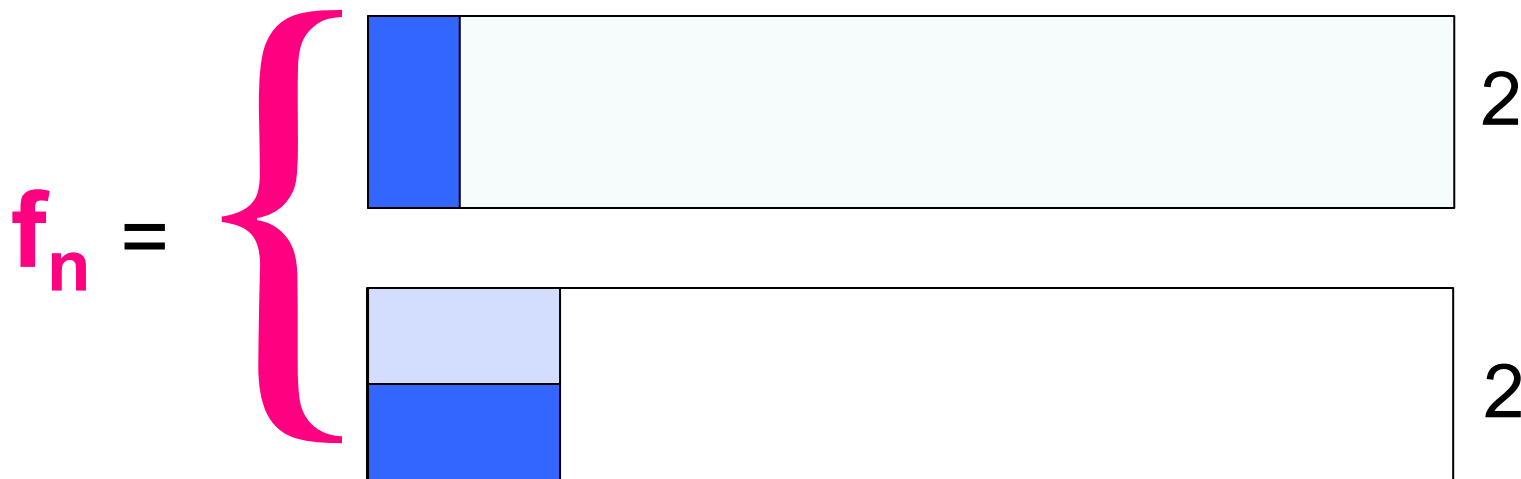
Building short walls



The number of walls equals:

$$f_n = 1, 1, 2, 3, 5, 8, 13, 21, \dots$$

n



n=0 :

n=1 :

n=2 :

n=3 :

n=4 :

Building short walls



The number of walls equals:

$$f_n = 1, 1, 2, 3, 5, 8, 13, 21, \dots$$

n

$$f_n = \left\{ \begin{array}{l} \text{[Blue block]} + f_{n-1} \\ \text{[Light blue block]} + \text{[Blue block]} + f_{n-2} \end{array} \right.$$

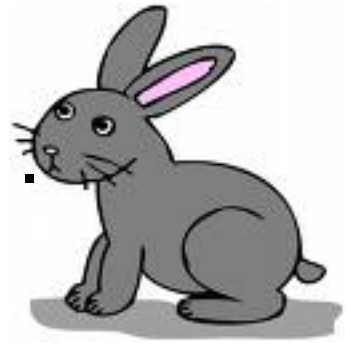
2

2

The Fibonacci Numbers

The number of walls equals:

$$f_n = 1, 1, 2, 3, 5, 8, 13, 21, \dots$$



$$f_n = \left\{ \begin{array}{l} \text{blue block} + \text{orange block } f_{n-1} \\ \text{light blue block} + \text{blue block} + \text{orange block } f_{n-2} \end{array} \right.$$

$$f_n = f_{n-1} + f_{n-2}, \quad f_0 = f_1 = 1$$

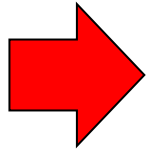
$$f_n = (\varphi^n + (1-\varphi)^n) / \sqrt{5},$$



where: $\varphi = \frac{1+\sqrt{5}}{2}$ (“golden ratio”)

Domino Tilings

Given a region R on the infinite chessboard, cover with non-overlapping 2×1 dominos.



Where is a tiling? Do any even exist?

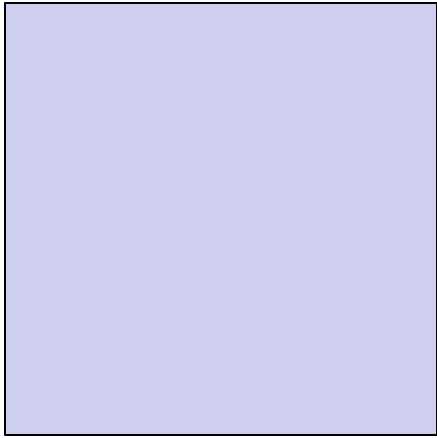
How many tilings are there?

What does a typical tiling look like?

When do we stop our algorithms?

Why do we care?

Where is a tiling? Do any exist?



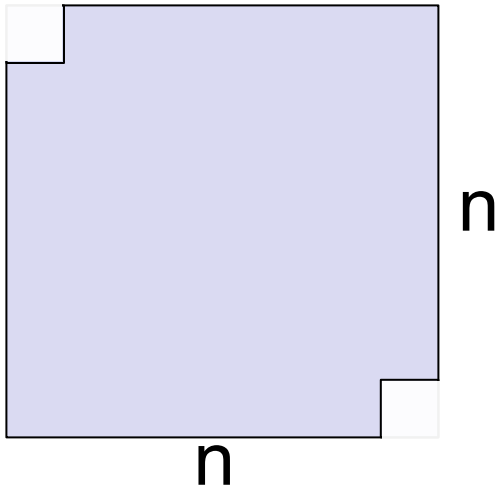
n

n

★ Only if n is even!

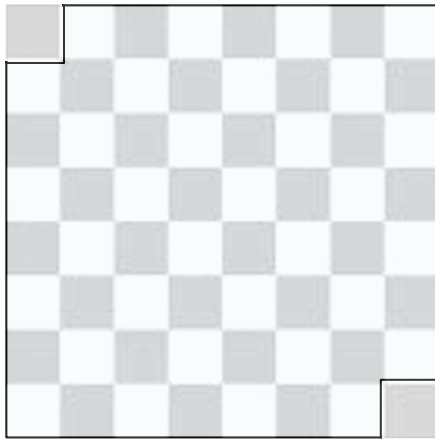
Where is a tiling? Do any exist?

- The Area of R must be even



Where is a tiling? Do any exist?

- The Area of R must be even



n

n

- ★ There must be an equal number of black and white squares.

Where is a tiling? Do any exist?

- The Area of R must be even
- With an equal number of white and black squares

Is this enough?

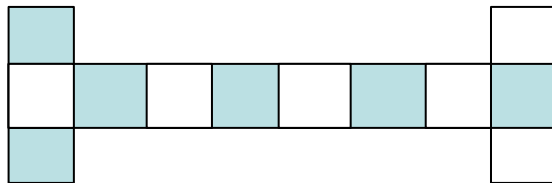


Where is a tiling? Do any exist?

- The Area of R must be even
- With an equal number of white and black squares

Is this enough?

⋮

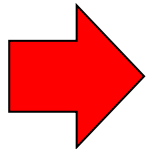


?

There is an **efficient algorithm** to decide if R is tileable and to find one if it is. [Thurston]

Domino Tilings

Where is a tiling? Do any even exist?



How many tilings are there?

What does a typical tiling look like?

When do we stop our algorithms?

Why do we care?

How many tilings are there?

Short 2 x n
walls



$$\approx \varphi^n$$

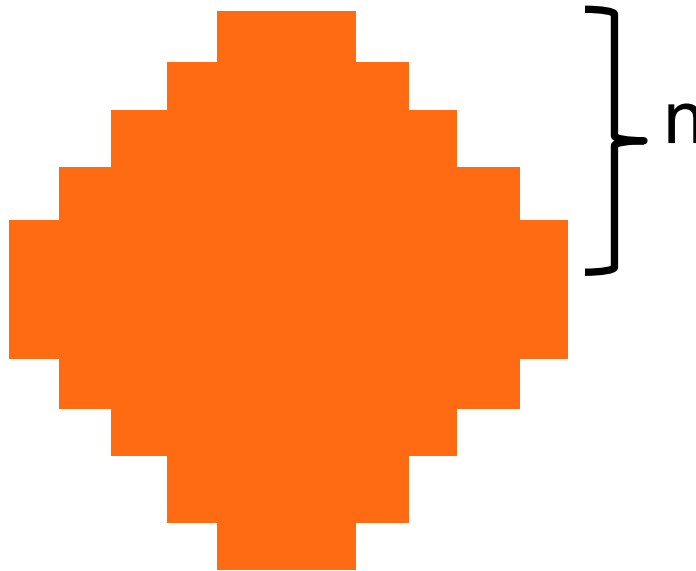
How many tilings are there?

Short 2 x n
walls



$$\approx \varphi^n$$

Aztec
Diamonds



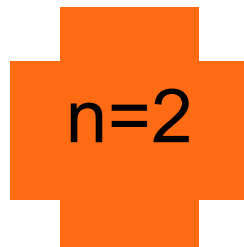
$$= 2^{n(n+1)/2}$$

[Elkies, Kuperberg,
Larson, Propp]

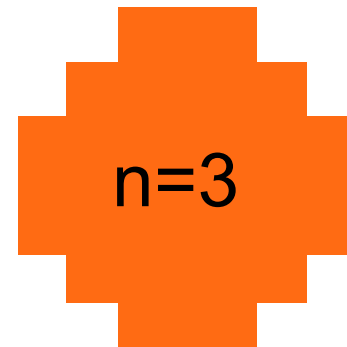
n=0 .



n=1



n=2



n=3

= 1

2

8

64

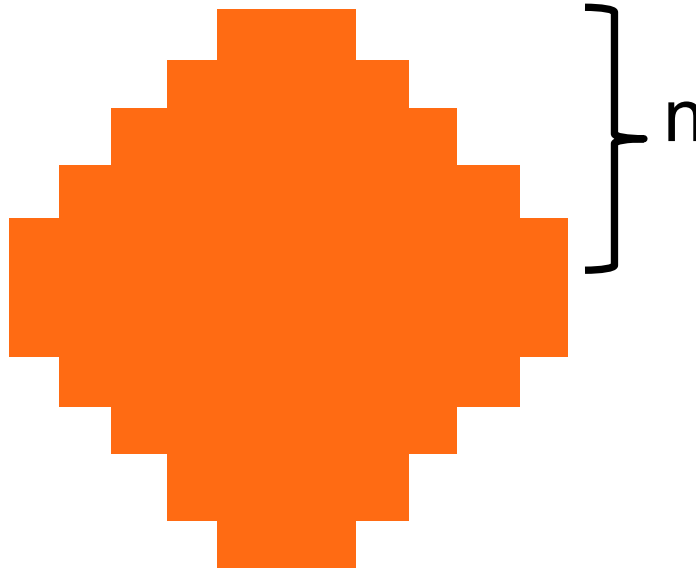
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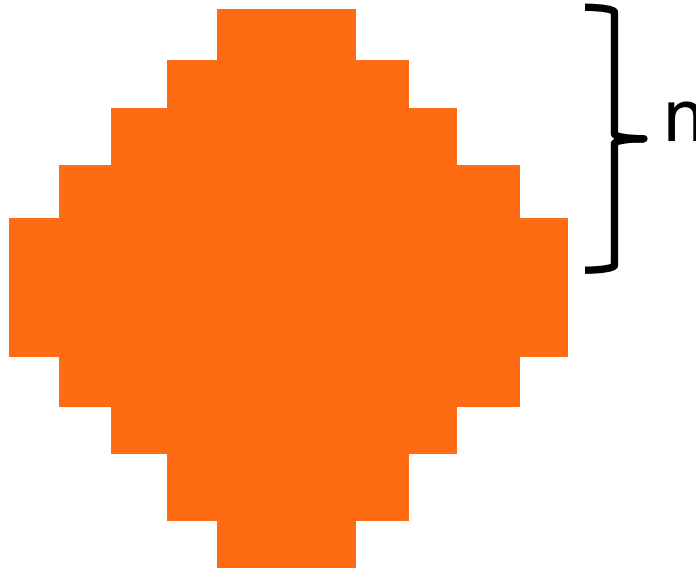
How many tilings are there?

Short 2 x n
walls



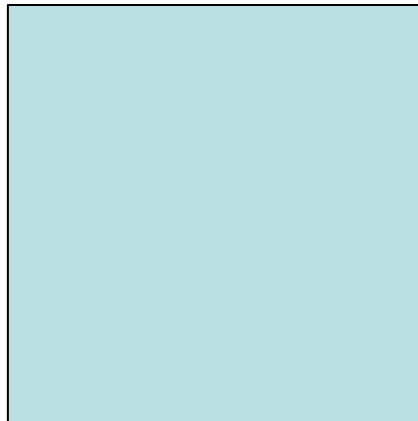
$$\approx \varphi^n$$

Aztec
Diamonds



$$= 2^{n(n+1)/2}$$

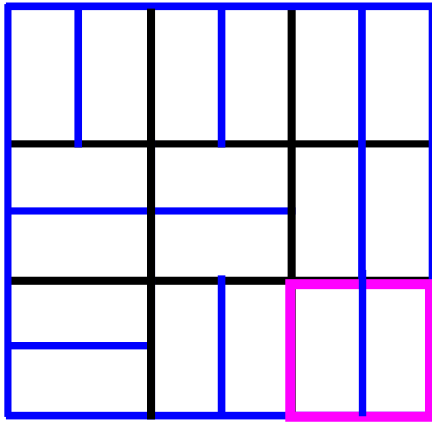
Square n x n
walls



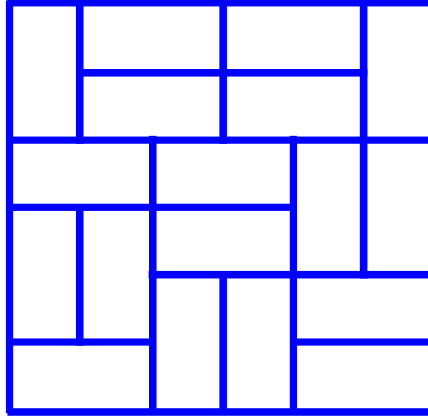
?

