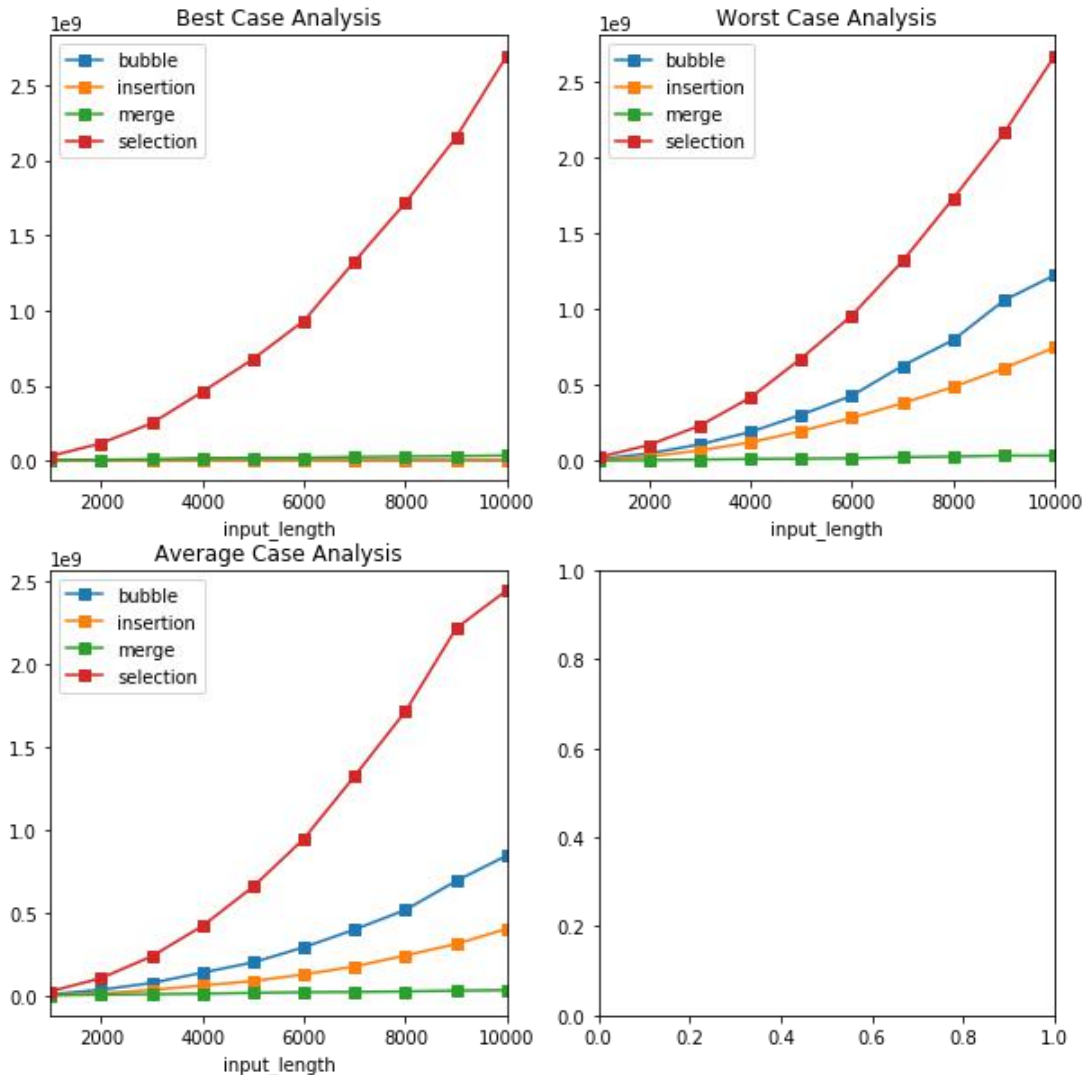


1. Empirical Analysis:

Best Case, Worst Case, Average Case comparisons for the 4 sorting algorithms.



2. Asymptotic analysis:

Sort Types	Best Case Performance	Average Case Performance	Worst Case Performance
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Merge Sort	$\theta(n \log(n))$	$\theta(n \log(n))$	$\theta(n \log(n))$
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$

Bubble Sort:

Analysis:

For Worst and Average case, there would be $n-1$ comparisons in the 1st pass, $n-2$ comparisons in the 2nd pass and so on.

$$T(n) = \sum_{k=1}^{n-1} k = n(n-1)/2 \implies O(n^2).$$

For Best case, the loop will run n times, $T(n) \implies O(n)$

Insertion Sort:

Analysis:

For Worst Case:

$$T(n) = \sum_{k=1}^n k = n(n+1)/2 \implies O(n^2).$$

For Average Case and Best Case:

$$T(n) = \sum_{k=1}^n k / 2 = n(n+1)/4 \implies O(n^2).$$

Merge Sort:

Analysis:

For Best, Average, and Worst cases the sort behaves in the same manner i.e input list is divided into 2 parts and solved recursively.

$$T(n) = 2T(n/2) + \theta(n)$$

Using the master theorem, we get

$$T(n) = \theta(n \log(n))$$

Selection Sort:

Analysis:

For Best, Average, and Worst cases the sort behaves in the same manner i.e the selection sort has 2 nested for loops, so the time complexity is as mentioned.

$$T(n) = \sum_{k=1}^{n-1} k = n(n-1)/2 \implies O(n^2)$$