CMBThe SML/NJ Bootstrap Compiler

(for SML/NJ version 110.35 and later)

User Manual

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Contents

1	Intr	oduction	2			
2	Basic usage					
	2.1	Requirements	2			
	2.2	Invoking the bootstrap compiler	3			
	2.3	Booting a new interactive system	4			
	2.4	Testing the newly booted system	4			
	2.5	Installing a newly booted system	4			
3	Diff	Ferences between CMB and CM	5			
	3.1	Code sharing	5			
	3.2	Init library	5			
		3.2.1 Linkage to runtime system	6			
	3.3	BOOTLIST file	6			
	3.4	PIDMAP file	6			
	3.5	Cross-compiling	6			
4	Stru	acture CMB	7			
5	File	naming—the role of path anchors	8			
	5.1	Anchoring requirement	8			
	5.2	Different anchor mappings at different times	8			
6	Scripts					
	6.1	The makeml script	9			
	6.2	The testml script	10			
	6.3	The installml script	10			
	6.4	The fixpt script	10			
	6.5	The alloross script	10			

1 Introduction

With the exception of the runtime system which is written in C, compiler and interactive system for Standard ML of New Jersey (SML/NJ) [AM91] are themselves implemented in Standard ML [MTHM97]. When the new Compilation and Library Manager (CM) [Blu00] was introduced, all ML code of SML/NJ was reorganized in such a way that the system itself works like any other stand-alone ML program compiled by SML/NJ.

However, there are a few important differences between compiling an ordinary ML program using CM and (re-)compiling the interactive system itself. These differences are handled by a special-purpose version of the compilation manager: the bootstrap compiler. Its interface is structure CMB, exported from library \$smlnj/cmb.cm.

This document describes how to use the bootstrap compiler and also explains how bootstrapping the SML/NJ compiler differs from compiling ordinary ML code.

2 Basic usage

2.1 Requirements

To be able to use the bootstrap compiler, one must first install SML/NJ (i.e., the interactive system that contains the compiler) as well as both ML-Yacc and ML-Lex.

It is further necessary to have all ML source code for the system available. If the basic installation did not install this source code, then one must now fetch all source archives and unpack them. (There is an option for the SML/NJ installer that lets it install all source trees automatically. However, by default this feature is turned off.)

The following list shows all required source packages. (Path names are shown relative to the SML/NJ installation directory.)

src/system This archive contains the sources for the SML Standard Basis library as well as lots of "glue" code. The glue is used for assembling a complete system from all its other parts. Directory src/system must be the current working directory at the time the bootstrap compiler is started.

src/MLRISC This is the implementation of the MLRISC backend (the low-level optimizer and code generator) used by SML/NJ.

src/cm This source tree contains most of CM's (and CMB's) implementation.

src/compiler This is the implementation of the frontend (parser, type checker, etc.) of the compiler as well as its high-level optimizer (FLINT).

src/ml-yacc This source tree contains the implementation of ML-Yacc and its library (\$/ml-yacc-lib.cm). Technically, the sources of ML-Yacc itself are not needed (provided a working executable for ML-Yacc has been installed), but the library sources are.

src/smlnj-lib This source tree hosts several sub-trees, each of which implements one of the libraries in the SML/NJ library collection. For bootstrap compilation, the following sub-trees are required:

src/smlnj-lib/Util This directory holds the sources for \$/smlnj-lib.cm.

src/smlnj-lib/PP This directory holds the sourcse for \$/pp-lib.cm, i.e., the pretty-printing library.

src/smlnj-lib/HTML This directory holds the sources for \$/html-lib.cm, a library for handling HTML files.
The need for this library arises from the fact that \$/pp-lib.cm statically depends on it. (The compiler does not actually use any services from this library.)