How many tilings are there?

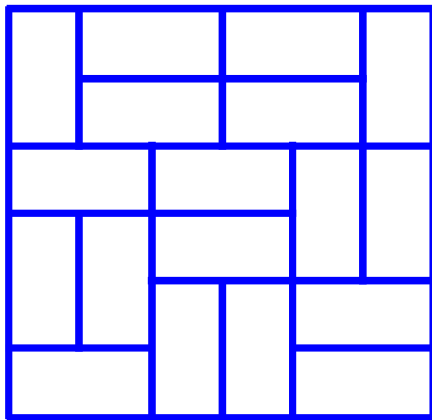
Square $n \times n$ walls



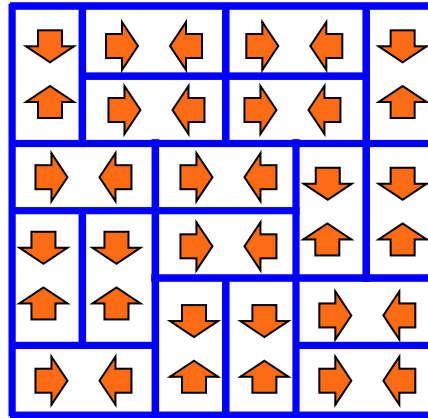
\wedge



$$2 \text{ (Area/4)} < \#$$



\wedge



$$\# < 4 \text{ (Area)}$$

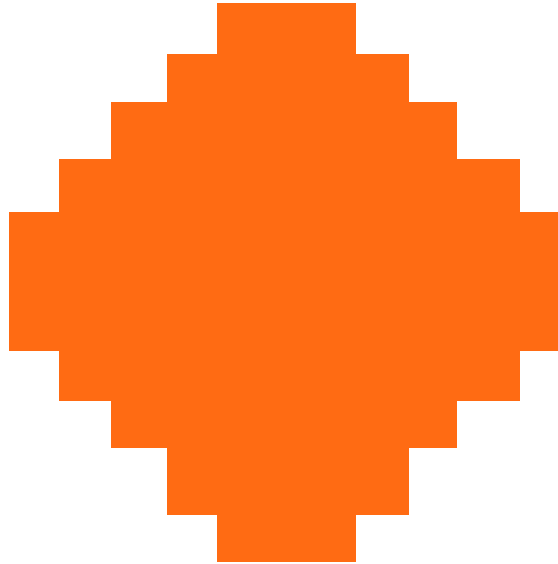
How many tilings are there?

Short 2 x n
walls



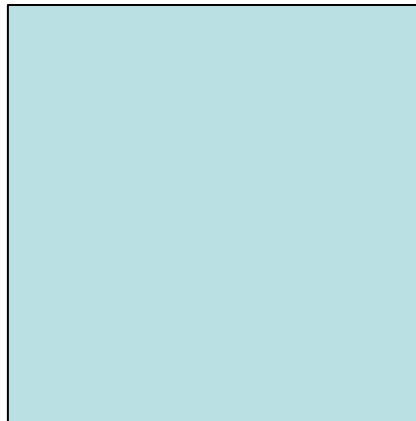
$$\# \approx \varphi^n$$

Aztec
Diamonds



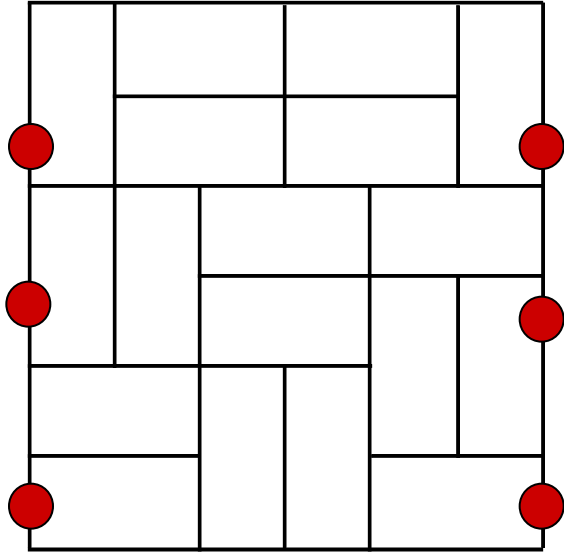
$$\# = 2^{n(n+1)/2}$$

Square n x n
walls



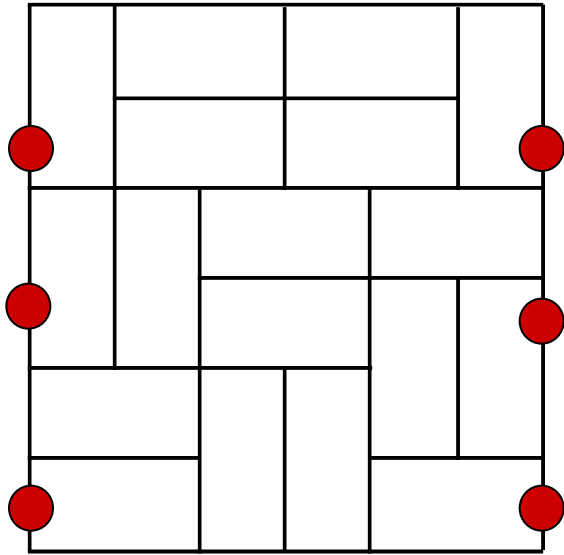
$$2^{\text{Area}/4} < \# < 4^{\text{Area}}$$

How many: An Algorithm

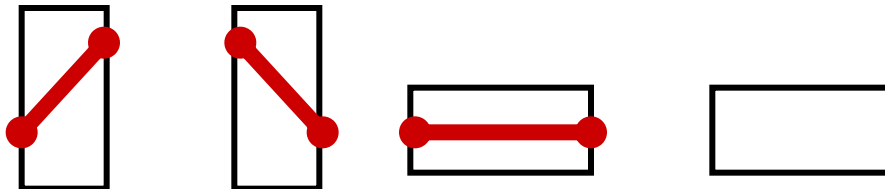


- Mark alternating vertical edges;

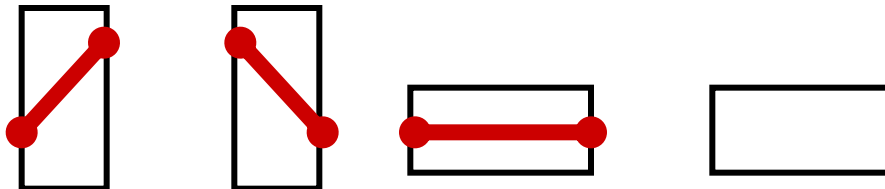
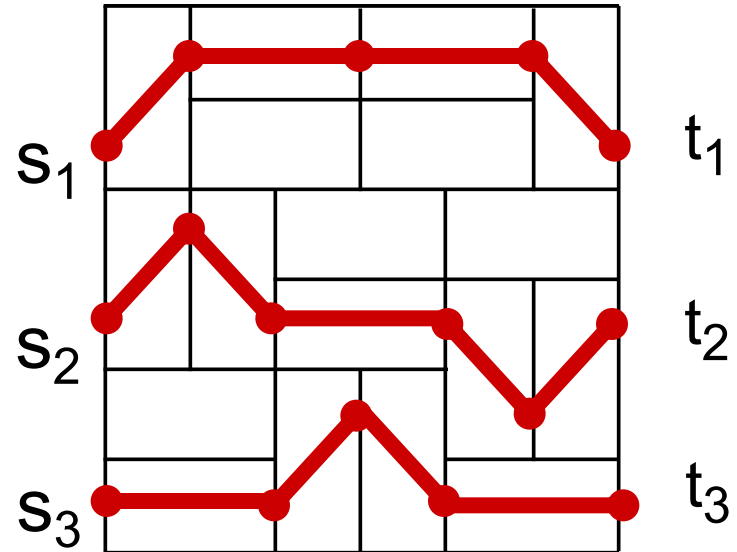
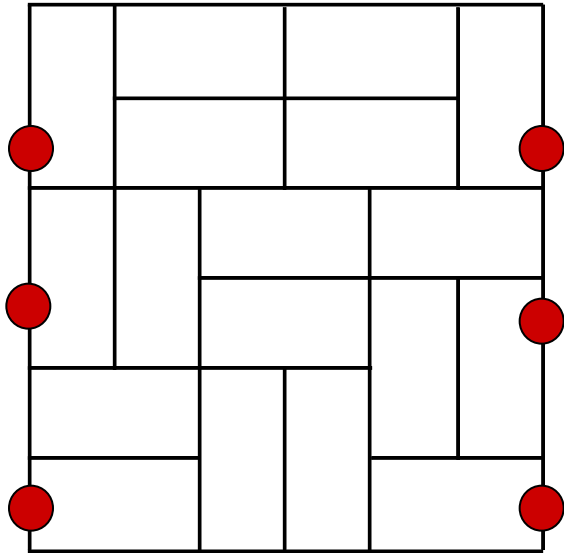
How many: An Algorithm



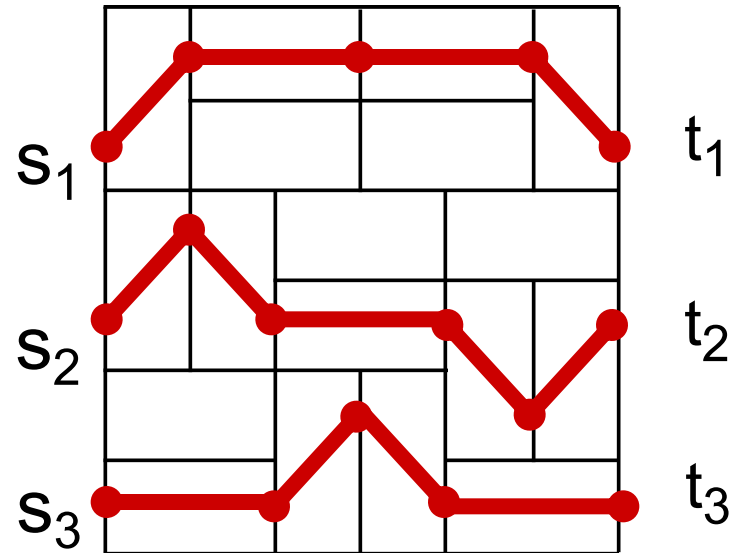
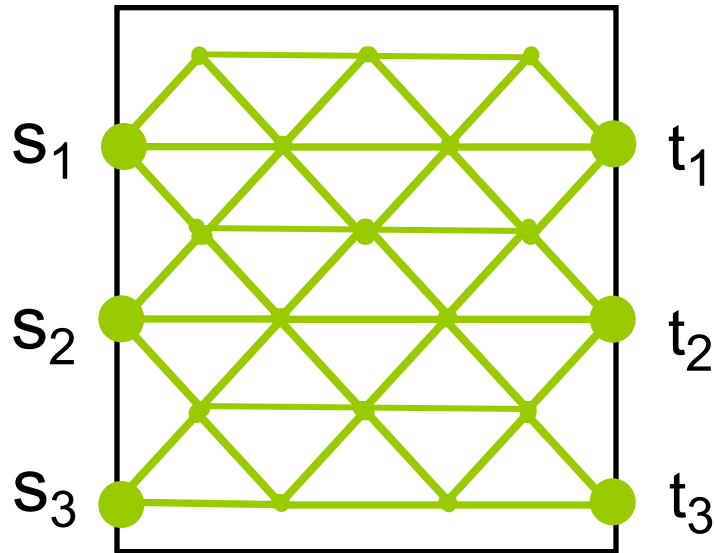
- Mark alternating vertical edges;
- Use marked tiles;
- Markings must line up!



How many: An Algorithm



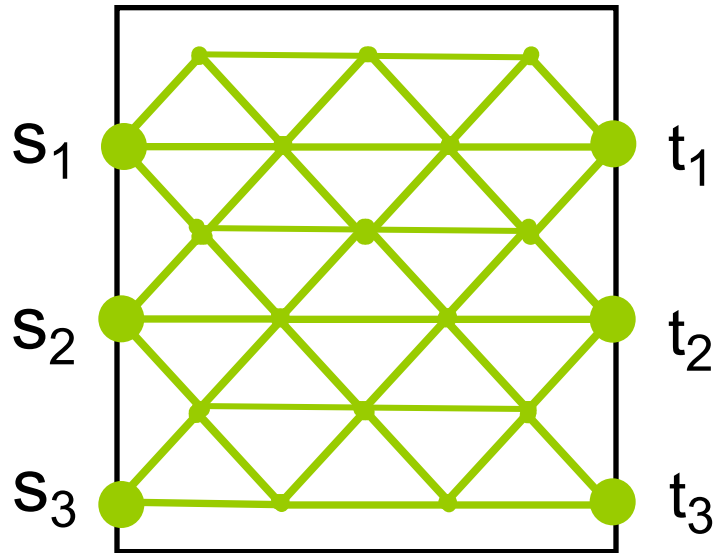
How many: An Algorithm



We want to count
non-intersecting sets of
paths from s_i to t_i .

[R.]

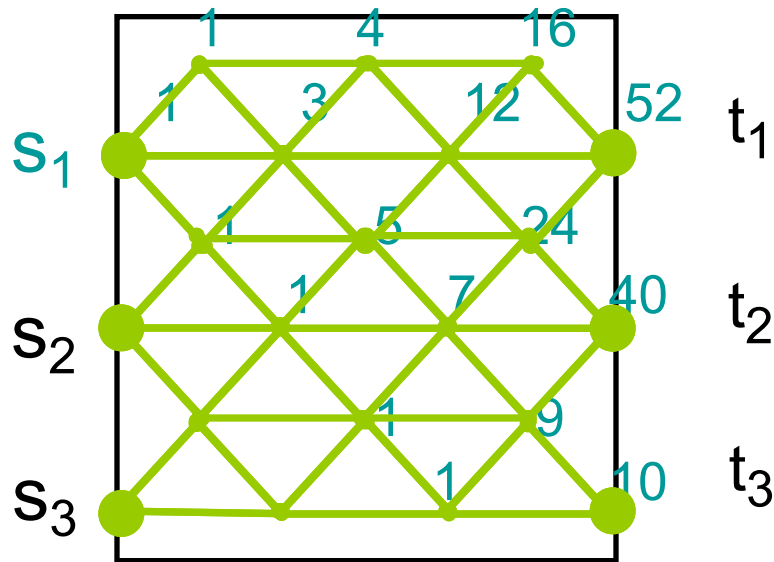
How many: An Algorithm



We want to count non-intersecting sets of paths from s_i to t_i .

