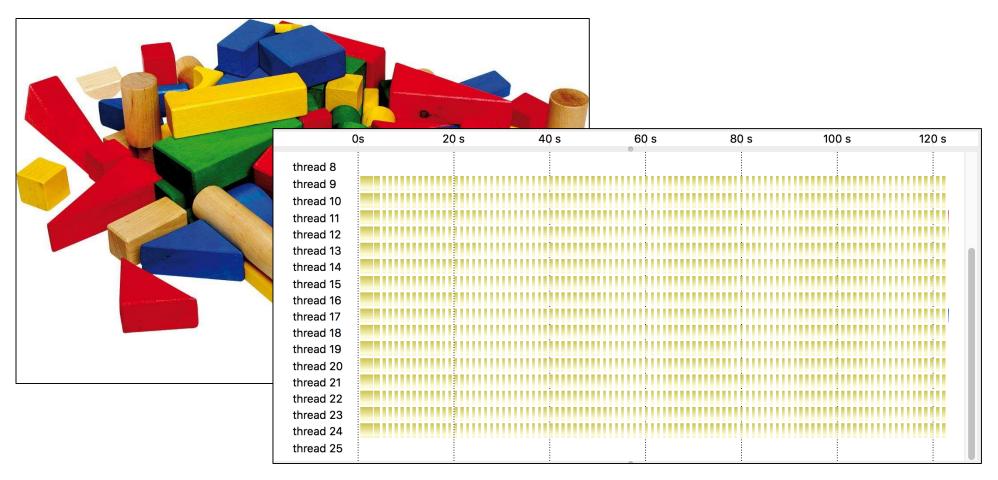
HPX

The C++ Standards Library for Concurrency and Parallelism

Hartmut Kaiser (hkaiser@cct.lsu.edu)

The Application Problems



HPX – A General Purpose Runtime System

- The C++ Standards Library for Concurrency and Parallelism
- Exposes a coherent and uniform, C++ standards-conforming API for ease of programming parallel, distributed, and heterogeneous applications.
 - Enables to write fully asynchronous code using hundreds of millions of threads.
 - · Provides unified syntax and semantics for local and remote operations.
 - Enables seamless data parallelism orthogonally to task-based parallelism
- HPX represents an innovative mixture of
 - · A global system-wide address space (AGAS Active Global Address Space)
 - Fine grain parallelism and lightweight synchronization
 - · Combined with implicit, work queue based, message driven computation
 - Support for hardware accelerators

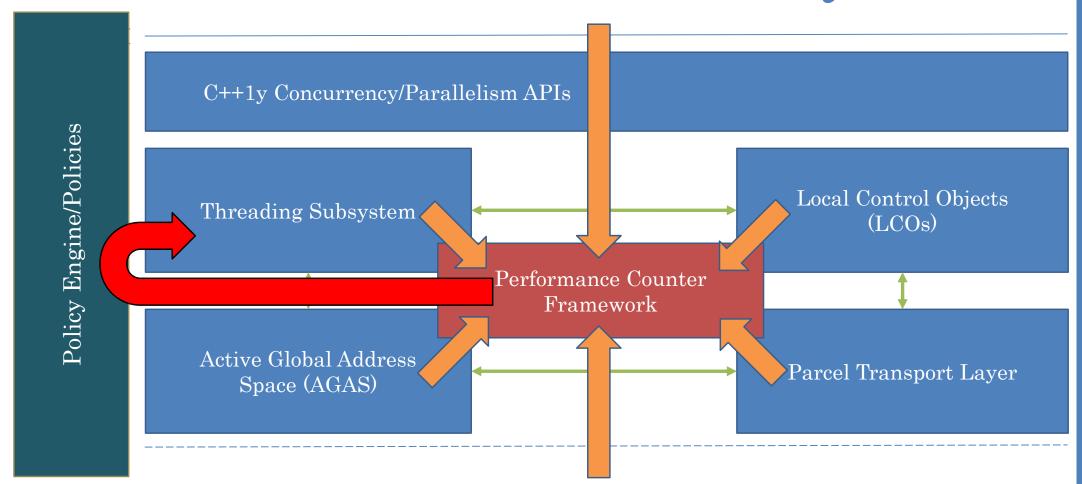


HPX – A C++ Standard Library

- Widely portable
 - Platforms: x86/64, Xeon/Phi, ARM 32/64, Power, BlueGene/Q
 - Operating systems: Linux, Windows, Android, OS/X
- Well integrated with compiler's C++ Standard libraries
- Enables writing applications which out-perform and out-scale existing applications based on OpenMP/MPI
 - http://stellar-group.org/libraries/hpx
 - https://github.com/STEllAR-GROUP/hpx/
- Is published under Boost license and has an open, active, and thriving developer community.
- · Can be used as a platform for research and experimentation



HPX – A C++ Standard Library



Programming Model

- Focus on the logical composition of data processing, rather than the physical orchestration of parallel computation
- Provide useful abstractions that shield programmer from low-level details of parallel and distributed processing
- Centered around data dependencies not communication patterns
- Make data dependencies explicit to system thus allows for auto-magic parallelization
- Basis for various types of higher level parallelism, such as iterative, fork-join, continuation-style, asynchronous, data-parallelism
- Enable runtime-based adaptivity while applying application-defined policies



Programming Model

- The consequent application of the Concept of Futures
 - · Make data dependencies explicit and visible to the runtime
- Implicit and explicit asynchrony
 - Transparently hide communication and other latencies
 - Makes over-subscription manageable
 - Uniform API for local and remote operation
 - · Local operation: create new thread
 - · Remote operation: send parcel (active message), create thread on behalf of sender
- Work-stealing scheduler
 - Inherently multi-threaded environment
 - · Supports millions of concurrently active threads, minimal thread overhead
 - Enables transparent load balancing of work across all execution resources inside a locality
- API is fully conforming with C++11/C++17 and ongoing standardization efforts



HPX – The API

• As close as possible to C++11/14/17 standard library, where appropriate, for instance

• std::thread

• std∷mutex

• std::future

• std∷async

• std::bind

• std::function

std∷tuple

• std∷any

• std∷cout

• std::for_each(par, ...), etc.

• std::experimental::task_block

hpx::thread

hpx∷mutex

hpx::future (including N4538, 'Concurrency TS')

hpx::async (including N3632)

hpx∷bind

hpx∷function

hpx∷tuple

hpx::any (N3508)

hpx∷cout

hpx::parallel::for_each (C++17)

hpx::parallel::task_block (Parallelism TS 2)



Control Model: How is parallelism achieved?

- Explicit parallelism:
 - Low-level: thread
 - Middle-level: async(), dataflow(), future::then()
- Higher-level constructs
 - Parallel algorithms (parallel::for_each and friends, fork-join parallelism for homogeneous tasks)
 - Asynchronous algorithms (alleviates bad effect of fork/join)
 - Task-block (fork-join parallelism of heterogeneous tasks)
 - Asynchronous task-blocks
 - Continuation-style parallelism based on composing futures (task-based parallelism)
 - Data-parallelism on accelerator architectures (vector-ops, GPUs)
 - Same code used for CPU and accelerators



Parallel Algorithms (C++17)

| adjacent difference | adjacent_find | all_of | any_of |
|---------------------|--------------------------|-----------------------|-------------------------|
| copy | copy_if | copy_n | count |
| count_if | equal | exclusive_scan | fill |
| fill_n | find | find_end | find_first_of |
| find_if | find_if_not | for_each | for_each_n |
| generate | generate_n | includes | inclusive_scan |
| inner product | inplace_merge | is_heap | is_heap_until |
| is_partitioned | is_sorted | is_sorted_until | lexicographical_compare |
| max_element | merge | min_element | minmax_element |
| mismatch | move | none_of | nth_element |
| partial_sort | partial_sort_copy | partition | partition_copy |
| reduce | remove | remove_copy | remove_copy_if |
| remove_if | replace | replace_copy | replace_copy_if |
| replace_if | reverse | reverse_copy | rotate |
| rotate_copy | search | search_n | set_difference |
| set_intersection | set_symmetric_difference | set_union | sort |
| stable_partition | stable_sort | swap_ranges | transform |
| uninitialized_copy | uninitialized_copy_n | $uninitialized_fill$ | uninitialized_fill_n |
| unique | unique_copy | | |



STREAM Benchmark

```
std::vector<double> a, b, c;
                              // data
// ... init data
auto a_begin = a.begin(), a_end = a.end(), b_begin = b.begin() ...;
// STREAM benchmark
parallel::copy(par, a_begin, a_end, c_begin);
                                                                  // copy step: c = a
                                                                  // scale step: b = k * c
parallel::transform(par, c begin, c end, b begin,
    [](double val) { return 3.0 * val; });
parallel::transform(par, a_begin, a_end, b_begin, b_end, c_begin, // add two arrays: c = a + b
    [](double val1, double val2) { return val1 + val2; });
parallel::transform(par, b_begin, b_end, c_begin, c_end, a_begin, // triad step: a = b + k * c
    [](double val1, double val2) { return val1 + 3.0 * val2; });
```

Dot-product: Vectorization

```
std::vector<float> data1 = {...};
std::vector<float> data2 = {...};
double p = parallel::inner product(
                                             // parallel and vectorized execution
   datapar,
   std::begin(data1), std::end(data1),
   std::begin(data2),
   0.0f,
   [](auto t1, auto t2) { return t1 + t2; }, // std::plus<>()
   [](auto t1, auto t2) { return t1 * t2; } // std::multiplies<>()
```

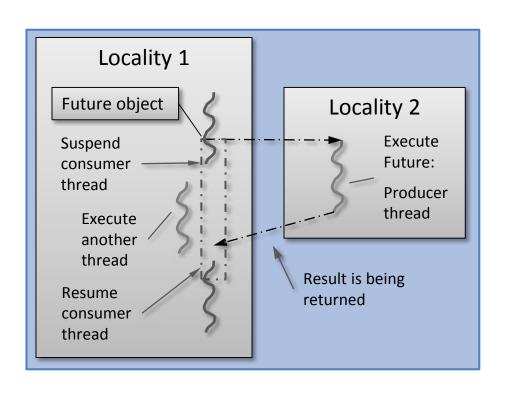
Control Model: How is synchronization expressed?

- Low-level (thread-level) synchronization: mutex, condition_variable, etc.
- Replace (global) barriers with finer-grain synchronization (synchronize of a 'as-need-basis')
 - Wait only for immediately necessary dependencies, forward progress as much as possible
- Many APIs hand out a future representing the result
 - Parallel and sequential composition of futures (future::then(), when_all(), etc.)
 - Orchestration of parallelism through launching and synchronizing with asynchronous tasks
- Synchronization primitives: barrier, latch, semaphore, channel,
 - Synchronize using futures



Synchonization with Futures

• A future is an object representing a result which has not been calculated yet



- Enables transparent synchronization with producer
- Hides notion of dealing with threads
- Makes asynchrony manageable
- Allows for composition of several asynchronous operations
- (Turns concurrency into parallelism)

What is a (the) Future?

