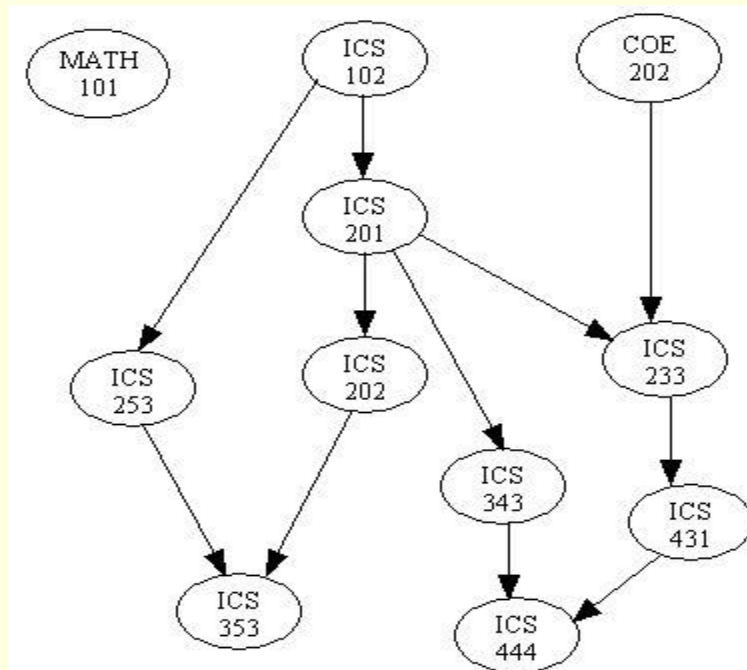


# Topological Sorting

*Kumkum Saxena*

# Introduction

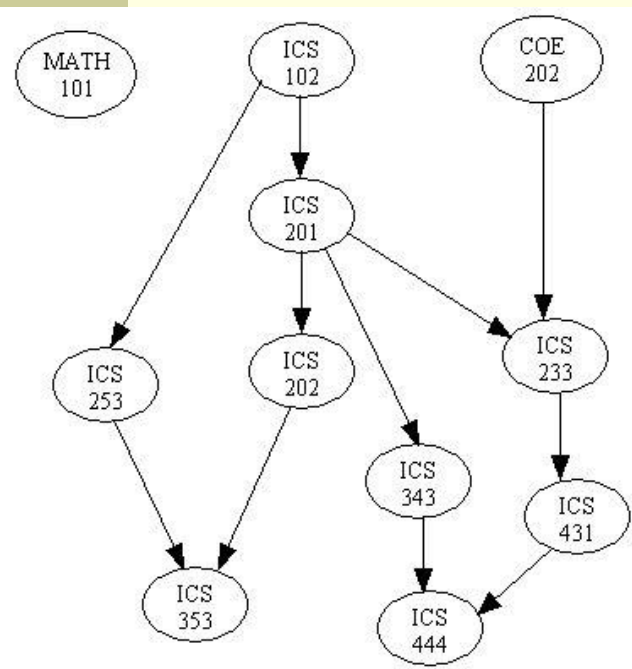
- There are many problems involving a set of tasks in which some of the tasks must be done before others.
- For example, consider the problem of taking a course only after taking its prerequisites.
- Is there any systematic way of linearly arranging the courses in the order that they should be taken?



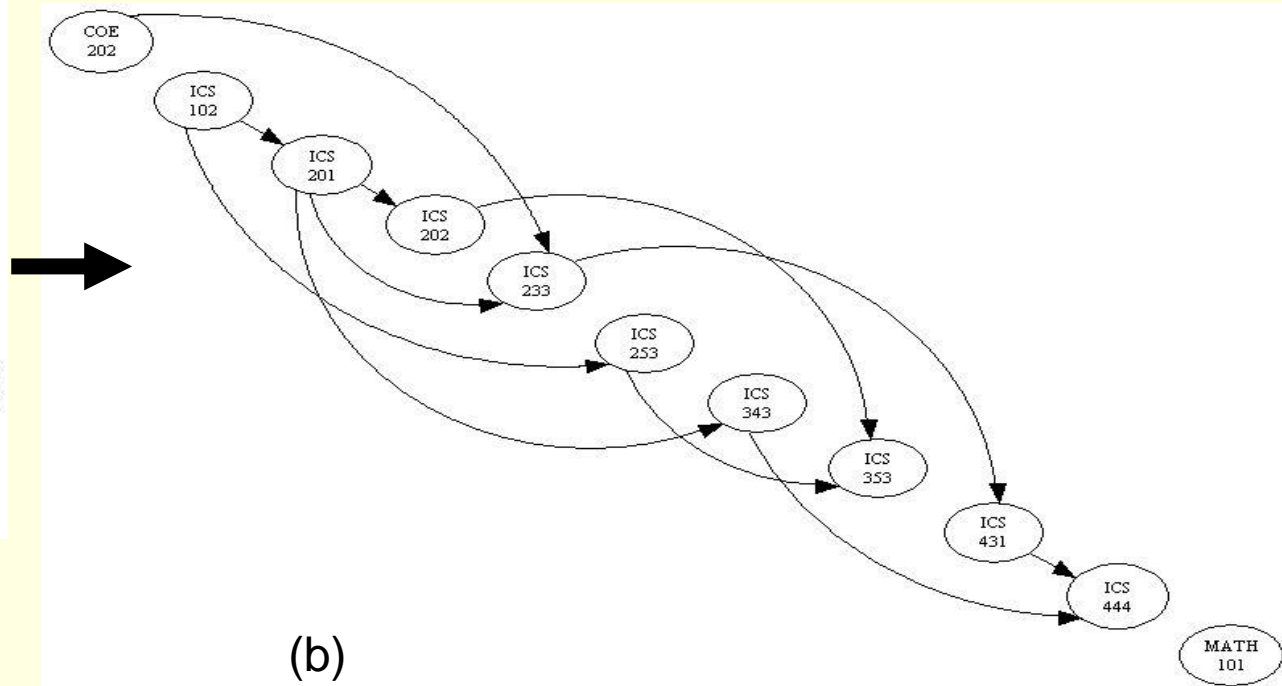
Yes! - Topological sort.