Let a_{ij} be the number of paths from s_i to t_j .

How many: An Algorithm

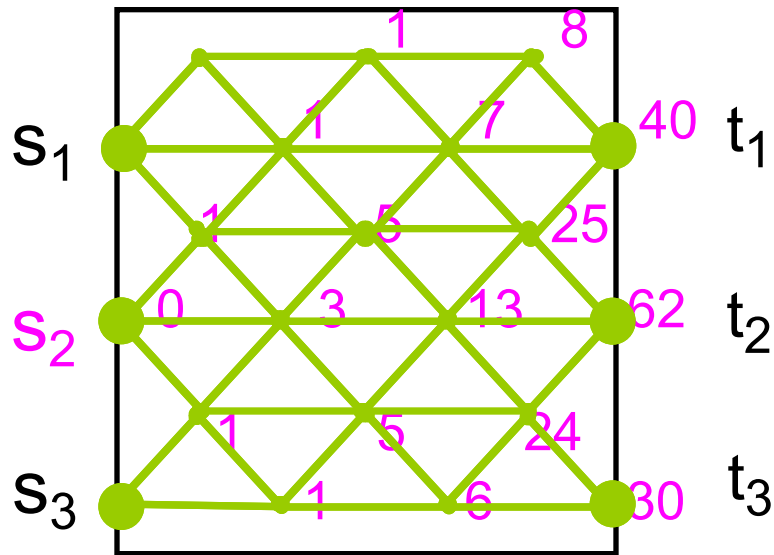


We want to count non-intersecting sets of paths from s_i to t_i .

Let a_{ij} be the number of paths from s_i to t_j .

52 40 10

How many: An Algorithm

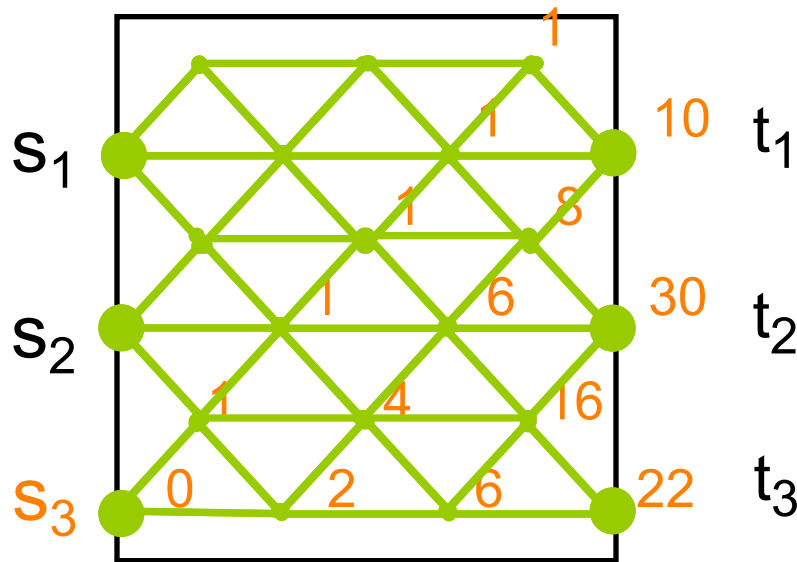


We want to count non-intersecting sets of paths from s_i to t_i .

Let a_{ij} be the number of paths from s_i to t_j .

52	40	10
40	62	30

How many: An Algorithm



We want to count non-intersecting sets of paths from s_i to t_i .

Let a_{ij} be the number of paths from s_i to t_j .

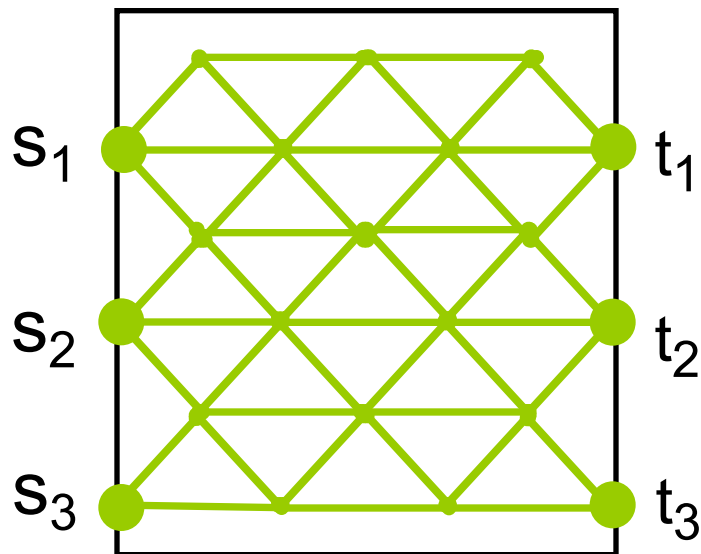
52 40 10

40 62 30

10 30 22

How many: An Algorithm

[Gessel, Viennot]



We want to count
non-intersecting sets of
paths from s_i to t_i .

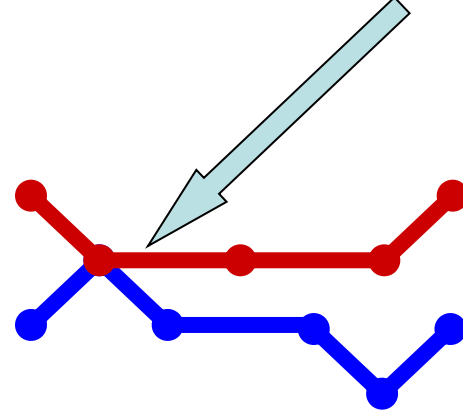
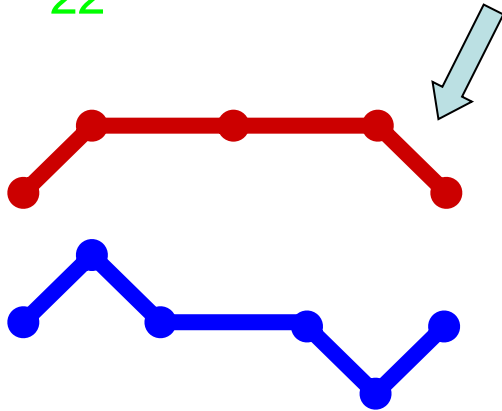
Let a_{ij} be the number of paths from s_i to t_j .

$$\text{Det} \begin{vmatrix} 52 & 40 & 10 \\ 40 & 62 & 30 \\ 10 & 30 & 22 \end{vmatrix} = 1728.$$

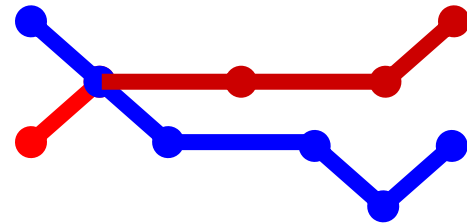
This is the number domino tilings!!

Proof sketch for two paths:

$a_{11} \times a_{22}$ counts what we want + extra stuff.



$a_{12} \times a_{21}$ also counts the extra stuff.



= #

Therefore $(a_{11} \times a_{22}) - (a_{12} \times a_{21})$ counts real tilings.

(This is the 2 x 2 determinant!)

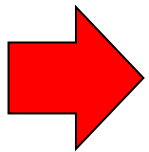
Domino Tilings

Where is a tiling? Do any even exist?

How many tilings are there?

What does a typical tiling look like?

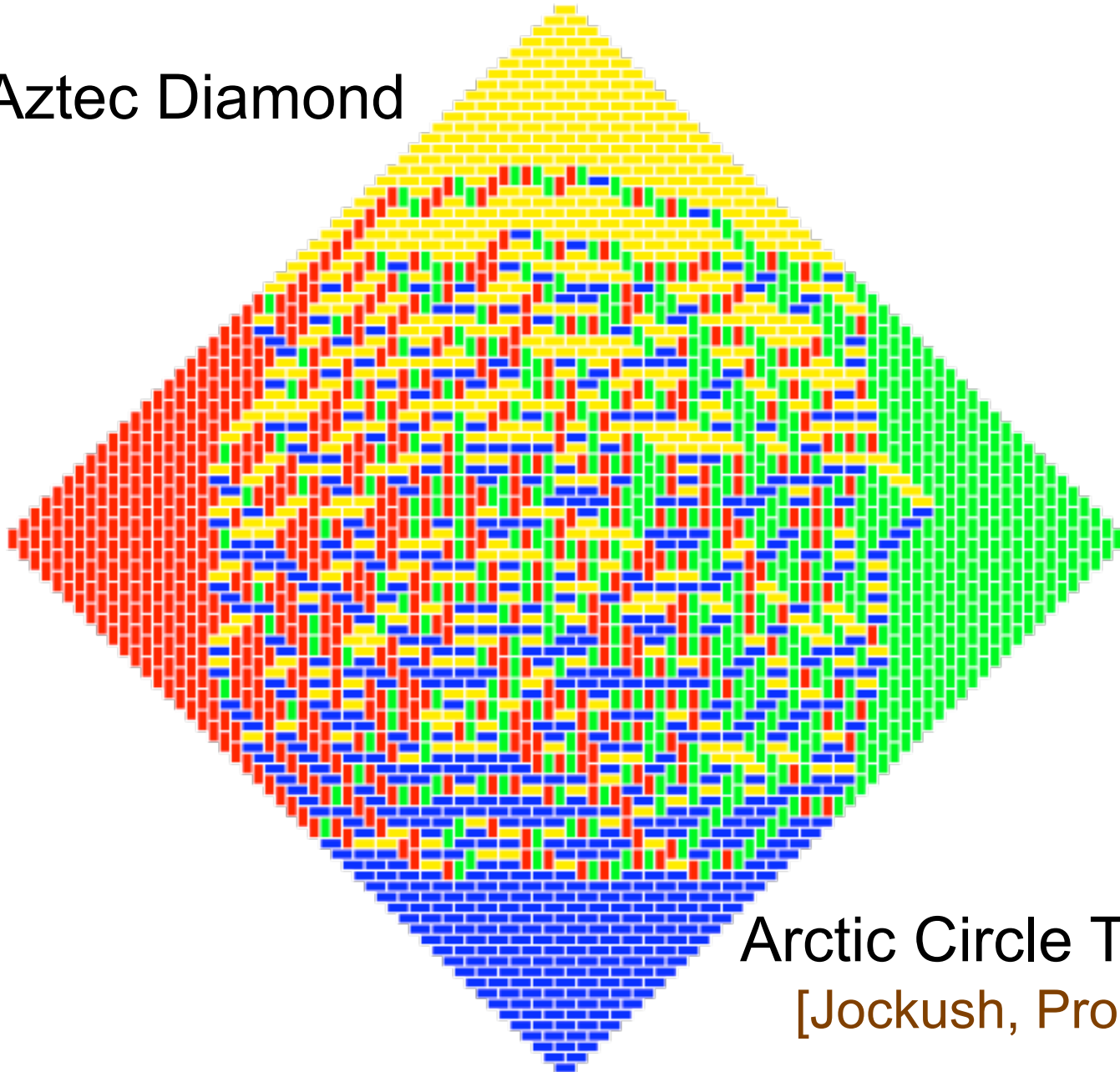
When do we stop our algorithms?



Why do we care?

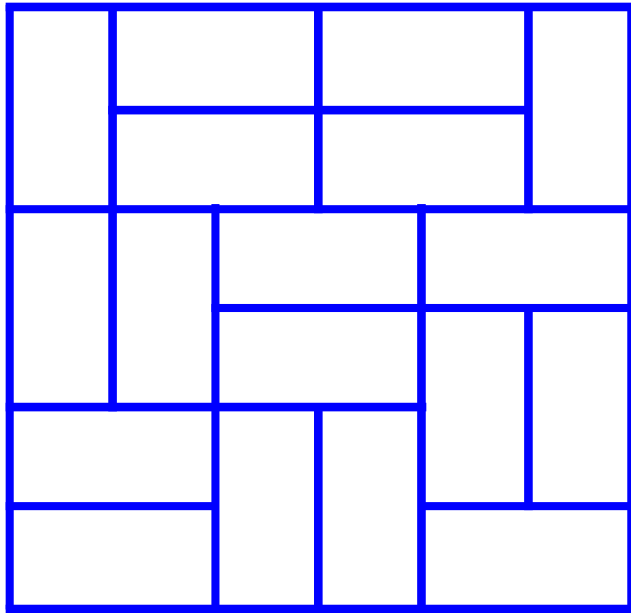
Why Mathematicians Care

The Aztec Diamond

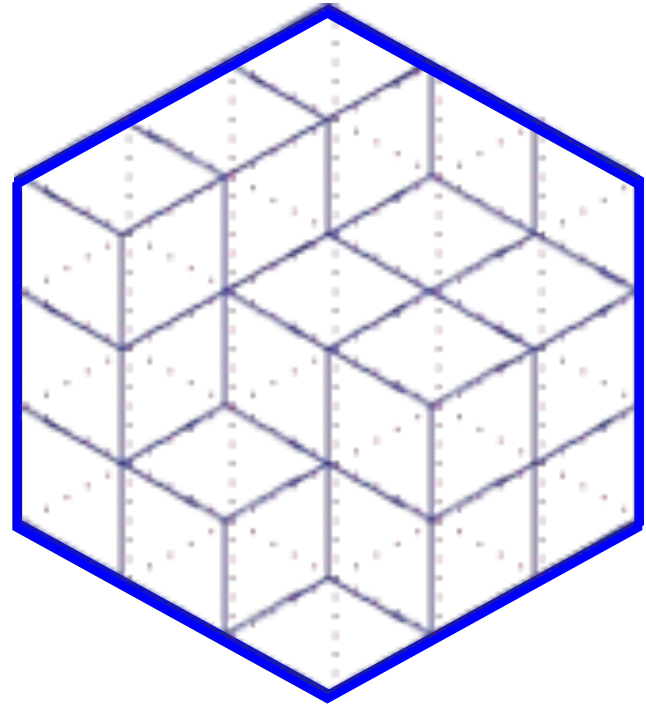


Arctic Circle Theorem
[Jockush, Propp, Shor]

What about tilings on lattices?