• Many ways to get hold of a future, simplest way is to use (std) async:

```
int universal_answer() { return 42; }

void deep_thought()
{
   future<int> promised_answer = async(&universal_answer);

   // do other things for 7.5 million years

   cout << promised_answer.get() << endl; // prints 42
}</pre>
```

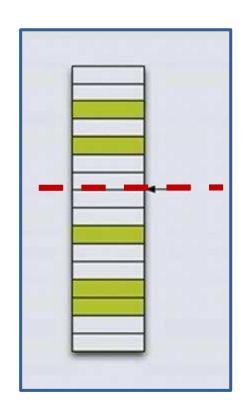
Data Model

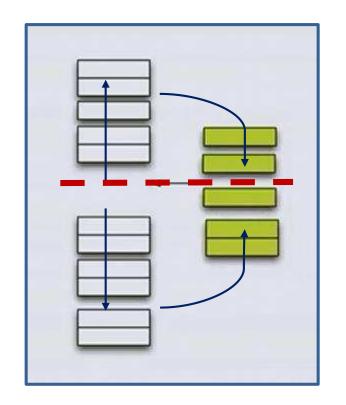
- AGAS essential underpinning for all data management
 - Foundation for syntactic semantic equivalence of local and remote operations
- Full spectrum of C++ data structures are available
 - Either as distributed data structures or for SPMD style computation
- Explicit data partitioning, manually orchestrated boundary exchange
 - Using existing synchronization primitives (for instance channels)
- Use of distributed data structures, like partitioned_vector
 - Use of parallel algorithms
 - Use of co-array like layer (FORTRAN users like that)
- Load balancing: migration
 - Move objects around in between nodes without stopping the application

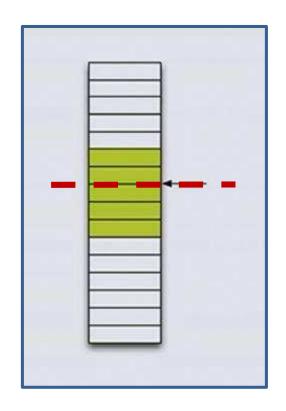


Small Example

Extending Parallel Algorithms







Sean Parent: C++ Seasoning, Going Native 2013



Extending Parallel Algorithms

• New algorithm: gather

```
template <typename BiIter, typename Pred>
pair<BiIter, BiIter> gather(BiIter f, BiIter l, BiIter p, Pred pred)
{
    BiIter it1 = stable_partition(f, p, not1(pred));
    BiIter it2 = stable_partition(p, l, pred);
    return make_pair(it1, it2);
}
```



Extending Parallel Algorithms

New algorithm: gather_async

```
template <typename Bilter, typename Pred>
future<pair<Bilter, Bilter>> gather_async(Bilter f, Bilter l, Bilter p, Pred pred)
{
    future<Bilter> f1 = parallel::stable_partition(par(task), f, p, not1(pred));
    future<Bilter> f2 = parallel::stable_partition(par(task), p, l, pred);
    return dataflow(
        unwrapping([](Bilter r1, Bilter r2) { return make_pair(r1, r2); }),
        f1, f2);
}
```

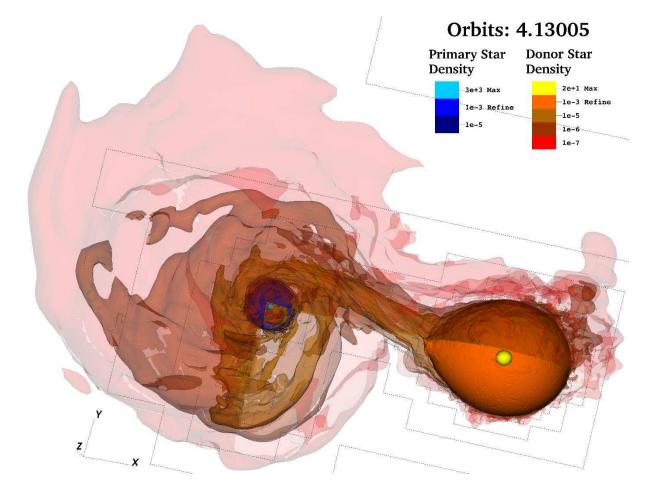
Extending Parallel Algorithms (await)

New algorithm: gather_async

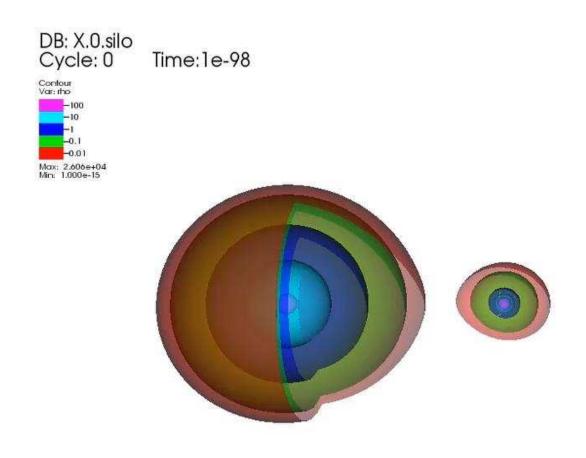
```
template <typename BiIter, typename Pred>
future<pair<BiIter, BiIter>> gather_async(BiIter f, BiIter l, BiIter p, Pred pred)
{
   future<BiIter> f1 = parallel::stable_partition(par(task), f, p, not1(pred));
   future<BiIter> f2 = parallel::stable_partition(par(task), p, l, pred);
   co_return make_pair(co_await f1, co_await f2);
}
```

Recent Results

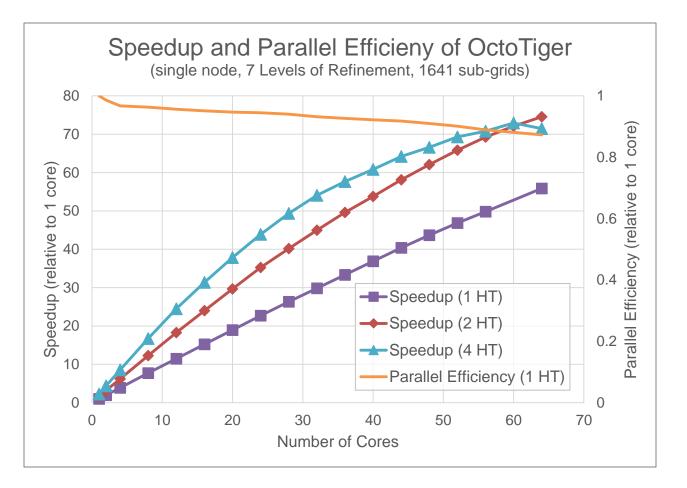
Merging White Dwarfs







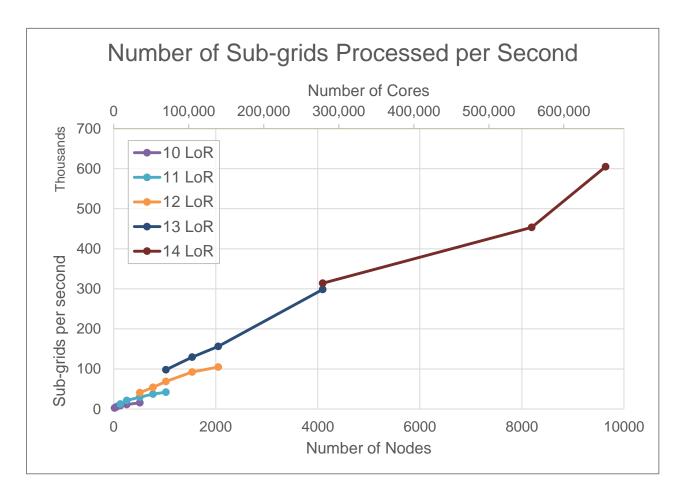
Adaptive Mesh Refinement



Cori II (NERSC)



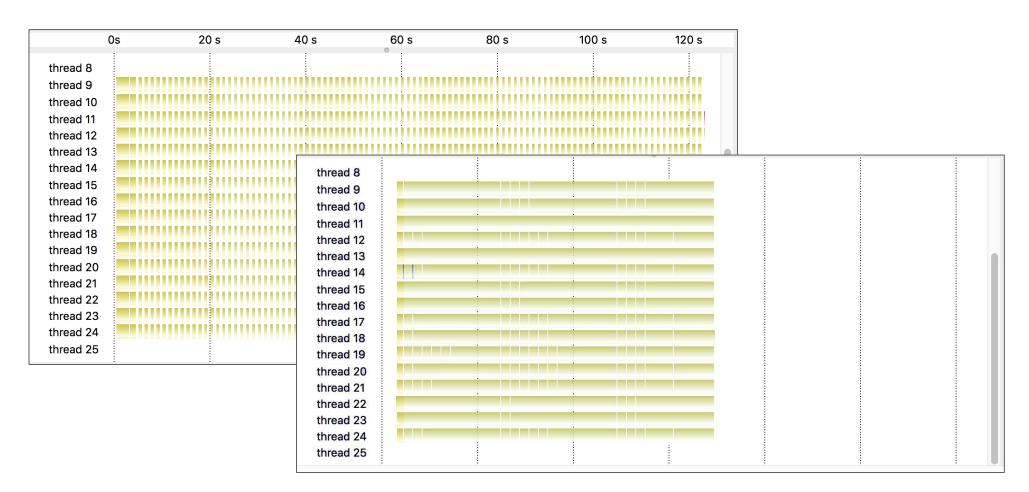
Adaptive Mesh Refinement



Cori II (NERSC)



The Solution to the Application Problem



The Solution to the Application Problem













HPX + Cling + Jupyter

This tutorial works in a special Jupyter notebook that can be used in one of two ways:

- From this website: https://hpx-jupyter.cct.lsu.edu (https://hpx-jupyter.cct.lsu.edu)
- From the docker image: stevenrbrandt/fedora-hpx-cling
- Normally, each cell should contain declarations, e.g. definitions of functions, variables, or #include statements.

```
```#include using namespace std;``
```

If you wish to process an expression, e.g. cout << "hello world\n" you can put .expr at the
front of the cell.</li>

```
```.expr cout << "hello, world\n";``
```

• Sometimes you will want to test a cell because you are uncertain whether it might cause a segfault or some other error that will kill your kernel. Othertimes, you might want to test a definition without permanently adding it to the current namespace. You can do this by prefixing your cell with .test. Whatever is calculated in a test cell will be thrown away after evaluation and will not kill your kernel.

```
```.test.expr int foo[5]; foo[10] = 1;```
```

## Docker Instructions

- · Frist, install Docker on your local resource
- Second, start Docker, e.g. sudo service docker start
- · Third, run the fedora-hpx-cling container, e.g.

```
```dockerpullstevenrbrandt/fe\overline{dora-hpx-cling} docker run -it -p 8000:8000 stevenrbrandt/fedora-hpx-cling```
```

After you do this, docker will respond with something like

```
http://0.0.0.0:8000/?token=5d1eb8a4797851910de481985a54c2fdc3be80280023bac5`
```

Paste that URL into your browser, and you will be able to interact with the notebook.

- Fourth, play with the existing ipynb files or create new ones.
- Fifth, save your work! This is an important step. If you simply quit the container, everything you did will be lost. To save your work, first find your docker image using docker ps.

```\$ docker ps CONTAINER ID IMAGE COMMAND CREATED STATUS PORTS NAMES 4f806b5f4fb3 stevenrbrandt/fedora-hpx-cling "/bin/sh -c 'jupyter " 11 minutes ago Up 11 minutes 0.0.0.0:8000->8000/tcp dreamy\_turing```

Once you have it (in this case, it's 4f806b5f4fb3), you can use docker cp to transfer files to or from your image.

```
```dockercp4f806b5f4fb3:/home/jup/HPX_by_example.ipynb. docker cp HPX_by_example.ipynb 4f806b5f4fb3:/home/jup```
```

```
In [1]: #include <hpx/hpx.hpp>
Out[1]:
```

```
In [2]: using namespace std;
using namespace hpx;
Out[2]:
```

What is a (the) Future?

Many ways to get hold of a future, simplest way is to use (std) async:

```
In [3]: int universal_answer() { return 42; }
    void deep_thought()
    {
        future<int> promised_answer = async(&universal_answer);
        // do other things for 7.5 million years
        cout << promised_answer.get() << endl; // prints 42
    }
Out[3]:</pre>
```

If we want to do something other than a declaration, use the ".expr" prefix.

