CS232 Exam 1 February 14, 2007

Name:					
Section:	10am	noon	2pm	4pm	(circle one)

- This exam has 6 pages, including this cover.
- There are three questions, worth a total of 100 points.
- The last two pages are a summary of the MIPS instruction set, calling convention, and hexadecimal notation, which you may remove for your convenience.
- No other written references or calculators are allowed.
- Write clearly and show your work. State any assumptions you make.
- You have 50 minutes. Budget your time, and good luck!

Question	Maximum	Your Score
1	40	
2	20	
3	40	
Total	100	

Question 1: Write a recursive MIPS function (40 points)

Here is a recursive function that multiplies two unsigned integers and returns the result as an unsigned integer. Translate multiply into a **recursive** MIPS assembly language function; iterative versions (*i.e.*, those with loops) will not receive full credit. You will not be graded on the efficiency of your code, but you must follow all MIPS conventions. Comment your code!!!

```
unsigned int
multiply(unsigned int x, unsigned int y) {
 if (y == 1) {
    return x;
 }
 unsigned temp = mul(x << 1, y>> 1);
 if ((y \& 1) == 1) {
    temp = temp + x;
 }
 return temp;
}
multiply:
   bne
          $a1, 1, not base case # check for the base case
   mov
          $v0, $a0
                                  # return x
   jr
          $ra
not base case:
   sub
          $sp, $sp, 12
                                # save x and y before shifting them
          $ra, 0($sp)
   SW
         $a0, 4($sp)
   SW
   sw
         $a1, 8($sp)
   sll $a0, $a0, 1
                                 # x << 1
         $a1, $a1, 1
                                 # v >> 1
   srl
   jal
         multiply
                                 # temp = multiply(...)
   lw
         $a1, 8($sp)
                                 # restore original y
          $a1, $a1, 1
                                 # y & 1
   and
          $a1, 1, skip if
                                 # skip if (y & 1) != 1
   bne
         $a0, 4($sp)
                                # restore original x
   lw
   add
          $v0, $v0, $a0
                                 \# temp += x
skip if:
          $ra, 0($sp)
                                # restore $ra
   1w
          $sp, $sp, 12
                                 # restore stack (same amount)
   add
          $ra
                                 # return temp
   jr
```

Question 2: Concepts (20 points)

Write a short answer to the following questions. For full credit, answers should not be longer than **two** sentences.

Part a) Write the body of a function that returns the second least significant byte (shown shaded in the picture) of a 32-bit integer using bit-wise logical operations and shifts. No control flow (e.g., loops/ifs) are needed. (10 points)

Part b) What is an abstraction layer? How does it relate to instruction set architectures (ISAs)? (5 points)

An abstraction layer is a way of hiding the implementation details of a particular set of functionality by providing an interface that separates the functionality from how it is implemented. An instruction set architecture (ISA) is the interface that abstracts the processor's funcationality from how it is implemented.

Part c) What is the "assembler temporary"? What is the motivation for it? (5 points)

The assembler temporary (\$at) is a register that, by convention, is not used by assembly programmers or compilers, so that it can be used as a temporary register for implementing pseudo-instructions that are implemented as multiple real assembly instructions. For example, blt \$a0, \$a1, Label is translated into the following:

Question 3: Understanding MIPS programs (40 points)

```
$t0, $a1, 1
knight:
             sub
pawn:
             blt
                     $t0, $0, rook
                     $t1, $t0, 4
             mul
                     $t1, $t1, $a0
             add
             lw
                     $v0, 0($t1)
             blt
                     $v0, $a2, bishop
                     $t0, $t0, 1
             sub
             j
                     pawn
             li
                     $v0, 0
rook:
bishop:
             jr
                     $ra
```

Part (a)

Translate the function knight above into a high-level language like C or Java. Your function header should list the types of any arguments and return values. Also, your code should be as concise as possible, without any gotos or pointer arithmetic. We will not deduct points for syntax errors unless they are significant enough to alter the meaning of your code. (30 points)

```
int
find_backwards(int arr[], int len, int min) {
    for (int i = len - 1 ; i >= 0 ; i = i - 1) {
        if (arr[i] < min) {
            return arr[i];
        }
    }
    return 0;
}</pre>
```

Part (b)

Describe briefly, in English, what this function does. (Note: It has nothing to do with chess.) (10 points)

The function walks backwards through an array of integers (potentially starting at the end), looking for an integer below the 3rd argument. The first one found is returned; if one is not found, the value 0 is returned.

