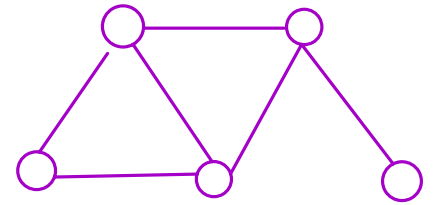


Graphs 1: Intro, DFS & Cycle Detection

Graph \rightarrow collection of nodes & edges.

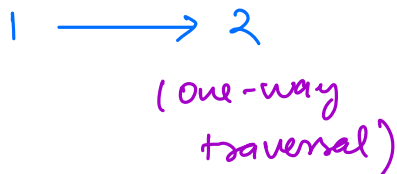


5 nodes
&
6 edges

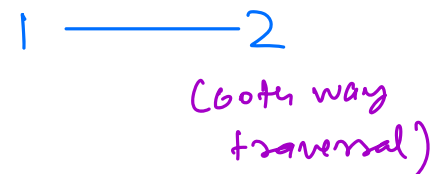
- \rightarrow Google maps
- \rightarrow Internet
- \rightarrow Social media

Properties

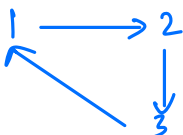
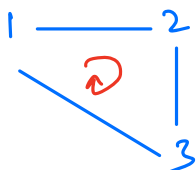
1. Directed (Digraph)



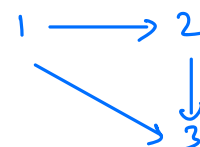
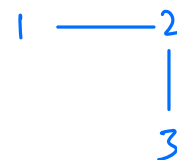
Undirected



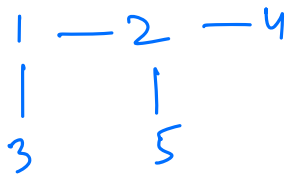
2. Cyclic



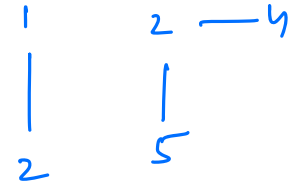
Acyclic



3. Connected



Dis-connected



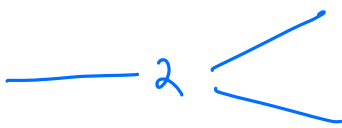
4. Weighted



Un-weighted

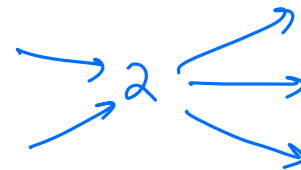


5. Degree



$$\text{degree}(2) = 3$$

total connected
edges



$$\text{in}(2) = 2$$

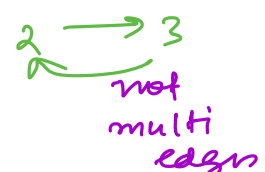
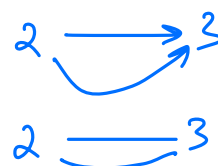
total incoming edges

$$\text{out}(2) = 3$$

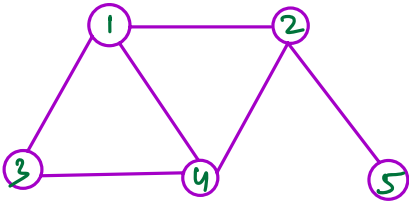
total outgoing edges

6. Simple graph

A connected graph without self-loops & multi-edges.



How to store graph



1. Adjacency Matrix

$mat[i][j] = 1$, $i \rightarrow j$
 $= 0$, else

	1	2	3	4	5
1	0	1	1	1	0
2	1	0	0	1	1
3	1	0	0	1	0
4	1	1	1	0	0
5	0	1	0	0	0

mat

// input $\rightarrow A[n][2]$

$A[i][0]$ & $A[i][1]$ represent the edge

int mat[N][N] = {0}

$V \rightarrow$ node/vertex

$E \rightarrow$ edge

for ($i=0$; $i < A.size()$; $i++$) {

u = $A[i][0]$

v = $A[i][1]$

$A = [[1,2], [1,3], [2,3], [1,5], \dots]$

mat[u][v] = 1;

mat[v][u] = 1; // not needed for Digraph

}

$$SC = O(V^2) \text{ or } O(N^2)$$

Advantage : Easy to update edges

Disadvantage : Space wastage if edges are less

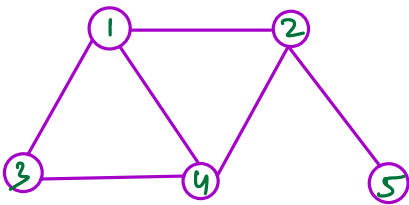
if $V = 10^5$

$$SC = O(V^2) = O(10^{10})$$

not possible to
store graph

2. Adjacency list

$graph[i] \rightarrow$ list of nodes connected to i .



