Recursion - 1

→ function calling itself

→ Problem is broken down into smaller sub-problems & using the solution of smaller subproblems, main problem is solved.

Question

find sum of first N natural no. using recursion.

$sum(N) = 1 + 2 + 3 + \ldots\ldots + N-1 + N$

$sum(5) = \underbrace{1 + 2 + 3 + 4}_{sum(4)} + 5$

$sum(5) = sum(4) + 5$

$sum(N) = sum(N-1) + N$

$\hookrightarrow sum(N-2) + N-1$

$\vdots$

Steps to write recursive code

1. Define the problem / function

   sum(N) = 1 + 2 + 3 + ... + N

2. Define recursive relation

   sum(N) = sum(N-1) + N

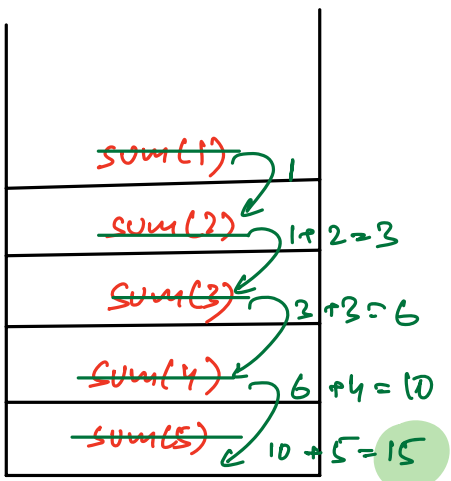
3. Define base case // simplest version of the problem

   sum(1) = 1

Code

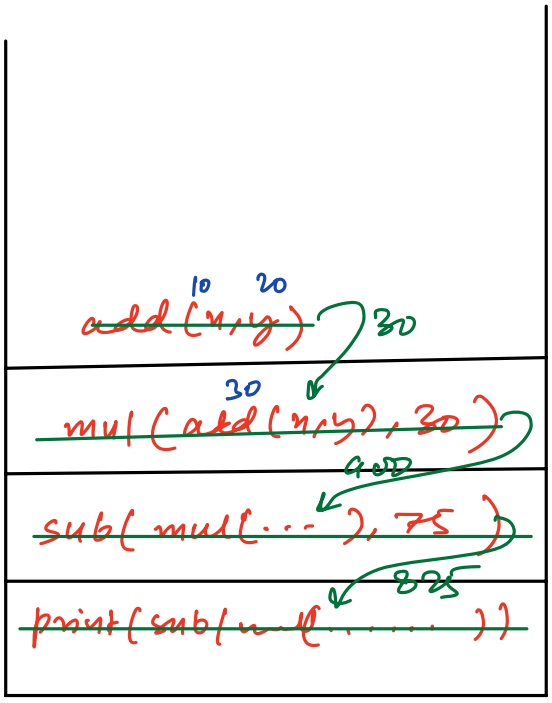
```
int sum(N) {
    //Base case
    if(N==1) return 1

    return sum(N-1) + N
}
```

sum(5)



function call tracing

```
int add( int x, int y){
    return x+y
}
int mul( int x, int y){
    return x*y
}
int sub( int x, int y){
    return x-y
}
main(){
   x=10, y=20
   print( sub( mul( add(x,y), 30), 75))
}
```



Call Stack

Output : 825

Question

Given a +ive integer N, find factorial of N.

N = 3 fact(3) = 1 * 2 * 3 = 6

N = 5 fact(5) = 1 * 2 * 3 * 4 * 5 = 120

Step 1

fact(N) = 1 * 2 * 3 * . . . . * N

Step 2

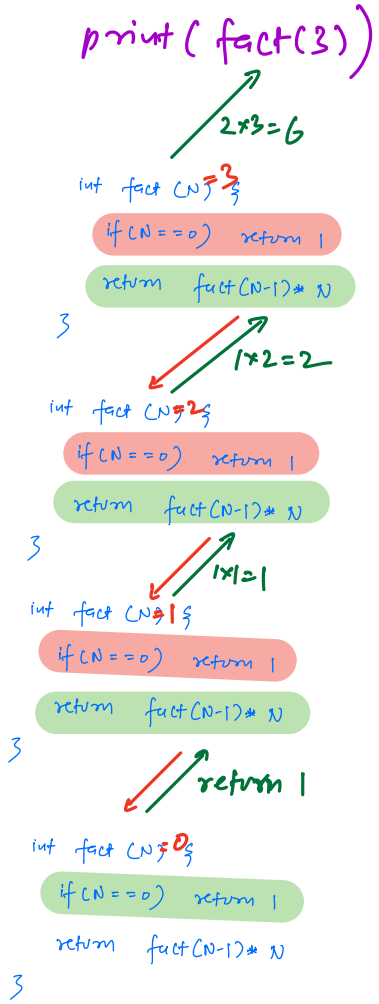
fact(N) = fact(N-1) * N

Step 3

fact(0) = 1

code

```
int fact (N) {
    if (N == 0)  return 1
    return  fact(N-1) * N
}
```



Question

Given N, print numbers from 1 to N in increasing order.

