

VLSI Architecture Design Course Final Exam Spring 1997

Student name:_____ **Student number**_____

date:_____

This exam contains three questions.

You must answer all questions.

The exam duration is 2:00 hours.

Please fill the answers ON THE EXAM forms. Additional support to your findings you should add in your exercise-book.

Good luck!

Question 1 (30%):

Given is an in order, single pipeline CPU, with Internal combined (Instruction/Data) L1 cache, which is physically tagged. The CPU internal operating frequency is 1GHz with its average Cycle Per Instruction is 2.0 (with perfect Branch prediction, perfect cache (L1)) . The CPU's external bus frequency is 100Mhz.

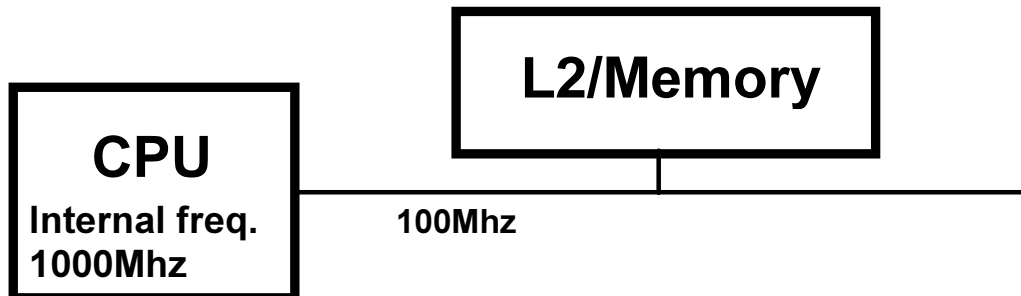


Figure 1

Branches occur on the average every 5 instructions, and the branch prediction accuracy is 90%, and the BTB hit rate is 100%. Average branch misprediction penalty is 20 internal cycles. The internal combined L1 cache hit rate (per access) is 95%. The access latency (for cache line) to L2 is 4 cycles (at a rate of 100Mhz), and the L2 hit rate is 100%. The machine's average instruction and data access is 1 access/instruction. TLB hit rate is 99% and its miss penalty is 15 cycles (at 100mhz).

Compute the average statistical CPI of this CPU:

- A. What is the CPI degradation due to wrong Branch prediction?

Answer A: _____ **Cycle/Instruction**

- B. What is the CPI degradation due to internal cache misses?

Answer B: _____ **Cycle/Instruction**

- C. What is the CPI degradation due TLB misses?

Answer C: _____ **Cycle/Instruction**

- D. What is the average statistical CPI of this CPU?

Answer D: _____ **Cycle/Instruction**

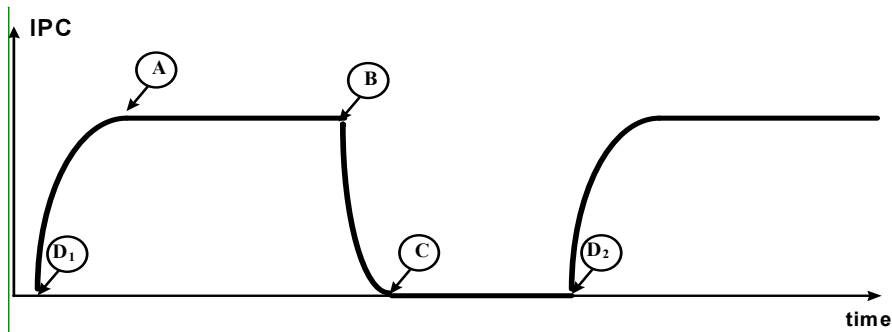
Question 2 (35%):

A CPU with >4 general purpose execution units, and on die Instruction and data caches is given. The CPU micro-architecture supports Out Of Order execution, and its internal to external frequency ratio is >20.

Following (Figure 2) is the CPU with Instruction Per Cycle (IPC) vs Time curve.

You should identify and explain what are the reasons (**what**) for the curve changes (see figure 2 points A="saturation", B="the drop", C="end of drop", and D="the ramp up") .

Answer below, for each point, what are the causes (**1=why**), the interval delay (**2=interval**), and the potential performance improvement (**3=solutions**).



Comment [OL1]:

Figure 2

A. 1. _____

A. 2. Interval from A to D₁: _____

A. 3. _____

B. 1. _____

B. 2. Interval from B to A: _____

B. 3. _____

C. 1. _____

C. 2. Interval from C to B: _____

C. 3. _____

D. 1. _____

D. 2. Interval from D₂ to C: _____

D. 3. _____

E. What is the cause of the Interval from one B to the next B?

Question 3 (35%):

This question will confront the issue of a partially decoded instruction cache in a variable instruction length machine.

Figure 3 depicts machine A with a 5 stage pipe, and 4 cycle latency to the L2 cache (i.e., Instruction L1 cache miss and L2 cache hit causes a 4 cycle penalty. Machine B has a shorter pipe, one pipe stage was reduced relative to machine A. The instruction length decode stage was removed, and the instruction length information is kept in the L1 instruction cache. On machine B, the length decode stage is between the instruction L1 cache and L2 cache, thus increasing the instruction L1 miss penalty to 5 cycles.

A) Propose a structure for the instruction L1 cache of machine B, with minimal area impact.

