

King Saud University College of Computer and Information Sciences Computer Science Department (CSC 281)

Group #5				
Name	ID	Section		
(Leader) Jomanah Almuhanna	442200133	41201		
Sarah Alajlan	442202314			
Leena Almeshal	442200408			
Majd Alqahtani	442202282			

Methods implementation and Sample run

Method1: Trial division

```
Algorithm TrialDevision(n)
Input: n an integer to test for primality
Output: boolean value
limit ← ceil the sqrt of n
isPrime ← true
loop for i \leftarrow 2 to limit do
   if n \mod i = 0 then
      isPrime ← false
    end if
end loop
if isPrime = false then
    loop for i \leftarrow 1 to n do
       if n \mod i = 0
        print i
       end if
    end loop
 return isPrime
end if
else return isPrime
•Java Implementation:
public static boolean TrialDevision(int n) {
int limit = ((int)Math.ceil(Math.sqrt(n)));
boolean isprime = true;
for(int i = 2; i \le limit; i++)
if(n\%i == 0)
isprime = false;
```

•Pseudocode:

```
if(isprime == false) {
    for (int i = 1; i <= n; i++)
    {
        if(n%i == 0)
            System.out.print(i+ ",");
    }
    return isprime;}
    else
        return isprime;
}</pre>
```

•Explanation:

We set the limit which is the square root of n and initialize a Boolean attribute 'isPrime', then a loop starting from 2 to the limit we calculated, we make sure it is not an even number because no even number is prime.

•Sample run:

```
Enter a value to test for primality:
341
Factors: 1,11,31,341,
Composite
Enter a value to test for primality:
Factors: 1,3,11,17,33,51,187,561,
Composite
   ------
Enter a value to test for primality:
Factors: 1,7,13,19,91,133,247,1729,
Composite
   ______
Enter a value to test for primality:
19973
probably prime
Enter a value to test for primality:
20051
```

Method2: Fermat Method

•Pseudocode:

```
Algorithm FermatMethod(k,n)
```

Input: *n* an integer to test for primality, *k* a parameter that determines the accuracy of the test.

output: Boolean value.

```
loop while k > 0
```

```
a ← random number between 2 and n-2
x ← a^n-1
if gcd(a,n) ≠ 1 then
    return false
end if
If x mod n ≠ 1 then
    return false
end if
k ← k-1
end loop
```

return true

```
 \begin{tabular}{l} public static boolean FermatMethod(long $k$ , long $n$) { } \\ while($k > 0$) { } \\ long $a = 2 + (int)Math.random()*(n-4); \\ long $x = modPow(a, n-1, n); $ \\ if(gcd(a \ , n)! = 1) \\ return false; \\ if($x \% n ! = 1) \\ return false; \\ k--; \\ } \\ return true; \\ \end{tabular}
```

•Explanation:

First we have a loop for the number of times we want to test a number for primality, inside of it we generate a random number 'a' with range [2.. n-2],then we need to calculate a^n-1 (mod n) we calculated a^n-1 by calling our implemented method 'modpow' and stored it in 'x', also we need to make sure there aren't any common devisors between 'a' and 'n' if it is proven false then we can take x mod n (which is a^n-1 (mod n)) and check the result if it is equal to 1 or not.

•Sample run:

```
Enter the number of times to test for primality:

10

Enter a value to test for primality:

341

probably prime

Enter the number of times to test for primality:

11

Enter a value to test for primality:

561

probably prime

Enter the number of times to test for primality:
```

```
15
Enter a value to test for primality:
1729
probably prime
Enter the number of times to test for primality:
10
Enter a value to test for primality:
19973
probably prime
   ------
   Enter the number of times to test for primality:
   Enter a value to test for primality:
   20051
   probably Prime
   ______
Enter the number of times to test for primality:
5
Enter a value to test for primality:
49597
probably prime
   ______
Enter the number of times to test for primality:
Enter a value to test for primality:
52721
probably prime
Enter the number of times to test for primality:
20
Enter a value to test for primality:
86069
probably prime
   -----
```

Method3: Miller-Rabin

•Pseudocode:

Algorithm millerRabinIsPrime(n,k)

Input: n an integer to test for primality, k a parameter that determines the accuracy of the test.