Inc(5) = 1 2 3 4 5

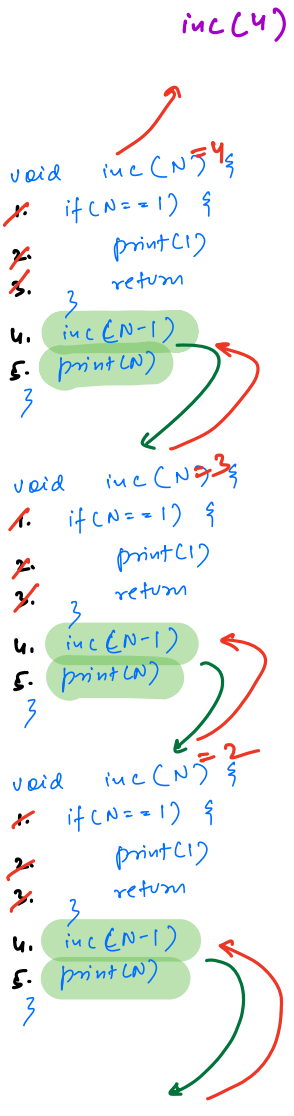
Inc(5) = Inc(4) 5

Inc(N) = Inc(N-1) N

Inc(1) = 1

Code

```
void inc(N) {
    if(N==1) {
        print(1)
        return
    }
    inc(N-1)
    print(N)
}
```



```
void inc(N) {
1.  if(N==1) {
2.      print(1)
3.      return
    }
4.  inc(N-1)
5.  print(N)
}
```

output: 1 2 3 4

Question

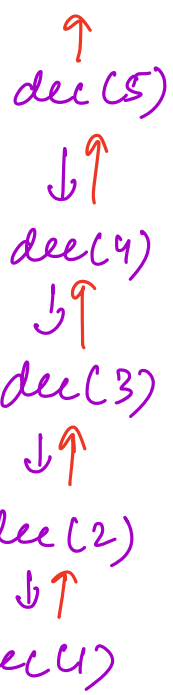
Given N, print numbers from 1 to N in decreasing order.

dec(5) = 5 4 3 2 1

dec(N) = N dec(N-1)

Code

```
void dec(N) {
    if(N==1) {
        print(1)
        return
    }
    print(N)
    dec(N-1)
}
```



output: 5 4 3 2 1

Time Complexity

O(no. of function calls * TC per function call)

Space Complexity

O(max depth of recursion tree / stack + space per function call)

for factorial

```
int fact(N) {
    if(N==0) return 1
    return fact(N-1)*N
}
```

fact(5)
↓
fact(4)
↓
fact(3)
↓
fact(2)
↓
fact(1)
↓
fact(0)

function calls = N+1

time per call = O(1)

TC = O(N)

max stack size = N

SC = O(N)

for increasing(N) / decreasing(N)

TC = O(N) SC = O(N)

Question

Given a +ive number N, find N^{th} fibonacci no.

N :	0	1	2	3	4	5	6	7
fib(N) :	0	1	1	2	3	5	8	13
						↑	↑	
						fib(5)	fib(6)	

fib(7) = fib(6) + fib(5)

