

#### Goals

- The goal is to understand:
  - Merits and limits of parallel computing
  - Parallel programming models (task / data parallelism)
  - Differences between shared and distributed memory systems



#### **Content**

What is parallel computing?

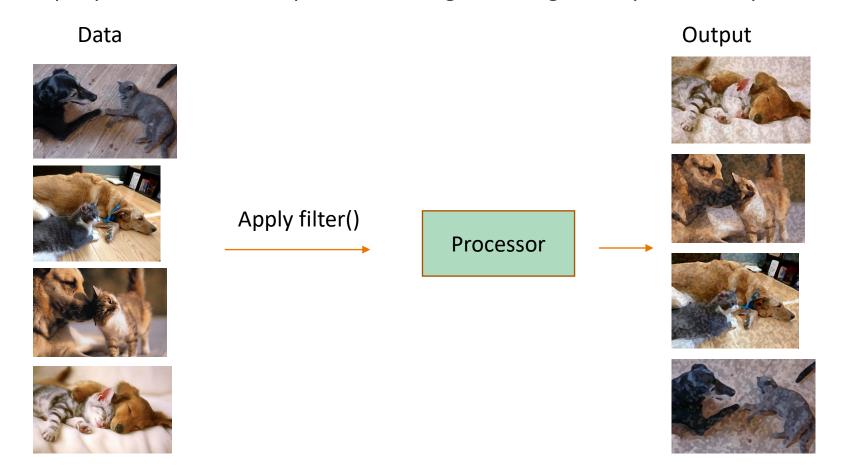
Parallel programming models

Different types of systems



# What is parallel computing?

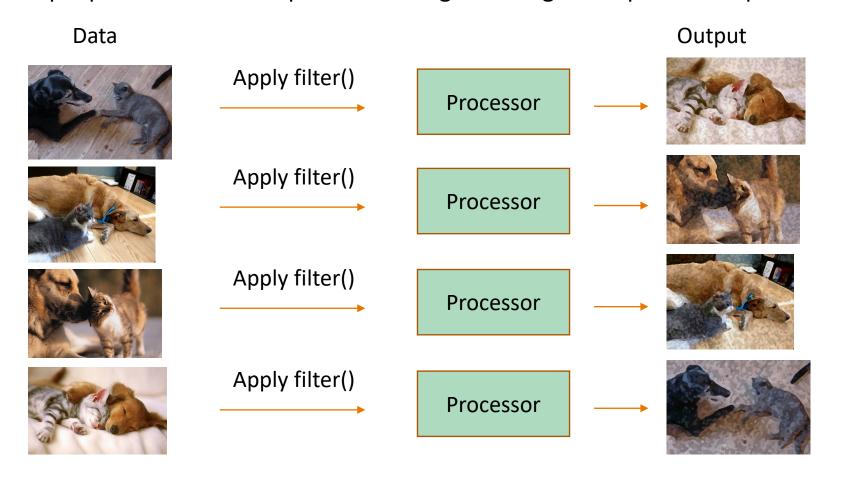
- Parallel computing
  - Multiple processors or computers working on a single computational problem





## What is parallel computing?

- Parallel computing
  - Multiple processors or computers working on a single computational problem



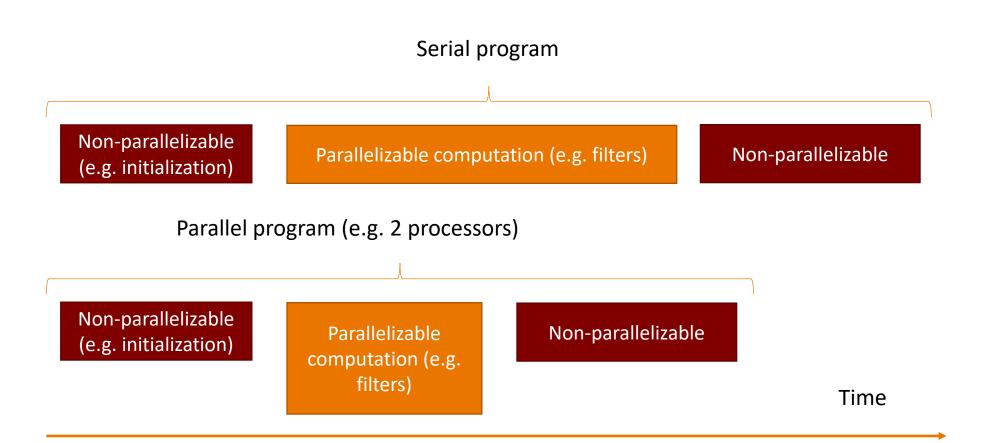


## Parallel computing: benefits

- Benefits
  - Solve computationally intensive problems (speedup)
  - Solve problems that don't fit a single memory (multiple computers)
- Requirements:
  - Problem should be divisible in smaller tasks