2.2 Invoking the bootstrap compiler

To invoke the bootstrap compiler, first one has to change the current working directory to src/system:

```
$ cd src/system
```

The next step is to start the interactive system and load the bootstrap compiler. This can be done in one of two ways:

1. Start the interactive system and then issue a CM.autoload command that causes the bootstrap compiler to be loaded. The resulting session could look like this:

```
$ sml
Standard ML of New Jersey ...
- CM.autoload "$smlnj/cmb.cm";
...
val it = true : bool
-
```

2. Start the interactive system and specify \$smlnj/cmb.cm on the command line:

```
$ sml '$smlnj/cmb.cm'
Standard ML of New Jersey ...
```

Note for frequent compiler hackers: The makeml script (see below) builds the the interactive system in such a way that \$smlnj/cmb.cm is already pre-registered for autoloading. Therefore, when using an interactive system built by makeml (as opposed to the original config/install.sh) there is no need for loading the bootstrap compiler explicitly.

At this point one can invoke the bootstrap compiler by simply issuing the command CMB.make():

```
- CMB.make ();
```

If CMB.make() does not run to successful completion, you do not have to start from the beginning. Instead, fix the problem at hand and re-issue CMB.make() without terminating the interactive session in between. This tends to be a lot faster than starting over.

This process can be repeated arbitrarily many times until CMB.make() is successful.

A successful run looks like this:

```
- CMB.make ();
...
New boot directory has been built.
val it = true : bool
-
```

The return value of true indicates success. This means that (as indicated by the message above the return value) a directory with stable libraries and some other special files that are needed for "booting" a new interactive system has been created.

The name of the boot directory depends on circumstances. A part of it can be chosen freely, other parts depend on current architecture and operating system. For example, on a Sparc system running some version of Unix, the default name of the directory is sml.boot.sparc-unix.

There is also a similarly-named em binfile directory which is used by CMB.make() itself but which is not required for the purpose of subsequent "boot" steps.

2.3 Booting a new interactive system

Once CMB.make() has completed its work successfully (indicated by its return value of true), the next step is to build a new heap image. To do so, issue the command makeml at the shell prompt:

```
$ ./makeml
...
Standard ML of New Jersey ...
./makeml: Heap image generated.
```

The makeml command generates a new heap image and also prepares stable libraries to be used by this image. Neither the heap image nor the libraries will at this time be installed for permanent use though. This means that invoking sml still starts the old system.

Again, the name of the generated heap image is variable and depends on programmer choice, current architecture, and current operating system. For example, on an Intel x86 Linux machine, the default would be sml.x86-linux.

2.4 Testing the newly booted system

To test-drive a newly booted system without installing it, issue the testml command at the shell prompt:

```
$ ./testml
```

This starts an interactive system in a way very similar to sml, but it uses the heap image and libraries from a previous run of makeml.

2.5 Installing a newly booted system

Once one is sure that a newly booted system is good enough to replace the old system, issuing the installml command will replace old heap image and old libraries with those generated by the previous run of makeml.

```
$ ./installml
```

This command will replace the system's heap image in ../../bin/.heap and its libraries in ../../lib. However, it will leave alone any unrelated libraries in ../../lib.

Sometimes changes to the compiler will render any previously installed libraries unusable. In this case one should erase them prior to issuing the installml command:

```
$ rm ../../lib/*
$ ./installml
```

Libraries that were installed as part of the SML/NJ installation process but which are unrelated to bootstrap compilation (e.g., \$/inet-lib.cm, CML, eXene) can be recovered (once they had been removed) by going back to the installation directory and issuing the config/install.sh command again:

```
$ cd ../..
$ config/install.sh
```

Since some changes to the compiler also render old binfiles unusable, one will occasionally have to remove those first (prior to re-running config/install.sh). Binfiles for libraries unrelated to bootstrapping are handled by CM (and not CMB), so the usual CM rules for locating them apply. (This means that in such a case the binfile for d/f.sml will be in d/CM/arch-os/f.sml where arch is a string describing the CPU architecture and os is a string describing the operating system kind. Example: x86-unix.)

¹The operating-system dependent portion of the name of a heap image is more discriminating than the operating-system dependent portion of bootor binfile directories. On the same Intel x86 Linux machine, the name of the boot directory is sml.boot.x86-unix.

3 Differences between CMB and CM

In this section we discuss why compiling the compiler is different from compiling other ML programs. Each of the following sub-sections focuses on one particular aspect.