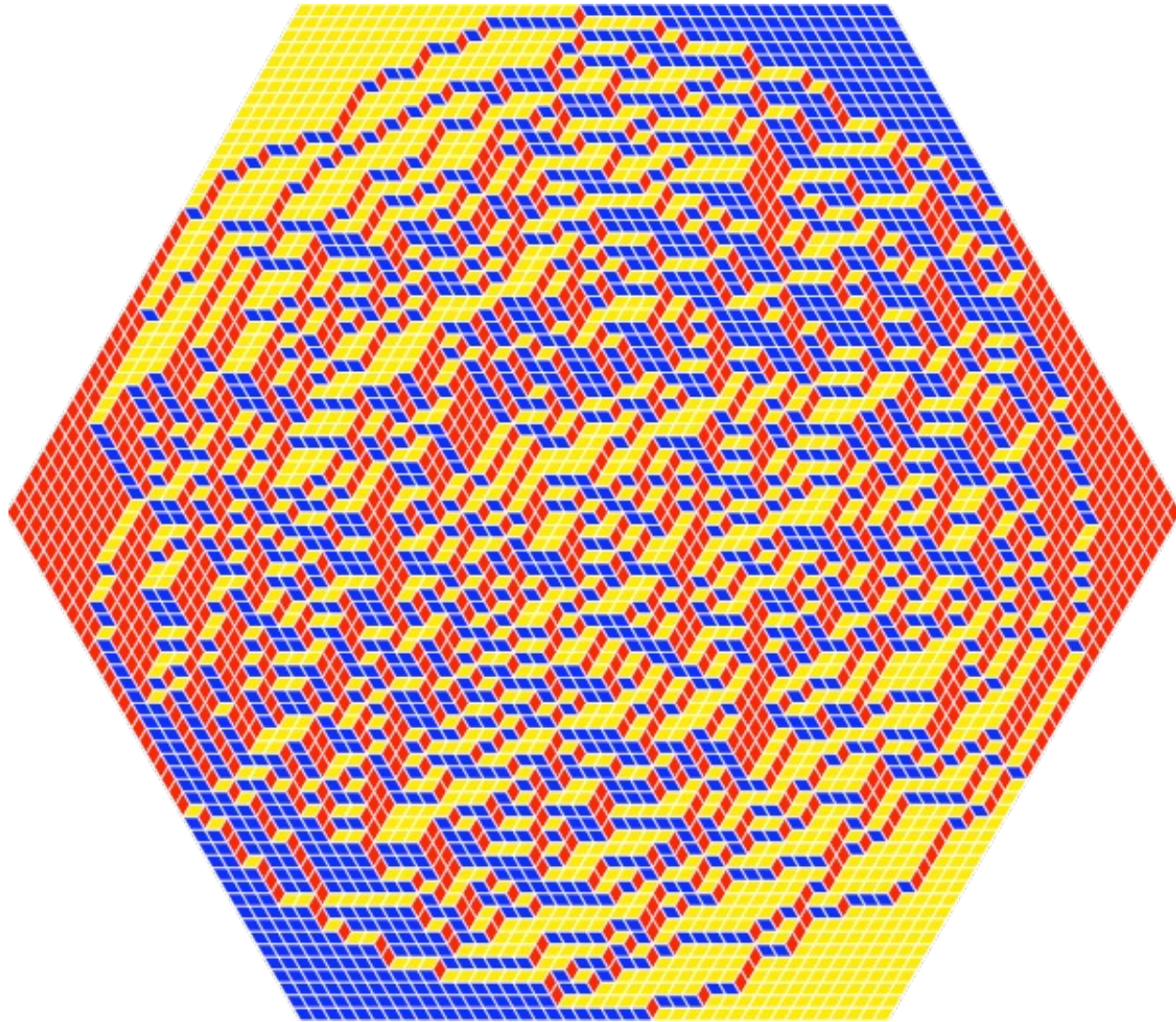


On the chessboard
“Domino tilings”

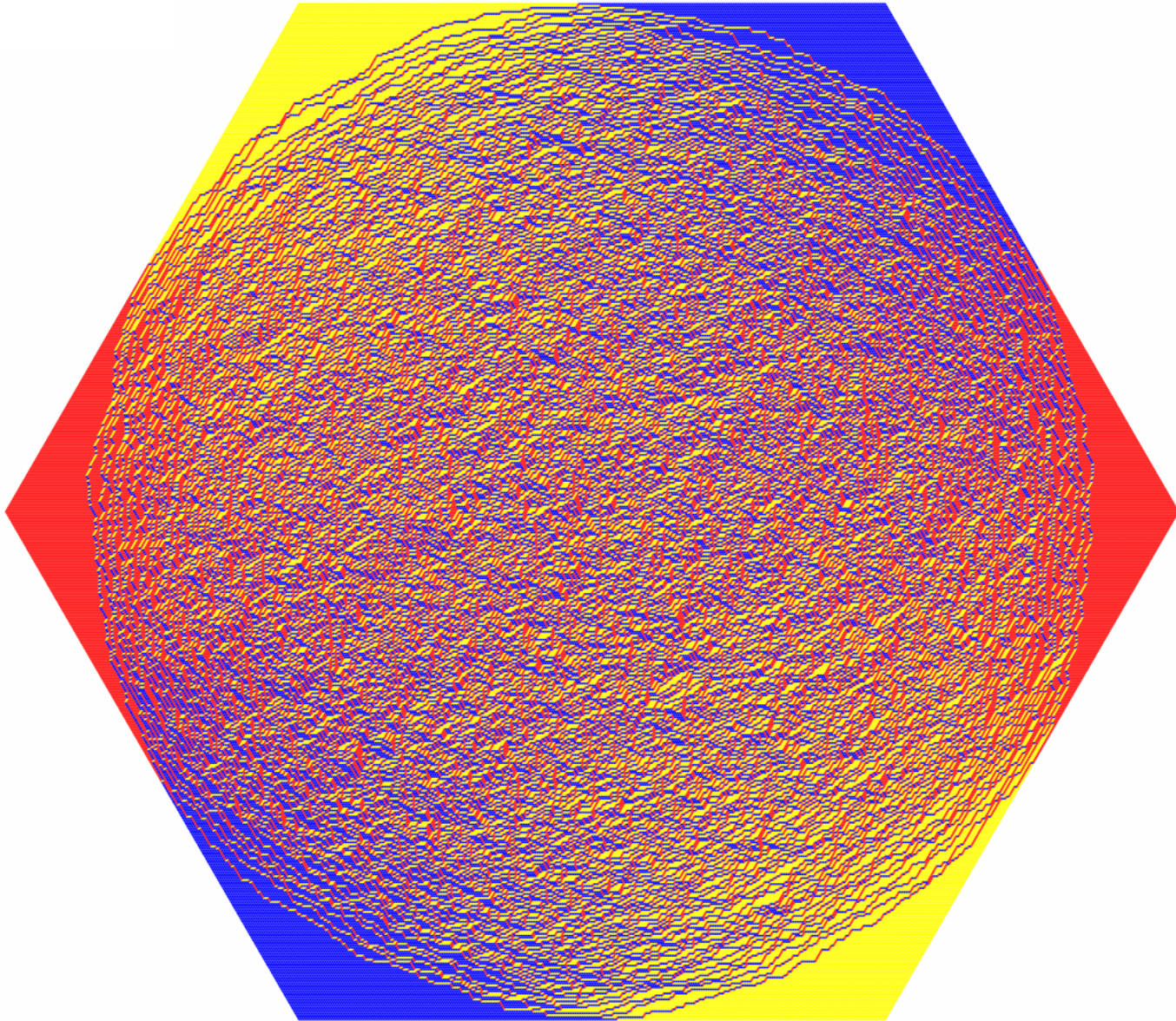


On the hexagonal lat.
“Lozenge tilings”
(little “cubes”)

Why Mathematicians Care



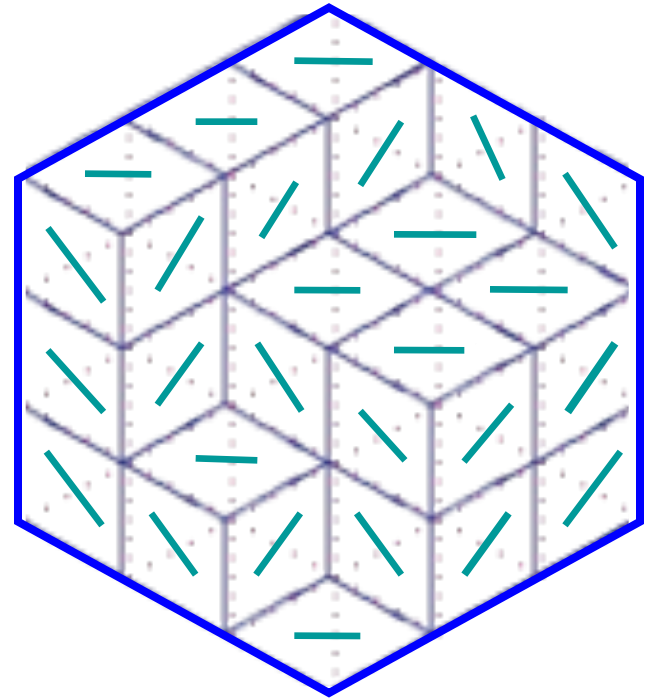
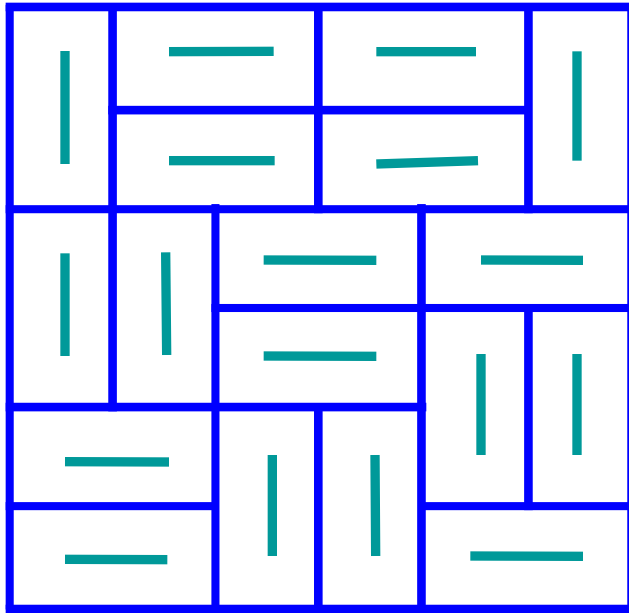
Why Mathematicians Care



Why do we care?

- **Mathematics**: Discover patterns
- **Chemistry, Biology**: Estimate probabilities
- **Physics**: Count and calculate other functions
to study a physical system
- **Nanotechnology**: Model growth processes

Why Physicists Care



“Dimer models”: diatomic molecules adhering to the surface of a crystal.

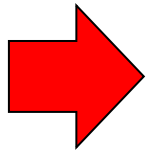
The count (“partition function”) determines:
specific heat, entropy, free energy, ...

What does “nature” compute?

Domino Tilings

Where is a tiling? Do any even exist?

How many tilings are there?

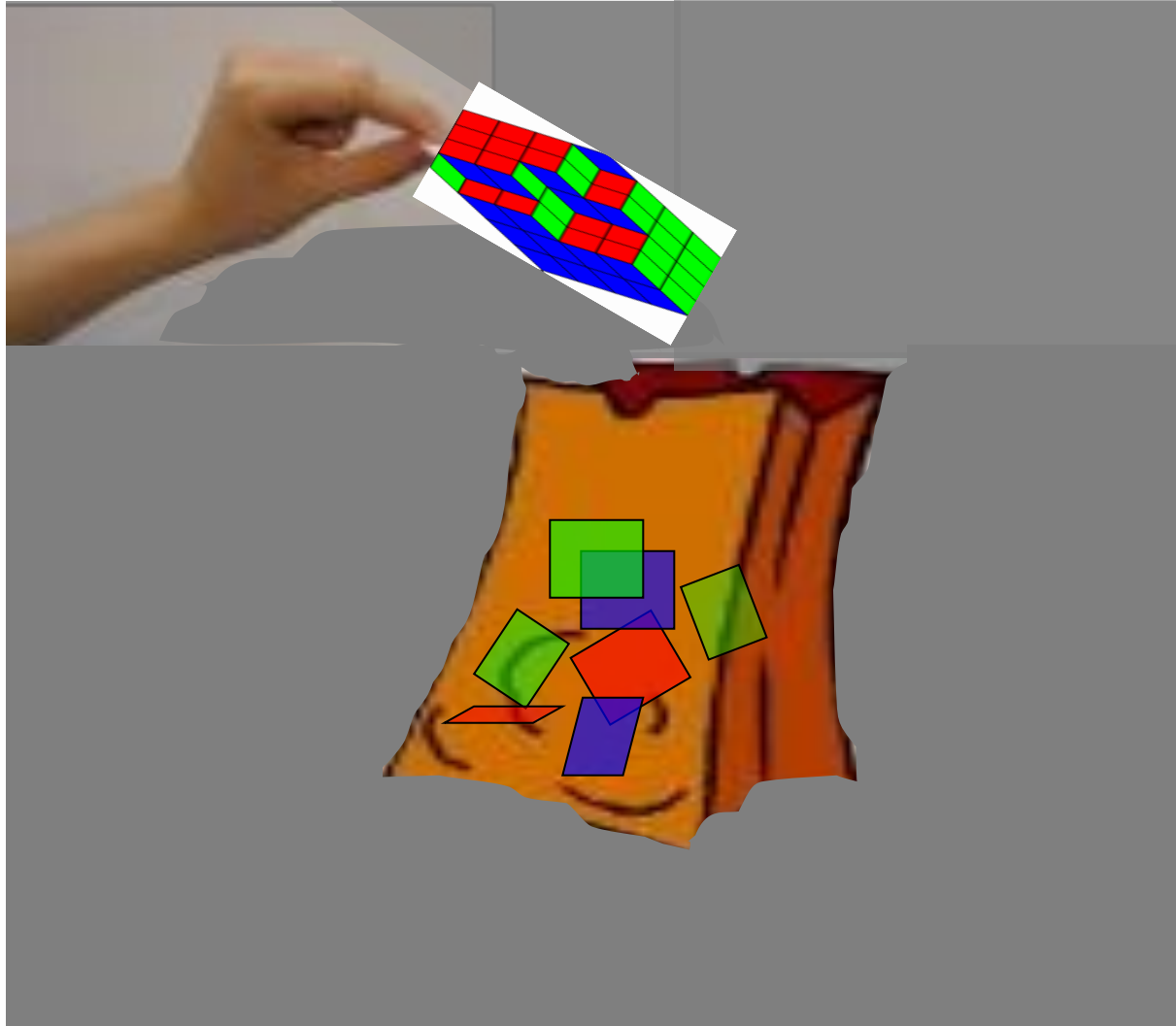


What does a typical tiling look like?