Compositional Facilities

```
In [7]: int do work(hpx::lcos::future<hpx::util::tuple<hpx::lcos::future<int>,
               hpx::lcos::future<std::basic string<char> > >& w) {
          // extract the value of the first argument.
          return hpx::util::get<0>(w.get()).get();
        }
        future<int> test_when_all()
        {
          future<int> future1 = async([]()->int { return 125; });
          future<string> future2 = async([]()->string { return string("hi");
        });
          auto all f = when all(future1, future2);
          future<int> result = all f.then(
             [](auto f)->int {
               return do_work(f);
          return result;
Out[7]:
In [8]:
        .test.expr cout << test when all().get() << endl;</pre>
        125
         Test
Out[8]:
```

Parallel Algorithms

HPX allows you to write loop parallel algorithms in a generic fashion, applying to specify the way in which parallelism is achieved (i.e. threads, distributed, cuda, etc.) through polcies.

```
In [9]: #include <hpx/include/parallel_for_each.hpp>
    #include <hpx/parallel/algorithms/transform.hpp>
#include <boost/iterator/counting_iterator.hpp>

Out[9]:
In [10]: vector<int> v = { 1, 2, 3, 4, 5, 6 };
Out[10]:
```

Transform

Here we demonstrate the transformation of a vector, and the various mechnanisms by which it can performed in parallel.

```
In [11]: | .expr
         // This parallel tranformation of vector v
         // is done using thread parallelism. An
         // implicit barrier is present at the end.
         parallel::transform (
           parallel::execution::par,
           begin(v), end(v), begin(v),
           [](int i) -> int
              return i+1;
           });
         for(int i : v) cout << i << ",";</pre>
         2,3,4,5,6,7,
Out[11]:
In [12]:
         .expr
         // This parallel tranformation of vector v
         // is done using thread parallelism. There
         // is no implicit barrier. Instead, the
         // transform returns a future.
         auto f = parallel::transform (
           parallel::par (parallel::execution::task),
           begin(v), end(v), begin(v),
           [](int i) -> int
              return i+1;
           });
           // work here...
         // wait for the future to be ready.
         f.wait();
         for(int i : v) cout << i << ",";</pre>
         3,4,5,6,7,8,
Out[12]:
In [13]: #include <hpx/include/parallel fill.hpp>
         #include <hpx/include/compute.hpp>
         #include <hpx/include/parallel executors.hpp>
Out[13]:
In [14]: auto host targets = hpx::compute::host::get local targets();
         typedef hpx::compute::host::block executor<> executor type;
         executor type exec(host targets);
Out[14]:
```

Other Algorithms

There are a great many algorithms. Here we demonstrate a handful of them.

```
In [16]:
          .expr
          std::vector<float> vd;
          for(int i=0;i<10;i++) vd.push back(1.f);</pre>
          parallel::fill(parallel::execution::par.on(exec),vd.begin(),vd.end(),
          0.0f);
Out[16]:
In [17]: #include <hpx/parallel/algorithms/reverse.hpp>
Out[17]:
In [18]:
         .expr
          std::vector<float> vd;
          for(int i=0;i<10;i++) vd.push back(1.f*i);</pre>
          parallel::reverse(parallel::par,vd.begin(),vd.end());
          for(int val : vd) cout << val << " ";</pre>
         9 8 7 6 5 4 3 2 1 0
Out[18]:
In [19]: #include <hpx/include/parallel minmax.hpp>
Out[19]:
In [20]:
          .expr
          std::vector<float> vd;
          for(int i=0; i<10; i++) vd.push back(1.f*rand());
          auto ptr = parallel::max element(parallel::par,vd,std::less<float>());
          for(float val : vd) cout << val << " ";</pre>
          cout << endl << *ptr << endl;</pre>
         8.02369e+08 1.63599e+09 1.60543e+09 3.22506e+08 4.34983e+08 1.87237e+0
         9 2.04466e+09 1.82667e+09 1.27975e+09 1.95976e+09
         2.04466e+09
Out[20]:
```

```
In [21]: #include <hpx/traits/is executor.hpp>
         #include <hpx/include/parallel executors.hpp>
Out[21]:
In [22]: int count_async = 0;
         struct test async executor
         {
             typedef hpx::parallel::parallel execution tag execution category;
             template <typename F, typename ... Ts>
             static hpx::future<typename hpx::util::result of<F&&(Ts&&...)>::ty
         pe>
             async execute(F && f, Ts &&... ts)
                  ++count async;
                  return hpx::async(hpx::launch::async, std::forward<F>(f),
                      std::forward<Ts>(ts)...);
             }
         };
Out[22]:
In [23]: // Note that the exact way to specify this trait for an executor is in
          flux
         // and the code here is tied to the specific version of HPX on the tes
         t machine.
         namespace hpx { namespace traits
         {
             template<>
             struct is two way executor<test async executor>
                : std::true type
             {};
         }}
Out[23]:
         .test.expr
In [24]:
         // This parallel tranformation of vector v
         // is done using using distributed parallelism.
         test async executor e;
         parallel::transform (
           parallel::execution::par.on(e),
           begin(v), end(v), begin(v),
           [](int i) -> int
             return i+1;
         cout << "count=" << count async << endl;</pre>
         count=3
          Test
```

Out[24]:

Let's Parallelize It - Adding Real Asynchrony

Here we take a step back. Instead of using a pre-designed parallel operation on a vector, we instead introduce task-level parallelism to an existing program.

Calculate Fibonacci numbers in parallel (1st attempt)

```
In [25]: uint64_t fibonacci(uint64_t n)
{
    // if we know the answer, we return the value
    if (n < 2) return n;
    // asynchronously calculate one of the sub-terms
    future<uint64_t> f = async(launch::async, &fibonacci, n-2);
    // synchronously calculate the other sub-term
    uint64_t r = fibonacci(n-1);
    // wait for the future and calculate the result
    return f.get() + r;
}
Out[25]:
In [26]: .expr cout << fibonacci(10) << endl;
55</pre>
```

Let's Parallelize It – Introducing Control of Grain Size

Parallel calculation, switching to serial execution below given threshold

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```
In [27]: const int threshold = 20;
         uint64_t fibonacci_serial(uint64_t n)
           if (n < 2) return n;
           uint64 t f1 = fibonacci serial(n-2);
           uint64 t f2 = fibonacci serial(n-1);
           return f1 + f2;
         }
         uint64 t fibonacci2(uint64 t n)
           if (n < 2) return n;
           if (n < threshold) return fibonacci serial(n);</pre>
           // asynchronously calculate one of the sub-terms
           future<uint64 t> f = async(launch::async, &fibonacci2, n-2);
           // synchronously calculate the other sub-term
           uint64 t r = fibonacci2(n-1);
           // wait for the future and calculate the result
           return f.get() + r;
         }
Out[27]:
In [28]:
         .expr cout << fibonacci2(22) << endl;</pre>
         17711
Out[28]:
```

Let's Parallelize It – Apply Futurization

Parallel way, futurize algorithm to remove suspension points

```
In [29]: future<uint64_t> fibonacci3(uint64_t n)
{
    if(n < 2) return make_ready_future(n);
    if(n < threshold) return make_ready_future(fibonacci_serial(n));

    future<uint64_t> f = async(launch::async, &fibonacci3, n-2);
    future<uint64_t> r = fibonacci3(n-1);

    return dataflow(
       [](future<uint64_t> f1, future<uint64_t> f2) {
        return f1.get() + f2.get();
       },
       f, r);
    }

Out[29]:
```

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Let's Parallelize It – Unwrap Argument Futures

```
In [31]: #include <hpx/util/unwrapped.hpp>
          using hpx::util::unwrapping;
          future<uint64 t> fibonacci4(uint64 t n)
            if(n < 2) return make ready future(n);</pre>
            if(n < threshold) return make ready future(fibonacci serial(n));</pre>
            future<uint64 t> f = async(launch::async, &fibonacci4, n-2);
            future<uint64 t> r = fibonacci4(n-1);
            return dataflow(
              unwrapping([](uint64_t f1, uint64_t f2) {
                return f1+f2;
              }),
              f, r);
          }
Out[31]:
          .expr cout << fibonacci4(22).get() << endl;</pre>
In [32]:
          17711
Out[32]:
```

Excercise: Parallelize a sort

Test what you've learned. See if you can speed up the quicksort program below by find a place to:

- 1. parallelize the code with async
- 2. use parallel transforms

```
In [33]: #include <unistd.h>
    #include <stdlib.h>
    #include <iostream>
    #include <vector>
    #include <functional>
    using namespace std;
    function<void(vector<int>&)> myqsort = [](vector<int>& v)->void {};
Out[33]:
```

11/14/2017 HPX_Training

Discussion

We want to define the myqsort function repeatedly, and call it recursively. This is hard to do in C++. So we define it as a std::function<>. There is a slight awkwardness to this. If you want to call myqsort() with an async function, you have to do it like this:

```
auto f = hpx::async([&arg](){ myqsort(arg); });
Not like this
     auto f = hpx::async(myqsort,arg);
   In [42]: | .test.expr
              myqsort = [](vector<int>& v)->void {
                if(v.size()<2) return;</pre>
                vector<int> pre, eq, post;
                int pivot = v[rand() % v.size()];
                for(int val : v) {
                  if(val < pivot) pre.push back(val);</pre>
                  else if(pivot < val) post.push back(val);</pre>
                  else eq.push_back(val);
                }
                myqsort(pre);
                mygsort(post);
                //for(int i=0;i<eq.size();i++) v[i+pre.size()] = eq[i];
                parallel::transform(parallel::par,
                  eq.begin(), eq.end(), v.begin()+pre.size(),[](int i) { return i;
              });
                for(int i=0;i<post.size();i++) v[i+pre.size()+eq.size()] = post[i];</pre>
                for(int i=0;i<pre.size();i++) v[i] = pre[i];</pre>
              };
              vector<int> vv{20};
              for(int i=0;i<20;i++) vv.push back(rand() % 100);</pre>
              for(int val : vv) cout << val << " ";</pre>
              cout << endl;</pre>
              myqsort(vv);
              for(int val : vv) cout << val << " ";</pre>
              cout << endl;</pre>
```

20 26 32 84 5 66 50 81 7 5 53 69 45 84 94 59 21 80 96 17 6 5 5 6 7 17 20 21 26 32 45 50 53 59 66 69 80 81 84 84 94 96

```
Test
```

Out[42]:

The Wave Equation

This problem file sets up a very simple physical system, a wave propagating in one dimension. It is a nice example because it requires several loops in sequence and presents an opportunity to practice creating and using HPX parallel algorithms.

```
In [1]: #include <hpx/hpx.hpp>
    #include <vector>
    #include <hpx/include/parallel_for_each.hpp>
    #include <hpx/parallel/algorithms/transform.hpp>
    #include <boost/iterator/counting_iterator.hpp>
Out[1]:
```

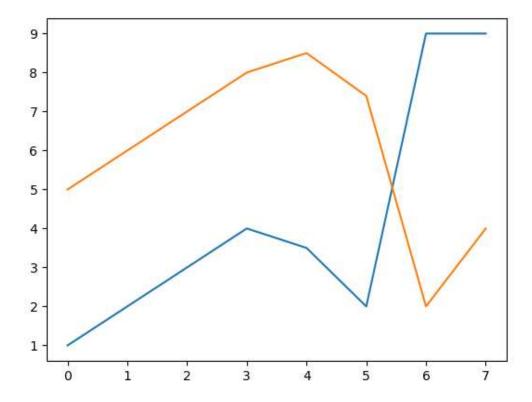
Basic Variables

We are going to evolve two variables, phi and psi on a grid described by N, dx, and x_0 . The system of equations we will solve is

```
\partial_t \phi = c \partial_x \psi \partial_t \psi = c \partial_x \phi In [2]:  \begin{array}{c} {\bf const} \ {\rm int} \ {\rm N=300;} \\ {\rm std::vector<double> \ phi(N),psi(N);} \\ {\bf const} \ {\rm double} \ {\rm dx=0.01, \ x0=-1.5;} \\ {\rm Out[2]:} \end{array}
```