# Definition of Topological Sort

- Topological sort is a method of arranging the vertices in a directed acyclic graph (DAG), as a sequence, such that no vertex appear in the sequence before its predecessor.
- The graph in (a) can be topologically sorted as in (b)



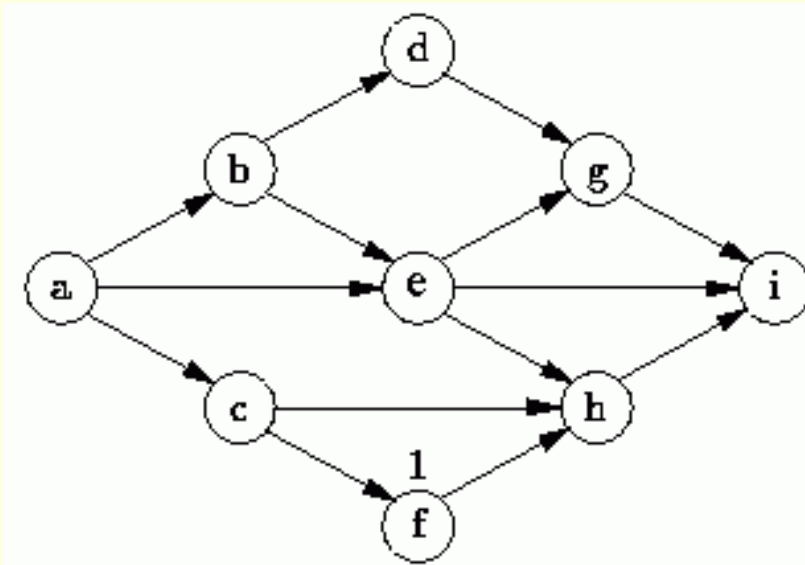
(a)



(b)

# Topological Sort is not unique

- Topological sort is not unique.
- The following are all topological sort of the graph below:



**s1 = {a, b, c, d, e, f, g, h, i}**

**s2 = {a, c, b, f, e, d, h, g, i}**

**s3 = {a, b, d, c, e, g, f, h, i}**

**s4 = {a, c, f, b, e, h, d, g, i}**  
**etc.**

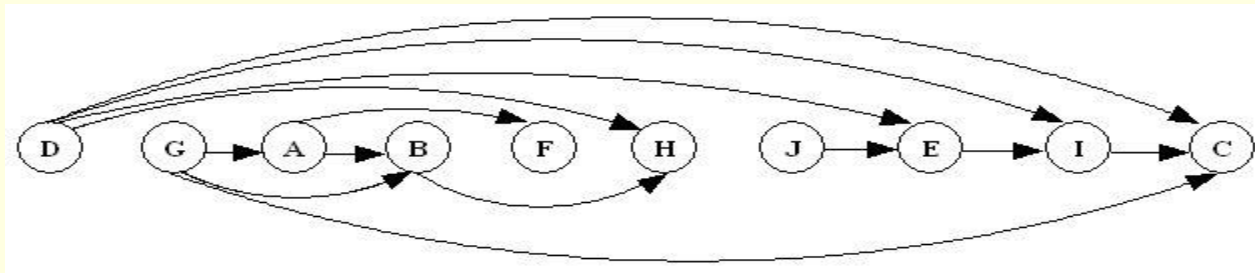
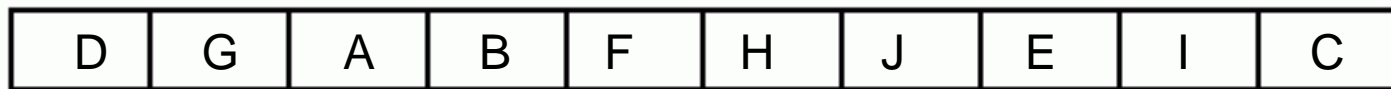
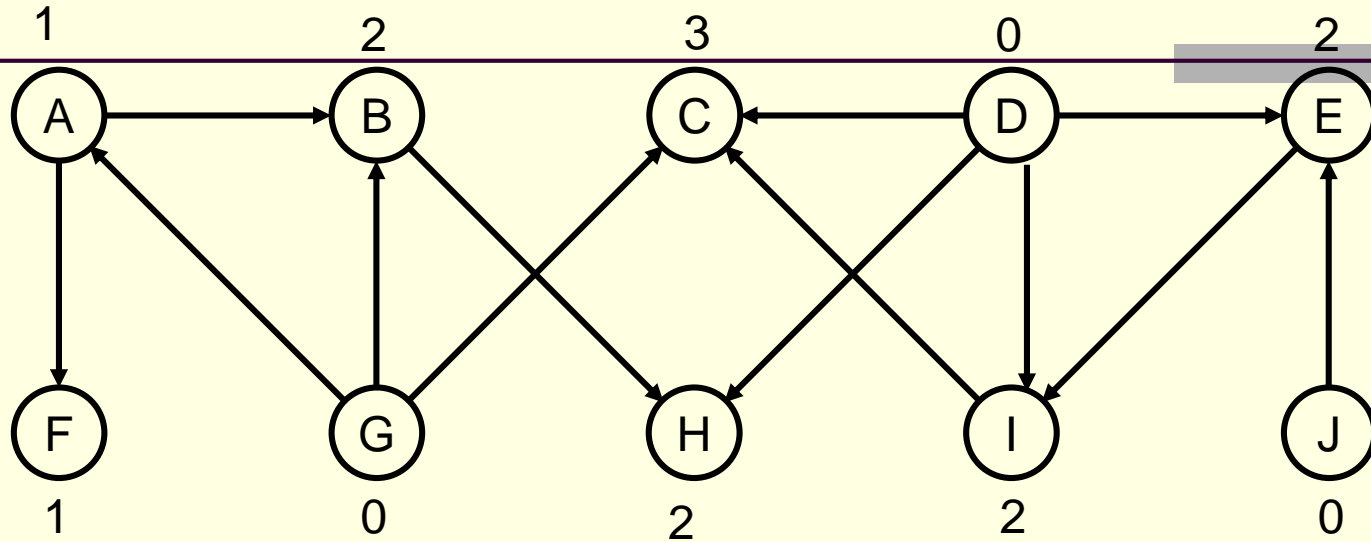
# Topological Sort Algorithm