output: Boolean value. If n<2 then return false End if if n=2 then return true End if If $n \mod 2 = 0$ then return false End if $m \leftarrow 0$ e **←** 0 **loop for** $i \leftarrow 1$ to n-1 do x **←**2^I If $(n-1) \mod x = 0$ then $m \leftarrow (n-1)/x$ e **←** i End if If $(n-1) \mod x \neq 0$ Break; End if **End loop** ver: loop for i ← 0 to k do a← random number between 2 and n-1 x ← a^m mod n if(x = 1 or x = n-1)Continue End if **loop for** $j \leftarrow 0$ to e **do** $x \leftarrow x^2 \mod n$ If x=1 then return false End if If x=n-1 then Continue ver End if **End loop** return false

End loop

```
public static boolean millerRabinIsPrime(int n, int k) {
if(n < 2)
return false;
if(n == 2)
return true;
if(n \% 2 == 0)
return false;
long m = 0;
long e = 0;
for(int i = 1; i < n-1; i++) {
int x = (int)Math.pow(2, i);
if((n-1)\%x == 0) {
m = (n-1)/x;
e = i;
}
if((n-1)\%x != 0)
break;
}
       for(int i = 0; i < k; i++) {
ver:
long a = 2 + (long)Math.random()*(n-3);
long x = modPow(a, m, n);
 if(x == 1 || x == n-1)
continue;
for(int j = 0; j < e; j++) {
x = modPow(x, 2, n);
if(x == 1)
return false;
if(x == n-1)
continue ver;
}
return false;
}
```

```
return true;
```

}

•Explanation:

Method will first handle the base cases if n<2 or if n is even, it will return false. It will search for an odd number m such that n-1 can be written as $m*2^e$ then it will check the following k times with a for loop: generates an a between 2 and n-1, compute $x = a^m \pmod n$ with our method 'modPow' if result is 1 or -1, it will continue the loop else it will do the following e times with a for loop: x = (x*x) % n if the result is 1 return false if -1 it will continue the outer loop if we continue through all the iterations it will return true at the end of the method meaning its prime else if it did not continue It will reach the end of the loop returning false meaning it didn't satisfy the condition.

•Sample run:

```
Enter the number of times to test for primality:
5
Enter a value to test for primality:
composite
Enter the number of times to test for primality:
10
Enter a value to test for primality:
composite
Enter the number of times to test for primality:
Enter a value to test for primality:
1729
composite
Enter the number of times to test for primality:
Enter a value to test for primality:
19973
probably prime
   -----
Enter the number of times to test for primality:
Enter a value to test for primality:
20051
probably prime
   -----
```

```
Enter the number of times to test for primality:

5
Enter a value to test for primality:
49597
probably prime
------
Enter the number of times to test for primality:
16
Enter a value to test for primality:
52721
probably prime
------
Enter the number of times to test for primality:
7
Enter a value to test for primality:
86069
probably prime
```

EXTRA METHODS

```
Algorithm modPow(a , b ,c)

Input a the base , b the exponent and c the modulus

Output a^b mod c

res \leftarrow 1

Loop for I \leftarrow 0 to b-1 do

res \leftarrow res*a

res \leftarrow res mod c

End loop

return res mod c
```

Algorithm gcd(a , b) Input a and b to calculate their gcd Output gcd of a and b

```
If a > b
  If b = 0
     return a
  End if
  Else
     return gcd(b, a mod b)
  End else
End if
Else
  If a = 0
     Return b
  End if
  Else
    Return gcd(a, b mod a)
  End else
End else
```

```
if(a > b) {

if(b == 0)
return a;

else return gcd(b, a % b);

}

else {
  if(a == 0)
  return b;
  else return gcd(a, b %a);
}}
```

Discussion

	Strength	Weakness	Accuracy
Trial Division	We can be certain if the number is prime or not Faster than the other two methods	We can't reach a high range of numbers	The most accurate out of the three methods
Fermat	We can reach higher range of numbers unlike trial divison	Its not that accurate. Pseudoprimes can pass the test. Faster when figuring out composite numbers but slower with primes comparing to the third method	A lot of pseudoprimes pass the test The method isn't reliable since it yields Different results for the same number

Miller- Rabin	Its more accurate than fermat and we can reach higher range of numbers Pseudoprimes are annonced as composites.		Few psudoprimes pass the test more accurate than fermat
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K = 10 Efficiency comparison:

n	Method 1	Method 2	Method 3
341	0.0004283	0.000743	0.0011614
561	0.000527	0.0007541	0.0009658
19973	0.0003745	0.0057174	0.0017754
20051	0.0003028	0.0048482	0.0022307