```
In [3]: #include <fstream>
         #include <sstream>
         #include <iostream>
         #include <vector>
         #include <content.hpp>
         /**
          * The following function plots an array of vectors
          * of doubles in a single overlapping plot.
         void plot vector(std::vector<std::vector<double>> v,std::string iname)
           // Store the data in a text file
           const char *fname = "data.txt";
           std::ofstream o(fname);
           for(int n=0; n<v.size(); n++) {
             const std::vector<double>& vv = v[n];
             for(int i=0;i<vv.size();i++) {</pre>
               if(i > 0) o << ' ';
               o << i;
             }
             o << std::endl;</pre>
             for(int i=0;i<vv.size();i++) {</pre>
               if(i > 0) o << ' ';
               0 << vv[i];</pre>
             }
             o << std::endl;</pre>
           }
           o.close();
           // Create a python script to run matplotlib
           std::ostringstream cmd;
           cmd << "import matplotlib\n";</pre>
           cmd << "matplotlib.use('Agg')\n";</pre>
           cmd << "import numpy as np\n";</pre>
           cmd << "import matplotlib.pyplot as plt\n";</pre>
           cmd << "f = np.genfromtxt('" << fname << "')\n";</pre>
           cmd << "plt.figure()\n";</pre>
           cmd << "for n in range(0,f.shape[0],2):\n";
           cmd << " plt.plot(f[n,:],f[n+1,:]) \n";
           cmd << "plt.savefig('" << iname << ".png')\n";</pre>
           cmd << "exit(0)\n";
           std::ofstream o2("p.py");
           o2 << cmd.str();
           o2.close();
           system("python3 p.py");
           // Create the html. The pid is added to prevent caching.
           // Note that the pid changes with every cell with the
           // current implementation of HPX/cling.
           std::ostringstream html;
           html << "<imq src='" << iname << ".pnq?pid=" << getpid() << "-" << r</pre>
         and() << "'>";
           std::string htmls = html.str();
           create content(htmls.c str());
         }
```

Out[3]:



```
Out[4]:
```

```
In [8]: // This version of plot_vector plots a single vector only.
void plot_vector(const std::vector<double>& v) {
    std::vector<std::vector<double>> vv;
    vv.push_back(v);
    plot_vector(vv,"vec");
}
```

Out[8]:

```
In [9]: // Apply boundary conditions. In this case, we are using periodic boun
         dary
         // conditions, i.e. if we move N-2 points to the right we come back to
          where we were.
         void boundary(std::vector<double>& vv) {
             const int n = vv.size();
             vv[0] = vv[n-2];
             vv[n-1] = vv[1];
         }
Out[9]:
In [10]: // The following are auxiliary variables which are required by our Run
         ge-Kutta
         // time integration scheme.
         std::vector<double> phi2(N), psi2(N), phi3(N), psi3(N);
         std::vector<double> k1 phi(N), k1 psi(N), k2 phi(N), k2 psi(N), k3 phi
         (N), k3 psi(N);
Out[10]:
         #include <hpx/include/parallel for each.hpp>
In [11]:
         #include <hpx/parallel/algorithms/transform.hpp>
         #include <boost/iterator/counting iterator.hpp>
Out[11]:
```

The Wave Equation Evolution Code

The sequence of loops below will perform an evolution of the wave equation, essentially, a sine wave that propagates unchanged to the right. The important thing is not understanding the physics, but in parallelizing the loops and recognizing the dependencies between them.

As an example, the first loop is already parallelized with for_each. Note: You should be able to find a place to use parallel::execution::task.

```
for(int i=0;i<N;i++) {
  // psi = A*cos(k*x + w*t)
  // phi = B*cos(k*x + w*t)
  // d(phi)/dt = c*d(psi)/dx
  // w^*B^*sin(k^*x + w^*t) = c^*k^*A^*sin(k^*x + w^*t)
  //B = -c*k*A/w
  // d(psi)/dt = c*d(phi)/dx
  // w^*A^*sin(k^*x + w^*t) = c^*k^*B^*sin(k^*x + w^*t)
  double x = x0 + i*dx;
  phi[i] = B*sin(k*x);
  psi[i] = A*sin(k*x);
*/
hpx::parallel::for each(
      hpx::parallel::par,
      boost::counting iterator<int>(0),
      boost::counting iterator<int>(N),
      [&](inti){
  double x = x0 + i*dx;
   phi[i] = B*sin(k*x);
   psi[i] = A*sin(k*x);
});
while(t < t end) {</pre>
 for(int i=1;i<N-1;i++) {
    k1 phi[i] = c*(psi[i+1]-psi[i-1])/(2*dx);
    k1 psi[i] = c*(phi[i+1]-phi[i-1])/(2*dx);
 }
 for(int i=1;i<N-1;i++) {
    phi2[i] = phi[i] + (1./3.)*dt*k1 phi[i];
    psi2[i] = psi[i] + (1./3.)*dt*k1 psi[i];
 // the boundary routines can be parallelized also
 boundary(phi2);
 boundary(psi2);
 for(int i=1;i<N-1;i++) {
    k2 phi[i] = c*(psi2[i+1]-psi2[i-1])/(2*dx);
    k2 psi[i] = c*(phi2[i+1]-phi2[i-1])/(2*dx);
 }
 for(int i=1;i<N-1;i++) {
    phi3[i] = phi[i] + (2./3.)*dt*k1_phi[i];
    psi3[i] = psi[i] + (2./3.)*dt*k1_psi[i];
  }
 boundary(phi3);
 boundary(psi3);
 for(int i=1;i<N-1;i++) {
    k3 phi[i] = c*(psi3[i+1]-psi3[i-1])/(2*dx);
   k3_psi[i] = c*(phi3[i+1]-phi3[i-1])/(2*dx);
  }
 for(int i=1;i<N-1;i++) {
    phi[i] = phi[i] + 0.5*dt*(k2 phi[i]+k3 phi[i]);
    psi[i] = psi[i] + 0.5*dt*(k2 psi[i]+k3 phi[i]);
  }
 boundary(phi);
 boundary(psi);
 t += dt;
 t plot += dt;
 if(t plot >= t every) {
```

Performance Analysis of HPX in Jupyter Notebooks using APEX

Kevin Huck

khuck@cs.uoregon.edu

http://github.com/khuck/xpress-apex

Download slides from: http://tau.uoregon.edu/SC17-HPX-APEX.pdf



UNIVERSITY OF OREGON

Install Docker image from USB

- Install Docker (if necessary)
- Insert USB key, open a terminal, navigate to key directory and:
- (sudo) docker load -i fedora-hpx-cling
- (sudo) docker pull stevenrbrandt/fedora-hpx-cling
- (sudo) docker run -it -p 8000:8000 stevenrbrandt/ fedora-hpx-cling

"sudo" may be necessary on some machines

Outline

- Introduction to APEX Autonomic Performance Environment for eXascale
- Motivation, overview, API
- Integration with HPX
- Building HPX with APEX
- APEX event listeners
- Postmortem analysis of HPX applications
 - Gnuplot/Python visualizations of APEX data
 - OTF2 output to Vampir
 - Profile output to TAU

APEX Measurement: Motivation

- Originally designed as a performance measurement library for distributed, asynchronous tasking models/runtimes
 - i.e. HPX, but there are others
- Why another measurement library?
 - "not invented here" mentality? Reinventing the wheel? No.
- New challenges:
 - Lightweight measurement (tasks <1ms)
 - High concurrency (both OS threads and tasks in flight)
 - Distinction between OS and runtime (HPX) thread context
 - Lack of a traditional call stack
 - Task dependency chain instead
 - Runtime controlled task switching
 - Dynamic runtime control

APEX Runtime Adaptation: Motivation

- Controlling concurrency
 - Energy efficiency
 - Performance
- Parametric variability
 - Granularity for this machine / dataset?
- Load Balancing
 - When to perform AGAS migration?
- Parallel Algorithms (for_each...)
 - Separate what from how
- Address the "SLOW(ER)" performance model

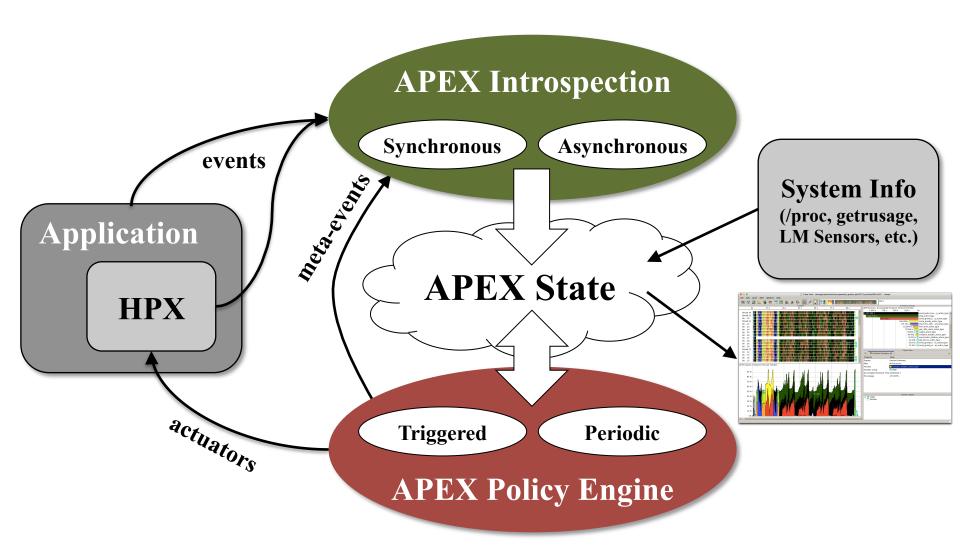
Introduction to APEX

- Performance <u>awareness</u> and performance <u>adaptation</u>
- Top down and bottom up performance mapping / feedback
 - Make node-wide resource utilization data and analysis, energy consumption, and health information available in real time
 - Associate performance state with policy for feedback control

APEX introspection

- OS: track system resources, utilization, job contention, overhead
- Runtime (HPX): track threads, queues, concurrency, remote operations, parcels, memory management
- Application timer / counter observation
- Post-mortem performance analysis
 - "secondary" goal, but has been useful
- Integrated with HPX performance counters

APEX architecture



APEX Introspection

- APEX collects data through "inspectors"
 - Synchronous uses an event API and event "listeners"
 - Initialize, terminate, new thread added to HPX runtime
 - Timer start, stop, yield*, resume* added to HPX task scheduler
 - Sampled value (counters from HPX)
 - Custom events (meta-events)
 - Asynchonous do not rely on events, but occur periodically
- APEX exploits access to performance data from lower stack components
 - "Health" data through other interfaces (/proc/stat, cpuinfo, meminfo, net/dev, self/status, lm_sensors, power*, etc.)

APEX Measurement API (subset)

Starting Timers:

```
apex::profiler* apex::start(const string &name,
   void** data ptr = 0LL);
apex::profiler* apex::start(const uint64 t
   address, void** data ptr = OLL);
apex::profiler* apex::resume(const string name,
   void** context = OLL);
apex::profiler* apex::resume(const uint64 t
   address, void** context = 0LL);
Stopping Timers:
void apex::stop(apex::profiler* p);
void apex::yield(apex::profiler* p);
Sampling a counter:
void apex::sample value (const string name,
   double value);
```

 Note about yield, resume – exist for accurate counting of "number of calls" in the face of preemption (usually handled by HPX scheduler):

| Timer Start Command | Timer Stop Command | Call count is increased by: |
|------------------------|-----------------------|-----------------------------|
| apex::start() | apex::stop() | 1 |
| apex::start() | apex::yield() | 0 |
| apex::resume() | apex::yield() | 0 |
| apex::resume() | apex::stop() | 0 |

APEX Measurement API example

```
void foo(int x) {
  // sample the argument, for example
  apex::sample counter("foo(x)", x);
  // start a timer
  apex::profiler* p = apex::start(&foo);
  /* do some work in function foo */
  // stop the current timer to wait on some asynchronous subtask
  apex::yield(p);
  /* wait on result from "subtask" */
  result = some future.get();
  // "resume" the APEX timer
 p = apex::start(&foo);
  /* do some more work in function foo */
  // stop the timer
  apex::stop(p);
```

HPX and APEX - Integration

- In HPX, all tasks scheduled by the thread scheduler are "automatically" timed – with some caveats
- HPX registered actions are automatically timed
- All threads/tasks are timed, attribution is the required user intervention
 - Asynchronous functions, direct actions are correctly attributed if wrapped with an

```
hpx::util::annotated function object.
```

See notebook examples for details

Annotation examples:

```
// Forward declaration of the action
uint64_t fibonacci_a(uint64_t n);
// This is to generate the required boilerplate we need for the remote
// invocation to work.