## Real problem: often partly parallelizable





## Real problem: often partly parallelizable

Non-parallelizable (e.g. initialization)

Parallel program (e.g. 50 processors)

Non-parallelizable von-parallelizable

Non-parallelizable (e.g. initialization)

Non-parallelizable

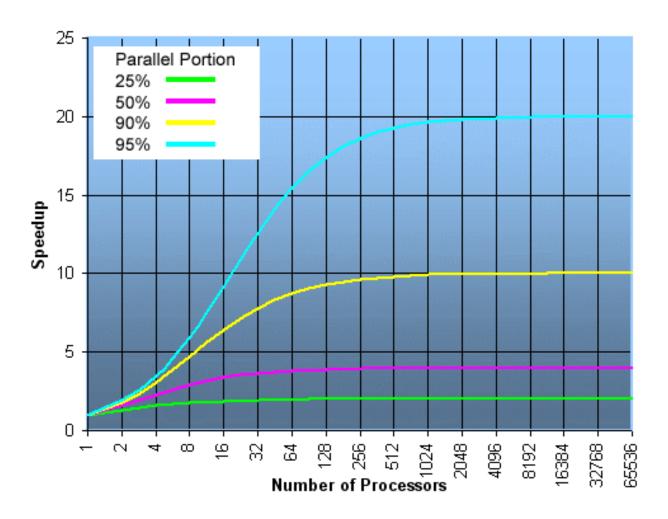
Time



#### Amdahl's Law

$$Speedup = \frac{1}{(1-p) + \frac{p}{N}}$$

N = number of processorsp = parallel fraction





## Amdahl's law describes strong scaling

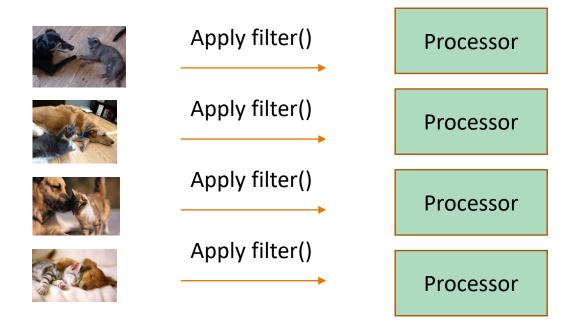
- Strong scaling
  - Variation of solution time with #processors for fixed total problem size
  - Possibilities to run the same problem in shorter time





# Amdahl's law describes strong scaling

- Strong scaling
  - Variation of solution time with #processors for fixed total problem size
  - Possibility to run the same problem in shorter time





#### Amdahl's law is theoretical...

- ... and reality is often worse!
- Further issues to tackle
  - Load balancing
  - Process scheduling
  - Communication overhead
  - File Input/Output
  - \_



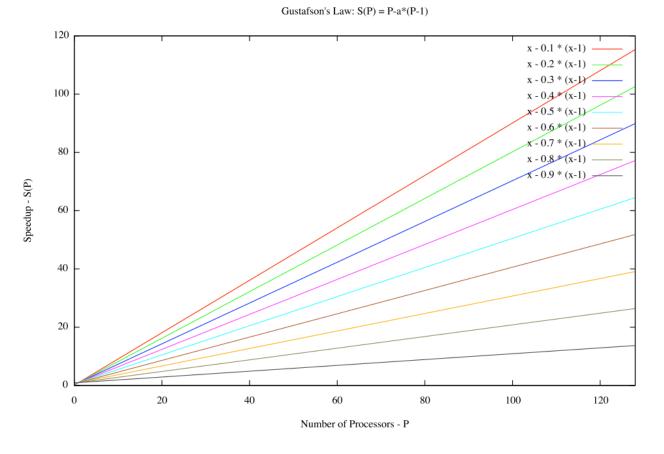
Do we have hope...???