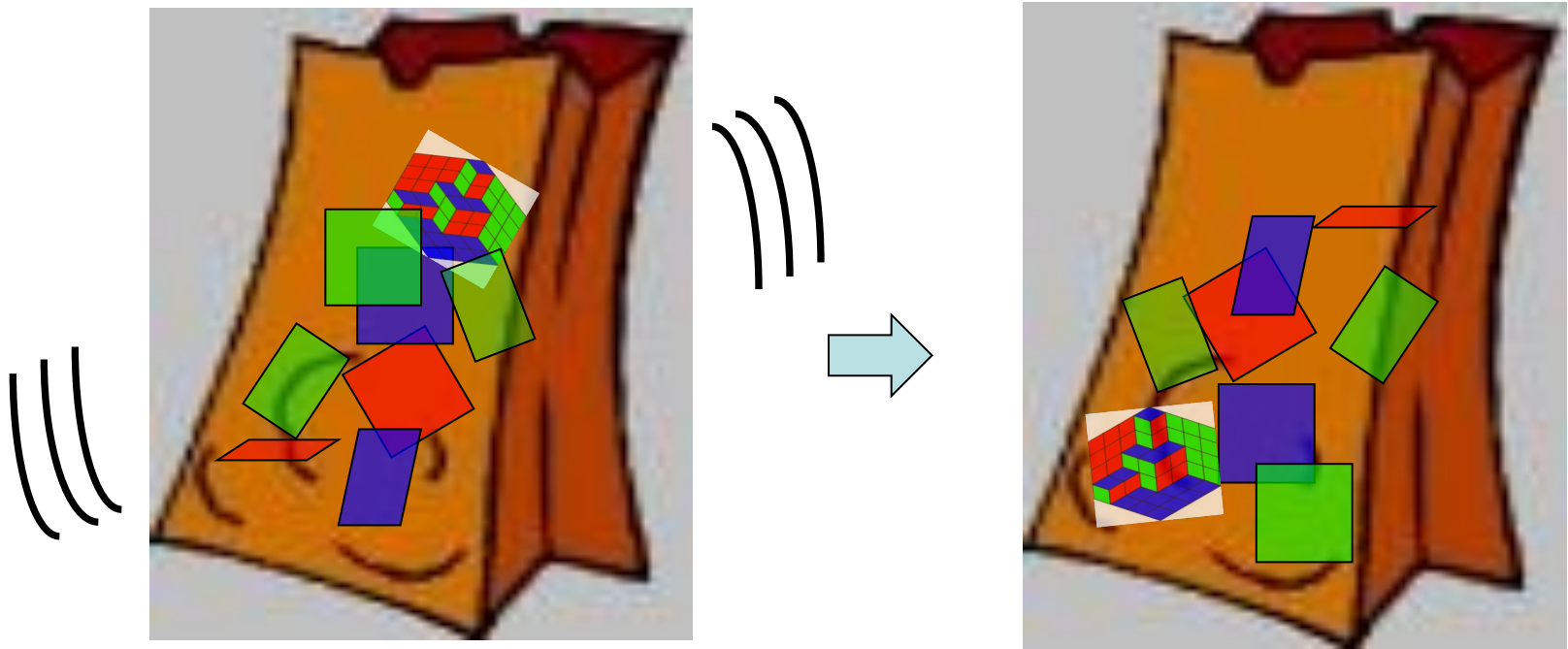
When do we stop our algorithms?

Why do we care?

What does a *typical* tiling look like?



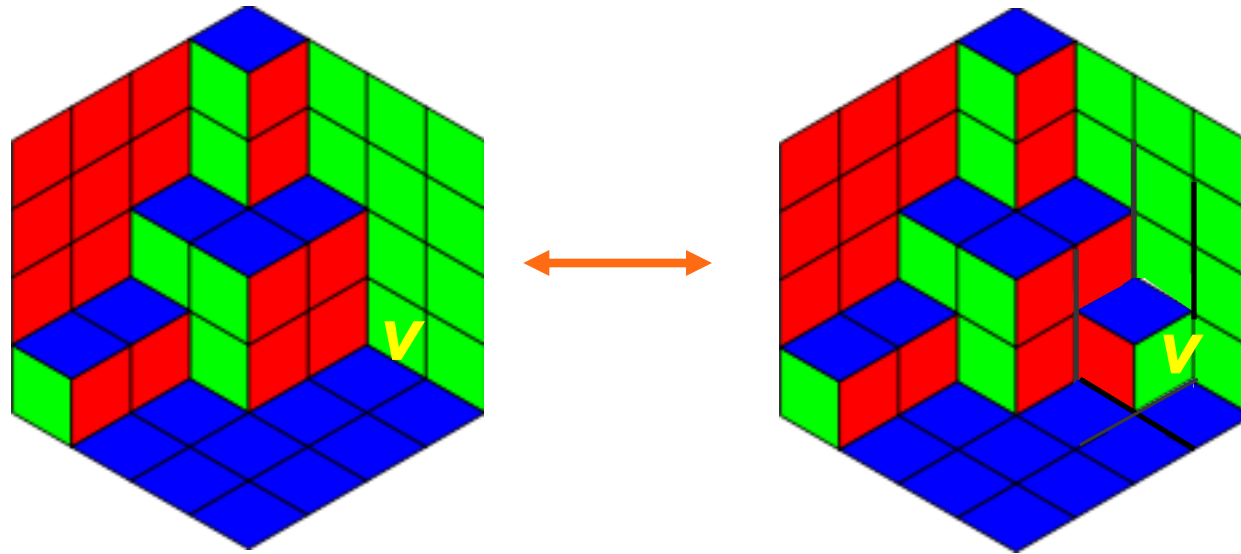
What does a *typical* tiling look like?



“Mix” them up!



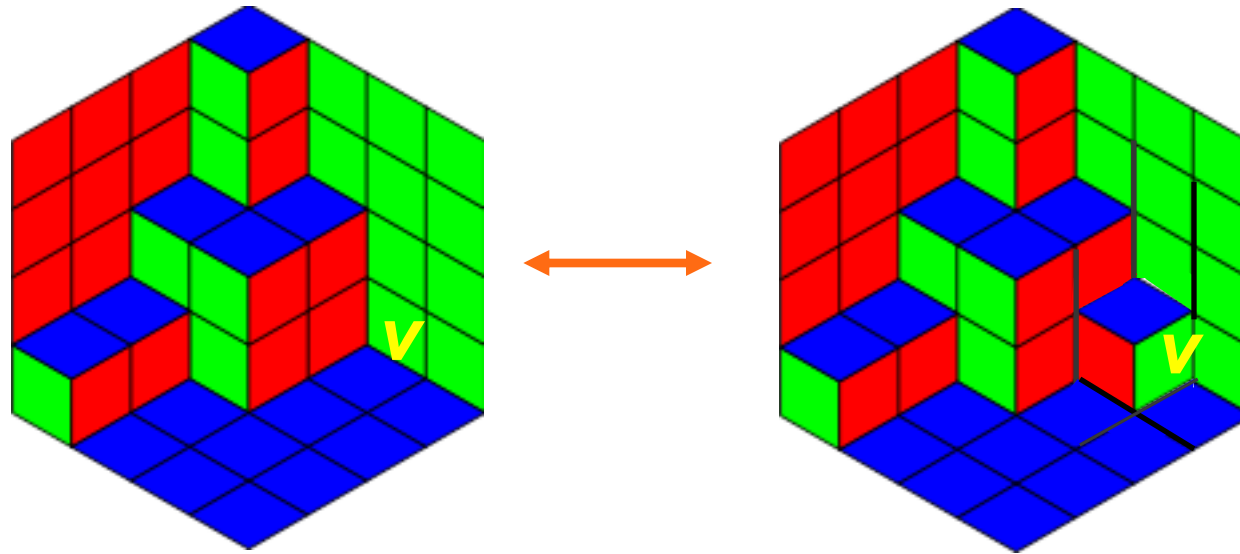
Markov chain for Lozenge Tilings



Repeat:

- Pick **v** in the lattice region;
- **Add / remove** the “cube”
at **v** w.p. $\frac{1}{2}$, if possible.

Markov chain for Lozenge Tilings



1. The state space is connected.
2. If we do this long enough, each tiling will be equally likely.
3. *How long* is “long enough” ?

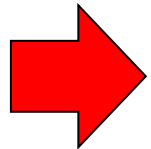


Domino Tilings

Where is a tiling? Do any even exist?

How many tilings are there?

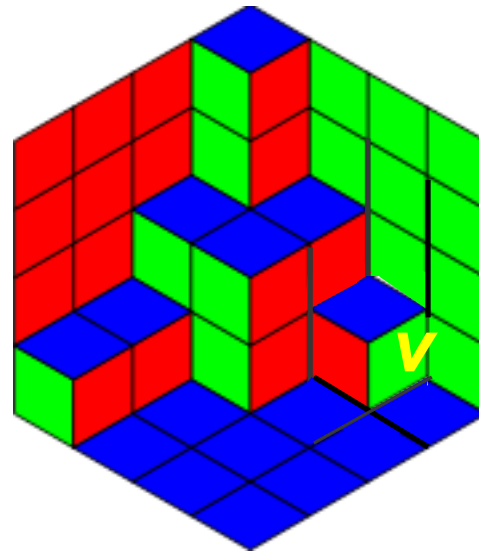
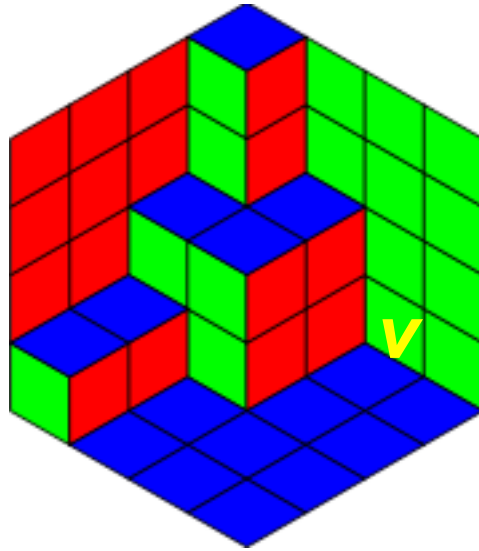
What does a typical tiling look like?



When do we stop our algorithms?

Why do we care?

When do we stop our algorithms?



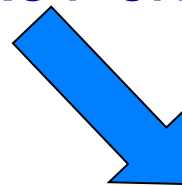
3. *How long* is “long enough” ?

Thm: The lozenge Markov chain is “rapidly mixing.”

[Luby, R., Sinclair]



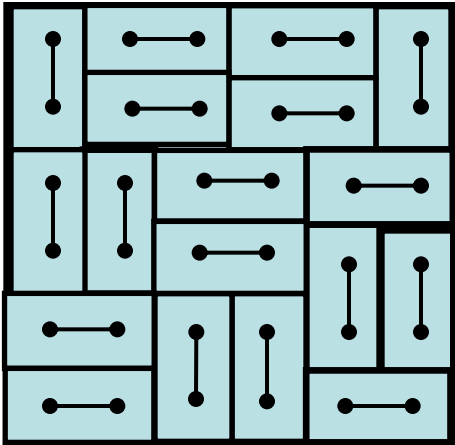
$2^n, 3^n, n^n, 10^{\sqrt{n}}, \dots$
(exponential)



$n^2, n \log n, n^{10}, \dots$
(polynomial)

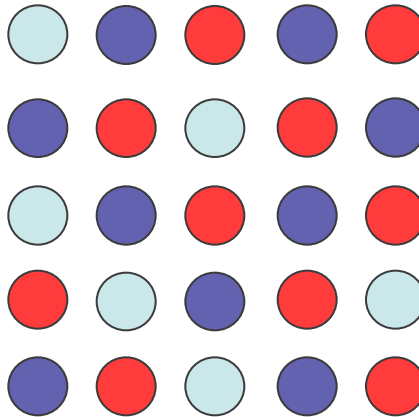


What about other models?



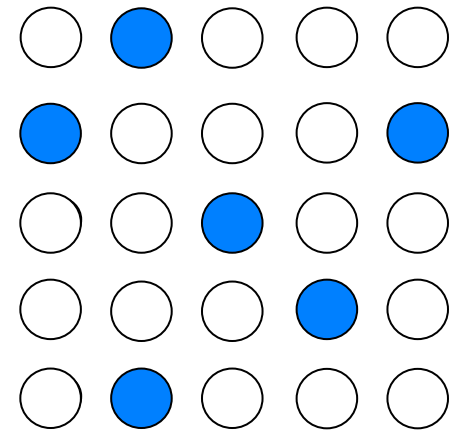
Dimer model

Domino tilings



Potts model

k-colorings

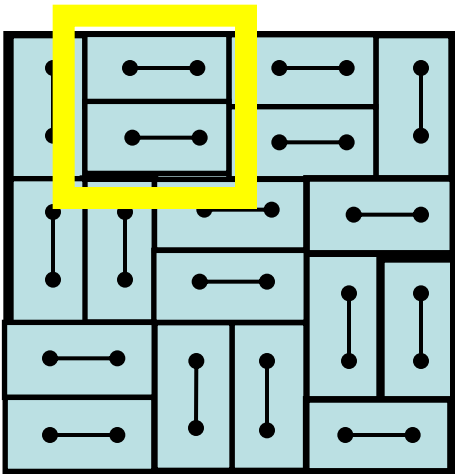


Hardcore model

Independent sets

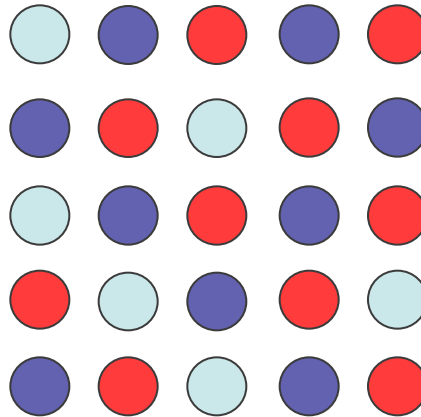
What about other models?