Answer A: _____

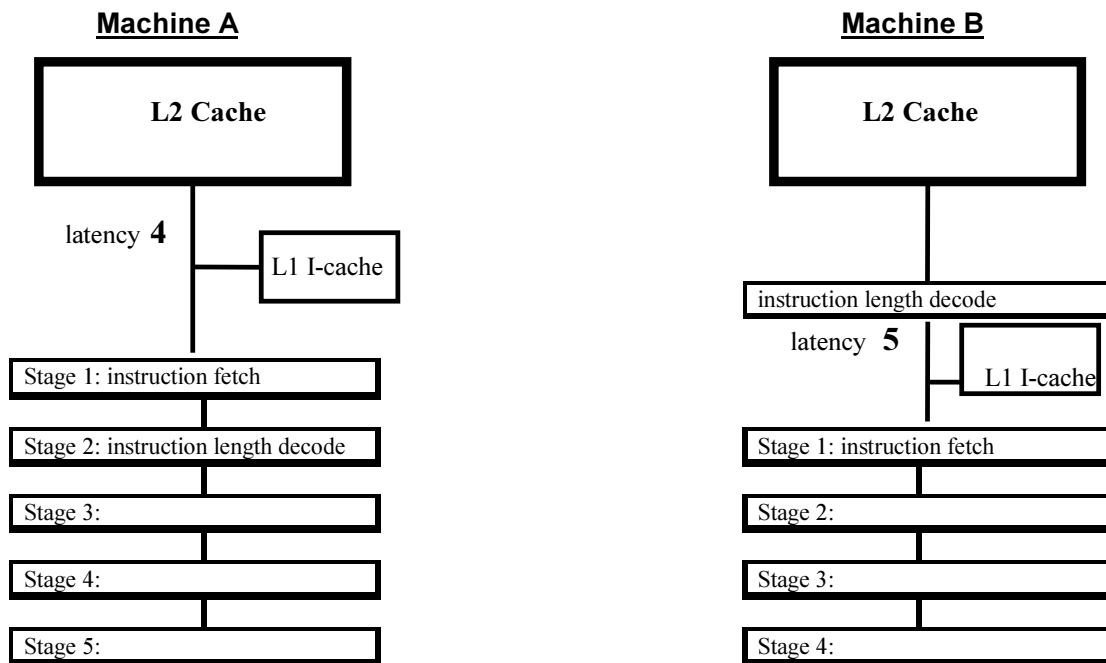


Figure 3 - Machine A has a length decode pipe stage, machine B has a partially decoded cache.

The following code section demonstrates a problem with machine B's pipeline for variable length Instruction Set Architectures (ISA). Assume value of R1 is 1 before this code begins. Assume a 16 byte instruction cache line size.

Address		Instruction	Comments
1F7h	Loop:	ADD R1, R1, #1	Increment R1 by 1
1F9h		AND R1, R1, #00000001	Mask R1 with the value 1
1FEh		JZ R1, Even	If R1=0 jump to label Even
200h		ADD R3, R3, #1	Increment R3 by 1
202h		JMP Odd	Jump to label Odd
204h	Even:	SUB R2, R2, #1	Decrement R2 by 1
206h	Odd:	JNZ R2, Loop	If R2!=0 jump to label Loop

Figure 4

B) Describe the problem of a length decoded instruction cache for variable length ISAs, as reflected by the above loop?

Answer B: _____

C) How many fetch requests are required in order to fetch all the instructions in the code above? Note the address of each instruction, and the address 200h begins on a new cache line.

Answer C: _____

Machine C (Figure 5) has a Trace cache instead of a partially decoded cache. The trace cache lines can contain up to 3 instructions per line. A trace has only one entry point - the beginning of the trace. Every backward taken branch ends the current trace (i.e. the instruction after a backward taken branch begins a new trace).

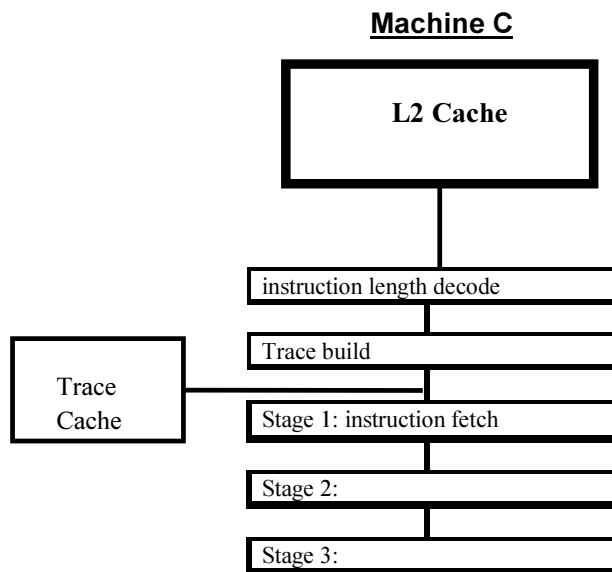


Figure 5 - Machine C has a Trace Cache

D) How would the code above (Figure 4) be mapped into the Trace Cache assuming the loop runs 10 times? Write the instructions into the TC lines below. Note the lines in which a new trace begins. Use as many lines that you need.

	Instruction1	Instruction2	Instruction3
Line1			
Line2			
Line3			
Line4			
Line5			