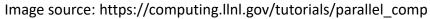


#### **Gustafson's Law**

$$Speedup = 1 + p \cdot (N - 1)$$

N = number of processorsp = parallel fraction







# Gustafson's law describes weak scaling

- Weak scaling
  - Variation of solution time with #processors for fixed problem sizeper processor
  - Possibilities to run a bigger problem in the same time



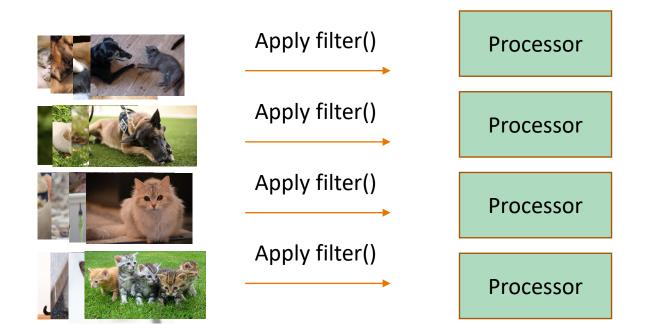
Apply filter()

Processor



# Gustafson's law describes weak scaling

- Weak scaling
  - Variation of solution time with #processors for fixed problem sizeper processor
  - Possibilities to run a bigger problem in the same time





# Strong scaling vs. weak scaling

Strong scaling: fixed work

$$Speedup = \frac{1}{(1-p) + \frac{p}{N}}$$

p=0.5

Weak scaling: fixed work per processor

$$Speedup = 1 + p \cdot (N - 1)$$









S=1







## Weak scaling is often the most relevant to HPC computations

- Some examples
  - Physics: "I can only run my fluid simulation a small domain / low resolution on my local PC"
  - Chemistry: "I can simulate a small molecule on my PC, but I want to simulate a big one"
- Common background
  - A big system / molecule increases the total work
  - Distributing larger work over multiple processors keeps the work per processor constant
  - ... this is weak scaling!



So now how can we program in parallel?



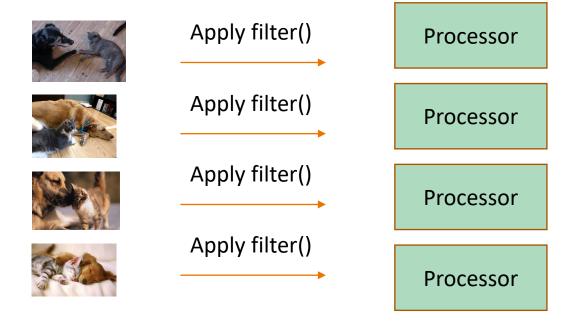
# Parallel programming models

- "How do I parallelize my code?"
- Two well-known programming models
  - Data parallelism
  - Task parallelism



# Parallel programming models: data parallelism

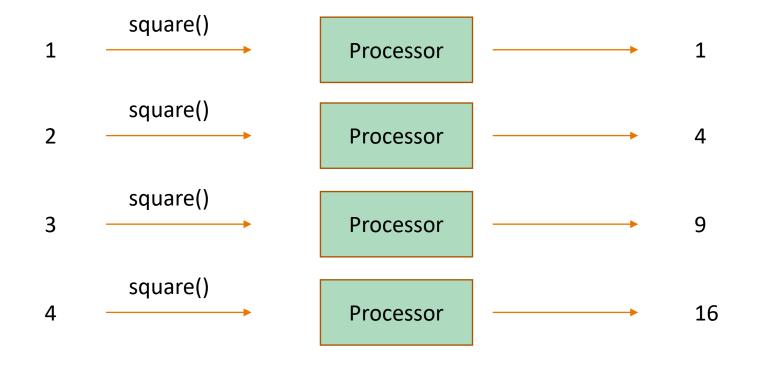
Each processor performs the same task on different data





# Parallel programming models: data parallelism

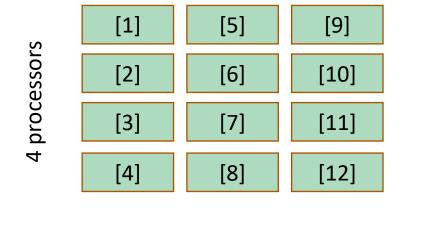
Each processor performs the same task on different data





## Parallel programming models: data parallelism

- Amount of parallelization depends on input data size
- Load balancing may be relatively easy
  - Same task on each data element
  - Approximately same time per element