3.1 Code sharing

CM keeps compiled code within the same directory tree that contains the corresponding ML source code. Thus, there is a fixed function that maps the names of source files to the names of corresponding binfiles and the names of CM description files (for libraries) to their corresponding stable files.

As a result, these files will be shared between programs that use the same libraries. Moreover, CM will let different programs that are loaded into an interactive session at the same time share their in-memory copies of common compiled modules. (There is also an issue of state-sharing, but this does not concern CMB because the bootstrap compiler only compiles code without linking it.)

Sharing of code is useful for ordinary usage, but when compiling the compiler itself, it is not desirable. During bootstrap compilation, it is often the case that several different versions of compiled code have to coexist. Some, or even all of these versions can differ significantly from those of the currently running system.

Therefore, CMB keeps binfiles and stable files in separate directory trees. The names of the directories where these trees are rooted at are constructed from three parts; the binfile directory's name is u.bin.arch-os and the stablefile directory's name is u.boot.arch-os. As mentioned before, arch is a string describing the current CPU architecture and os is a string describing the current operating system kind. Component u is a string that can be selected freely when the bootstrap compiler is invoked. When using CMB.make it defaults to sml, otherwise it is the argument given to CMB.make'. (The u component is kept variable to make it possible to keep and use several compiled versions of the system at the same time.)

The auxiliary script makeml (which is responsible for bootstrapping a new system) also accepts a parameters to select u. If the parameter is missing, it defaults to sml (in accordance with CMB .make's behavior).

3.2 Init library

The *init library* (\$smlnj/internal/init.cmi) is a library that is used implicitly by all programs. This library is "special" in several ways and cannot be described using an ordinary CM description file. It is the bootstrap compiler's responsibility to properly prepare a stable version.

Ordinary programs (those managed by CM) do not have to worry about the special aspects of how to construct this library; they just have to be able to use its stable version.

There are several reasons why the library cannot be described as an ordinary CM library:

- The library exports the *pervasive environment* which normally is imported implicitly by every compilation unit. Within the init library, no pervasive environment is available yet.
- One binding in the above-mentioned pervasive environment is a binding for structure _Core. The symbol _Core is not a legal SML identifier, and the bootstrap compiler has to take special action to create a binding for it anyway.
- One of the compilation units in this library is merely a placeholder which at link time has to be replaced by the SML/NJ runtime system (which is written in C).

3.2.1 Linkage to runtime system

The ML source file dummy.sml (located in directory src/system/smlnj/init) contains a carefully constructed module whose signature matches that of the runtime system's binary API. This file is being compiled as part of constructing the init library, but it has been marked specially as *runtime system placeholder*. During compilation, this file pretends to *be* the runtime system; other modules that use the runtime system "think" they are using dummy.sml.

At link time (i.e., bootstrap time—when makeml is run), the boot loader will ignore dummy.sml and use the actual runtime system in its place. This trick makes it possible that (from the point of view of all other modules) using services from the runtime system appears to be no different than using services from an ordinary ML compilation unit.

3.3 BOOTLIST file

Linking of SML/NJ programs involves executing the code of each of the concerned compilation units. The code of each compilation unit is technically a closed function; all its imports have been turned into arguments and all exports have been turned into return values.

For ordinary programs, this process is under control of CM; CM will take care of properly passing the exports of one compilation unit to the imports of the next.

When booting a stand-alone program, though, there is no CM available yet. Thus, executing module-level code and passing exports to imports has to be done by the (bare) runtime system. The runtime system understands enough about the layout of binfiles and library files so that it can do that—provided there is a special *bootlist* file that contains instructions about which modules to load in what order.

The bootlist mechanism is not restricted to building SML/NJ. Ordinary ML code can also be turned into stand-alone programs, and as far as the runtime system is concerned, the mechanisms are the same. The bootlist file used by such ordinary stand-alone ML programs will be constructed by CM; only in the case of bootstrapping SML/NJ itself it will be constructed by CMB.

The name of the bootlist file is BOOTLIST, and it is located at the root of the directory tree that contains stable files (i.e., its name is u, boot, arch-os/BOOTLIST).

