

Solutions - Midterm Exam

(February 17th @ 7:30 pm)

Presentation and clarity are very important! Show your procedure!

PROBLEM 1 (20 PTS)

- Compute the result of the following operations. The operands are signed fixed-point numbers. The result must be a signed fixed point number. For the division, use $x = 5$ fractional bits.

1.010001 + 1.011	1001.1101 - 1.011101	0.010101 + 01.11111
10.101 × 0.10011	1.011 × 1.0101	10.10010 ÷ 0.101

$$\begin{array}{r}
 \begin{array}{c} \text{C}_8=1 \\ \text{C}_7=1 \\ \text{C}_6=0 \\ \text{C}_5=1 \\ \text{C}_4=0 \\ \text{C}_3=0 \\ \text{C}_2=0 \\ \text{C}_1=0 \\ \text{C}_0=0 \end{array} \\
 \begin{array}{r} 1\ 1.0\ 1\ 0\ 0\ 0\ 1\ + \\ 1\ 1.0\ 1\ 1\ 0\ 0\ 0\ 0 \\ \hline 1\ 0.1\ 0\ 1\ 0\ 0\ 1 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{C}_8=0 \\ \text{C}_7=0 \\ \text{C}_6=1 \\ \text{C}_5=1 \\ \text{C}_4=0 \\ \text{C}_3=0 \\ \text{C}_2=0 \\ \text{C}_1=0 \\ \text{C}_0=0 \end{array} \\
 \begin{array}{r} 1\ 0\ 0\ 1.1\ 1\ 0\ 1\ 0\ 0\ - \\ 1\ 1\ 1\ 1.0\ 1\ 1\ 1\ 0\ 1 \\ \hline 1\ 0\ 1\ 0.0\ 1\ 0\ 1\ 1\ 1 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{C}_8=0 \\ \text{C}_7=0 \\ \text{C}_6=1 \\ \text{C}_5=1 \\ \text{C}_4=1 \\ \text{C}_3=1 \\ \text{C}_2=0 \\ \text{C}_1=0 \\ \text{C}_0=0 \end{array} \\
 \begin{array}{r} 0\ 0\ 0\ 0.0\ 1\ 0\ 1\ 0\ 1\ + \\ 0\ 0\ 1.1\ 1\ 1\ 1\ 1\ 0 \\ \hline 0\ 1\ 0.0\ 1\ 0\ 0\ 1\ 1 \end{array}
 \end{array}$$

$$\begin{array}{r}
 10.101 \times \rightarrow \begin{array}{r} 01.011 \times \\ 0.10011 \\ \hline 1\ 0\ 0\ 1\ 1 \\ 1\ 0\ 1\ 1 \\ \hline 1\ 0\ 0\ 1\ 1 \\ 1\ 0\ 0\ 1\ 1 \\ \hline 0\ 0\ 0\ 0\ 0 \\ 1\ 0\ 0\ 1\ 1 \\ \hline 0\ 1\ 1\ 0\ 1\ 0\ 0\ 0\ 1 \\ \hline 0.1\ 1\ 0\ 1\ 0\ 0\ 0\ 1 \\ \hline 1.0\ 0\ 1\ 0\ 1\ 1\ 1\ 1 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 1.011 \times \rightarrow \begin{array}{r} 0.101 \times \\ 0.1011 \\ \hline 1\ 0\ 1\ 1 \\ 1\ 0\ 1 \\ \hline 1\ 0\ 1\ 1 \\ 0\ 0\ 0\ 0 \\ \hline 1\ 0\ 1\ 1 \\ \hline 0\ 0\ 1\ 1\ 0\ 1\ 1\ 1 \\ \hline 0.0\ 1\ 1\ 0\ 1\ 1\ 1 \end{array}
 \end{array}$$

- ✓ $\frac{10.10010}{0.101}$: To unsigned (numerator) and then alignment, $a = 4$: $\frac{01.0111}{0.101} = \frac{01.0111}{00.1010} = \frac{010111}{001010} \equiv \frac{10111}{1010}$

$$\begin{array}{r}
 0001001001 \\
 1010 \overline{) 1011100000} \\
 \underline{1010} \\
 1100 \\
 \underline{1010} \\
 10000 \\
 \underline{1010} \\
 110
 \end{array}$$

Append $x = 5$ zeros: $\frac{1011100000}{1010}$
Integer Division:

$$\begin{array}{l}
 Q = 1001001, R = 110 \\
 \rightarrow Qf = 10.01001 (x = 5)
 \end{array}$$

Final result (2C): $\frac{01.01110}{1.011} = 2C(010.01001) = 101.10111$

PROBLEM 2 (10 PTS)

- Represent these numbers in Fixed Point Arithmetic (signed numbers). Use the FX format [12 4].

✓ -16.375

$+16.375 = 010000.011$

$\Rightarrow -16.375 = 11101111.1010$

✓ 32.3125

$+32.3125 = 00100000.0101$

- Complete the table for the following fixed point formats (signed numbers): (6 pts.)

