

Class starts @ 7:07 am

DSA: Lab Session on Heaps and Greedy

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DSA SME @ Scaler.

↳ Real time Problem Solving Situations.

↳ Solving problems.



Structure:

Pick a problem



Quick read



Try to solve it → (Doubt, app dis)



Discussion of Solutions

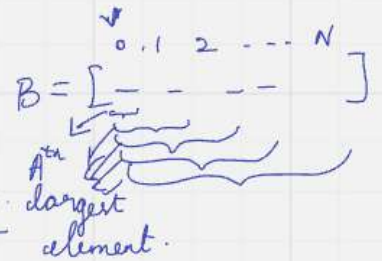
Problems.

1. A^{th} largest element
2. Merge K sorted linked lists
3. Flipkart's bla bla.
4. Distribute candies.

7:33 am

Question 1) A^{th} Largest Element

Given an array B of length N , find the A^{th} largest element of every window that starts from index 0, and ends at all indices.

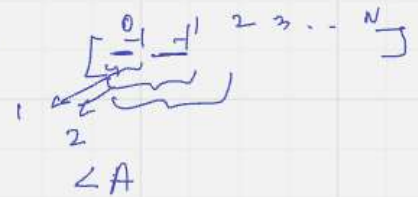


(i) $s = 0$ ✓

(ii) $0 \leq e < N$

(iii) If window size is less than A , return -1

$A = 3$



Example 1)

$res = [A^{\text{th}} \quad A^{\text{th}} \quad A^{\text{th}} \quad A^{\text{th}}]$

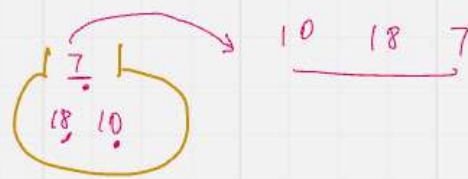
$A = 3$ ✓ 3rd

$B = [\text{---} \text{---} \text{---} \text{---} \text{---} \text{---}]$
 $0 \quad 1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6$
 $B = [10 \quad 18 \quad 7 \quad 5 \quad 16 \quad 19 \quad 3]$
 $res = [-1 \quad -1 \quad 7 \quad 7 \quad 10 \quad 16 \quad 16]$

$A = 3$

HINT 1 ★

Min Heap



HINT 2 ★

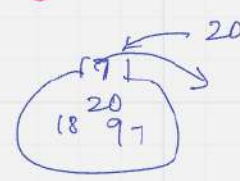
maintain the size of the heap at A .

Pseudocode)

```

int[] solve (int[] B, int A) {
    int N = B.length;
    int[] res = new int[N];
    res.fill(-1); // N

    PriorityQueue<Integer> q = new PriorityQueue<>();
    for (int i = 0; i < N; i++) {
        q.add(B[i]);
        if (q.size() < A) continue;
        if (q.size() > A) q.remove();
        res[i] = q.peek();
    }
    return res;
}
    
```

$-1 \quad -1 \quad -1$
 $A = 3$
 $2 \log(N)$
 $s = -A$


$TC \Rightarrow O(N \log N)$
 $SC \Rightarrow \cancel{O(N)} \quad O(K)$

Try to not use the auxiliary array.

because the question has demanded the auxiliary array.

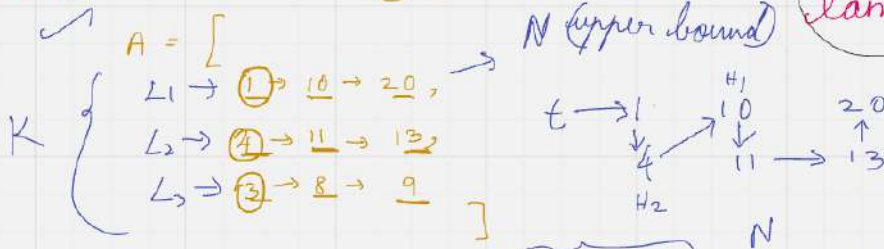
Question 2) Merge K sorted linked lists

Given K sorted linked list, merge them together as a single sorted linked list.