$$graph[1] \rightarrow \{2, 3, 4\}$$

$$graph[2] \rightarrow \{1, 4, 5\}$$

$$graph[3] \rightarrow \{1, 4\}$$

$$graph[4] \rightarrow \{1, 2, 3\}$$

$$graph[5] \rightarrow \{2\}$$

```
list<int> graph[V];
```

```
for (i=0; i < A.size(); i++) {
```

```
    u = A[i][0]
```

```
    v = A[i][1]
```

```
    graph[u].add(v);
```

```
    graph[v].add(u);
```

```
}
```

$$SC = O(V + E)$$

#nodes ↑
#edges ↗

$$\left. \begin{array}{l} V \leq 10^5 \\ E \leq 10^5 \end{array} \right\} \text{store in adjacency list}$$

if each node is connected to every other node \Rightarrow

each node will have $V-1$ edges.

$$\begin{aligned} \text{total edges} &= V \times (V-1) \\ &= O(V^2) \end{aligned}$$

$$E = O(V^2)$$

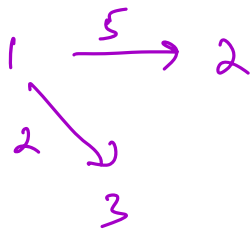
for weighted graph - how to store

Adjacency Matrix

$$\begin{aligned} \text{mat}[i][j] &= w & , i \xrightarrow{w} j \\ &= 0 & , \text{else} \end{aligned}$$

Adjacency list

$\text{graph}[i]$ = list of (node, weight) connected to i

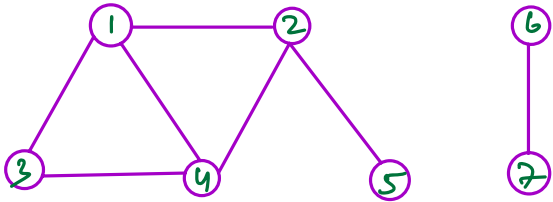


$$\text{graph}[1] = \{ (2, 5), (3, 2), \dots \}$$

Graph traversals

pre order
in order
post order } dfs
level order → bfs

1. Depth first Search (DFS)



1 → 2, 3, 4
2 → 1, 4, 5
3 → 1, 4
4 → 1, 2, 3
5 → 2
6 → 7
7 → 6

visit[V]
if, visit[i] = false

for (i = 0 to V-1) { → O(V)

if (visit[i] == false)
dfs(i)

}

void dfs(u) {

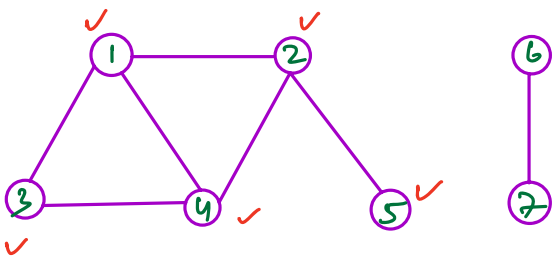
visit[u] = true;
print(u)

for (int v : graph[u]) {

if (!visit[v])
dfs(v)

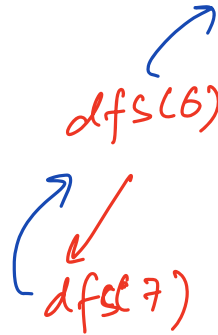
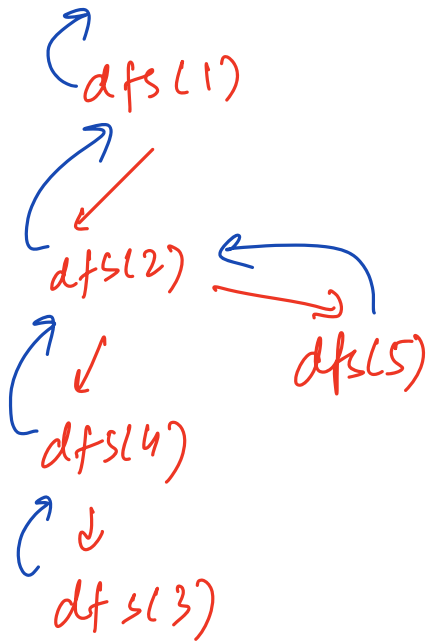
}