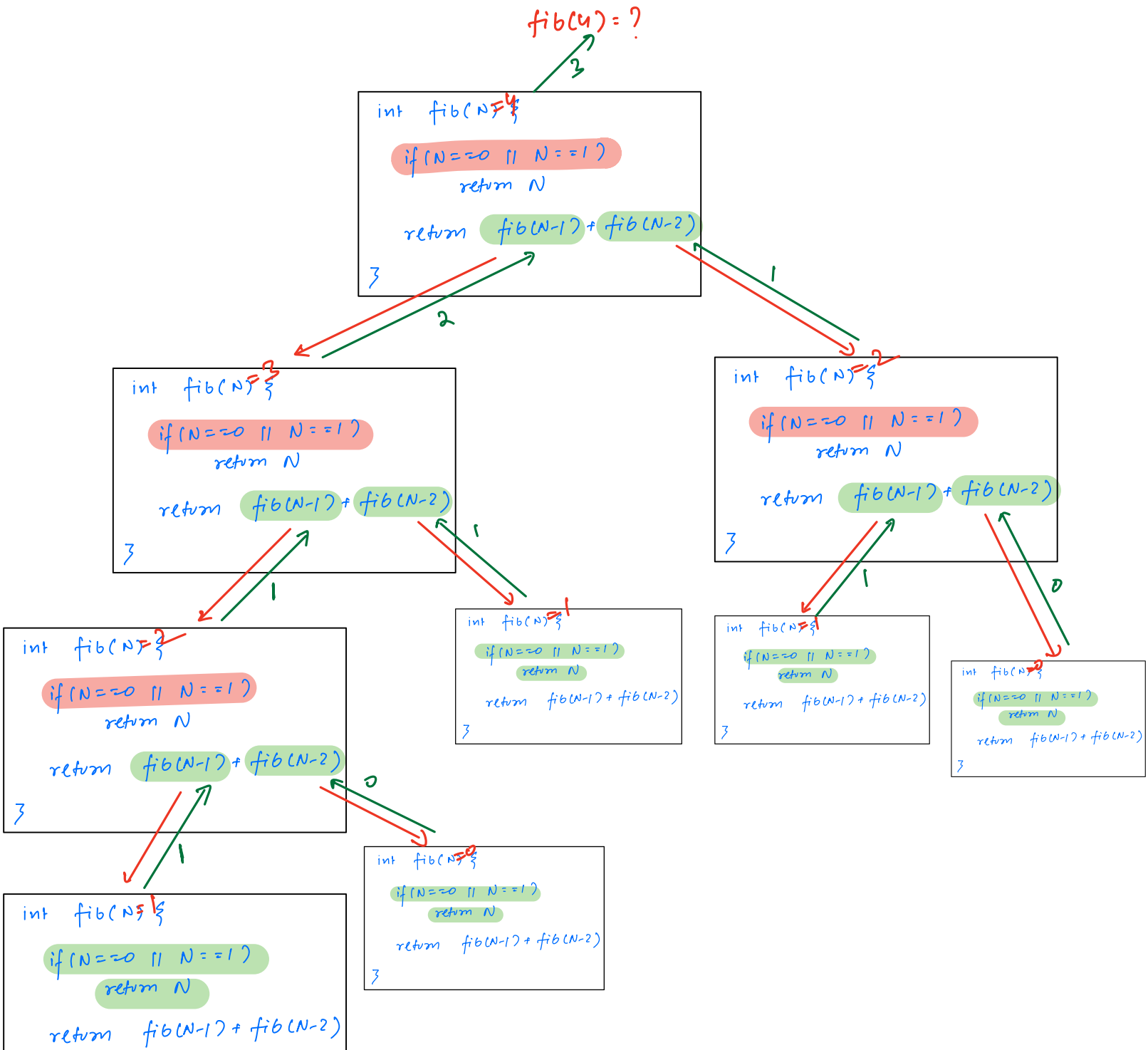
1. fib(N) → N^{th} fibonacci no.
2. fib(N) = fib(N-1) + fib(N-2)
3. fib(0) = 0
   fib(1) = 1

fib(1) = fib(0) + fib(-1)

