

Solutions - Midterm Exam

(February 16th @ 7:30 pm)

Presentation and clarity are very important! Show your procedure!

PROBLEM 1 (20 PTS)

- Compute the result of the following operations. The operands are signed fixed-point numbers. The result must be a signed fixed-point number. For the division, use $x = 5$ fractional bits.

1.0111 + 1.101001	1.010101 - 1000.0101	01.11111 + 0.10001
10.101 × 1.01101	0.111 × 1.0101	10.101 ÷ 0.101

$$\begin{array}{r}
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \end{array} + \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 1 & 1 & 0 & 1 & 0 & 0 & 1 & \end{array} \\
 \hline
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 1 & 0 & 0 & 0 & 1 & 0 & 1 & \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 \end{array} - \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \end{array} \\
 \hline
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \end{array}
 \end{array}
 \rightarrow
 \begin{array}{r}
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 \end{array} + \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 \end{array} \\
 \hline
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{array} + \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{array} \\
 \hline
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccccccc} 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \end{array}
 \end{array}$$

$$\begin{array}{r}
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 10.101 \times \\ 1.01101 \end{array} \rightarrow \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 01.011 \times \\ 0.10011 \end{array} \rightarrow \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 1 & 0 & 0 & 1 & 1 \times \\ & 1 & 0 & 1 & 1 \end{array} \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 1 & 0 & 0 & 1 & 1 \\ & 1 & 0 & 0 & 1 & 1 \\ & 0 & 0 & 0 & 0 & 0 \\ & 1 & 0 & 0 & 1 & 1 \\ \hline & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 1 \end{array} \\
 \downarrow \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 0.1 & 1 & 0 & 1 & 0 & 0 & 0 & 1 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 0.111 \times \\ 1.0101 \end{array} \rightarrow \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 0.111 \times \\ 0.1011 \end{array} \rightarrow \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 1 & 0 & 1 & 1 \times \\ & 1 & 1 & 1 \end{array} \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 1 & 0 & 1 & 1 \\ & 1 & 0 & 1 & 1 \\ & 1 & 0 & 1 & 1 \\ \hline & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 \end{array} \\
 \downarrow \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 0.1 & 0 & 0 & 1 & 1 & 0 & 1 \end{array} \\
 \downarrow \\
 \begin{array}{c} \text{1's} \\ \text{0's} \end{array} \begin{array}{cccc} 1.0 & 1 & 1 & 0 & 0 & 1 & 1 \end{array}
 \end{array}$$

- ✓ $\frac{10.1010}{0.101}$: To unsigned (numerator) and then alignment, $a = 3$: $\frac{01.0110}{0.101} = \frac{01.011}{00.101} \equiv \frac{1011}{101}$

$$\begin{array}{r}
 001000110 \\
 101 \overline{) 101100000} \\
 \underline{101} \\
 01000 \\
 \underline{101} \\
 110 \\
 \underline{101} \\
 10
 \end{array}$$

Append $x = 5$ zeros: $\frac{101100000}{101}$
Integer Division:

$$Q = 1000110, R = 10 \\
 \rightarrow Qf = 10.00110 (x = 5)$$

$$\text{Final result (2C): } \frac{10.1010}{0.101} = 2C(010.0011) = 101.1101$$

PROBLEM 2 (10 PTS)

- Represent these numbers in Fixed Point Arithmetic (signed numbers). Use the FX format [12 4].

✓ -16.125

✓ 19.25

+16.125 = 010000.001 \Rightarrow -16.125 = 11101111.1110

+19.25 = 00010011.0100

- Complete the table for the following fixed point formats (signed numbers): (6 pts.)

Integer bits	Fractional Bits	FX Format	Range	Resolution
8	6	[14 6]	$[-2^7, 2^7 - 2^{-6}]$	2^{-6}
6	4	[10 4]	$[-2^5, 2^5 - 2^{-4}]$	2^{-4}

PROBLEM 3 (40 PTS)

- Perform the following 32-bit floating point operations. For fixed-point division, use 4 fractional bits. Truncate the result when required. Show your work: how you got the significand and the biased exponent bits of the result. Provide the 32-bit result.