3.4 PIDMAP file

The last file to be loaded by the bootstrap process contains module-level code which will trigger the self-initialization of the interactive system—including CM. One job of CM is to manage sharing of link-time state (i.e., dynamic state created by module-level code at link time). Link-time state of a module used by the interactive system should be shared with any program using the same module. The file u.boot.arch-os/PIDMAP contains information that enables CM to relate existing link-time state to particular library modules and also to identify any link-time state that will never be shared and which can therefore be dropped. It is CMB's responsibility to construct the PIDMAP file.

3.5 Cross-compiling

Several different versions of the bootstrap compiler can coexist—each being responsible for targeting another CPU-OS combination. Structure CMB is the default bootstrap compiler that targets the current system; it is exported from \$smlnj/cmb.cm. The following table lists the names of other structures—those corresponding to various cross-compilers. All these structures share the same signature.

The table also shows the names of libraries that the structures are exported from as well as those *arch* and *os* strings that are used to name binfile- and stablefile-directory.

library	structure	architecture	OS	arch	os
\$smlnj/cmb.cm	CMB	current	current		
\$smlnj/cmb/current.cm					
\$smlnj/cmb/alpha32-unix.cm	Alpha32UnixCMB	Alpha	Unix	alpha32	unix
\$smlnj/cmb/hppa-unix.cm	HPPAUnixCMB	HP-PA	Unix	hppa	unix
\$smlnj/cmb/ppc-macos.cm	PPCMacOSCMB	Power-PC	Mac-OS	ppc	macos
\$smlnj/cmb/ppc-unix.cm	PPCUnixCMB	Power-PC	Unix	ppc	unix
\$smlnj/cmb/sparc-unix.cm	SparcUnixCMB	Sparc	Unix	sparc	unix
\$smlnj/cmb/x86-unix.cm	X86UnixCMB	Intel x86	Unix	x86	unix
\$smlnj/cmb/x86-win32.cm	X86Win32CMB	Intel x86	Win32	x86	win32
\$smlnj/cmb/all.cm	all of the above (except CMB)				

As an example, consider targeting a Sparc/Unix system. The first step is to load the library that exports the corresponding cross-compiler:

```
CM.autoload "$smlnj/cmb/sparc-unix.cm";
```

Once this is done, run the equivalent of CMB.make:

```
SparcUnixCMB.make ();
```

This will recompile the compiler, producing object code for a Sparc. Binfiles will be stored under sml.bin.sparc-unix and stable libraries under sml.boot.sparc-unix.

4 Structure CMB

This section describes the signature of structure CMB. Since structures representing cross-compilers have the same signature, everything said here applies (*mutatis mutandis*) to them as well.

The primary function to invoke the bootstrap compiler is CMB.make':

```
val make' : string option -> bool
```

This (re-)compiles the interactive system's entire source tree, constructing stable versions for all libraries involved. In the process, binfiles are placed under directory u. bin. arch-os where arch and os are strings describing target architecture and target OS, respectively. The string u is the optional argument to CMB.make'. If set to NONE, it defaults to "sml".

An alternative equivalent to invoking CMB.make' with NONE is to use CMB.make:

```
val make : unit -> bool
```

CMB—like CM—maintains a lot of internal state to speed up repeated invocations. (Between sessions, much of this state is preserved in those binfile- and stablefile-directories. However, reloading is still quite a bit more expensive than directly using existing in-core information.)

Information that CMB keeps in memory can be completely erased by issuing the CMB.reset command:

```
val reset : unit -> unit
```

After a CMB.reset (), the next CMB.make (or CMB.make') will have to re-load everything from the file system.

CMB has its own registry of "CM identifiers"—named values that can be queried by using the conditional compilation facility. This registry is initialized according to CM's rules. Of course, initial values are not based on current architecture and OS but on those of the target system. To explicitly set or erase the values of specific variables, one can use CMB.symval (which acts in a way analogous to CM.symval):

5 File naming—the role of path anchors

Under normal operation of CM, the mapping from path names to the files they denote is supposed to be a fixed one. The path anchor mechanism is merely a means of configuring this mapping according to the actual filesystem layout.