- Pick a 2×2 square;



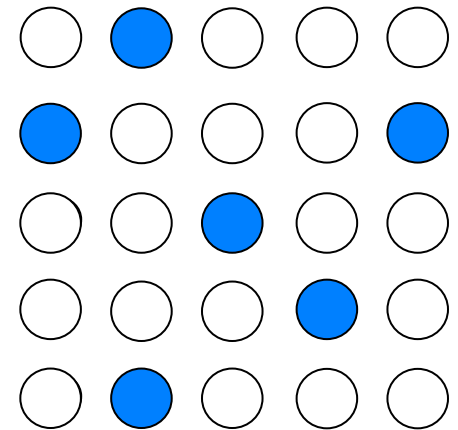
Dimer model

Domino tilings



Potts model

k-colorings

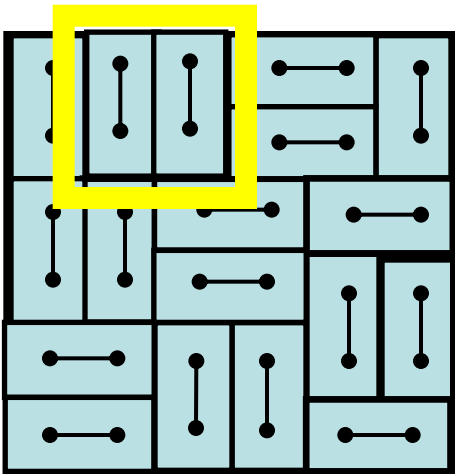


Hardcore model

Independent sets

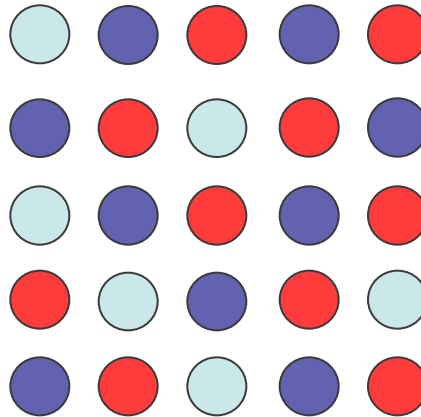
What about other models?

- Pick a 2×2 square;
- Rotate, if possible;



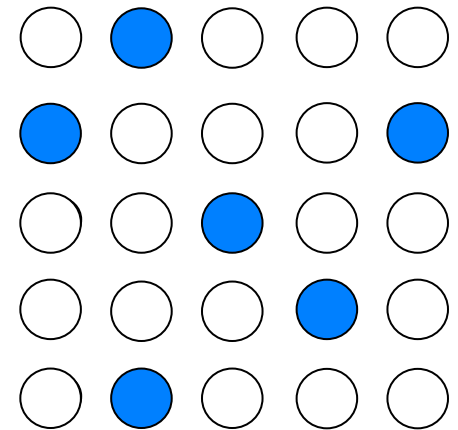
Dimer model

Domino tilings



Potts model

k-colorings

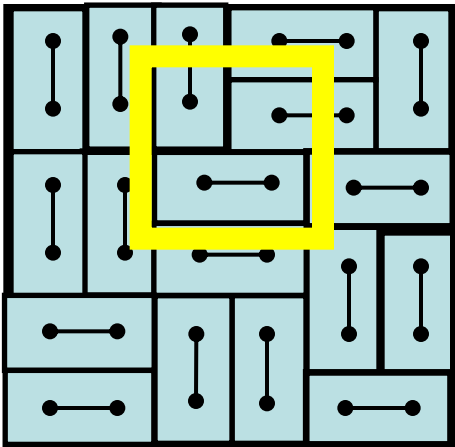


Hardcore model

Independent sets

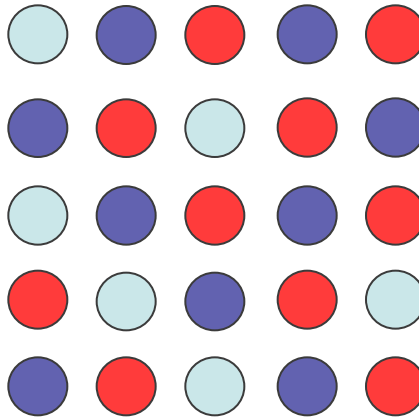
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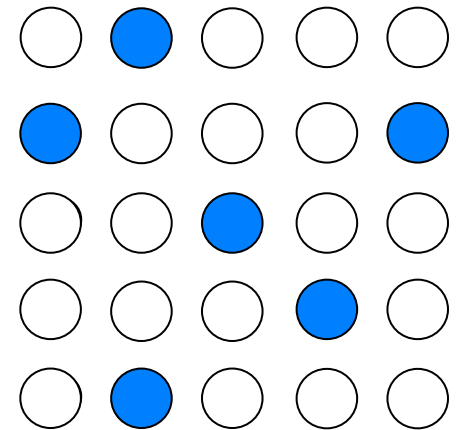
Dimer model

Domino tilings



Potts model

k-colorings

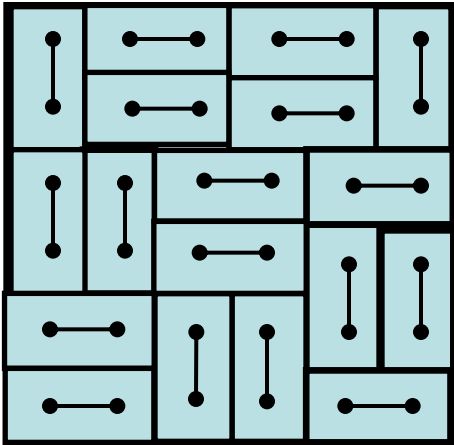


Hardcore model

Independent sets

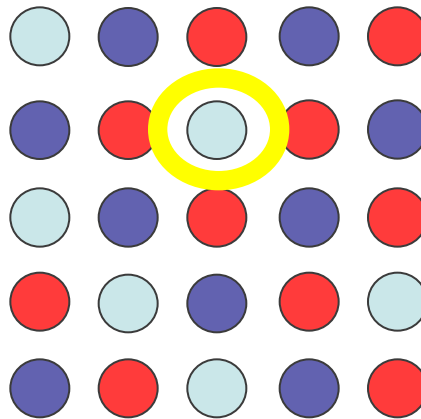
What about other models?

- Pick a vtx and a color; 



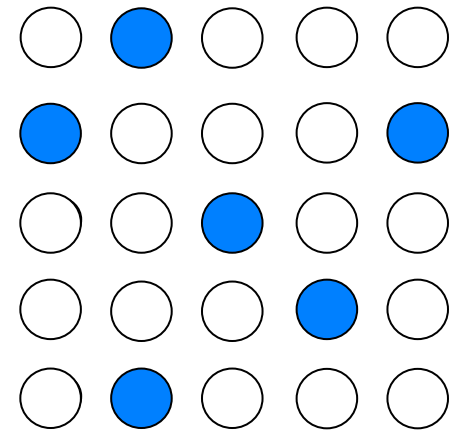
Dimer model

Domino tilings



Potts model

k-colorings

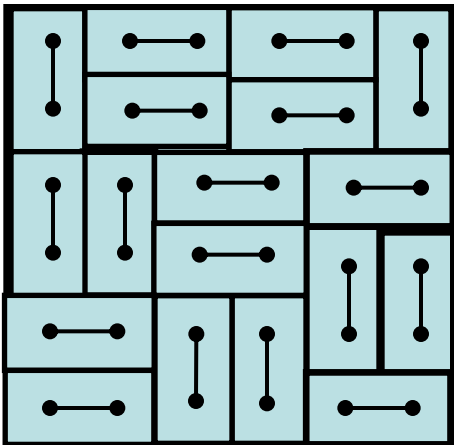


Hardcore model

Independent sets

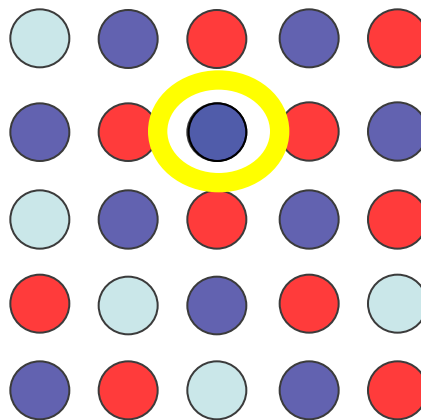
What about other models?

- Pick a vtx and a color; 
- Recolor, if possible;



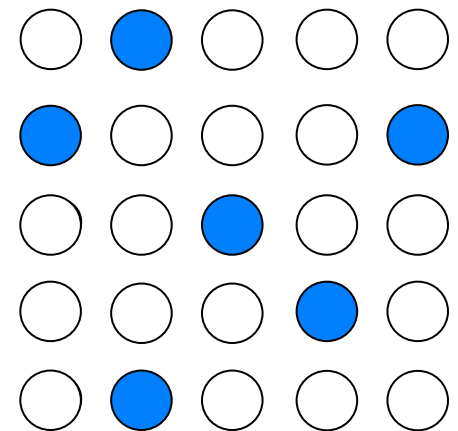
Dimer model

Domino tilings



Potts model


k-colorings

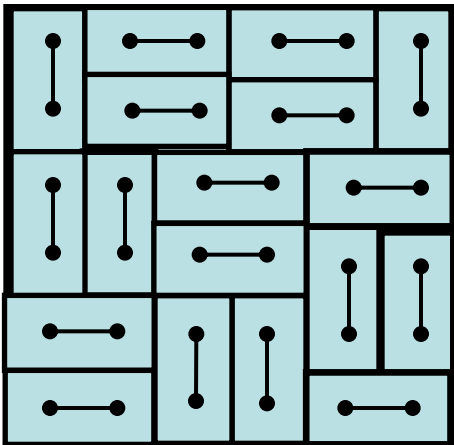


Hardcore model

Independent sets

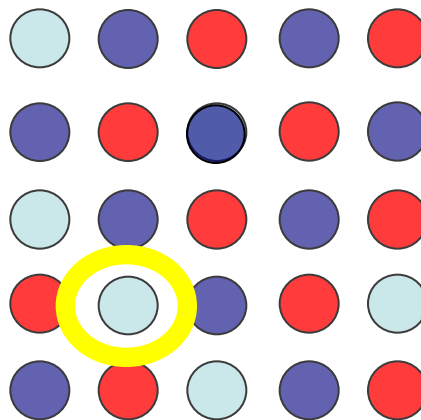
What about other models?

- Pick a vtx and a color; 
- Recolor, if possible;
- Otherwise do nothing.



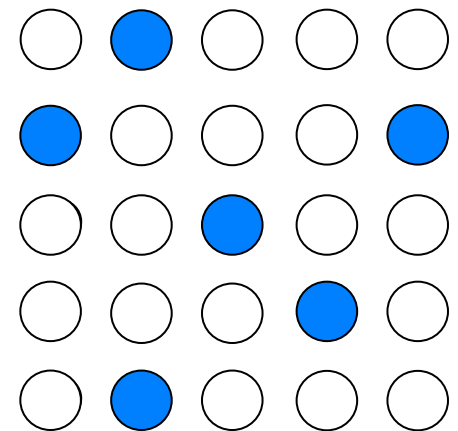
Dimer model

Domino tilings



Potts model

k-colorings

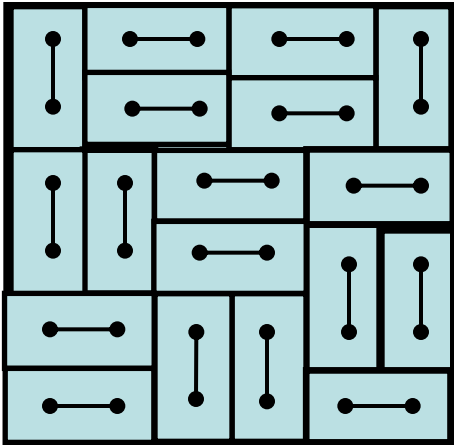


Hardcore model

Independent sets

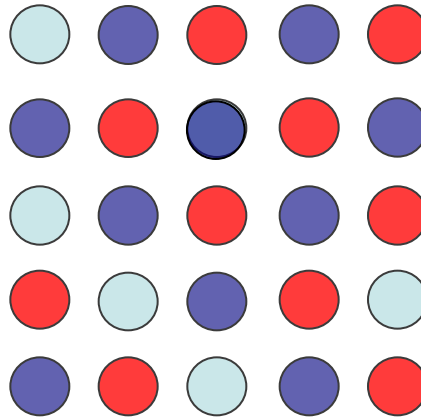
What about other models?

- Pick a vtx v and a bit b ;



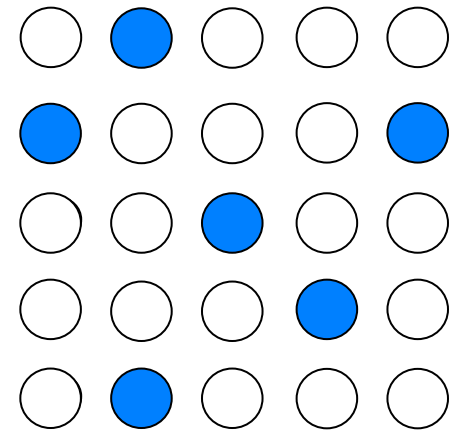
Dimer model

Domino tilings



Potts model

k-colorings

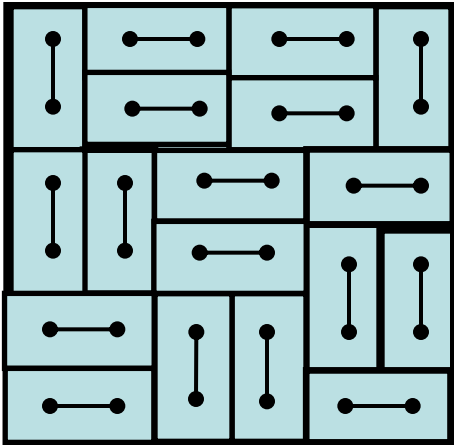


Hardcore model

Independent sets

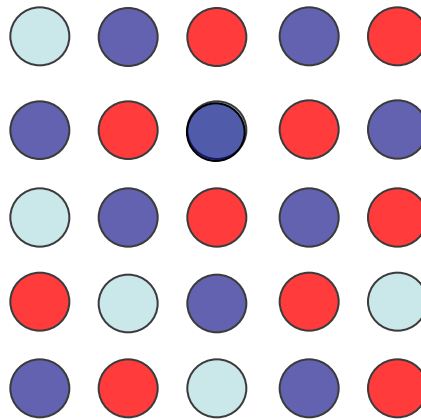
What about other models?