HPX_PLAIN_ACTION(fibonacci_a, fibonacci_action);
/* ... */
future<uint64_t> nl = async(fib, locality_id, n - l);
```

```
using namespace hpx::util;
future<uint64_t> f = async(launch::async,
    annotated_function(unwrapping(&fibonacci3), "fibonacci3"), n-2);
```

APEX Event Listeners

Profiling listener

- Start event: input name/address, get timestamp, return profiler handle
- Stop event: get timestamp, put profiler object in a queue for back-end processing, return
- Sample event: put the name & value in the queue
- Asynchronous consumer thread: process profiler objects and samples to build statistical profile (in HPX, processed/scheduled as a thread/task)
- Optional: screen/CSV output, build task scatterplot, build taskgraph, etc.

TAU Listener (postmortem analysis)

Synchronously passes all measurement events to TAU to build an offline profile

OTF2 Listener (postmortem analysis)

Synchronously passes all measurement events to libotf2 for trace analysis

Concurrency listener (postmortem analysis)

- Start event: push timer ID on stack
- Stop event: pop timer ID off stack
- Asynchronous consumer thread: periodically log current timer for each thread, output report at termination

APEX Policy Listener

- Policies are rules that decide on outcomes based on observed state
 - Triggered policies are invoked by introspection API events
 - Periodic policies are run periodically on asynchronous thread
- Polices are registered with the Policy Engine
 - Applications, runtimes, and/or OS register callback functions
- Callback functions define the policy rules
 - "If x < y then..."
- Enables runtime adaptation using introspection data
 - Feedback and control mechanism
 - Engages actuators across stack layers
 - Could also be used to involve online auto-tuning support*
 - Active Harmony http://www.dyninst.org/harmony

APEX Policy API (subset)

```
apex event type
   apex::register custom event(const
   string &name);
apex::custom event(apex event type, void*
   event data);
apex tuning session handle
   apex::setup custom tuning(std::functio
   n<double(void) > metric,
   apex event type event type, int
   num inputs, long** inputs, long* mins,
   long* maxs, long* steps);
```

APEX Policy API example

See last example in notebook:

```
apex::register_custom_event() in
  setup_counters()
```

```
apex::get_profile() (as potential
    exercise) in get_counter_value()
```

```
apex::setup_custom_tuning() in
   do_1d_solve_repart()
```

Building HPX with APEX

```
$ cmake <usual HPX settings>...
-DHPX_WITH_APEX=TRUE \
-DAPEX_WITH_ACTIVEHARMONY=TRUE/FALSE \
-DAPEX_WITH_PAPI=TRUE/FALSE \
-DAPEX_WITH_MSR=TRUE/FALSE \
-DAPEX_WITH_OTF2=TRUE/FALSE \
```

\$ make (as usual) i.e. "make –j 8 core examples tests"

APEX environment variables

APEX DISABLE: 0 APEX THROTTLE CONCURRENCY: 0 APEX SUSPEND: 0 APEX THROTTLING MAX THREADS: 4 APEX PAPI SUSPEND: 0 APEX THROTTLING MIN THREADS: 1 APEX PROCESS ASYNC STATE: 1 APEX THROTTLE ENERGY: 0 APEX THROTTLE ENERGY PERIOD: 1000000 APEX TAU:0 APEX THROTTLING MAX WATTS: 300 APEX POLICY: 1 **APEX_MEASURE_CONCURRENCY: 0** APEX THROTTLING MIN WATTS: 150 APEX MEASURE CONCURRENCY PERIOD: 1000000 APEX PTHREAD WRAPPER STACK SIZE: 0 **APEX SCREEN OUTPUT: 1** APEX OMPT REQUIRED EVENTS ONLY: 0 **APEX PROFILE OUTPUT: 0** APEX OMPT HIGH OVERHEAD EVENTS: 0 **APEX CSV OUTPUT: 0** APEX PIN APEX THREADS: 1 APEX TASKGRAPH OUTPUT: 0 APEX_TASK_SCATTERPLOT: 1 APEX PROC CPUINFO: 0 APEX POLICY DRAIN TIMEOUT: 1000 APEX PROC MEMINFO: 0 **APEX PAPI METRICS:** APEX PROC NET DEV:0 **APEX PLUGINS:** APEX PROC SELF STATUS: 0 APEX PLUGINS PATH:./ APEX PROC SELF IO:0 APEX OTF2:0 APEX PROC STAT: 1 APEX OTF2 ARCHIVE PATH: OTF2 archive APEX PROC PERIOD: 1000000 APEX OTF2 ARCHIVE NAME: APEX

Note about environment variables...

- Some batch submission systems do not pass environment variables to the application without considerable effort
- APEX will read environment variables from \$PWD/ apex.conf file
- All variables also have API calls for getting/setting
 - Note: some are only effective at startup because they change/set the APEX configuration
 - i.e. APEX_DISABLE, APEX_TAU, APEX_PAPI_METRICS, ...

Post-mortem output examples

Example Screen Output

Elapsed time: 2.40514 seconds

Cores detected: 4

Worker Threads observed: 5

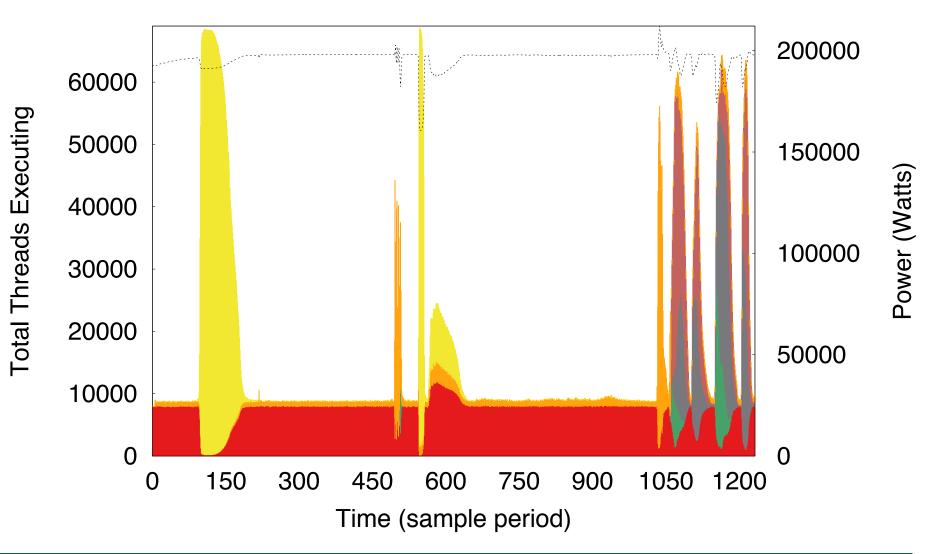
Available CPU time: 9.62057 seconds

```
: #calls | mean | total |
Timer
                                                                      % total
                           someThread(void*) :
                                                2 1.39e+00 2.79e+00
                                                                       28,965
                                  foo(int): 131072 2.00e-05 2.63e+00
                                                                       27,306
             bar(int, apex::profiler**, void**) : 131072 9.69e-06 1.27e+00
                                                                       13.198
                      someUntimedThread(void*): 2 1.12e+00 2.24e+00
                                                                       23.271
                                     main: 1 1.40e+00 1.40e+00
                                                                       14.548
                                 APEX MAIN: 1 2.41e+00 2.41e+00
                                                                      100.000
```

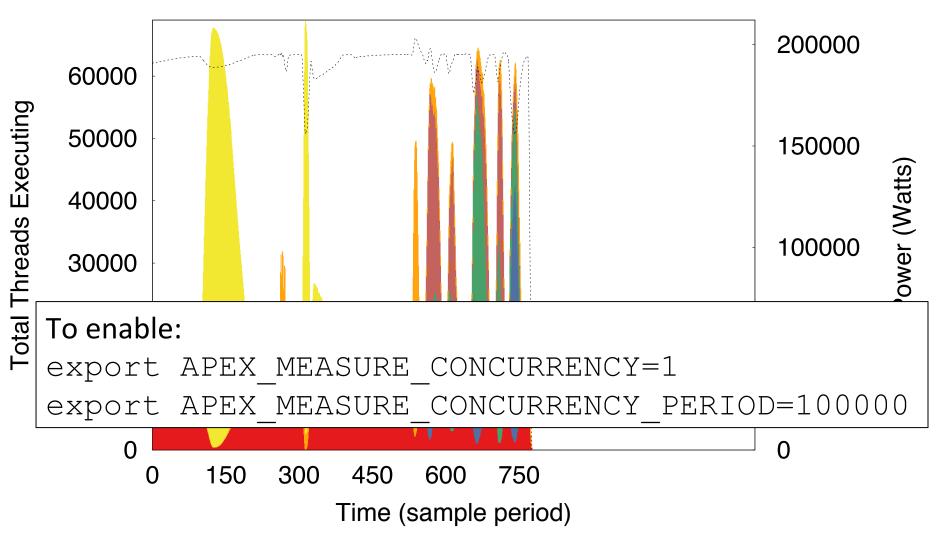
Total timers : 262,150