✓ 40D00000 + C2EA0000	✓ 50A90000 - 4F480000	✓ 80200000 × 7AB80000	✓ FB380000 ÷ 48C00000
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- ✓ $X = 40D00000 + C2EA0000$:

40D00000: 0100 0000 1101 0000 1000 0000 0000 0000

$$e + \text{bias} = 10000001 = 129 \rightarrow e = 129 - 127 = 2$$

Significand = 1.101

$$40D00000 = 1.101 \times 2^2$$

C2EA0000: 1100 0010 1110 1010 0000 0000 0000 0000

$$e + \text{bias} = 10000101 = 133 \rightarrow e = 133 - 127 = 6$$

Significand = 1.110101

$$C2EA0000 = -1.110101 \times 2^6$$

$$X = 1.101 \times 2^2 - 1.110101 \times 2^6 = + \frac{1.101}{2^4} \times 2^6 - 1.110101 \times 2^6 = (0.0001101 - 1.110101) \times 2^6$$

To subtract these numbers, we first convert to 2C:

$$R = 0.0001101 - 0.110101 = 0.0001101 + 10.001011 \text{ (2C addition)}$$

The result in 2C is: $R = 10.0100011$, $-R = 01.1011101$

* Note that you can also do unsigned subtraction: $X = -(1.110101 - 0.0001101) \times 2^6$

For floating point, we need to convert to sign-and-magnitude:

$$\Rightarrow R(SM) = -1.1011101$$

$$X = -1.1011101 \times 2^6, e + \text{bias} = 6 + 127 = 133 = 10000101$$

$$X = 1100 0010 1101 1101 0000 0000 0000 0000 = C2DD0000$$

$$\begin{array}{r} \begin{array}{cccccccc} \text{b}_{31} & \text{b}_{30} & \text{b}_{29} & \text{b}_{28} & \text{b}_{27} & \text{b}_{26} & \text{b}_{25} & \text{b}_{24} \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{array} + \\ \begin{array}{cccccccc} 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \end{array} \\ \hline \begin{array}{cccccccc} 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{array} \end{array}$$

- ✓ $X = 50A90000 - 4F480000$:

50A90000: 0101 0000 1010 1001 0000 0000 0000 0000

$$e + \text{bias} = 10100001 = 161 \rightarrow e = 161 - 127 = 34$$

Significand = 1.0101001

$$50A90000 = 1.0101001 \times 2^{34}$$

4F480000: 0100 1111 0100 1000 0000 0000 0000 0000

$$e + \text{bias} = 10011110 = 158 \rightarrow e = 158 - 127 = 31$$

Significand = 1.1001

$$4F480000 = 1.1001 \times 2^{31}$$

$$X = 1.0101001 \times 2^{34} - 1.1001 \times 2^{31} = 1.0101001 \times 2^{34} - \frac{1.1001}{2^3} \times 2^{34}$$

$$X = (1.0101001 - 0.0011001) \times 2^{34} \text{ (unsigned subtraction)}$$

$$X = 1.001 \times 2^{34}, e + \text{bias} = 34 + 127 = 161 = 10100001$$

$$X = 0101 0000 1001 0000 0000 0000 0000 0000 = 50900000$$

$$\begin{array}{r} \begin{array}{cccccccc} \text{b}_{31} & \text{b}_{30} & \text{b}_{29} & \text{b}_{28} & \text{b}_{27} & \text{b}_{26} & \text{b}_{25} & \text{b}_{24} \\ 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \end{array} - \\ \begin{array}{cccccccc} 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 \end{array} \\ \hline \begin{array}{cccccccc} 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \end{array} \end{array}$$

- ✓ $X = 80200000 \times 7AB80000$:

80200000: 1000 0000 0010 0000 0000 0000 0000 0000

$$e + \text{bias} = 00000000 = 0 \rightarrow \text{Denormal number} \rightarrow e = -126$$

Significand = 0.01

$$80200000 = -0.01 \times 2^{-126}$$

7AB80000: 0111 1010 1011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110101 = 245 \rightarrow e = 245 - 127 = 118$$

Significand = 1.0111

$$7AB80000 = 1.0111 \times 2^{118}$$

$$X = (-0.01 \times 2^{-126}) \times (1.0111 \times 2^{118}) = -0.010111 \times 2^{-8} = -1.0111 \times 2^{-10}$$

$$e + \text{bias} = -10 + 127 = 117 = 01110101$$

$$X = 1011 1010 1011 1000 0000 0000 0000 0000 = BAB80000$$

- ✓ $X = FB380000 \div 48C00000$:

FB380000: 1111 1011 0011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110110 = 246 \rightarrow e = 246 - 127 = 119$$

Significand = 1.0111

$$FB380000 = -1.0111 \times 2^{119}$$

48C00000: 0100 1000 1100 0000 0000 0000 0000 0000

$$e + \text{bias} = 10010001 = 145 \rightarrow e = 145 - 127 = 18$$

Significand = 1.1

$$48C00000 = 1.1 \times 2^{18}$$

$$X = -\frac{1.0111 \times 2^{119}}{1.1 \times 2^{18}} = -\frac{1.0111}{1.1} \times 2^{101}$$