- One way to find a topological sort is to consider in-degrees of the vertices.
- The first vertex must have in-degree zero -- every DAG must have at least one vertex with in-degree zero.
- The Topological sort algorithm is:

```
int topologicalOrderTraversal( ){  
    int numVisitedVertices = 0;  
    while(there are more vertices to be visited){  
        if(there is no vertex with in-degree 0)  
            break;  
        else{  
            select a vertex v that has in-degree 0;  
            visit v;  
            numVisitedVertices++;  
            delete v and all its emanating edges;  
        }  
    }  
  
    return numVisitedVertices;  
}
```

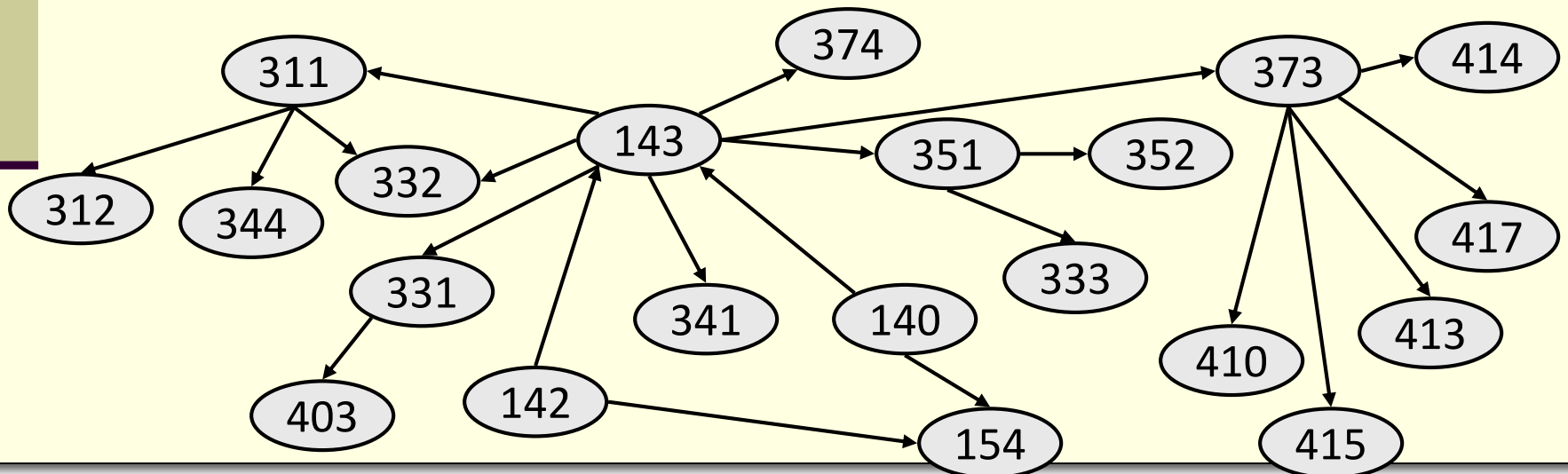
# Topological Sort Example

## ■ Demonstrating Topological Sort.



# Ordering a graph

- Suppose we have a directed acyclic graph (DAG) of courses, and we want to find an order in which the courses can be taken.
  - Must take all prereqs before you can take a given course. Example:
    - [142, 143, 140, 154, 341, 374, 331, 403, 311, 332, 344, 312, 351, 333, 352, 373, 414, 410, 417, 413, 415]
    - There might be more than one allowable ordering.
  - How can we find a valid ordering of the vertices?