- Pick a vtx v and a bit b ;
- If $b=1$, try to add v



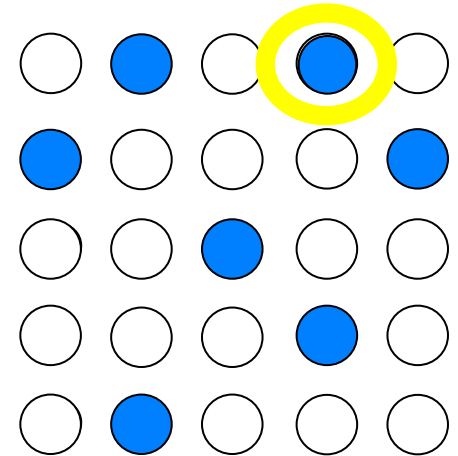
Dimer model

Domino tilings



Potts model

k-colorings

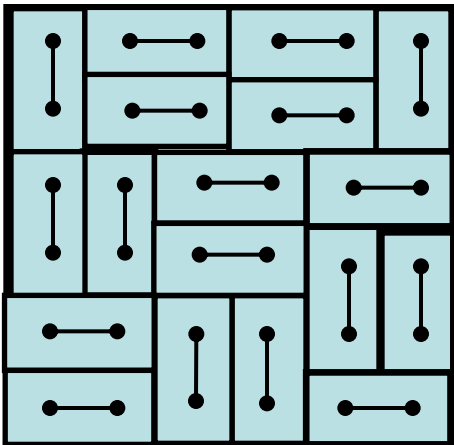


Hardcore model

Independent sets

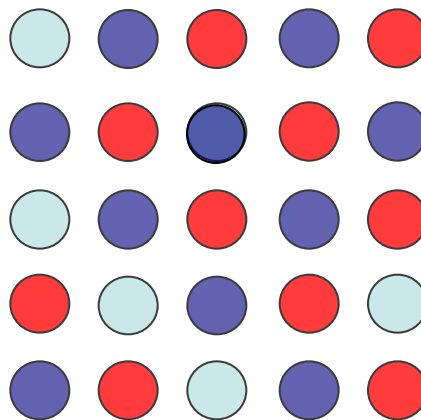
What about other models?

- Pick a vtx v and a bit b ;
- If $b=1$, try to **add** v ;
- If $b=0$, try to **remove** v ;



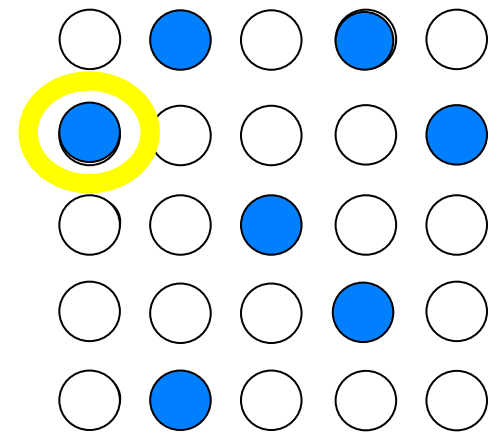
Dimer model

Domino tilings



Potts model

k -colorings

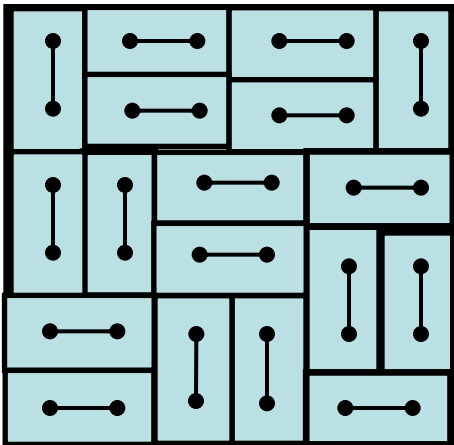


Hardcore model

Independent sets

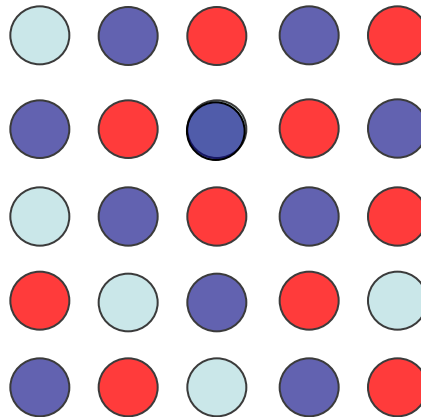
What about other models?

- Pick a vtx v and a bit b ;
- If $b=1$, try to **add** v ;
- If $b=0$, try to **remove** v ;
- O.w. do nothing.



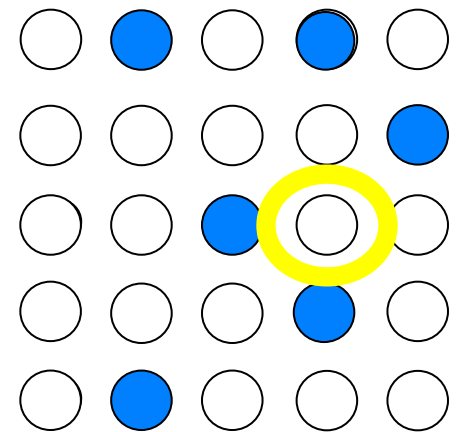
Dimer model

Domino tilings



Potts model

k-colorings

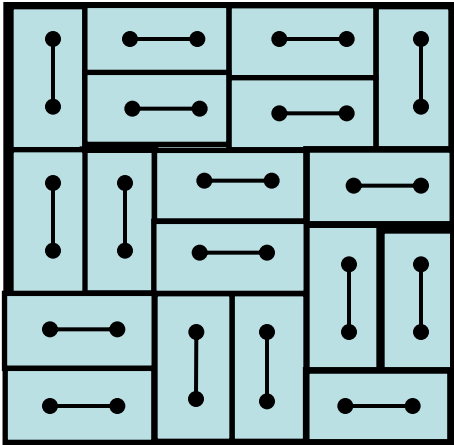


Hardcore model

Independent sets

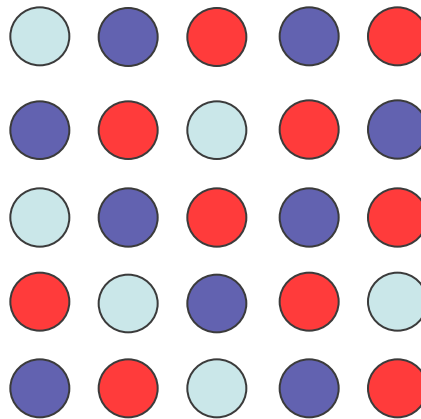
What about other models?

Thm: **All** of these chains are rapidly mixing.



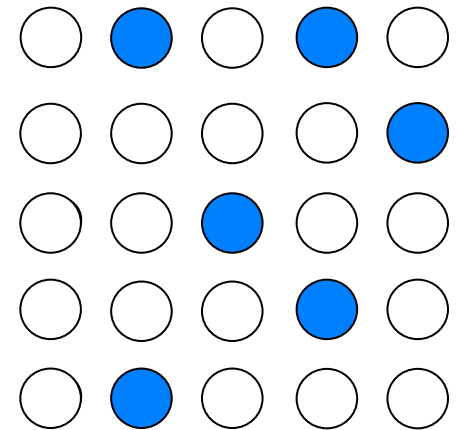
Dimer model

Domino tilings



Potts model

k-colorings



Hardcore model

Independent sets

HOWEVER . . .

Three-colorings:

The local chain is **fast** for 3-colorings in **2-d** [LRS]

but **slow** for 3-colorings in sufficiently **high dimension**.

[Galvin, Kahn, R, Sorkin], [Galvin, R]

Independent Sets

The local chain is **fast** for **sparse** Ind Sets in 2-d

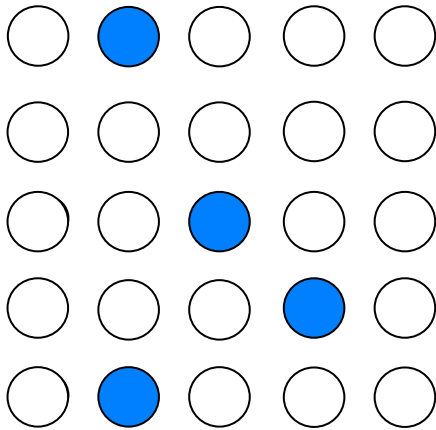
[Luby, Vigoda], ..., [Weitz]

but **slow** for **dense** Ind Sets.

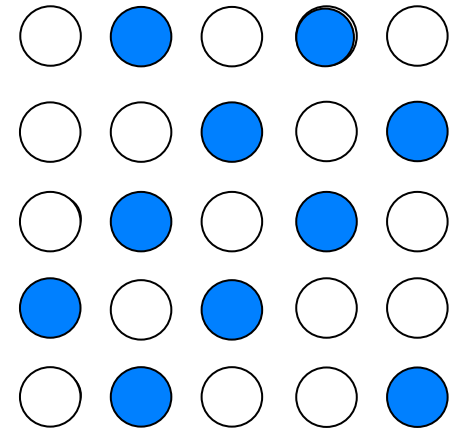
[R.]

Weighted Independent Sets

Sparse



Dense



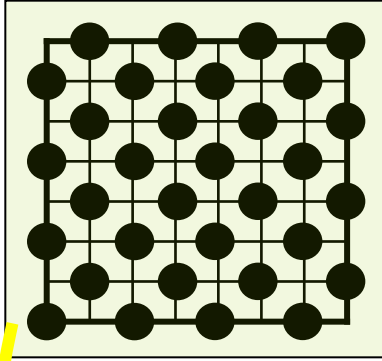
Fast

Phase Transition

Slow

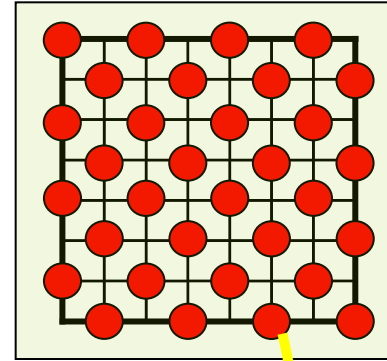
Why?

“Even”

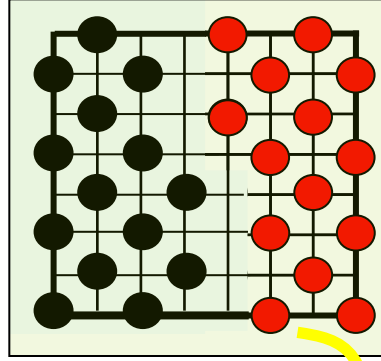


$$\lambda^{n^2/2}$$

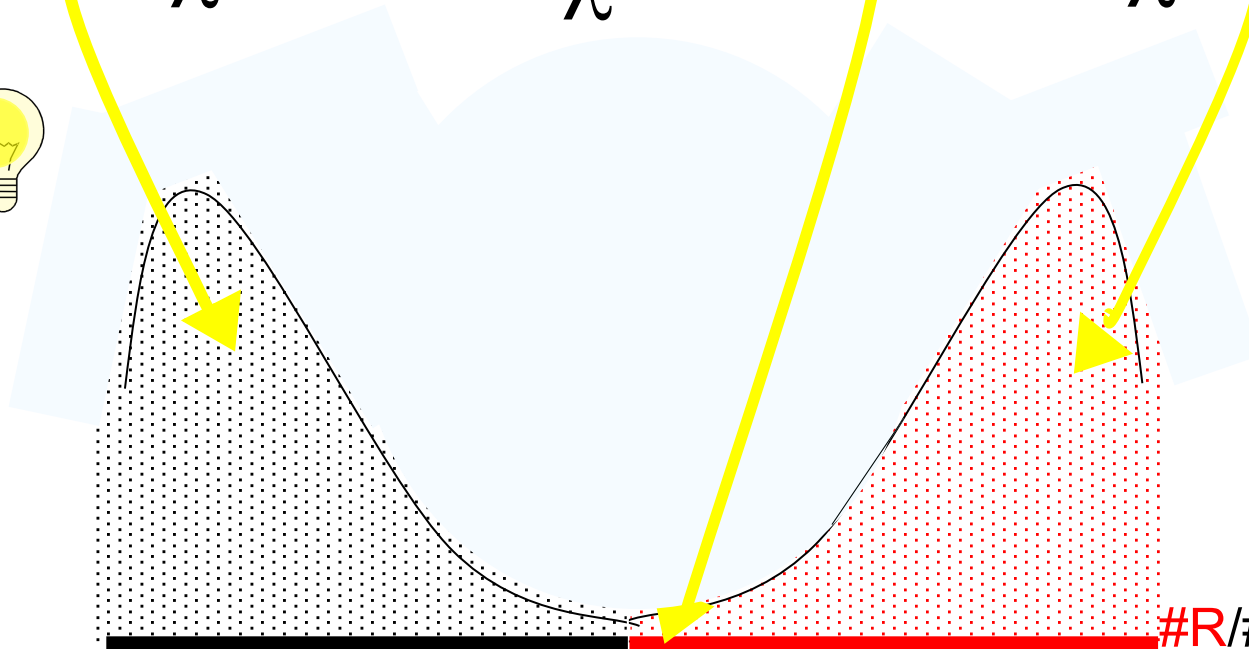
“Odd”