Integer bits	Fractional Bits	FX Format	Range	Resolution
6	3	[9 3]	$[-2^5, 2^5 - 2^{-3}]$	2^{-3}
8	5	[13 5]	$[-2^7, 2^7 - 2^{-5}]$	2^{-5}

PROBLEM 3 (40 PTS)

- Perform the following 32-bit floating point operations. For fixed-point division, use 4 fractional bits. Truncate the result when required. Show your work: how you got the significand and the biased exponent bits of the result. Provide the 32-bit result.

✓ C1500000 + 436A0000	✓ D0A90000 - CF480000	✓ 80400000 × 7AB80000	✓ FBB80000 ÷ 49400000
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- ✓ $X = C1500000 + 436A0000$:

C1500000: 1100 0001 0101 0000 1000 0000 0000 0000

$$e + \text{bias} = 10000010 = 130 \rightarrow e = 130 - 127 = 3$$

$$C1500000 = -1.101 \times 2^3$$

Significand = 1.101

436A0000: 0100 0011 0110 1010 0000 0000 0000 0000

$$e + \text{bias} = 10000110 = 134 \rightarrow e = 134 - 127 = 7$$

$$436A0000 = 1.110101 \times 2^7$$

Significand = 1.110101

$$X = -1.101 \times 2^3 + 1.110101 \times 2^7 = -\frac{1.101}{2^4} \times 2^7 + 1.110101 \times 2^7$$

$$X = (-0.0001101 + 1.110101) \times 2^7$$

To subtract these unsigned numbers, we first convert to 2C:

$$R = 01.110101 - 0.0001101 = 01.110101 + 1.1110011$$

The result in 2C is: $R = 01.1011101$

For floating point, we need to convert to sign-and-magnitude:

$$\Rightarrow R(SM) = +1.1011101$$

* You can also do unsigned subtraction: $X = (1.110101 - 0.0001101) \times 2^7$

$$X = 1.1011101 \times 2^7, e + \text{bias} = 7 + 127 = 134 = 10000110$$

$$X = 0100 0011 0101 1101 0000 0000 0000 0000 = 435D0000$$

$$\begin{array}{r} \begin{array}{cccccccc} \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} \\ 0 & 1.1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1.1 & 1 & 1 & 0 & 0 & 1 & 1 \end{array} + \\ \hline 0 & 1.1 & 0 & 1 & 1 & 1 & 0 & 1 \\ \begin{array}{cccccccc} \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} \\ 1.1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0.0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \end{array} - \\ \hline 1.1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \end{array}$$

- ✓ $X = D0A90000 - CF480000$:

D0A90000: 1101 0000 1010 1001 0000 0000 0000 0000

$$e + \text{bias} = 10100001 = 161 \rightarrow e = 161 - 127 = 34$$

$$D0A90000 = -1.0101001 \times 2^{34}$$

Significand = 1.0101001

CF480000: 1100 1111 0100 1000 0000 0000 0000 0000

$$e + \text{bias} = 10011110 = 158 \rightarrow e = 158 - 127 = 31$$

$$CF480000 = -1.1001 \times 2^{31}$$

Significand = 1.1001

$$X = -1.0101001 \times 2^{34} + 1.1001 \times 2^{31} = -1.0101001 \times 2^{34} + \frac{1.1001}{2^3} \times 2^{34}$$

$$X = -(1.0101001 - 0.0011001) \times 2^{34} \text{ (unsigned subtraction)}$$

$$X = -1.001 \times 2^{34}, e + \text{bias} = 34 + 127 = 161 = 10100001$$

$$X = 1101 0000 1001 0000 0000 0000 0000 0000 = D0900000$$

$$\begin{array}{r} \begin{array}{cccccccc} \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} & \overline{1} \\ 1.0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0.0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \end{array} - \\ \hline 1.0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \end{array}$$

- ✓ $X = 80400000 \times 7AB80000$:

80400000: 1000 0000 0100 0000 0000 0000 0000 0000

$$e + \text{bias} = 00000000 = 0 \rightarrow \text{Denormal number} \rightarrow e = -126$$

$$80400000 = -0.1 \times 2^{-126}$$

Significand = 0.1

7AB80000: 0111 1010 1011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110101 = 245 \rightarrow e = 245 - 127 = 118$$

$$7AB80000 = 1.0111 \times 2^{118}$$

Significand = 1.0111

$$X = (-0.1 \times 2^{-126}) \times (1.0111 \times 2^{118}) = -0.10111 \times 2^{-8} = -1.0111 \times 2^{-9}$$

$$e + \text{bias} = -9 + 127 = 118 = 01110110$$

$$X = 1011 1011 0011 1000 0000 0000 0000 0000 = BB380000$$

- ✓ $X = FBB80000 \div 49400000$:

FBB80000: 1111 1011 1011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110111 = 247 \rightarrow e = 247 - 127 = 120$$

$$FBB80000 = -1.0111 \times 2^{120}$$

Significand = 1.0111

49400000: 0100 1001 0100 0000 0000 0000 0000 0000

$$e + \text{bias} = 10010010 = 146 \rightarrow e = 146 - 127 = 19$$

$$49400000 = 1.1 \times 2^{19}$$

Significand = 1.1

$$X = -\frac{1.0111 \times 2^{120}}{1.1 \times 2^{19}} = -\frac{1.0111}{1.1} \times 2^{101}$$