8:19 am

Example 1)

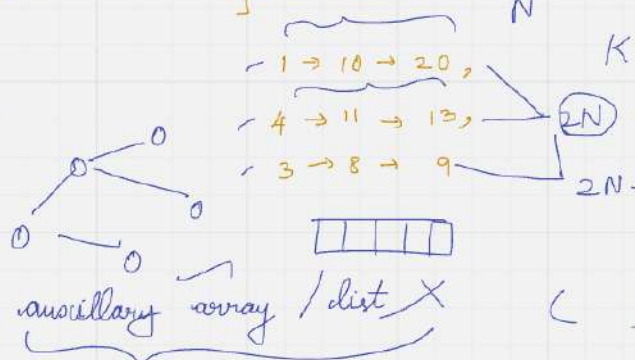
$NK (\log K)$



comparator \rightarrow \rightarrow (Pair)
 object / class \rightarrow \rightarrow (Pair)
 Parse Int ("234") \rightarrow
 slicing \rightarrow
 accessing DS \rightarrow
 lambda expressions.

HINT 1: min Heap

HINT 2 ✱:
 $\min(P_1, P_2)$
 $\min(P_1, P_2, P_3, \dots, P_k)$



8:20 am

Pseudocode)

```
ListNode solve(List<ListNode> nodes) {
```

```
    PriorityQueue<ListNode> q = new PriorityQueue<>((a,b) -> a.val - b.val);
```

```
    ListNode dummy = new ListNode(-1);
```

```
    ListNode temp = dummy;
```

```
    while (!q.isEmpty()) {
```

```
        ListNode smallNode = q.remove();
```

```
        temp.next = smallNode;
```

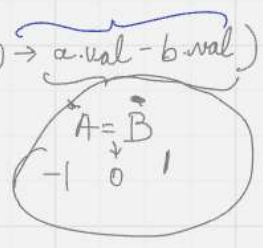
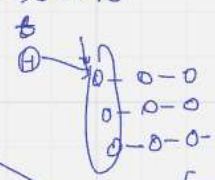
```
        temp = temp.next;
```

```
        if (smallNode.next != null) q.add(smallNode.next);
```

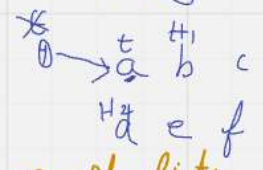
```
    }
```

```
    return dummy.next;
```

```
}
```



{add all first elements in the pq}



$K \rightarrow$ no. of lists

$N \rightarrow$ length of each list

TC $\Rightarrow O(NK \log K)$
 SC $\Rightarrow O(K)$

Break till 8:45

9:03

Question 3) Flipkart's Challenge.

You are given with 2 arrays A and B of size N. Here A[i] is the expiry time of a particular product, and B[i] is the profit margin of that product if sold.

Your task is to maximize the profit by selling the products before the expiry time.

- (i) Selling a product takes 1 unit of time.
- (ii) A product cannot be sold if the current time \geq expiry time.

$$P_0, P_1, P_2, \dots, P_N$$

$$A = [ET_0, ET_1, ET_2, \dots, ET_N]$$

$$B = [PM_0, PM_1, PM_2, \dots, PM_N]$$

$$ct \geq ET$$

$$P_0$$

$$ET = 2$$

$$PM = 1000$$

$$ct = 2, 3, 4, \dots$$

Example 1)