}



T	F	T	F	T	F	T
1	2	3	4	5	6	7

order $\rightarrow 1, 2, 4, 3, 5, 6, 7$



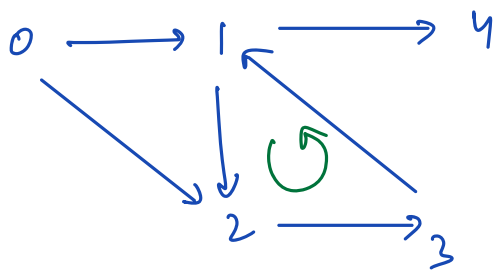
$$TC = O(V + E)$$

$$SC = O(V)$$

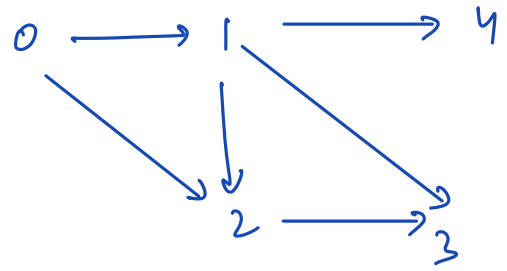
↑ visit array + recursion stack

Question

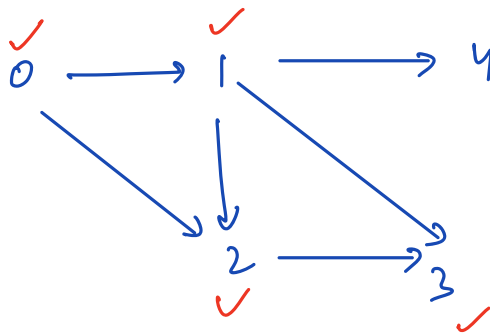
Check if the directed simple graph has a cycle.



ans = true



ans = false



dfs(0)



dfs(1)

→ cycle is detected since 3 is already visited

↓
dfs(2)

which is wrong

↓
dfs(3)

so we can't

use visit array to detect cycle.

Solution : cycle is present if we are visiting a node again which already is in current path.

visit[V] , path[V]

for i, visit[i] = false

for i, path[i] = false

for (i = 0 to V-1) { $\rightarrow O(V)$

if (visit[i] == false) {

if (dfs(i) == true) {

print("cycle is detected");

break;

}

}

}

bool dfs(u) {

visit[u] = true;

path[u] = true;

for (int v : graph[u]) {

if (path[v] == true)

return true;

if (!visit[v]) {

if (dfs(v) == true) {

return true;

}

}

}

path[u] = false;

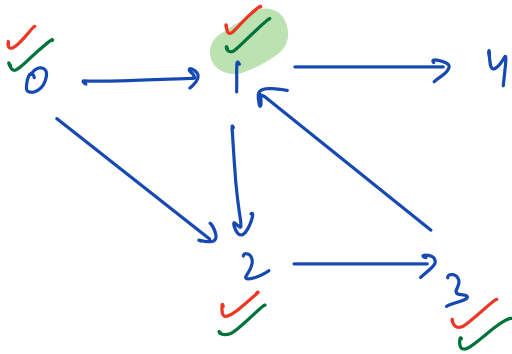
return false

}

$TC = O(V+E)$

$SC = O(V)$

✓ → visit
✓ → path



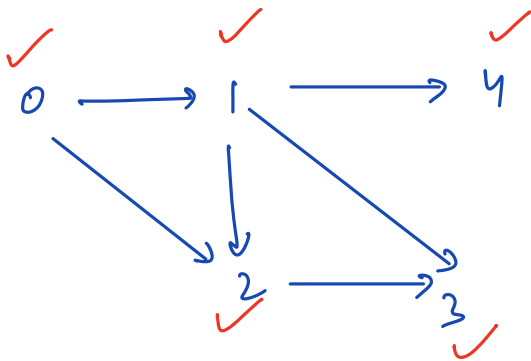
0 → 1, 2

1 → 2, 4

2 → 3

3 → 1

4 →



0 → 1, 2

1 → 2, 4, 3

2 → 3

3 →

4 →

NO CYCLE