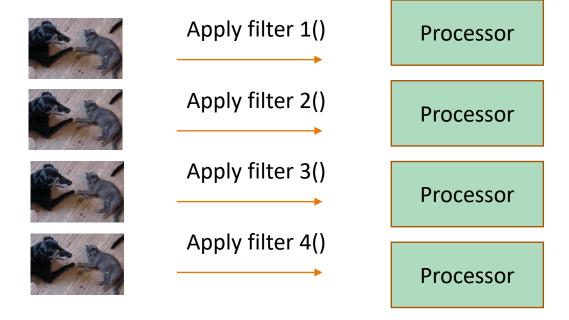


Time

SURF SARA

# Parallel programming models: task parallelism

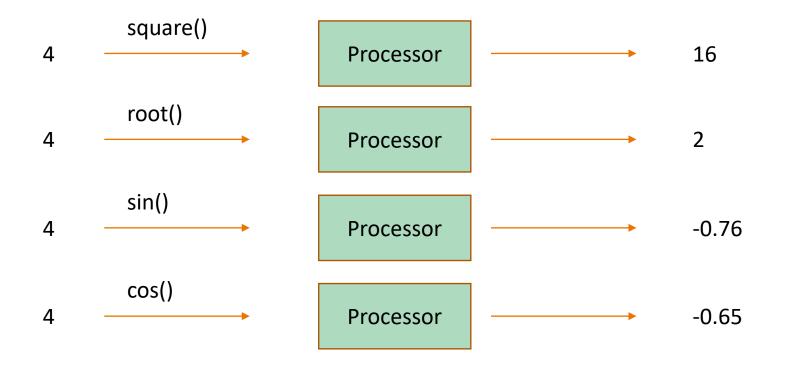
Each processor performs a different task on the same data





# Parallel programming models: task parallelism

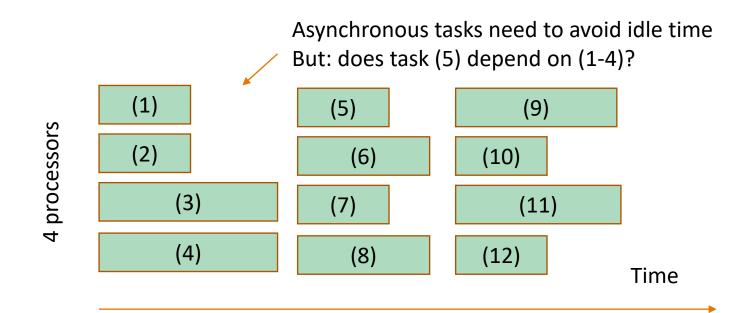
Each processor performs a different task on the same data





## Parallel programming models: task parallelism

- Amount of parallelization depends on the number of tasks
- Load balancing can be very difficult
  - Heterogeneous tasks may be executed over the same data
  - Each task may take a very different amount of time



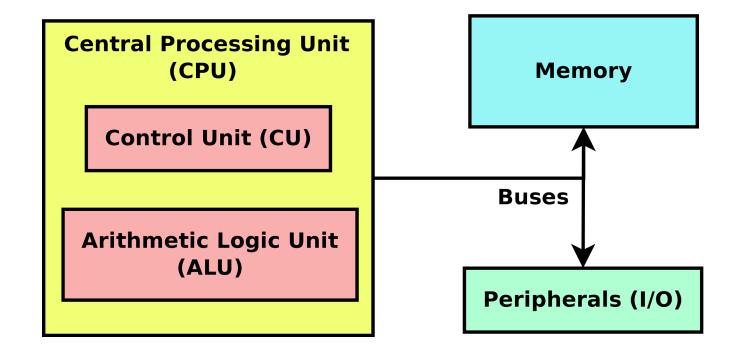


Theory is essential to understand concepts...

but parallel programming also requires understanding of the systems!



