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II)

Recurrence relation  $T(n) = T(n-1) + n$

Decreasing func.

• void Test(int n) —  $T(n)$  time.

```

    {
        if (n > 0) — 1
        {
            for (i = 0; i < n; i++)
            {
                printf(".\n", n); — n
            }
            Test(n - 1); —  $T(n-1)$ 
        }
    }
  
```

∴ Recurrence relation —

$$T(n) = T(n-1) + 2n + 2$$

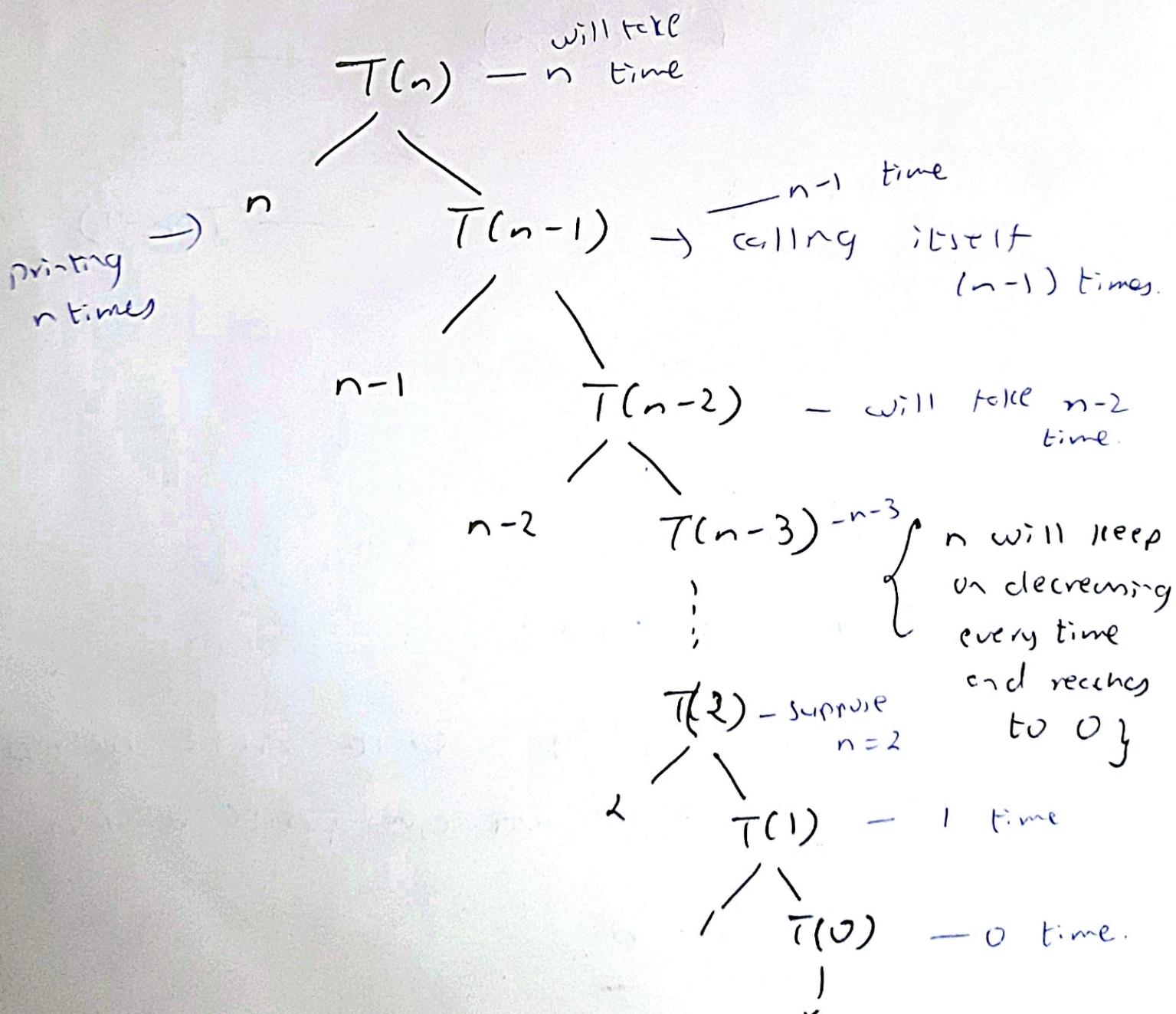
In order to simplify, since  $2n+2$  belongs to linear class, we can write only  $n$ .

$$\therefore \underline{T(n) = T(n-1) + n}$$

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + n & n>0 \end{cases} \quad \textcircled{1} \quad \text{Recurrence relation}$$

So we have prepared recurrence relation for a given program or algorithm.

Now solving  $\textcircled{1}$  using Recursion Tree



(3)

Summing all

$$\begin{aligned} T(n) &= 0 + 1 + 2 + \dots + n-2 + n-1 + n \\ &= \frac{n(n+1)}{2} \end{aligned}$$

$$\therefore T(n) = \frac{n(n+1)}{2}$$

$\rightarrow = \underline{\underline{\Theta(n^2)}} \text{ or } O(n^2)$

Solved using

Recurrence tree



$\rightarrow$  Solving above by Substitution method -

$$\rightarrow T(n) = \begin{cases} 1 & n=0 \\ T(n-1)+n & n>0 \end{cases}$$

Soln

$$T(n) = T(n-1) + n$$

Since  $\Rightarrow T(n) = T(n-1) + n$

$\therefore T(n-1) = T(n-2) + n-1$

(In place of  $n$ , put  $n-1$ )

So substitute  $T(n-1)$  to  $T(n-2) + n-1$  (4)

$$\therefore T(n) = \underline{\underline{[T(n-2) + n-1]}} + n - (1)$$

$$T(n) = T(n-2) + (n-1) + n - (2)$$

Substitute for this also.

$$\therefore T(n-2) = \underline{\underline{T(n-3) + n-2}}$$

After substituting

$$\begin{aligned} T(n) &= [T(n-3) + n-2] + (n-1) + n \\ &= T(n-3) + \underline{n-2} + n-1 + n - (3) \end{aligned}$$

So after substitution for 2 times, we have an idea of sequence. So if we do it for  $k$  times then -

∴  $k$  times

similar to  $n-2$   
as  $n-2$  is 1 less than  
 $n-3$ .  
so it is  
 $k-1$ .

$$\begin{aligned} T(n) &= T(n-k) + (n-(\underline{k-1})) + \\ &\quad (n-(k-2)) + \dots + (\cancel{(k)}) + n. - (4) \end{aligned}$$

(J)

→ Assume  $n-k=0$

$$\therefore n=k.$$

→ In (4) substitute  $n=k$ .

$$\therefore T(n) = T(n-n) + (n-n+1) + \\ (n-n+2) + \dots + (n-n+n) + n$$

$$\therefore T(n) = T(0) + 1 + 2 + 3 + \dots + \underbrace{(n)}_{\text{sum of } n \text{ nos.}} + n$$

$$\therefore T(n) = 1 + \frac{n(n+1)}{2} \\ = \underline{\underline{\Theta(n^2)}}$$

Same as recursive tree method.

III)

Consider an example —

$$\underline{\underline{T(n) = T(n-1) + \log n.}}$$

- Decreasing func.

(6)

→ `void Test(int n)` — will take  $T(n)$  time

```

{
    if (n > 0)
    {
        for (i=1; i < n; i = i*2)
        {
            printf(".d", i);
        }
    }
    Test(n-1);
}

```

→  $T(n) = T(n-1) + \log n$

∴ Recurrence relation is -

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + \log n & n>0 \end{cases}$$

→ Solving using Tree method -