```
To enable:
    export APEX_SCREEN_OUTPUT=1
    or, call apex::dump(bool reset);
```

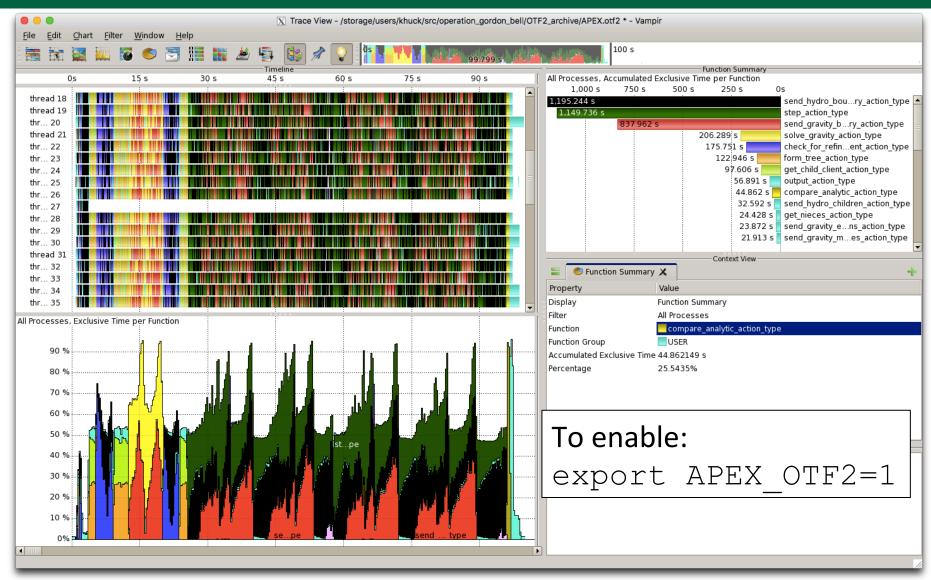
Concurrency View (before fix)



Concurrency View (after fix)

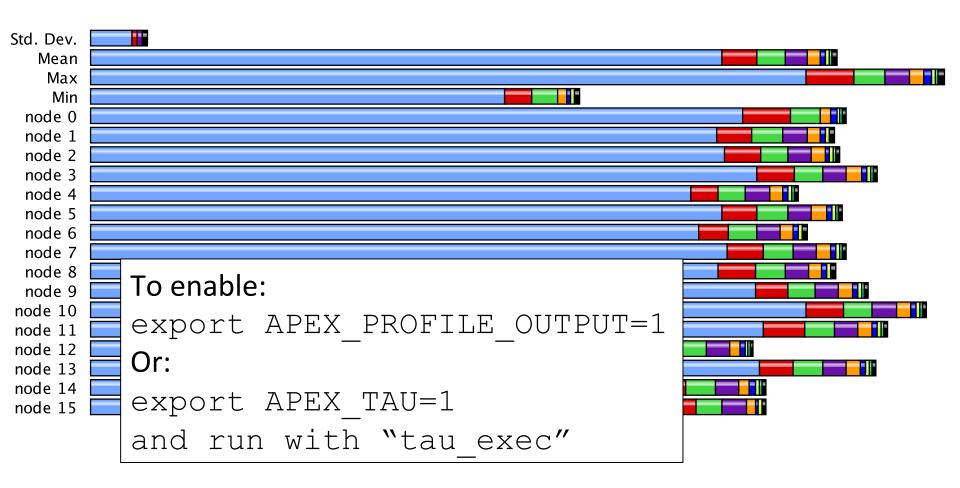


OTF2 View in Vampir

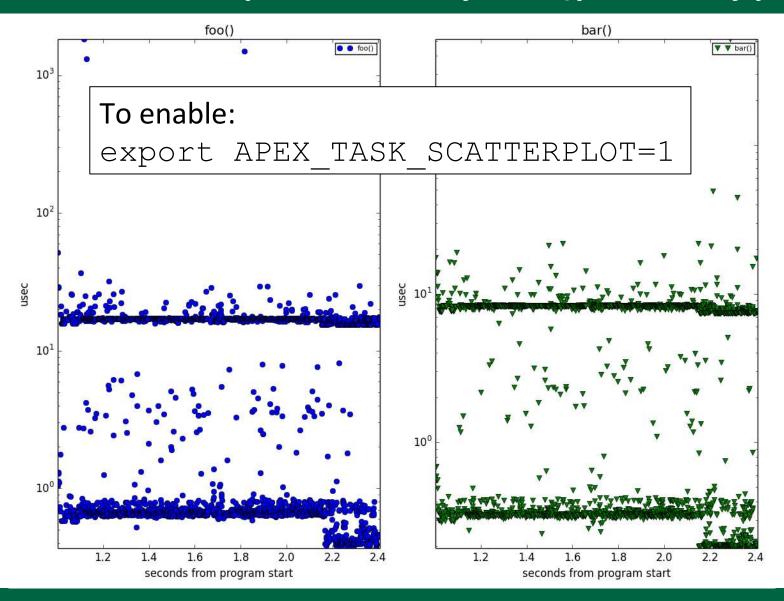


Profile View in TAU ParaProf

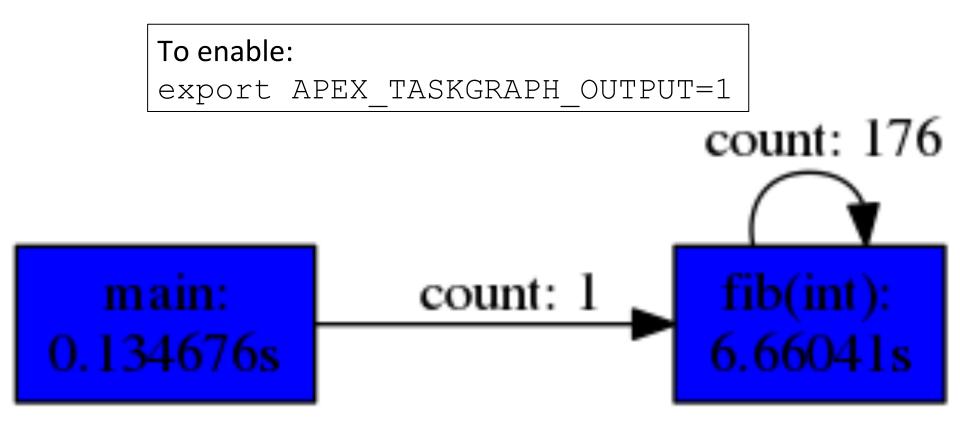
Metric: TIME Value: Exclusive



Task Scatterplot Analysis (prototype)



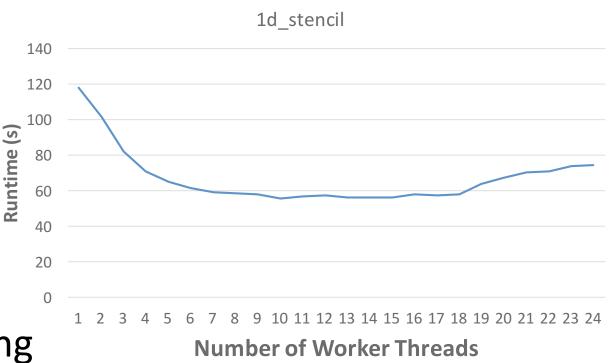
Taskgraph View (prototype)



HPX + APEX Policy Examples

Example: HPX+APEX

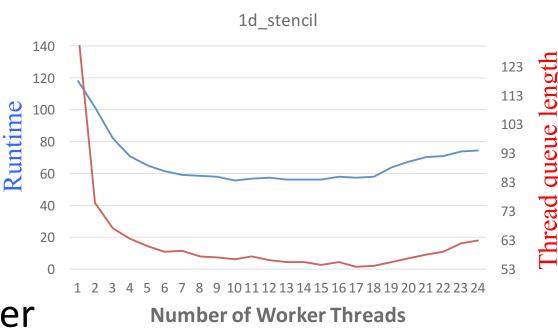
- Heat diffusion
- 1D stencil code
- Data array partitioned into chunks
- 1 node with no hyperthreading



 Performance increases to a point with increasing worker threads, then decreases

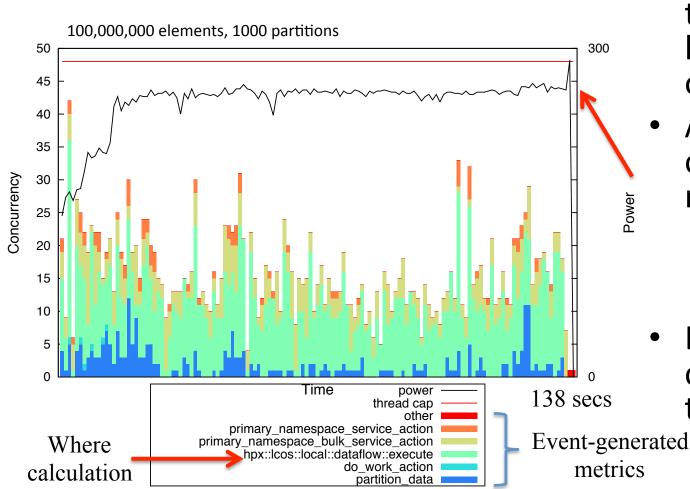
Concurrency & Performance

Region of maximum performance correlates with thread queue length runtime performance counter



- Represents # tasks currently waiting to execute
- Could do introspection on this to control concurrency throttling policy (see next example)

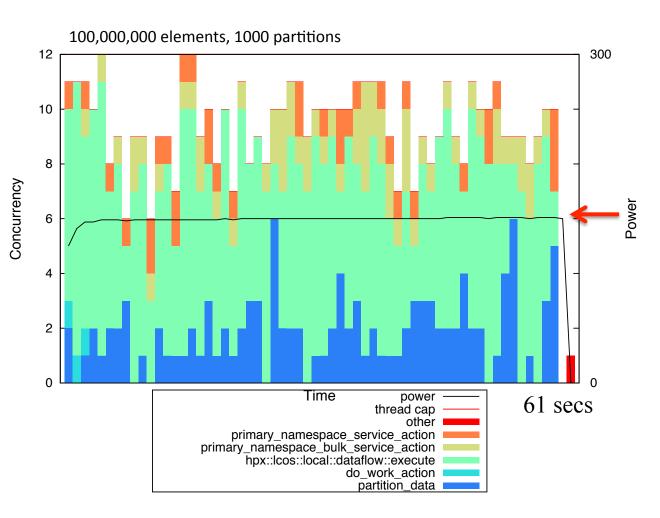
1d stencil 4 Baseline



- 48 worker threads (with hyperthreading, on Edison)
 - Actual concurrency much lower
 - Implementation is memory bound
- Large variation in concurrency over time
 - Tasks waiting on prior tasks to complete

takes place

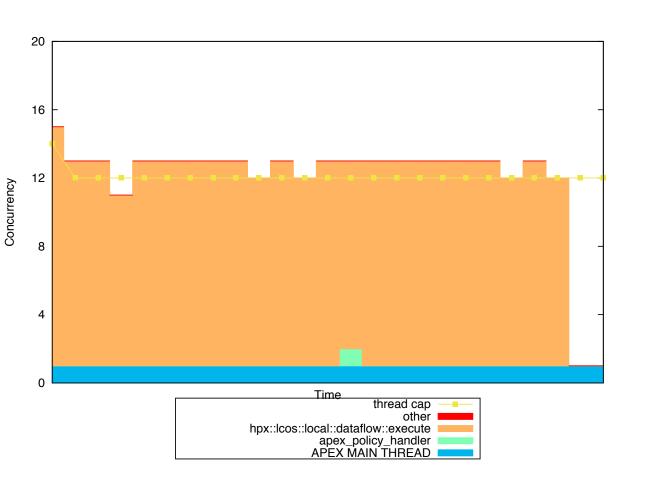
1d_stencil w/optimal # of Threads



- 12 worker threads on Edison
- Greater proportion of threads kept busy
 - Less

 interference
 between active
 threads and
 threads waiting
 for memory
- Much faster
 - 61 sec. vs 138 sec.

1d_stencil Adaptation with APEX

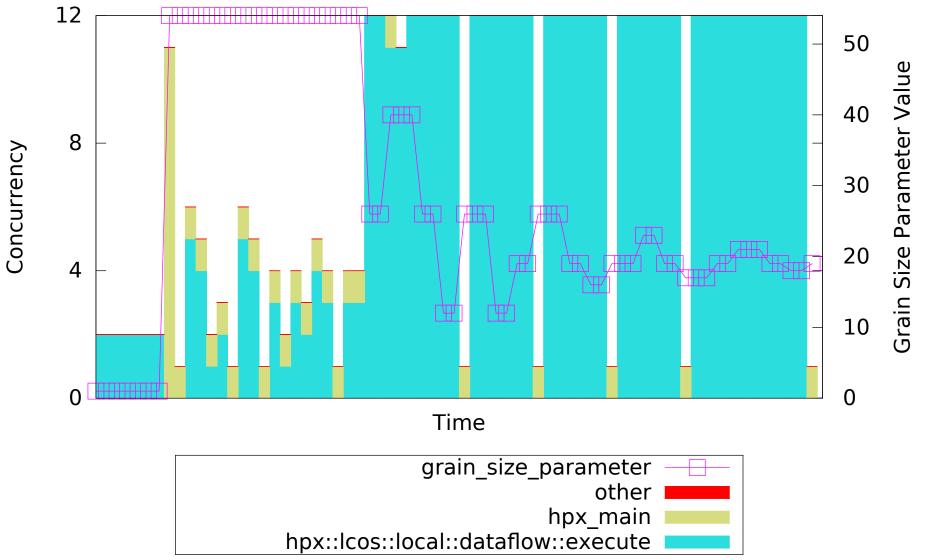


- Initially 32 worker threads
- ActiveHarmony searches for minimal thread queue length
- Quickly converges on 12

Adapting Block Size

- Is 1000 partitions of 100000 cells the best partitioning?
- Parametric studies say "no".
- Can we modify the example to repartition as necessary to find better performance?

1d_stencil: adapting block size



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HPX + Cling + Jupyter

This tutorial works in a special Jupyter notebook that can be used in one of two ways:

- From this website: https://hpx-jupyter.cct.lsu.edu (https://hpx-jupyter.cct.lsu.edu)
- From the docker image: stevenrbrandt/fedora-hpx-cling
- Normally, each cell should contain declarations, e.g. definitions of functions, variables, or #include statements.