## A computer is...





# A memory is...

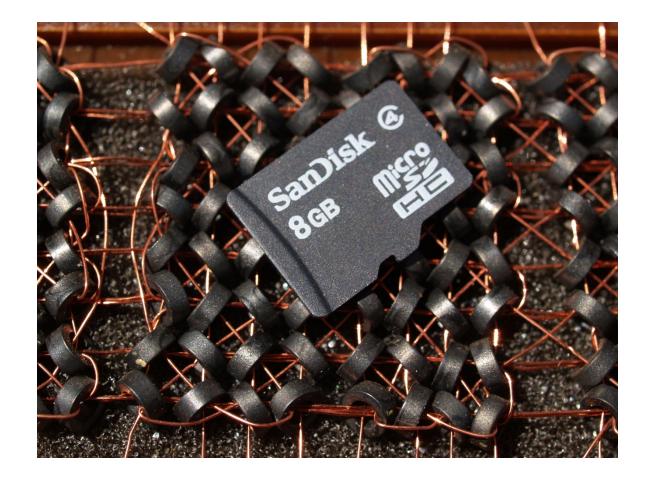
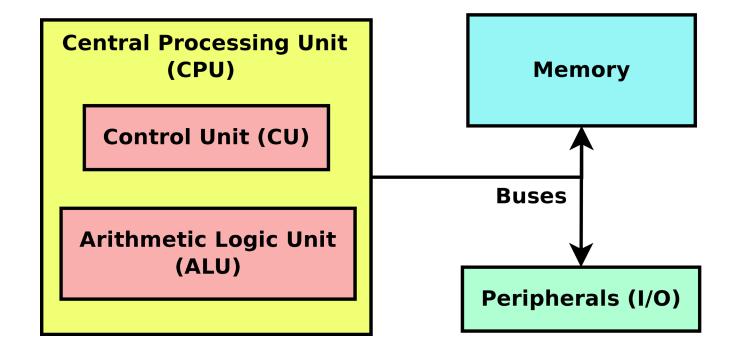


Image source: https://upload.wikimedia.org/wikipedia/commons/c/c0/8 bytes vs. 8Gbytes.jpg



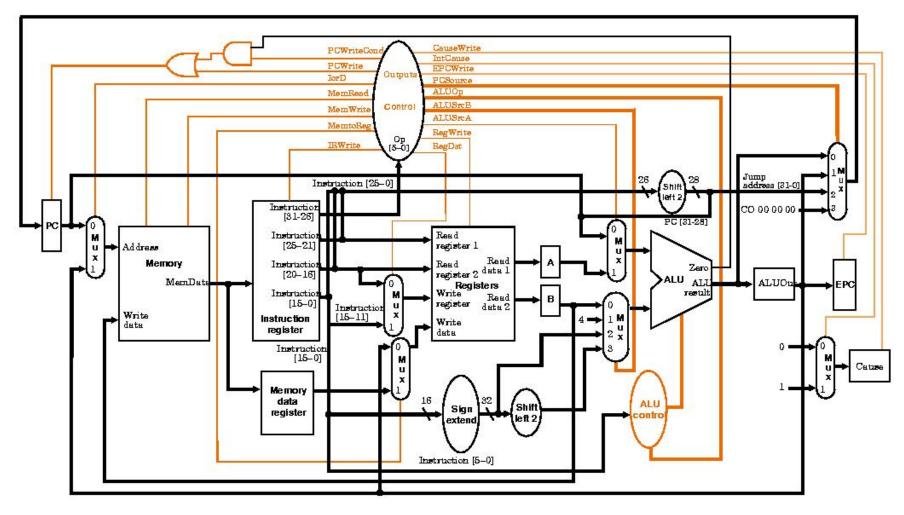
## A computer is...





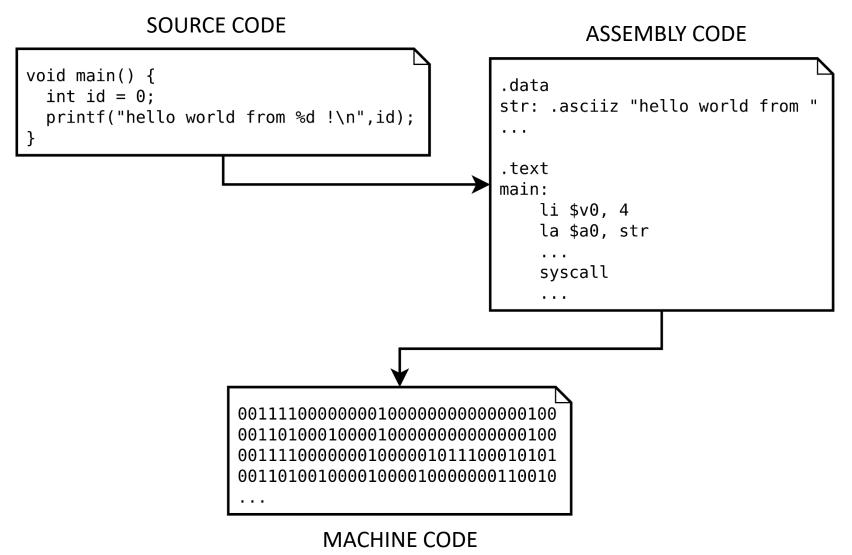
#### A CPU is...