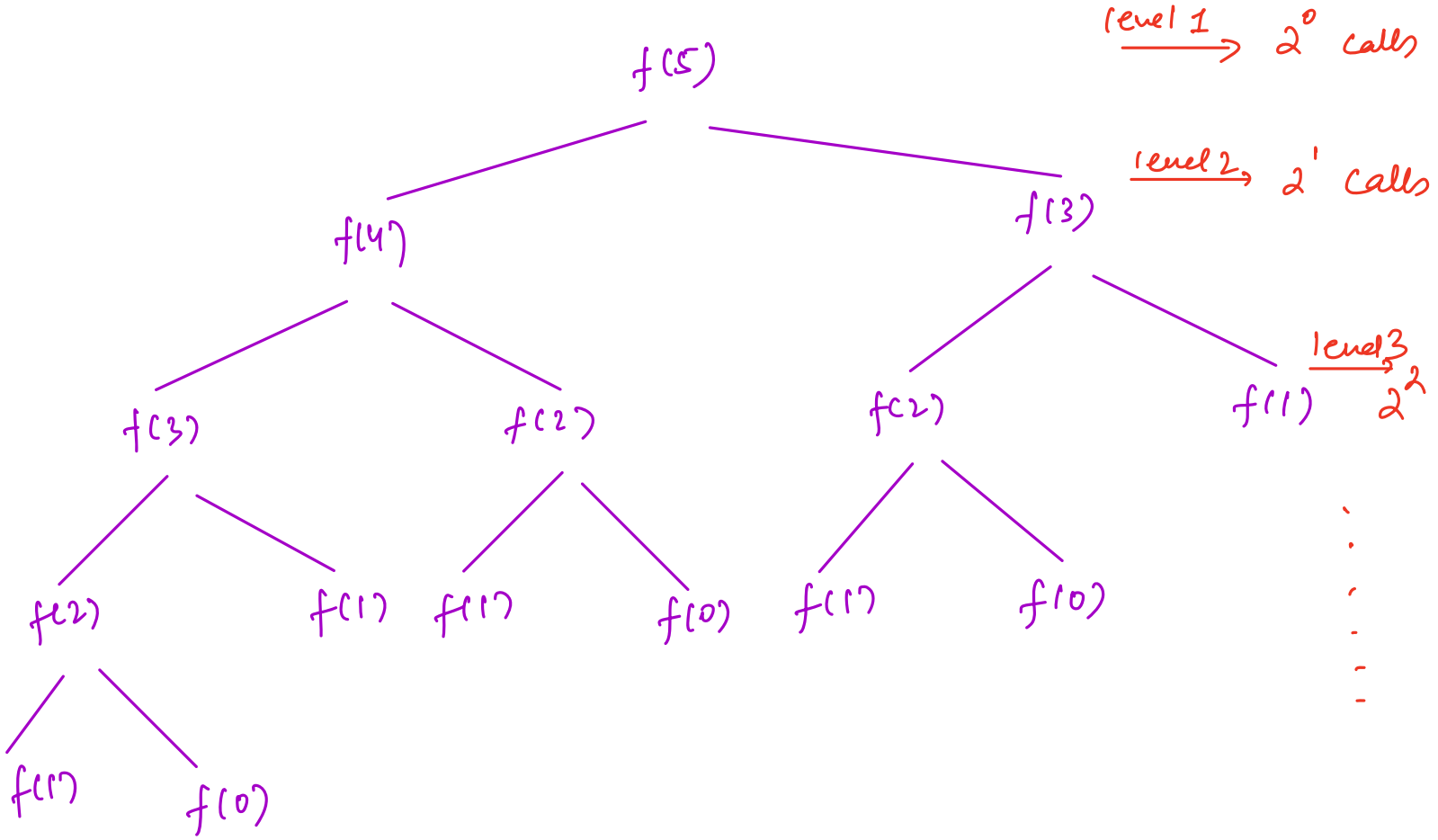
fib(2) = fib(1) + fib(0)

Code

```
int fib(N) {
    if (N == 0 || N == 1)
        return N
    return fib(N-1) + fib(N-2)
}
```



Time Complexity



there is $\sim 2^x$ calls at each level

total levels = N

total f^n calls $= 2^0 + 2^1 + 2^2 + \ldots + 2^{N-1}$

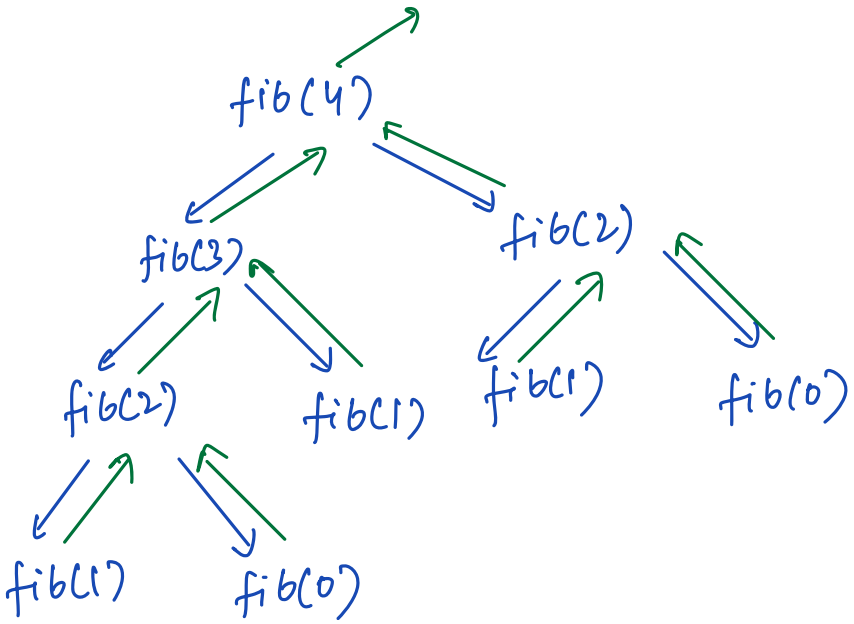
$= 2^0 \left(\frac{2^N - 1}{2 - 1} \right)$

$= 2^N - 1$

$\left[\text{sum} = a\left(\frac{r^n - 1}{r - 1} \right) \right.$

$$TC = O(2^N)$$

Space complexity



max size = 3

$$SC = O(N)$$