```
```#include using namespace std;``
```

If you wish to process an expression, e.g. cout << "hello world\n" you can put .expr at the
front of the cell.</li>

```
```.expr cout << "hello, world\n";``
```

• Sometimes you will want to test a cell because you are uncertain whether it might cause a segfault or some other error that will kill your kernel. Othertimes, you might want to test a definition without permanently adding it to the current namespace. You can do this by prefixing your cell with .test. Whatever is calculated in a test cell will be thrown away after evaluation and will not kill your kernel.

```
```.test.expr int foo[5]; foo[10] = 1;```
```

## Docker Instructions

- · Frist, install Docker on your local resource
- Second, start Docker, e.g. sudo service docker start
- Third, run the fedora-hpx-cling container, e.g.

```
```dockerpullstevenrbrandt/fedora-hpx-cling docker run -it -p 8000:8000 stevenrbrandt/fedora-hpx-cling```
```

After you do this, docker will respond with something like

```
http://0.0.0.0:8000/?token=5d1eb8a4797851910de481985a54c2fdc3be80280023bac5`
```

Paste that URL into your browser, and you will be able to interact with the notebook.

- Fourth, play with the existing ipynb files or create new ones.
- Fifth, save your work! This is an important step. If you simply quit the container, everything you did will be lost. To save your work, first find your docker image using docker ps.

```\$ docker ps CONTAINER ID IMAGE COMMAND CREATED STATUS PORTS NAMES 4f806b5f4fb3 stevenrbrandt/fedora-hpx-cling "/bin/sh -c 'jupyter " 11 minutes ago Up 11 minutes 0.0.0.0:8000->8000/tcp dreamy\_turing```

Once you have it (in this case, it's 4f806b5f4fb3), you can use docker cp to transfer files to or from your image.

```
```dockercp4f806b5f4fb3:/home/jup/HPX_by_example.ipynb. docker cp HPX_by_example.ipynb 4f806b5f4fb3:/home/jup```
```

Measuring HPX performance

HPX is integraged with APEX, "An Autonomic Performance Environment for eXascale". APEX serves two primary roles in HPX - to measure the HPX runtime and application tasks, and to use introspection of those measurements to control behavior.

Measurement

APEX provides an API for measuring actions within a runtime. The API includes methods for timer start/stop, as well as sampled counter values. APEX is designed to be integrated into a runtime, library and/or application and provide performance introspection for the purpose of runtime adaptation. While APEX can provide rudimentary post-mortem performance analysis measurement, there are many other performance measurement tools that perform that task much better (such as TAU http://tau.uoregon.edu (http://tau.uoregon.edu). That said, APEX includes an event listener that integrates with the TAU measurement system, so APEX events can be forwarded to TAU and collected in a TAU profile and/or trace to be used for post-mortem performance anlaysis. That process is not covered in this tutorial, but for more information, see http://github.com/khuck/xpress-apex (http://github.com/khuck/xpress-apex).

Runtime Adaptation

APEX provides a mechanism for dynamic runtime feedback and control, either for autotuning or adaptation to changing environment. The infrastruture that provides the adaptation in APEX is the Policy Engine, which executes policies either periodically or triggered by events. The policies have access to the performance state as observed by the APEX introspection API. APEX is integrated with Active Harmony (http://www.dyninst.org/harmony (http://www.dyninst.org/harmony)) to provide dynamic search for autotuning.

Fibonacci example - what's the performance?

Using the first fibonacci implementation from the HPX introduction, let's examine the performance. To get a simple text report of performance from HPX in a regular program, you would set the APEX_SCREEN_OUTPUT environment variable to a postitive number (i.e. "1"). In the Jupyter notebook, we will use the apex::dump(bool reset); method instead. Because the HPX environment is continuously running in the Jupyter kernel, we also need to reset the timers before executing our test.

First, include the HPX header for our example, and declare some useful namespaces.

```
In [ ]: #include <hpx/hpx.hpp>
using namespace std;
using namespace hpx;
```

Next, we will define the first implementation of the fibonacci algorithm from the previous presentation.

```
In []: uint64_t fibonacci(uint64_t n)
{
    // if we know the answer, we return the value
    if (n < 2) return n;
    // asynchronously calculate one of the sub-terms
    future<uint64_t> f = async(launch::async, &fibonacci, n-2);
    // synchronously calculate the other sub-term
    uint64_t r = fibonacci(n-1);
    // wait for the future and calculate the result
    return f.get() + r;
}
```

And we will execute that definition:

We get some useful information, but what is task_object::apply? That is the HPX runtime executing asynchronous tasks. To get a useful label for that function (and distinguish the different task types), we will use hpx::util::annotated_function. Note carefully that we will add the hpx::util namespace, and also that we will change the name to fibonacci2, to distinguish from the previous definition (that still exists in this compilation unit). We will use that renaming pattern throughout this tutorial.

(markdown padding for proper formatting)

Next, we will execute the fibonacci2 method with the annotated function:

```
In [ ]:
         .expr
         apex::reset(0L);
         cout << fibonacci2(22) << endl;</pre>
         apex::dump(true);
```

That's better, but not quite right. Note carefully that the fibonacci2 synchronous events are double-counting -- the synchronous timer for (n) includes the time spent computing both (n-1) and (n-2). Let's see what happens with the example that uses a serial cutoff:

```
In [ ]: const int threshold = 10;
         uint64 t fibonacci serial(uint64 t n)
           if (n < 2) return n;
           uint64 t f1 = fibonacci serial(n-2);
           uint64 t f2 = fibonacci serial(n-1);
           return f1 + f2;
         }
         future<uint64 t> fibonacci3(uint64 t n)
           if(n < 2) return make ready future(n);</pre>
           if(n < threshold) return make ready future(fibonacci serial(n));</pre>
           future<uint64_t> f = async(launch::async, annotated function(unwrapp
         ing(&fibonacci3), "fibonacci3"), n-2);
           future<uint64 t> r = fibonacci3(n-1);
           return dataflow(
             [](future<uint64 t> f1, future<uint64 t> f2) {
               return f1.get() + f2.get();
             },
             f, r);
         }
In [ ]:
         .expr
         apex::reset(false);
         cout << fibonacci3(22).get() << endl;</pre>
```

```
apex::dump(true);
```

Note that the timers are only around the calls where 10 > n >= 22. We are now under-counting, because the synchronous executions aren't measured. If we were interested in a careful evaluation of the serial execution, we would include a third function with a timer that is called once. That function would then recursively call fibonacci serial. Feel free to implement that version as an exercise.

What if we want to compare the different methods? We can put them all into one compilation unit, and execute them in the same cell for a clearer comparison:

```
In []: .expr
    apex::reset(0L);
    cout << fibonacci(22) << endl;
    cout << fibonacci2(22) << endl;
    cout << fibonacci3(22).get() << endl;
    apex::dump(true);</pre>
```

Well, there is some useful data in there - but we need to do some aggregation. We could add them up, or we could wrap the top level calls with a timer and get the wall clock time for each approach:

Heat Equation Example

```
In [ ]: #include <hpx/include/parallel_algorithm.hpp>
       #include <boost/range/irange.hpp>
       #include <boost/format.hpp>
       #include <cstddef>
       #include <cstdint>
       #include <iostream>
       #include <memory>
       #include <utility>
       #include <vector>
       #include <stdexcept>
       #include <string>
       void print_time_results(
           std::uint32_t num_localities
         , std::uint64_t num_os_threads
         , std::uint64_t elapsed
         , std::uint64 t nx
         , std::uint64 t np
         , std::uint64_t nt
         , bool header
           )
       {
           if (header)
```

```
std::cout << "Localities,OS_Threads,Execution_Time_sec,"</pre>
               "Points per Partition, Partitions, Time Steps\n"
            << std::flush;
   std::string const locs str = boost::str(boost::format("%u,") % num
localities);
   std::string const threads str = boost::str(boost::format("%lu,") %
num os threads);
   std::string const nx str = boost::str(boost::format("%lu,") % nx);
   std::string const np str = boost::str(boost::format("%lu,") % np);
   std::string const nt str = boost::str(boost::format("%lu ") % nt);
   std::cout << ( boost::format("%-6s %-6s %.14g, %-21s %-21s %-21s \n
")
           % locs str % threads str % (elapsed / 1e9) %nx str % np st
r
           % nt str) << std::flush;
}
void print_time_results(
   std::uint64 t num os threads
  , std::uint64_t elapsed
  , std::uint64 t nx
   std::uint64 t np
  , std::uint64 t nt
  , bool header
{
   if (header)
       std::cout << "OS Threads,Execution Time sec,"</pre>
               "Points per Partition, Partitions, Time Steps\n"
            << std::flush;
   std::string const threads str = boost::str(boost::format("%lu,") %
num_os threads);
   std::string const nx_str = boost::str(boost::format("%lu,") % nx);
   std::string const np str = boost::str(boost::format("%lu,") % np);
   std::string const nt str = boost::str(boost::format("%lu ") % nt);
   std::cout << ( boost::format("%-21s %.14g, %-21s %-21s %-21s\n")
           % threads str % (elapsed / 1e9) %nx str % np str
           % nt str) << std::flush;
}
void print_time_results(
   std::uint64_t num_os_threads
  , std::uint64 t elapsed
  , std::uint64 t nx
   std::uint64 t nt
   bool header
{
   if (header)
       std::cout << "OS Threads,Execution Time sec,"</pre>
               "Grid Points, Time Steps\n"
```

```
<< std::flush;
   std::string const threads str = boost::str(boost::format("%lu,") %
num os threads);
   std::string const nx str = boost::str(boost::format("%lu,") % nx);
   std::string const nt str = boost::str(boost::format("%lu ") % nt);
   std::cout << ( boost::format("%-21s %10.12s, %-21s %-21s\n")
           % threads str % (elapsed / 1e9) %nx str % nt str) << std::
flush;
}
// Command-line variables
bool header = true; // print csv heading
               // heat transfer coefficient
double k = 0.5;
double dt = 1.;
                  // time step
double dt = 1.; double dx = 1.;
                  // grid spacing
inline std::size t idx(std::size t i, int dir, std::size t size)
{
   if(i == 0 \&\& dir == -1)
       return size-1;
   if(i == size-1 \&\& dir == +1)
       return 0;
   HPX ASSERT((i + dir) < size);</pre>
   return i + dir;
}
// Our partition data type
struct partition data
{
public:
   partition data(std::size t size)
     : data (new double[size]), size (size)
   {}
   partition_data(std::size_t size, double initial_value)
     : data (new double[size]),
       size (size)
   {
       double base value = double(initial value * size);
       for (std::size t i = 0; i != size; ++i)
           data [i] = base value + double(i);
   }
   partition data(partition data && other)
     : data_(std::move(other.data_))
     , size_(other.size )
   {}
   double& operator[](std::size t idx) { return data [idx]; }
```

```
double operator[](std::size t idx) const { return data [idx]; }
   std::size t size() const { return size ; }
private:
   std::unique_ptr<double[]> data_;
   std::size t size ;
};
std::ostream& operator<<(std::ostream& os, partition data const& c)
   os << "{";
   for (std::size t i = 0; i != c.size(); ++i)
       if (i != 0)
           os << ", ";
       os << c[i];
   os << "}";
   return os;
}
struct stepper
   // Our data for one time step
   typedef hpx::shared future<partition data> partition;
   typedef std::vector<partition> space;
   // Our operator
   static double heat(double left, double middle, double right)
       return middle + (k*dt/(dx*dx)) * (left - 2*middle + right);
   }
   // The partitioned operator, it invokes the heat operator above on
all
   // elements of a partition.
   static partition data heat part(partition data const& left,
       partition_data const& middle, partition_data const& right)
   {
       apex_wrapper profiler("partition_data::heat_part", OL);
       std::size t size = middle.size();
       partition data next(size);
       next[0] = heat(left[size-1], middle[0], middle[1]);
       for(std::size t i = 1; i != size-1; ++i)
       {
           next[i] = heat(middle[i-1], middle[i], middle[i+1]);
       }
       next[size-1] = heat(middle[size-2], middle[size-1], right[0]);
       return next;
   }
```