$$\lambda^{n^2/2}$$



$$\lambda^{(n^2/2 - n/2)}$$



#R/#B

Summary

Where

How

What

When

Why

Summary



Where

How

What

When

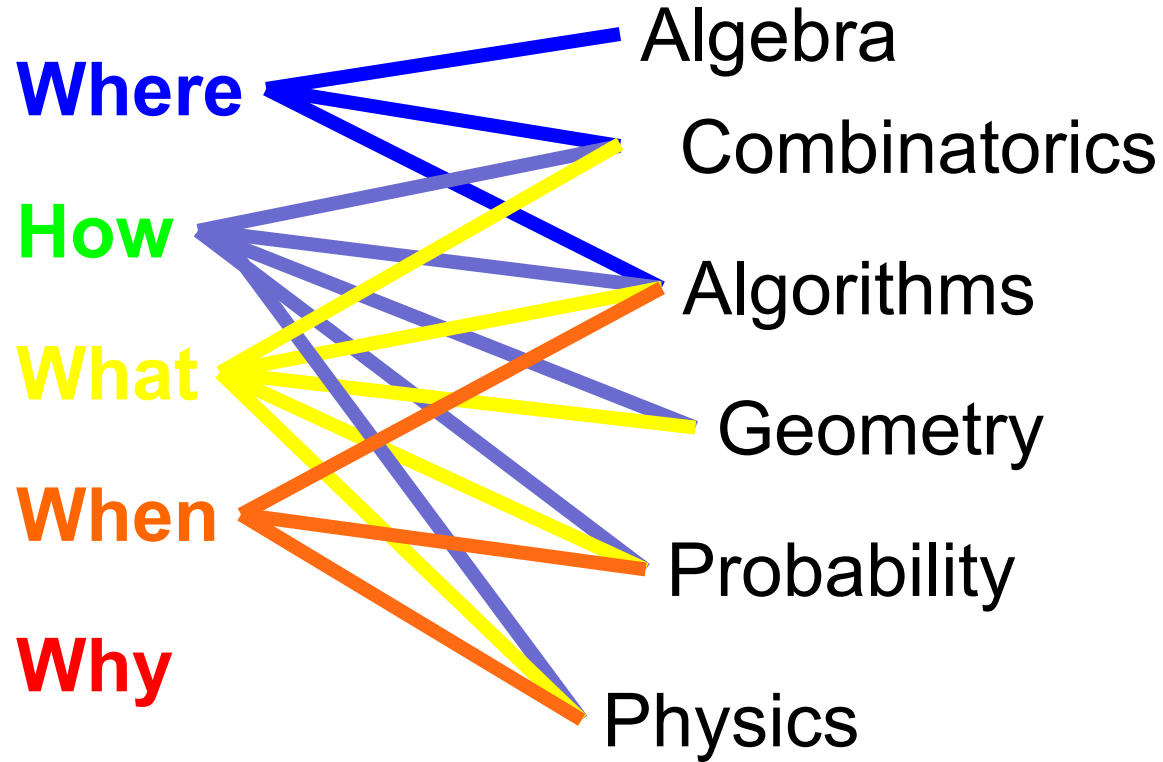
Why

Algebra

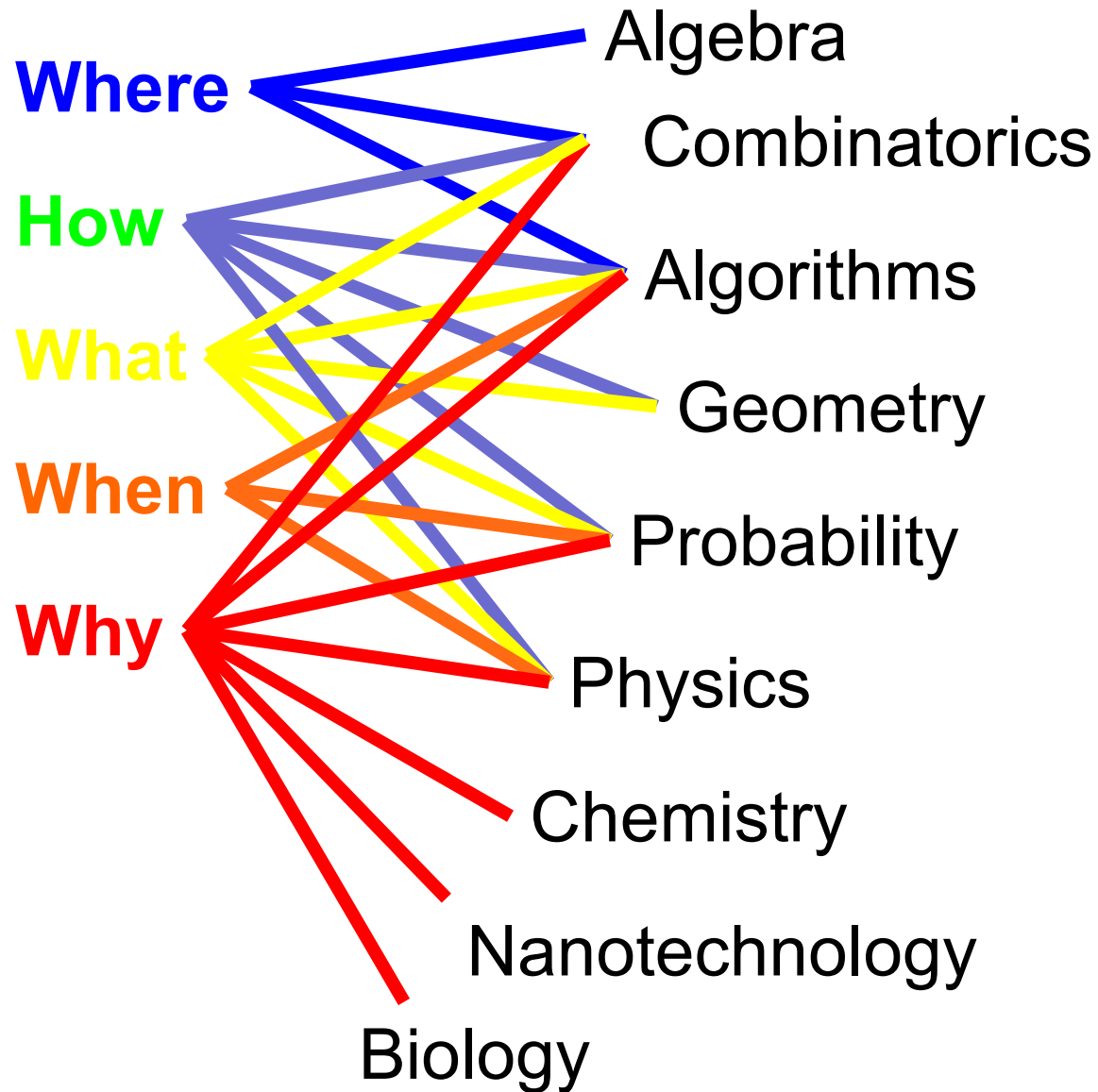
Combinatorics

Algorithms

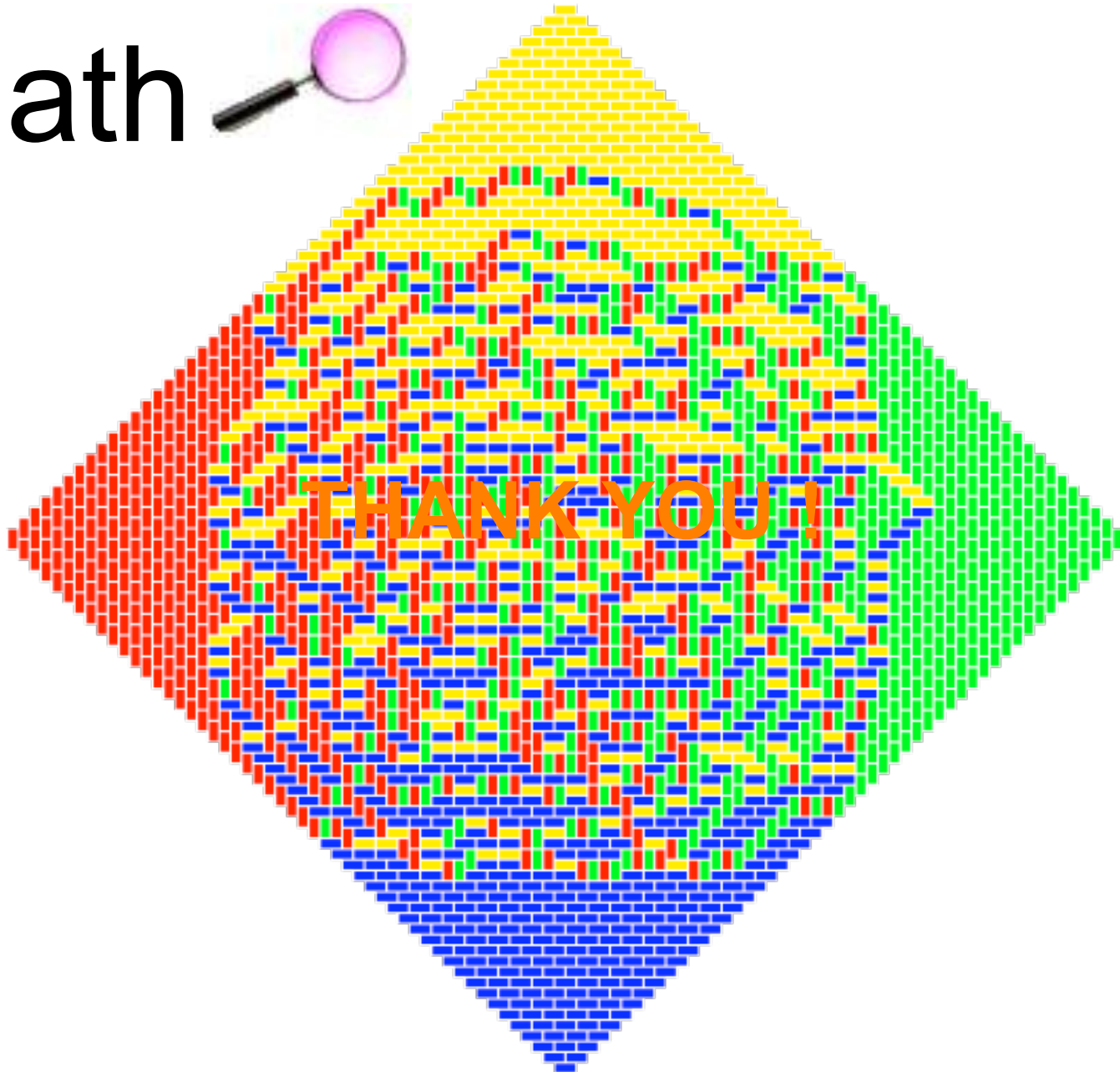
Summary



Summary



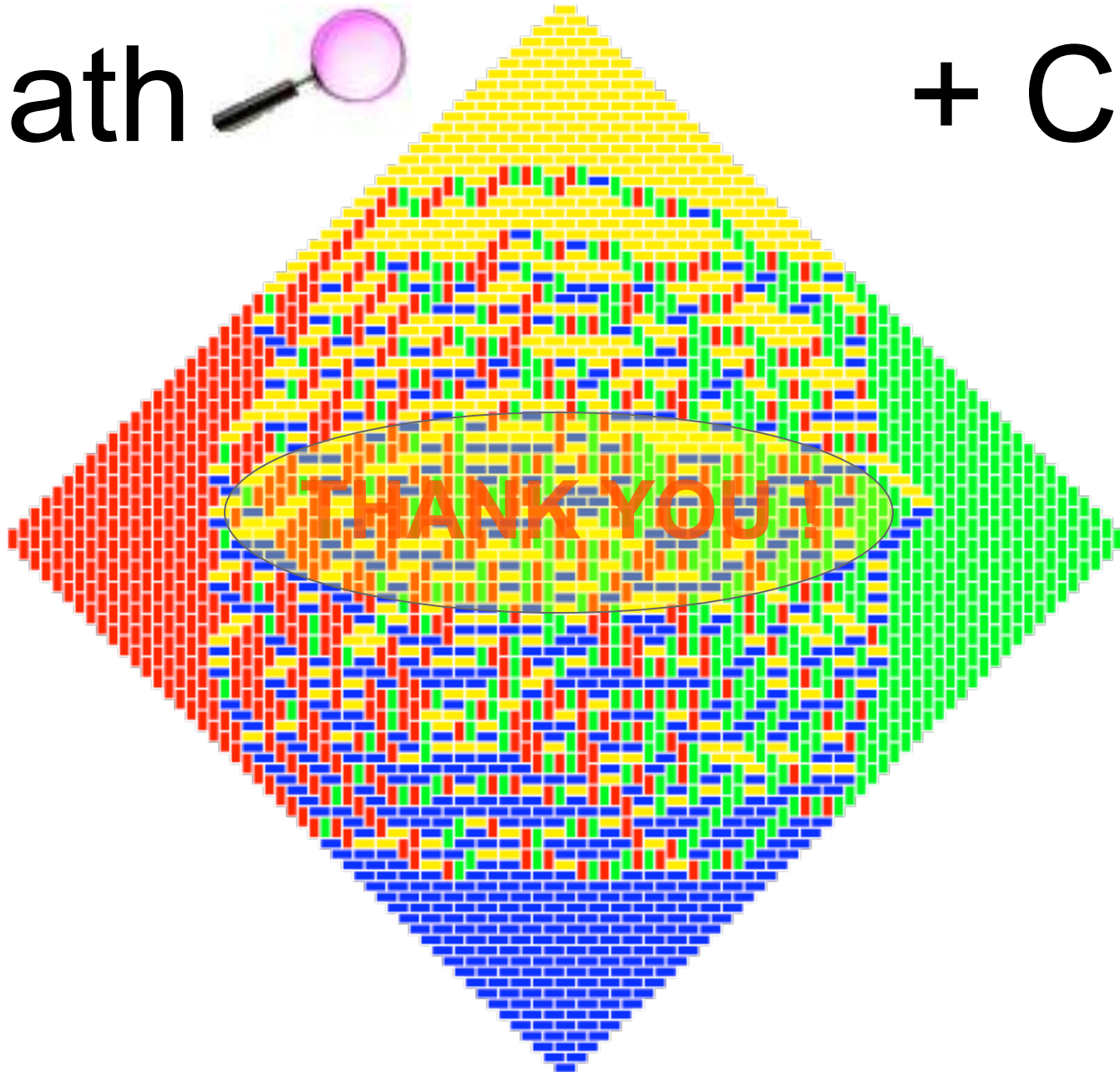
Math



Math



+ CS



Math



+ CS

+

Biology,
Chemistry,

THANK YOU !

Math



+ CS

THANK YOU !

+

Biology,
Chemistry,

+

Physics,
Nanotechnology,

...