Alignment:

$$\frac{1.0111}{1.1} = \frac{1.0111}{1.1000} = \frac{10111}{11000}$$

Append $x = 4$ zeros: $\frac{101110000}{11000}$

Integer division
 $Q = 1111 \rightarrow Qf = 0.1111$

$$\begin{array}{r} 000001111 \\ 11000 \overline{) 101110000} \\ \underline{11000} \\ 101100 \\ \underline{11000} \\ 101000 \\ \underline{11000} \\ 100000 \\ \underline{11000} \\ 1000 \end{array}$$

Thus: $X = -0.1111 \times 2^{101} = -1.111 \times 2^{100}$
 $e + bias = 100 + 127 = 227 = 11100011$

$X = 1111 \ 0001 \ 1111 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 = F1F00000$

PROBLEM 4 (30 PTS)

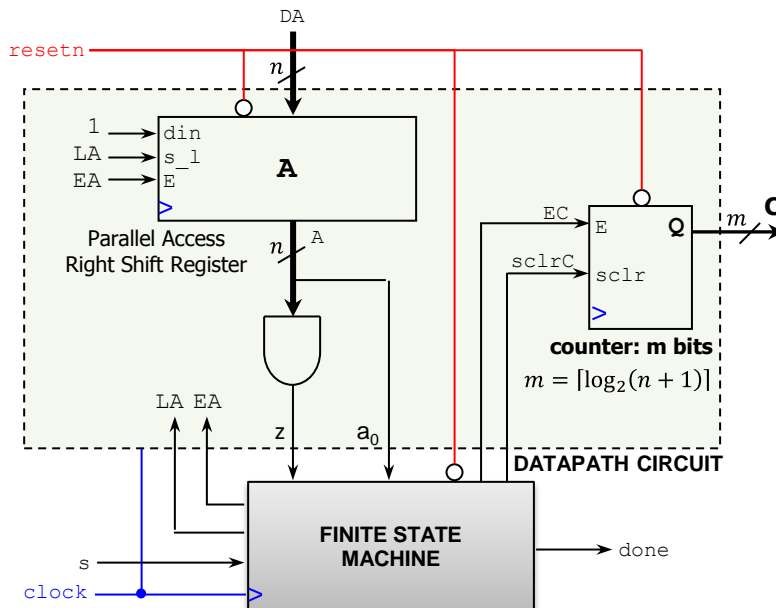
- “Counting 0’s” Circuit: It counts the number of bits in register A with a ‘0’ value. The digital system is depicted below.
- ✓ Example: for $n = 8$: if $A = 00110010$, then $C = 0101$.
- ✓ The behavior (on the clock tick) of the generic components is as follows:

m -bit counter (modulo- $n+1$): If $E=0$, the count stays.

```
if E = 1 then
  if sclr = 1 then
    Q ← 0
  else
    Q ← Q+1
  end if;
end if;
```

n -bit Parallel access shift register: If $E=0$, the output is kept.

```
if E = 1 then
  if s_l = '1' then
    Q ← D
  else
    Q ← shift in 'din' (to the right)
  end if;
end if;
```

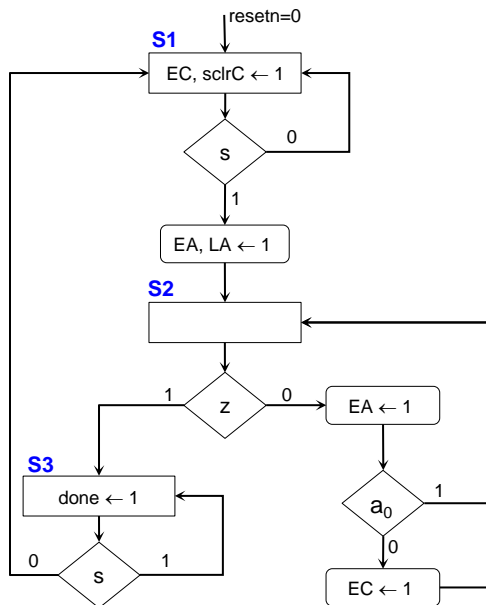


ALGORITHM

```
C ← 0
while A ≠ 11...1 (2^n-1)
  if a_0 = 0 then
    C ← C + 1
  end if
  right shift A
end while
```

- Sketch the Finite State Machine diagram (in ASM form) given the algorithm (for $n=8$, $m=4$). (18 pts.)
 - The process begins when s is asserted, at this moment we capture DA on register A . Then, we shift A one bit at a time.
 - The process ends when $A = 2^n - 1$ (i.e., when $z=1$). The signal $done$ is asserted when we finish counting.
 - As A is being shifted: we need to increase the count C every time $a_0 = 0$.

Finite State Machine:



- Complete the timing diagram ($n=8$, $m=4$). A is represented in hexadecimal format, while C is in binary format (12 pts.)