```
// do all the work on 'np' partitions, 'nx' data points each, for
 'nt'
    // time steps, limit depth of dependency tree to 'nd'
   hpx::future<space> do work(std::size t np, std::size t nx, std::si
ze_t nt,
        std::uint64 t nd)
    {
        using hpx::dataflow;
        using hpx::util::unwrapping;
        // U[t][i] is the state of position i at time t.
        std::vector<space> U(2);
        for (space& s: U)
            s.resize(np);
        // Initial conditions: f(0, i) = i
        std::size t b = 0;
        auto range = boost::irange(b, np);
        using hpx::parallel::execution::par;
        hpx::parallel::for each(par, std::begin(range), std::end(range)
),
            [&U, nx](std::size t i)
                U[0][i] = hpx::make_ready_future(partition_data(nx, do
uble(i)));
        );
        // limit depth of dependency tree
        hpx::lcos::local::sliding_semaphore sem(nd);
        auto Op = hpx::util::annotated function(unwrapping(&stepper::h
eat_part),
                                                 "stepper::heat part");
        // Actual time step loop
        for (std::size t t = 0; t != nt; ++t)
            space const& current = U[t % 2];
            space& next = U[(t + 1) % 2];
            for (std::size t i = 0; i != np; ++i)
                next[i] = dataflow(
                    Op, current[idx(i, -1, np)], current[i], current[i
dx(i, +1, np)]);
            }
            // every nd time steps, attach additional continuation whi
ch will
            // trigger the semaphore once computation has reached this
point
            if ((t % nd) == 0)
                next[0].then(
                    [&sem, t](partition &&)
```

```
{
                       // inform semaphore about new lower limit
                       sem.signal(t);
                   });
           }
           // suspend if the tree has become too deep, the continuati
on above
           // will resume this thread once the computation has caught
up
           sem.wait(t);
       }
       // Return the solution at time-step 'nt'.
       return hpx::when all(U[nt % 2]);
   }
};
std::uint64_t np{10};
                          // Number of partitions.
   std::uint64 t nx{10};
                         // Number of grid points. (local x dimensi
on of each partition)
   std::uint64 t nt{45};
                         // Number of steps.
   std::uint64 t nd{10};
                         // Max depth of dep tree.
*/
void do_1d_solve(std::uint64_t np, std::uint64_t nx,
                std::uint64 t nt, std::uint64 t nd, bool results)
{
   header = false;
   // Create the stepper object
   stepper step;
   // Measure execution time.
   std::uint64 t t = hpx::util::high resolution clock::now();
   // Execute nt time steps on nx grid points and print the final sol
ution.
   hpx::future<stepper::space> result = step.do_work(np, nx, nt, nd);
   stepper::space solution = result.get();
   hpx::wait_all(solution);
   std::uint64 t elapsed = hpx::util::high resolution clock::now() -
t;
   // Print the final solution
   if (results)
   {
       for (std::size t i = 0; i != np; ++i)
           std::cout << "U[" << i << "] = " << solution[i].get() << s
td::endl;
   }
   std::uint64 t const os thread count = hpx::get os thread count();
```

```
print_time_results(os_thread_count, elapsed, nx, np, nt, header);

return;
}

In []: .expr
apex::reset(0L);
do_ld_solve(100, 100, 450, 100, false);
apex::dump(true);
```

Heat Equation with an APEX policy

```
In [ ]: #include <hpx/hpx init.hpp>
        #include <hpx/hpx.hpp>
        #include <hpx/include/parallel_algorithm.hpp>
        #include <hpx/include/performance counters.hpp>
        #include <boost/range/irange.hpp>
        #include <boost/format.hpp>
        #include <algorithm>
        #include <cstddef>
        #include <cstdint>
        #include <iostream>
        #include < limits >
        #include <memory>
        #include <string>
        #include <utility>
        #include <vector>
        #include <apex_api.hpp>
        #include <boost/shared_array.hpp>
        using hpx::naming::id_type;
        using hpx::performance counters::get counter;
        using hpx::performance_counters::stubs::performance_counter;
        using hpx::performance counters::counter value;
        using hpx::performance_counters::status_is_valid;
        static bool counters initialized = false;
        static std::string counter name = "/threads{locality#0/total}/idle-rat
        static apex event type end iteration event = APEX CUSTOM EVENT 1;
        static hpx::naming::id type counter id;
        void setup counters() {
            try {
                id_type id = get_counter(counter_name);
                // We need to explicitly start all counters before we can use
         them. For
```

```
// certain counters this could be a no-op, in which case start
will return
        // 'false'.
        performance counter::start(hpx::launch::sync, id);
        std::cout << "Counter " << counter name << " initialized " <<
id << std::endl;</pre>
        counter value value = performance counter::get value(hpx::laun
ch::sync, id);
        std::cout << "Counter value " << value.get value<std::int64 t>
() << std::endl;</pre>
        counter id = id;
        end_iteration_event = apex::register_custom_event("Repartitio")
n");
        counters initialized = true;
    catch(hpx::exception const& e) {
        std::cerr << "ld stencil 4 repart: caught exception: "</pre>
            << e.what() << std::endl;
        counter id = hpx::naming::invalid id;
        return;
   }
}
double get_counter_value() {
    if (!counters initialized) {
        std::cerr << "get counter value(): ERROR: counter was not init</pre>
ialized"
            << std::endl;
        return false;
    }
    try {
        counter value value1 =
            performance counter::get value(hpx::launch::sync, counter
id, true);
        std::int64 t counter value = value1.get value<std::int64 t>();
        std::cout << "counter value " << counter value << std::endl;</pre>
        return (double)(counter value);
    }
    catch(hpx::exception const& e) {
        std::cout << "get_counter_value(): caught exception: " << e.wh</pre>
at()
            << std::endl;
        return (std::numeric limits<double>::max)();
    }
// Our partition data type
struct partition data2
{
public:
    partition data2(std::size t size)
      : data (new double[size]), size (size)
    {}
```

```
partition data2(std::size t size, double initial value)
     : data (new double[size]),
       size_(size)
   {
       double base value = double(initial value * size);
       for (std::size_t i = 0; i != size; ++i)
           data_[i] = base_value + double(i);
   }
   partition data2(std::size t size, const double * other)
        : data (new double[size]),
        size_(size)
   {
       for(std::size_t i = 0; i != size; ++i) {
           data [i] = other[i];
       }
   }
   partition_data2(partition_data2 && other)
     : data (std::move(other.data ))
     , size_(other.size_)
   {}
   double& operator[](std::size_t idx) { return data_[idx]; }
   double operator[](std::size_t idx) const { return data_[idx]; }
   void copy into array(double * a) const
   {
       for(std::size t i = 0; i != size(); ++i) {
           a[i] = data [i];
       }
   }
   std::size_t size() const { return size_; }
private:
   std::unique_ptr<double[]> data_;
   std::size_t size_;
};
std::ostream& operator<<(std::ostream& os, partition data2 const& c)</pre>
{
   os << "{";
   for (std::size t i = 0; i != c.size(); ++i)
   {
       if (i != 0)
           os << ", ";
       os << c[i];
   os << "}";
   return os;
}
```

```
struct stepper2
    // Our data for one time step
    typedef hpx::shared future<partition data2> partition;
    typedef std::vector<partition> space;
   // Our operator
   static inline double heat(double left, double middle, double right
)
   {
        return middle + (k*dt/dx*dx) * (left - 2*middle + right);
    }
   // The partitioned operator, it invokes the heat operator above on
all
    // elements of a partition.
    static partition data2 heat part(partition data2 const& left,
        partition data2 const& middle, partition data2 const& right)
    {
        std::size t size = middle.size();
        partition data2 next(size);
        if(size == 1) {
            next[0] = heat(left[0], middle[0], right[0]);
            return next;
        next[0] = heat(left[size-1], middle[0], middle[1]);
        for(std::size t i = 1; i < size-1; ++i)</pre>
            next[i] = heat(middle[i-1], middle[i], middle[i+1]);
        next[size-1] = heat(middle[size-2], middle[size-1], right[0]);
        return next;
   }
   // do all the work on 'np' partitions, 'nx' data points each, for
   // time steps
   hpx::future<space> do_work(std::size_t np, std::size_t nx, std::si
ze_t nt,
        boost::shared array<double> data)
    {
        using hpx::dataflow;
        using hpx::util::unwrapping;
        // U[t][i] is the state of position i at time t.
        std::vector<space> U(2);
        for (space& s: U)
            s.resize(np);
        if (!data) {
            // Initial conditions: f(0, i) = i
            std::size t b = 0;
```

```
auto range = boost::irange(b, np);
           using hpx::parallel::execution::par;
           hpx::parallel::for each(
               par, std::begin(range), std::end(range),
               [&U, nx](std::size t i)
                   U[0][i] = hpx::make ready future(
                       partition data2(nx, double(i)));
               }
           );
       else {
           // Initialize from existing data
           std::size t b = 0;
           auto range = boost::irange(b, np);
           using hpx::parallel::execution::par;
           hpx::parallel::for each(
               par, std::begin(range), std::end(range),
               [&U, nx, data](std::size_t i)
               {
                   U[0][i] = hpx::make ready future(
                       partition_data2(nx, data.get()+(i*nx)));
           );
       }
       auto Op = hpx::util::annotated function(unwrapping(&stepper2
::heat_part),
                                              "stepper2::heat part"
);
       // Actual time step loop
       for (std::size t t = 0; t != nt; ++t)
           space const& current = U[t % 2];
           space\& next = U[(t + 1) % 2];
           for (std::size_t i = 0; i != np; ++i)
               next[i] = dataflow(
                       hpx::launch::async, Op,
                       current[idx(i, -1, np)], current[i], current[i
dx(i, +1, np)
                   );
           }
       }
       // Return the solution at time-step 'nt'.
       return hpx::when all(U[nt % 2]);
   }
};
/*
   std::uint64 t np{10};
                          // Number of partitions.
   std::uint64 t nx{10};
                          // Number of grid points. (local x dimensi
```

```
on of each partition)
    std::uint64_t nt{45}; // Number of steps.
    std::uint64 t nr{10}; // Number of runs to tune
*/
void do 1d solve repart(std::uint64 t nx,
                 std::uint64 t nt, std::uint64 t nr, bool results)
{
    setup counters();
    /* Number of partitions dynamically determined
                                                       */
    header = false;
    std::uint64 t const os thread count = hpx::get os thread count();
   // Find divisors of nx
    std::vector<std::uint64 t> divisors;
    // Start with os thread count so we have at least as many
    // partitions as we have HPX threads.
    for(std::uint64_t i = os_thread_count; i < std::sqrt(nx); ++i) {</pre>
        if(nx \% i == 0) {
            divisors.push back(i);
            divisors.push back(nx/i);
        }
    }
   // This is not necessarily correct (sqrt(x) does not always evenly
divide x)
   // and leads to partition size = 1 which we want to avoid
   //divisors.push back(static cast<std::uint64 t>(std::sqrt(nx)));
    std::sort(divisors.begin(), divisors.end());
    if(divisors.size() == 0) {
        std::cerr << "ERROR: No possible divisors for " << nx</pre>
            << " data elements with at least " << os_thread_count
            << " partitions and at least two elements per partition."
            << std::endl;
        return hpx::finalize();
   }
    //std::cerr << "Divisors: ";
    //for(std::uint64 t d : divisors) {
   //
          std::cerr << d << " ";
   //}
   //std::cerr << std::endl;
   // Set up APEX tuning
   // The tunable parameter -- how many partitions to divide data int
0
    long np index = 1;
    long * tune params[1] = { 0L };
    long num params = 1;
    long mins[1] = \{0\};
    long maxs[1] = { (long)divisors.size() };
    long steps[1] = \{1\};
    tune params[0] = &np index;
    apex::setup custom tuning(get counter value, end iteration event,
num params,
            tune params, mins, maxs, steps);
```