Example of the MIPS architecture: control unit lines in orange, data lines in black



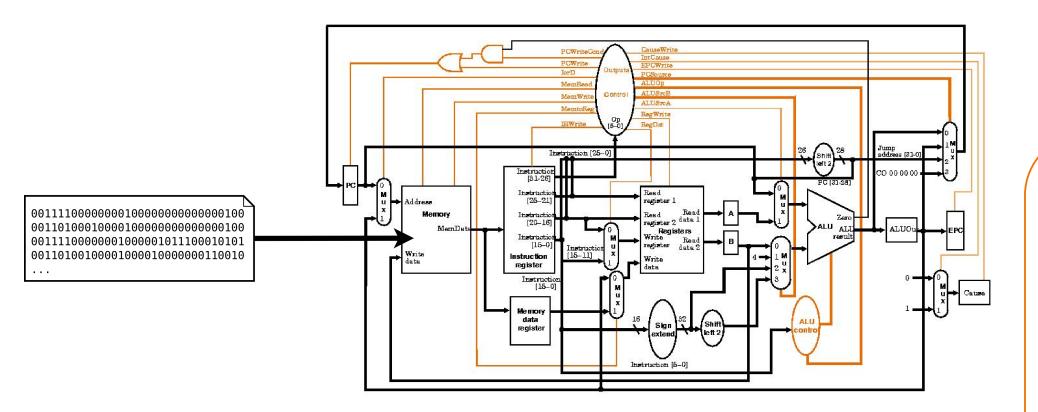


### A program becomes a process...

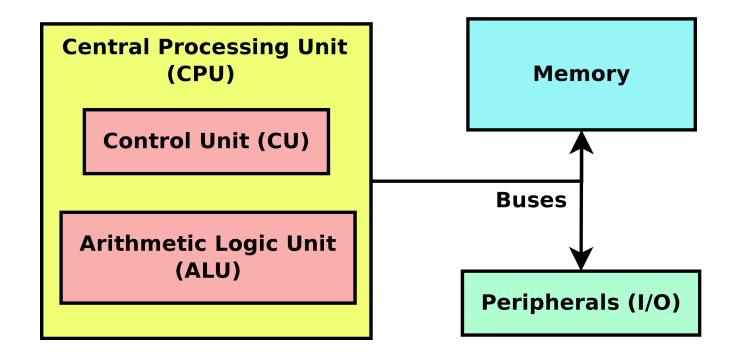




## A program becomes a process...

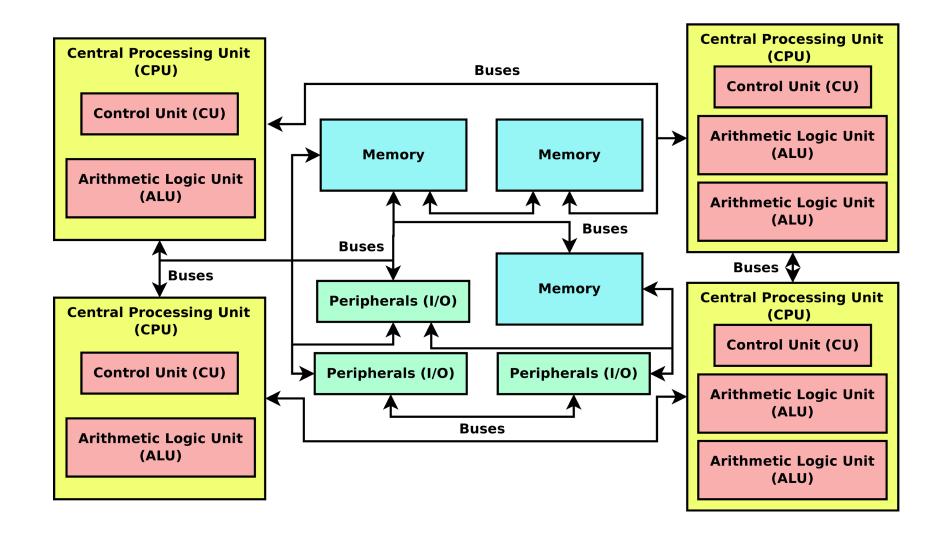






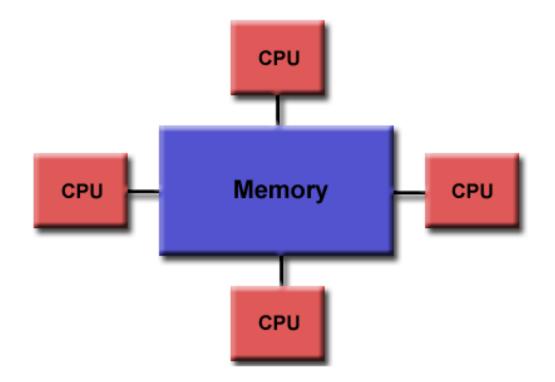


## Different types of computers architectures (!!1!!!!11!!eleven!!)



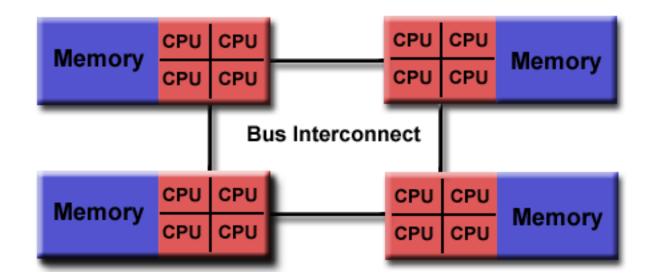


- Shared memory: all processors access the same memory (e.g. typical laptop)
  - Early model: uniform memory access for every processor (UMA)



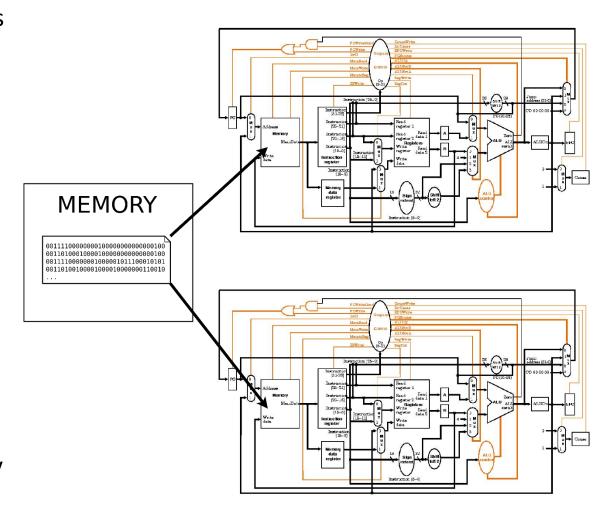


- Shared memory: all processors access the same memory (e.g. typical server)
  - General model: non-uniform memory access for every processor with protocol for cache coherency (ccNUMA)





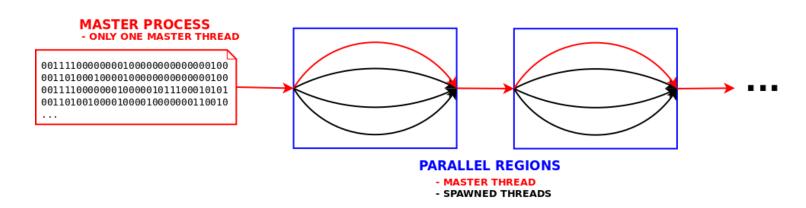
- Shared memory: all processors access the same memory
  - The programmer is seeing a single unified memory
  - Different sequences of execution (threads) run on the same process
  - Different memory modules may be used, but only one logical memory space is addressed
  - Communication between processors is done implicitly





#### Many threads in a process: OpenMP

- A process creates lightweight instances of itself (threads) that are coordinated for simultaneous execution
  - A thread shares the program code and data section with all other threads inside the same process
  - The parallel computing approach is called fork-join
    - A sequential program begins with one process (that is, only one thread)
    - A parallel region is defined by creating (spawning) threads from the process and destroying the original thread remains





... and this is where we are going to start!