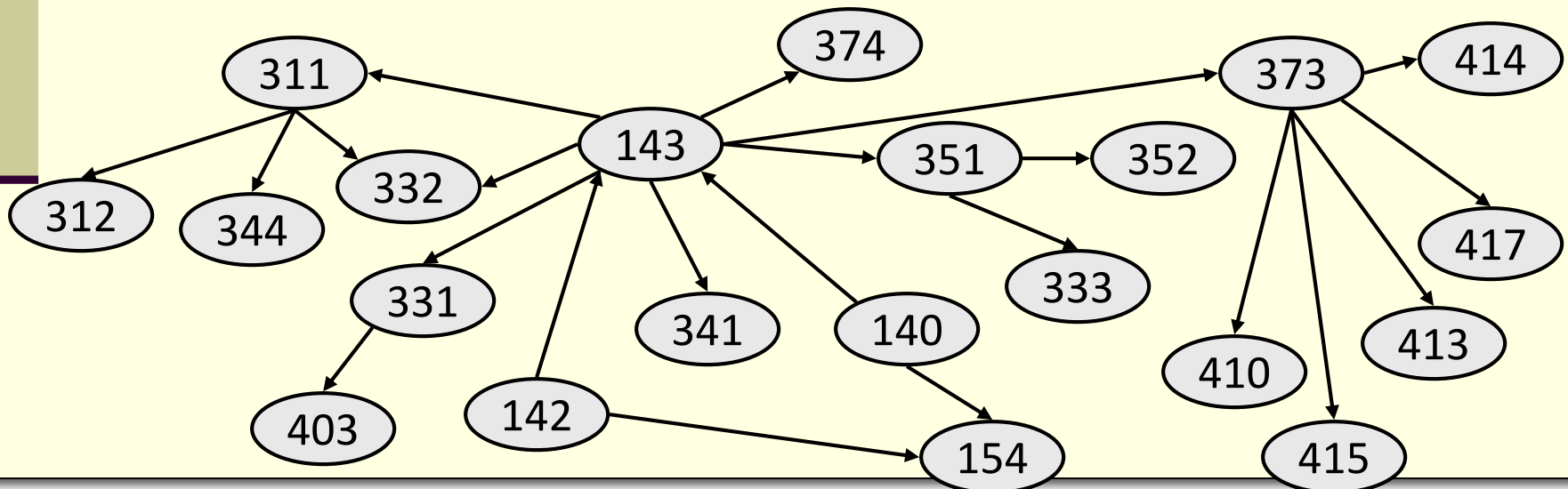


# Topological Sort

- **topological sort:** Given a digraph  $G = (V, E)$ , a total ordering of  $G$ 's vertices such that for every edge  $(v, w)$  in  $E$ , vertex  $v$  precedes  $w$  in the ordering.

Examples:

- determining the order to recalculate updated cells in a spreadsheet
- finding an order to recompile files that have dependencies
  - (any problem of finding an order to perform tasks with dependencies)

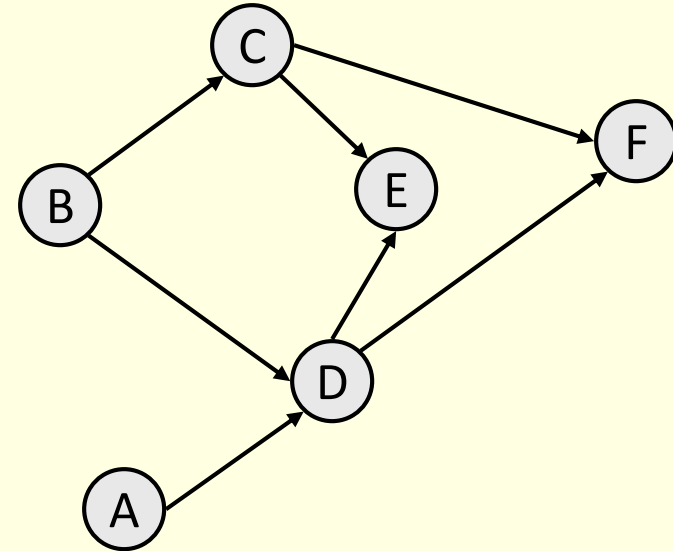




# Topo sort example

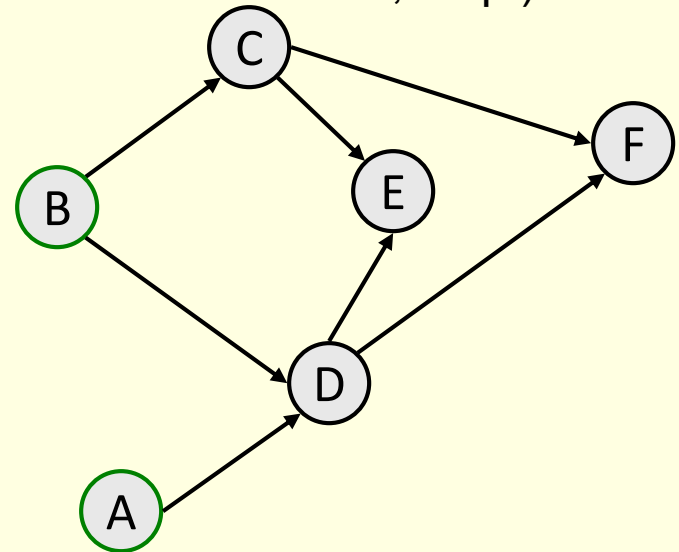
■ How many valid topological sort orderings can you find for the vertices in the graph below?