The bootstrap compiler and the associated boot procedure, however, use path anchors extensively for other purposes. In particular, the above-mentioned mapping will vary over time.

The "casual" compiler hacker does not actually need to remember all the details because these details are mostly taken care of automatically by CMB and its various scripts (makeml, testml, installml). But it is useful to have a rough idea of what rôles are being played by various files and directories that are being used or created.

5.1 Anchoring requirement

During bootstrap compilation it is necessary that all pathnames be either anchored or relative to another anchored pathname. This rule guarantees that every path name can be mapped to different locations in the filesystem at different times. (Technically, this restriction is necessary only for library description files.)

5.2 Different anchor mappings at different times

We can distinguish between four different anchor configurations which will be in effect at different times:

compilation: At the time CMB.make runs, the anchor configuration is taken from file pathconfig.² This configuration maps anchors to their respective directories in the actual source tree. At the same time the policy which determines the names of binfiles and stablefiles is modified in such a way that all binfiles will be created in directory u.bin.arch-os and all stablefiles will be created in u.boot.arch-os. In particular, if a is the root anchor name controlling the path leading to an ML source file $a/\cdots/file$ sml, then the corresponding binfile will be created as $a.bin.arch-os/a/\cdots/CM/arch-os/file$.sml. Similarly, the stablefile for the library description $a/\cdots/file$ cm controlled by root anchor a gets written to $a.boot.arch-os/a/\cdots/CM/arch-os/file$.cm. (Notice how the anchor name a is being incorporated into the resulting pathnames.)

boot: At bootstrap time (e.g., when the makeml script is invoked), CM scans the contents of u.boot.arch-os, and for each directory name a it finds there it maps the anchor a to that directory. The result is that the location for library description a/c.../file.cm is considered to be a.bin.arch-os/a/.../file.cm, which means that—under the a.bin.arch-os/a/.../cm/arch-os/file.cm. This is precisely where CMB.make created it. (Of course, the libary description file itself does not exist, but this is not a problem because the stablefile does.)

The resulting mapping for anchors is the one that is being saved to the newly generated heap image.

test: After generating a new heap image $v.arch-osname^3$, the makem1 script will copy the hierarchy of directories under u.boot.arch-os to a new directory v.lib and generate a path configuration file v.lib/pathconfig for it. Files in this hierarchy will be re-created as hard links. The prefix v can be chosen at the time makem1 is invoked (see section 6.1)—just like v can be chosen at the time CMB.make is invoked. The default for v is sml. Copying the directory prevents any future CMB.make from clobbering the libraries that belong to the newly generated heap image.

Running the testml script will start the runtime system, instruct it to load heap image v.arch-osname, and arrange for the path configuration file v.lib/pathconfig to take effect. This will cause any anchor a to be resolved within the v.lib hierarchy. Thus, the new system can be tested in a way that is completely independent from any existing installation. (Again, v can be specified as a parameter to testml.)

²All relative pathnames shown here are relative to src/system.

³As mentioned earlier, we distinguish between *os* which describes the *kind* of operating system and *osname* which is the name of a particular operating system. For example, on a Linux machine *osname* is linux while *os* is still unix.

install: When testing has been satisfactory, the new system can be installed permamently, replacing the old heap image and its old libraries with their newly created counterparts. To do so, one should run the installml script. It will move the heap image v.arch-osname to ../../bin/.heap/sml.arch-osname and all stablefiles from v.lib to their usual place under ../../lib. It then proceeds to edit ../../lib/pathconfig (the main path configuration file of the installation) to reflect the new situation. (Library files in ../../lib that have nothing to do with the bootstrap process will remain untouched.)

6 Scripts

This section gives a detailed description for each script, its function, and its options. All scripts described here are located in the src/system directory.⁴

6.1 The makeml script

This script is used after a successful run of CMB.make (or CMB.make'). It *links* the compiler and its interactive system, forming a corresponding heap image. In addition to that, it prepares a new directory containing stable CM libraries to be used with the new image.

