K180268

- St. @ The problem in code is that the Threads will headlock/ interlock among them, as they will acquie back A, book B, respectively, but execute at the same time. So each thread needs both of lock A and back B to be free.
 - The code can be solved, by removing either both A completely or bock B completely. Now the threads voill execute in a way that the one which comes first will be dealt with first, and after it gus executed the other will resume. This is also needed to be done as they acquire the data same as a and bo have sure that the books are always taken in same order by threads, so deadlock doesn't occur. Lock ordering will also be effective. In

92: while (true) {
wait (fork [(i+1) 202])

// eat now

signal (fork[(i+1)702]) 3

93: asg pages = 2^8 page size = $0a68 = asg \times a^{10} = a^{12}$ frames = $1a8 = a^7$ 8+18 = asg bits

Physical address space bits = asg bits

Squents are created, where each has a different virtual address space that directly corresponds to process. Example can be taken, that a program may be having many segments in the logical address space, such as different libraries, mapping on to physical memory is done by limit and base registers. The segments, now, may not be contiguous recessarily, but the exact amount of space that it needs is miportant, therefore no witered gragmentation problem.

(2) Multitheaded process provides more speed as compared to processes without thread, because it miscoses performance geatly because CPU can divide simultaneous work aworgst nulty be the ends. The sum of this spent on all processers is higher b/c extra work of coordinating multiple threads. It allows shared resource mage like memory, while it has to do additional synchronization work b/w the ends. but finishes faster b/c simultaneous work being done, i.e. work on threads of diff type at save time.