- [A, B, C, D, E, F], [A, B, C, D, F, E],
- [A, B, D, C, E, F], [A, B, D, C, F, E],
- [B, A, C, D, E, F], [B, A, C, D, F, E],
- [B, A, D, C, E, F], [B, A, D, C, F, E],
- [B, C, A, D, E, F], [B, C, A, D, F, E],
- ...



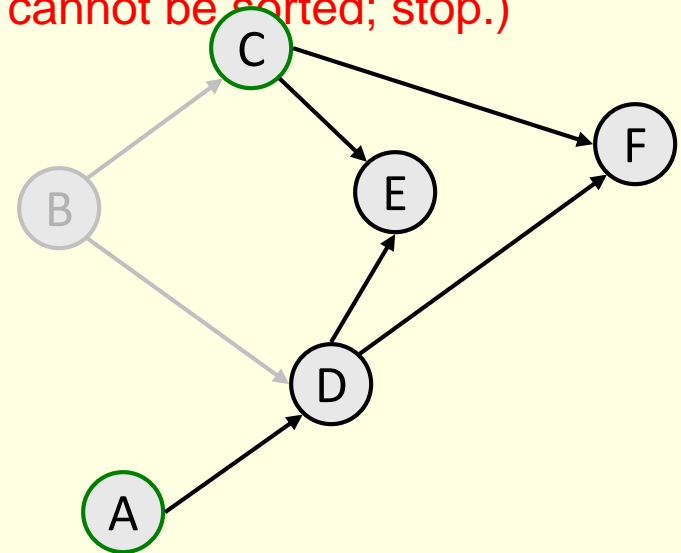
# Topo sort: Algorithm 1

- function topologicalSort():
  - *ordering* := { }.
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - *ordering* +=  $v$  .



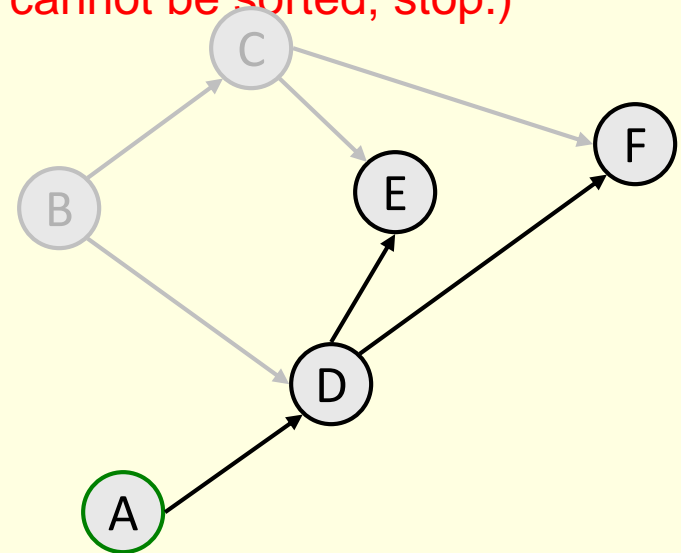
# Topo sort example

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
- $ordering = \{ B \}$



# Topo sort example

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C \}$



# Topo sort example

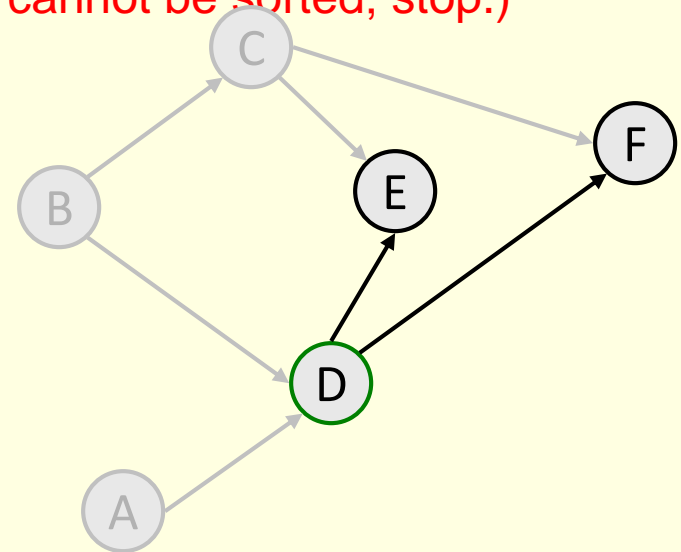
function topologicalSort():

- $ordering := \{ \}$ .

- Repeat until graph is empty:

- Find a vertex  $v$  with in-degree of 0 (no incoming edges).
  - (If there is no such vertex, the graph cannot be sorted; stop.)
- Delete  $v$  and all of its outgoing edges from the graph.
- $ordering += v$ .

- $ordering = \{ B, C, A \}$



# Topo sort example

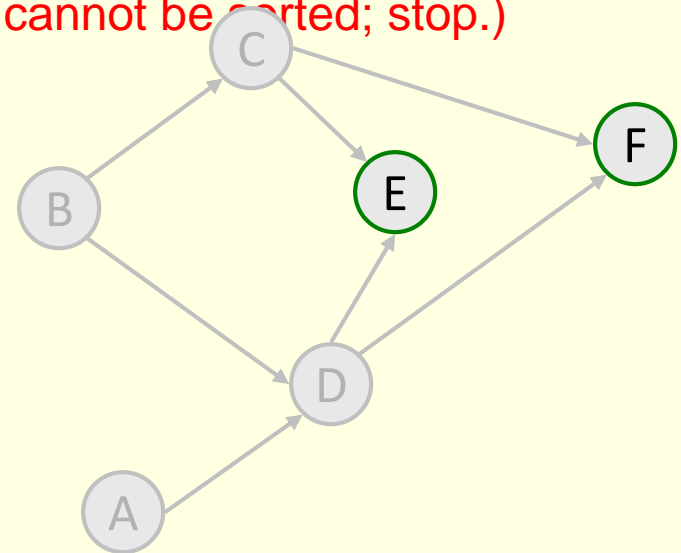
■ function topologicalSort():

■  $ordering := \{ \}$ .