One way of thinking about this is to view makem1 as a function mapping a bootfile directory b to a pair consisting of a heap image v. arch-osname and a library directory v. lib. The strings b and v are optional parameters; the defaults are sml.boot.arch-os for b and sml for v.

The script accepts a number of options as follows:

- **-o** v specifies a v other than the default sml.
- **-boot** b specifies a b other than the default sml.boot.arch-os.
- **-quiet** instructs makeml to greatly reduce its diagnostic output. In particular, the names of files being linked are not shown. The default for this can be controlled via a boolean-valued environment variable MAKEML_VERBOSITY. If the variable is not set, then the default is true (meaning "verbose").
- -verbose is the opposite of -quiet.
- -rebuild u puts makeml into a different mode: After loading the executable section of all binfiles and linking them, it will not read any static environments, will not initialize the usual interactive system and will not produce a heap image. Instead, it internally invokes the equivalent of CMB.make' (SOME "u"), thus recompiling everything again. When recompilation is complete, makeml stops; the newly-built system must be linked using a separate explicit invocation of makeml. Notice that u.boot.arch-os must be different from b.
- **-rebuildlight** u is the same as <code>-rebuild</code> u except that the symbol <code>LIGHT</code> will be defined (using <code>CMB.symval</code>) for the duration of the compilation. The effect of this is that no cross-compilers will be built (which can save considerable time). Alternative names for **-rebuildlight** are **-light** and **-lightrebuild**.
- **-bare** causes makeml to build a reduced version of the system without the compilation manager CM included. This is useful for people who are interested in an interactive system only.
- **-run** r selects the executable for the SML/NJ runtime system. The default is ../../bin/.run/run.arch-osname.
- -alloc a specifies the size of the SML/NJ garbage-collector's allocation area. The default depends on the current machine architecture.

The most common usage is to simply say . /makeml without any arguments, taking advantage of the defaults as described above

⁴If the current directory . is not on the shell's search path, then each script must be invoked using an explicit ./-prefix. Example: ./makeml.

6.2 The testml script

The testml script launches a newly-built and newly-linked interactive system. It requires a previous run of the makeml script because it uses the heap image v.arch-osname and libraries in v.lib.

If no arguments are given, v defaults to sml. If at least one argument is given, then v is assumed to be the first one. All other arguments are passed to the initialization routine of the interactive system just like normal arguments to the sml command would.

This means that one *must* specify a v unless there are no command-line arguments at all.

6.3 The installml script

The installml script moves a heap image v.arch-osname to ../../bin/.heap/sml.arch-osname. Moreover, it moves the libraries under v.lib to ../../lib and updates ../../lib/pathconfig accordingly.

The script takes one optional argument which specifies v. If no argument is given, v defaults to sml.

6.4 The fixpt script

One good test of whether a new version of the compiler is working properly is to see if it compiles to a *fixed point*. For this, the compiler is compiled n times with the result of the (k-1)st compilation being responsible for compiling the kth version. A fixed point is reached if two consecutive compilations produce identical results.

The fixpt script automates the task of compiling to a fixed point. It internally uses the sml command for the first compilation and the makeml script with its -rebuild parameter for all subsequent runs. This produces a series of bin-directories u.bin.arch-os, u1.bin.arch-os, u2.bin.arch-os,... and boot-directories u.boot.arch-os, u1.boot.arch-os,... where u is a common "stem" that is used for naming the whole series.

The fixpt script accepts the following options:

- **-iter** n limits the number of iterations to n. The default for n is 3. If no fixed point is found after n iterations, then fixpt terminates with an error message.
- **-base** u selects the "stem" u (see above). The default is sml.
- -light causes LIGHT to be defined during compilation. (See the discussion of makeml's -rebuildlight option.)

Failure to reach a fixed point after 2 iterations usually indicates some serious problem within the compiler.

6.5 The allcross script

Finally, the allcross script is handy for building bin- and boot-files for all architectures. The script currently takes no arguments and compiles (and cross-compiles) for 6 supported combinations of *arch* and *os*: alpha32-unix, hppa-unix, ppc-unix, sparc-unix, x86-unix, and x86-win32.

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