Alignment:

$$\frac{1.0111}{1.1} = \frac{1.0111}{1.1000} = \frac{10111}{11000}$$

Append $x = 4$ zeros: $\frac{101110000}{11000}$

Integer division
 $Q = 1111 \rightarrow Qf = 0.1111$

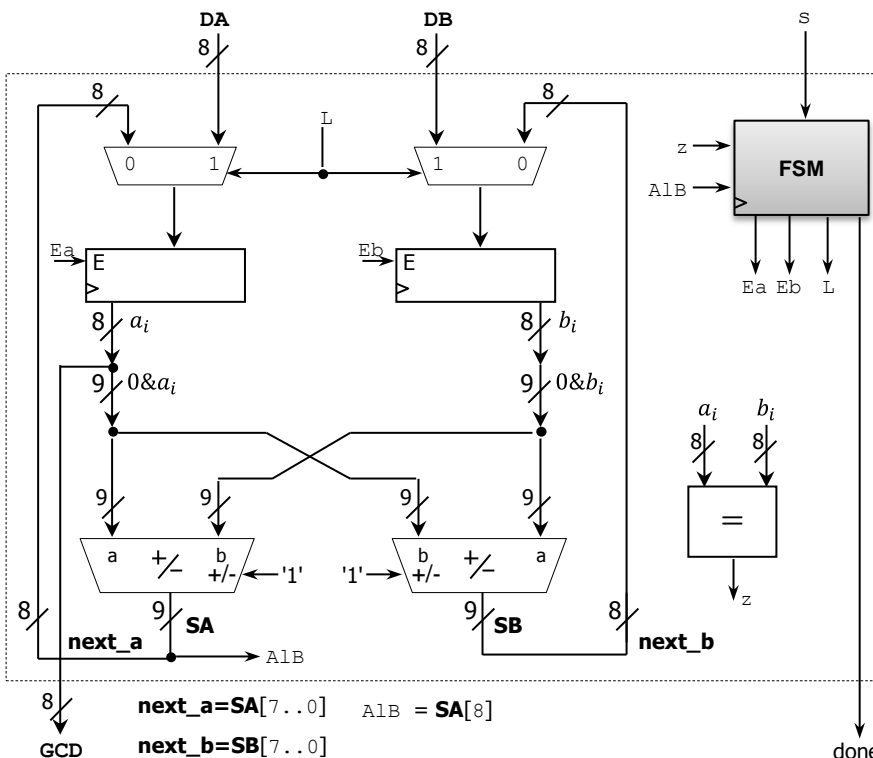
Thus: $X = -0.1111 \times 2^{101} = -1.111 \times 2^{100}$
 $e + \text{bias} = 100 + 127 = 227 = 11100011$

$X = 1111 \ 0001 \ 1111 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 = F1F00000$

PROBLEM 4 (30 PTS)

- Greatest Common Divisor (GCD):** This circuit computes the GCD of two n -bit unsigned numbers (A, B). For example:

✓ A = 216, B = 192 → GCD = 24.	✓ A = 132, B = 72 → GCD = 12.	✓ A = 169, B = 63 → GCD = 1.
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- The digital system is depicted below (FSM + Datapath) for $n = 8$. This iterative circuit is based on Euclid's GCD algorithm.
 - ✓ Input Data: DA, DB Output data: GCD



✓ $z=1$ when $a_i = b_i$, else $z=0$.

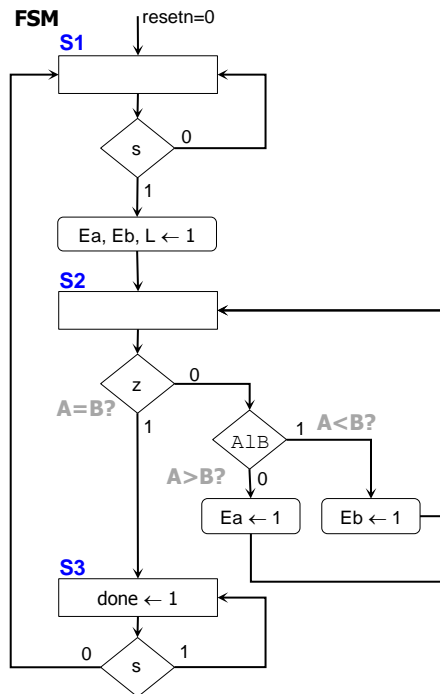
Sequential Algorithm

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a,b: unsigned integers
while a ≠ b
  if a > b
    a ← a-b
  else
    b ← b-a
end while
return a
  
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- Sketch the Finite State Machine diagram (in ASM form) given the sequential algorithm (for $n = 8$). (18 pts.)
 - The process begins when s is asserted, at this moment we capture DA and DB on register a_i and b_i (respectively). Then the process continues by updating a_i and b_i and it is concluded when $a_i = b_i$. The signal done is asserted when the result is computed and appears on output GCD .

Finite State Machine:



- Complete the timing diagram where $n = 8$. DA and DB are provided as unsigned decimals. You can provide a_i and b_i as unsigned decimals. (12 pts.)