■ Repeat until graph is empty:

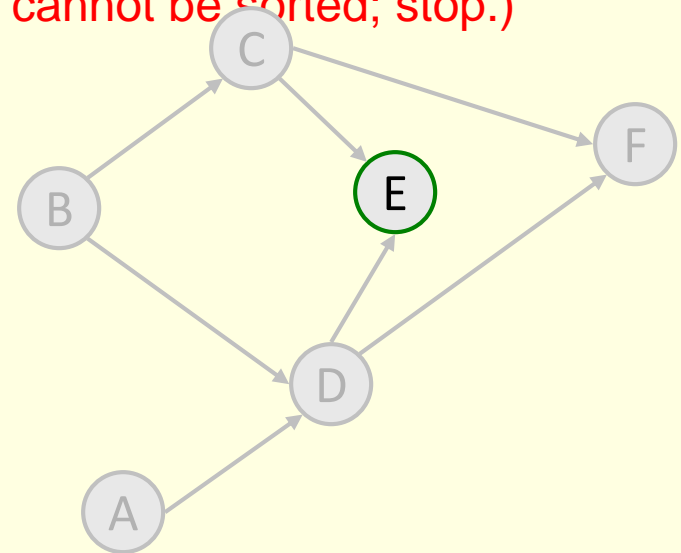
- Find a vertex  $v$  with in-degree of 0 (no incoming edges).
  - (If there is no such vertex, the graph cannot be sorted; stop.)
- Delete  $v$  and all of its outgoing edges from the graph.
- $ordering += v$ .

■  $ordering = \{ B, C, A, D \}$



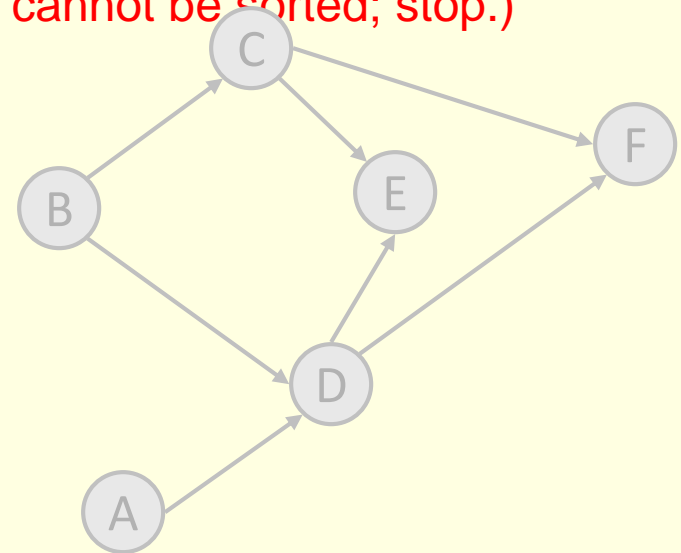
# Topo sort example

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
- $ordering = \{ B, C, A, D, F \}$



# Topo sort example

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
- $ordering = \{ B, C, A, D, F, E \}$





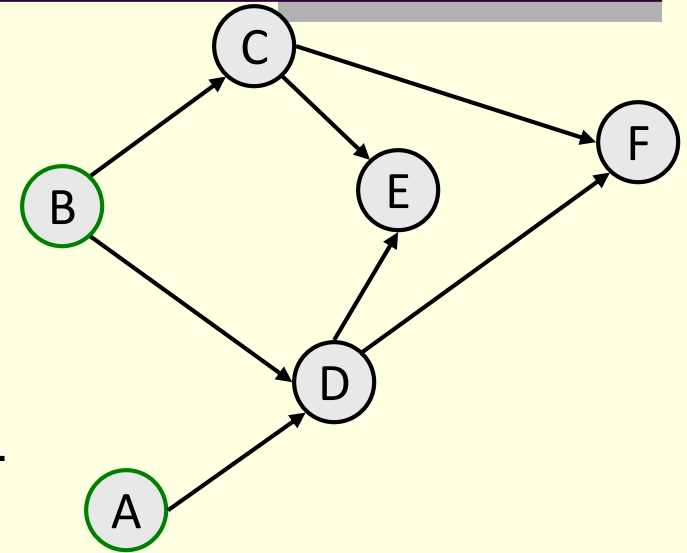
# Revised algorithm

■ We don't want to literally delete vertices and edges from the graph while trying to topological sort it; so let's revise the algorithm:

- $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
- $queue := \{\text{all vertices with in-degree} = 0\}.$
- $ordering := \{ \}.$
- Repeat until queue is empty:
  - Dequeue the first vertex  $v$  from the queue.
  - $ordering += v.$
  - Decrease the in-degree of all  $v$ 's neighbors by 1 in the  $map$ .
  - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- If all vertices are processed, success.  
Otherwise, there is a cycle.