MIPS instructions

These are some of the most common MIPS instructions and pseudo-instructions, and should be all you need. However, you are free to use *any* valid MIPS instructions or pseudo-instruction in your programs.

Category	Example Instruction		Meaning	
Arithmetic	add	\$t0, \$t1, \$t2	\$t0 = \$t1 + \$t2	
	sub	\$t0, \$t1, \$t2	\$t0 = \$t1 - \$t2	
	addi	\$t0, \$t1, 100	\$t0 = \$t1 + 100	
	mul	\$t0, \$t1, \$t2	$$t0 = $t1 \times $t2$	
Logical	and or sll srl sra	\$t0, \$t1, \$t2 \$t0, \$t1, \$t2 \$t0, \$t1, \$t2 \$t0, \$t1, \$t2 \$t0, \$t1, \$t2 \$t0, \$t1, \$t2	\$t0 = \$t1 & \$t2 (Logical AND) \$t0 = \$t1 \$t2 (Logical OR) \$t0 = \$t1 << \$t2 (Shift Left Logical) \$t0 = \$t1 >> \$t2 (Shift Right Logical) \$t0 = \$t1 >> \$t2 (Shift Right Arithmetic)	
Register Setting	move	\$t0, \$t1	\$t0 = \$t1	
	li	\$t0, 100	\$t0 = 100	
Data Transfer	lw	\$t0, 100(\$t1)	\$t0 = Mem[100 + \$t1] 4 bytes	
	lb	\$t0, 100(\$t1)	\$t0 = Mem[100 + \$t1] 1 byte	
	sw	\$t0, 100(\$t1)	Mem[100 + \$t1] = \$t0 4 bytes	
	sb	\$t0, 100(\$t1)	Mem[100 + \$t1] = \$t0 1 byte	
Branch	beq	\$t0, \$t1, Label	if $(\$t0 = \$t1)$ go to Label	
	bne	\$t0, \$t1, Label	if $(\$t0 \neq \$t1)$ go to Label	
	bge	\$t0, \$t1, Label	if $(\$t0 \geq \$t1)$ go to Label	
	bgt	\$t0, \$t1, Label	if $(\$t0 > \$t1)$ go to Label	
	ble	\$t0, \$t1, Label	if $(\$t0 \leq \$t1)$ go to Label	
	blt	\$t0, \$t1, Label	if $(\$t0 \leq \$t1)$ go to Label	
Set	slt	\$t0, \$t1, \$t2	if (\$t1 < \$t2) then \$t0 = 1 else \$t0 = 0	
	slti	\$t0, \$t1, 100	if (\$t1 < 100) then \$t0 = 1 else \$t0 = 0	
Jump	j	Label	go to Label	
	jr	\$ra	go to address in \$ra	
	jal	Label	\$ra = PC + 4; go to Label	

The second source operand of the arithmetic, logical, and branch instructions may be a constant.

Register Conventions

The *caller* is responsible for saving any of the following registers that it needs, before invoking a function.

\$t0-\$t9 \$a0-\$a3 \$v0-\$v1

The *callee* is responsible for saving and restoring any of the following registers that it uses.

\$s0-\$s7 \$ra

Pointers in C:

Declaration: either char *char ptr -or- char char array[] for char c

Dereference: c = c array[i] -or- c = *c pointer

Take address of: $c_pointer = &c$

Hexadecimal Notation

C and languages with a similar syntax (such as C++, C# and Java) prefix hexadecimal numerals with "0x", e.g. "0x5A3". The leading "0" is used so that the parser can simply recognize a number, and the "x" stands for hexadecimal.

Hex	Bin	Dec
0	0000	0
1	0001	1
2	0010	2
3	0011	3
4	0100	4
5	0101	5
6	0110	6
7	0111	7
8	1000	8
9	1001	9
A	1010	10
В	1011	11
С	1100	12
D	1101	13
Е	1110	14
F	1111	15