$$A = [1, 3, 3, 3, 2]$$

$$B = [5, 6, 9, 3, 1]$$

$$ct \geq 3$$

$$ct = 0, 1, 2, 3$$

$$P = 0 + 5 + 6 + 9$$

$$€ 20 \Rightarrow \text{max?}$$

HINT 1 ★ : Sorting

HINT 2 ★ : XDP \approx LIS ~~DP~~

HINT 3 ★ : Sort based on expiry time

HINT 4 ★ : minHeap

Pseudocode)

```
1 public class Solution {
2     class Product {
3         int expiry;
4         int profit;
5
6         Product(int expiry, int profit) {
7             this.expiry = expiry;
8             this.profit = profit;
9         }
10    }
11    int mod = (int)1e9+7;
12
13    public int solve(int[] A, int[] B) {
14        int N = A.length;
15        List<Product> products = new ArrayList<>();
16
17        for (int i = 0; i < N; i++) {
18            products.add(new Product(A[i], B[i]));
19        }
20
21        Collections.sort(products, (a, b) -> a.expiry - b.expiry);
22
23        int totalProfit = 0;
24        int curTime = 0;
25        PriorityQueue<Integer> pastProfits = new PriorityQueue<>();
26
27        for (Product product : products) {
28            if (product.expiry <= curTime) {
29                if (product.profit <= pastProfits.peek()) continue;
30                totalProfit -= pastProfits.remove();
31                curTime += product.profit;
32            }
33
34            totalProfit += product.profit;
35            totalProfit %= mod;
36            curTime += product.profit;
37            pastProfits.add(product.profit);
38        }
39
40        return totalProfit;
41    }
42 }
43
44 }
```

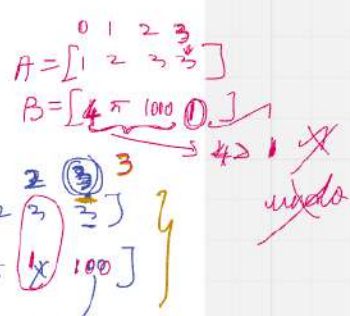
} obj



int compare(-, -) {

N log N

(N log N)



TC $\Rightarrow O(N \log N)$
SC $\Rightarrow O(N)$

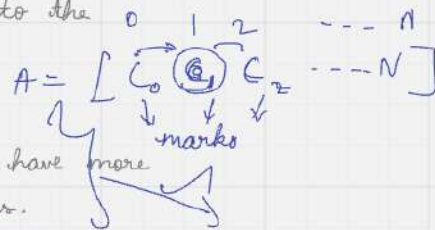
Question 4) Distribute Candies

N children are standing in a line.
Each child is assigned with a rating.
Your task is to distribute candies to the children in such a way that:

- (i) Every child ~~has~~ ^{must have} at least 1 candy.
- (ii) Children with higher ratings, should have more candies than their immediate neighbours.

Find the minimum number of candies you have to distribute to satisfy the above criteria.

$C_1 > C_0$ $C_1 > C_2$
* more cand...



Example 1)

$$A = \begin{bmatrix} 1 & 5 & 2 & 1 \end{bmatrix}$$

$$res = \begin{bmatrix} 1 & * & * & 1 \\ & 2 & 2 & \\ & 3 & & \end{bmatrix}$$

$$1 + 3 + 2 + 1 \Rightarrow 7$$

$$res = \begin{bmatrix} - & - & - & - \\ - & - & - & - \end{bmatrix}$$

$$\begin{bmatrix} 1 & 5 & 2 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & * & * & 1 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 2 & 3 \end{bmatrix}$$

Pseudocode)

int solve (int [] A) { ✓

```
int N = A.length;
```

```
int[] res = new int[N];
```

```
over-fill(1);
```

left \rightarrow right

```
for (int i = 1; i < N; i++) {
```

if ($A[i] > A[i-1]$)

$res[i] = res[i-1] + 1 \rightarrow$ stringent (greedy)

right \rightarrow left

```
for (int i = N-2; i > -1; i--) {
```

If $(A[i] > A[i+1])$ {

4. $(res[i] \leq res[i+1])$ $res[i] = res[i+1] + 1$ → stringent

$res[i] \Rightarrow candies$

```
return sum(res);
```

$$\text{int}[\][\] \Rightarrow \left[\begin{array}{cc} \text{EF} & \text{PM} \\ \text{EF} & \text{PM} \end{array} \right]$$
$$A = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \\ 3 & 3 & 3 \end{bmatrix}$$

$B = \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}}_{I_3}$

L J

↓ ↓ ↓ ↓

۱۲۳۴

$$A[i] > A[i+1]$$
$$res[i] = res[i+1] +$$

~~res[i] > res[i+1]~~

$$TC = O(N) //$$
$$S = O(N)$$

$[1 \quad 3 \quad 2 \quad 5]$

$\text{if } (i == N-1) \text{ skip}$

res $[1 \quad * \quad 1 \quad *]$

\times immediate neighbors.