# Topo sort example 2

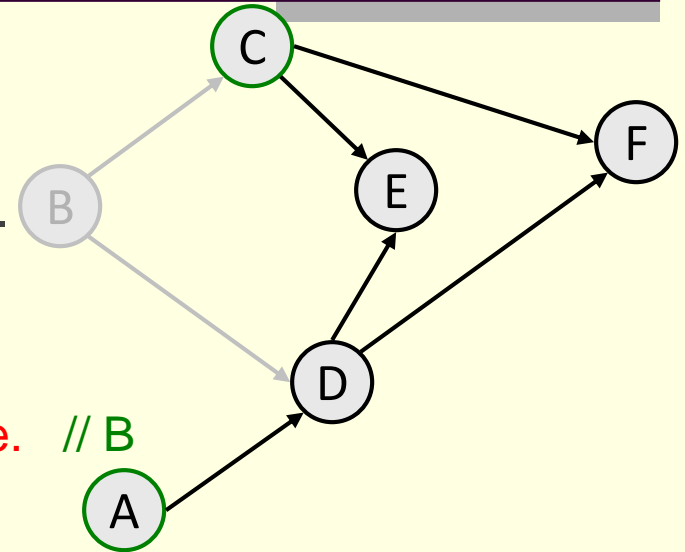
- function topologicalSort():
  - *map* := {each vertex  $\rightarrow$  its in-degree}.
  - *queue* := {all vertices with in-degree = 0}.
  - *ordering* := { }.
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue.
    - *ordering* +=  $v$ .
    - Decrease the in-degree of all  $v$ 's neighbors by 1 in the *map*.
    - *queue* += {any neighbors whose in-degree is now 0}.



- map := { A=0, B=0, C=1, D=2, E=2, F=2 }
- queue := { B, A }
- ordering := { }

# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // B
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // C, D neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$



- $map := \{ A=0, B=0, \mathbf{C=0}, \mathbf{D=1}, E=2, F=2 \}$
- $queue := \{ A, \mathbf{C} \}$
- $ordering := \{ \mathbf{B} \}$

# Topo sort example 2

function topologicalSort():

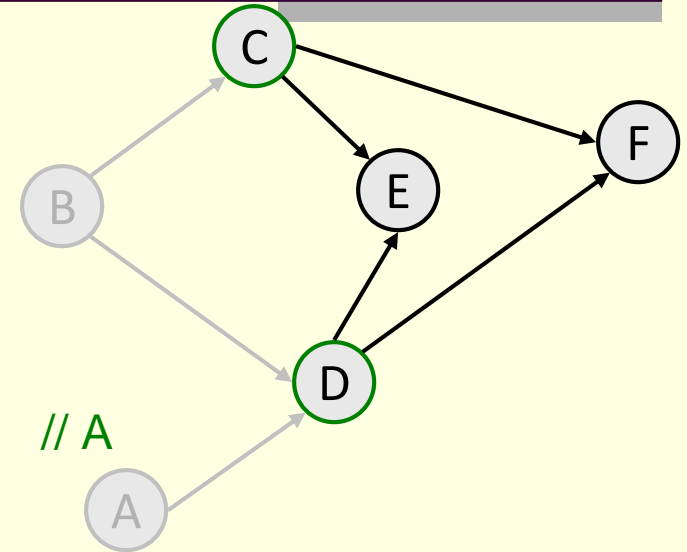
*map* := {each vertex  $\rightarrow$  its in-degree}.

*queue* := {all vertices with in-degree = 0}.

*ordering* := { }.

Repeat until queue is empty:

- Dequeue the first vertex *v* from the queue. // A
- *ordering* += *v*.
- Decrease the in-degree of all *v*'s // D  
neighbors by 1 in the *map*.
- *queue* += {any neighbors whose in-degree is now 0}.



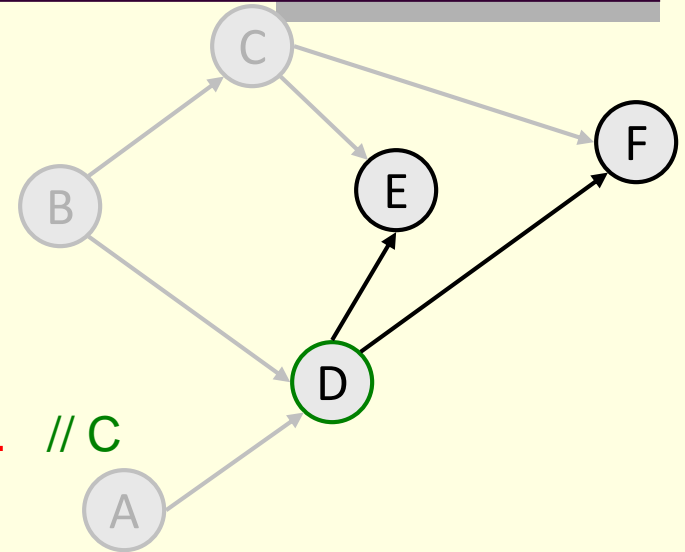
*map* := { A=0, B=0, C=0, **D=0**, E=2, F=2 }

*queue* := { C, **D** }

*ordering* := { B, **A** }

# Topo sort example 2

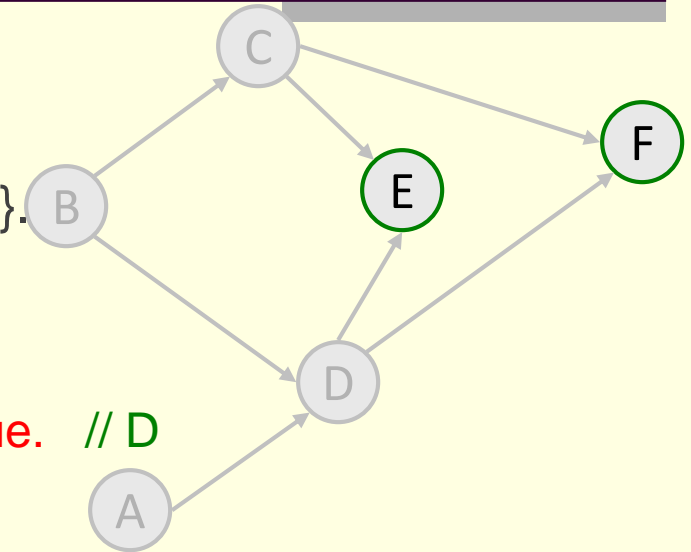
- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // C
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // E, F  
neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$



- $map := \{ A=0, B=0, C=0, D=0, E=1, F=1 \}$
- $queue := \{ D \}$
- $ordering := \{ B, A, \mathbf{C} \}$

# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // D
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // F, E neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$

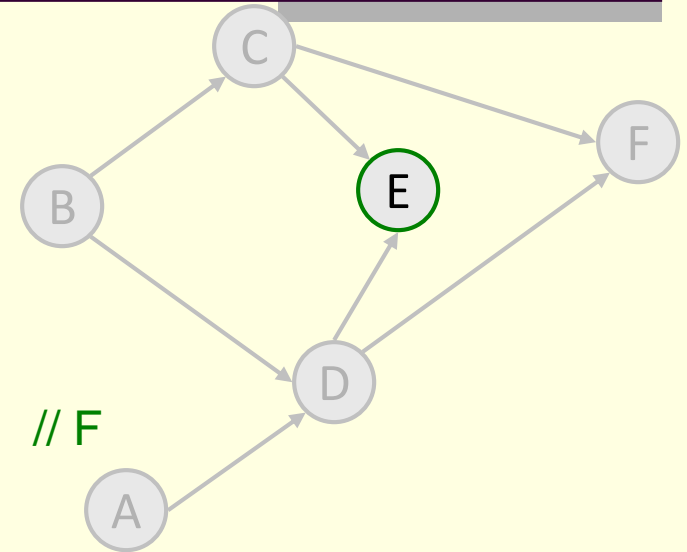


- $map := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue := \{ \mathbf{F}, \mathbf{E} \}$
- $ordering := \{ B, A, C, \mathbf{D} \}$

# Topo sort example 2

function topologicalSort():

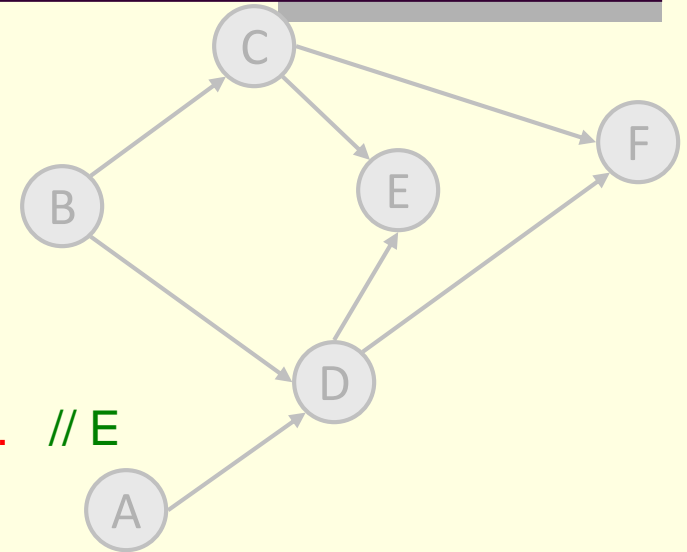
- $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
- $queue := \{\text{all vertices with in-degree} = 0\}.$
- $ordering := \{ \}.$
- Repeat until queue is empty:
  - Dequeue the first vertex  $v$  from the queue. // F
  - $ordering += v.$
  - Decrease the in-degree of all  $v$ 's // none neighbors by 1 in the  $map$ .
  - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$



- $map := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue := \{ E \}$
- $ordering := \{ B, A, C, D, F \}$

# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // E
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // none neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$



- $map := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue := \{ \}$
- $ordering := \{ B, A, C, D, F, E \}$



# Topo sort runtime

- What is the runtime of our topological sort algorithm?
  - (with an "adjacency map" graph internal representation)
  - function topologicalSort():
    - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$  //  $O(V)$
    - $queue := \{\text{all vertices with in-degree} = 0\}.$
    - $ordering := \{\}$ .
    - Repeat until queue is empty: //  $O(V)$ 
      - Dequeue the first vertex  $v$  from the queue. //  $O(1)$
      - $ordering += v.$  //  $O(1)$
      - Decrease the in-degree of all  $v$ 's neighbors by 1 in the  $map$ . //  $O(E)$  for all passes
      - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- Overall:  $O(V + E)$  ; essentially  $O(V)$  time on a sparse graph (fast!)