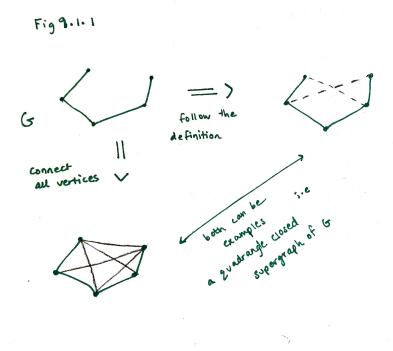
## CS 202, PSET 9

## 9.1.1

Let A be the set of all quadrangle closed supergraphs of G. We want to find the smallest graph in this set and show it's a unique closure.

The first thing we do is show that set A is non-empty:

A simple proof of this is that, given any graph, we can use this method: connect all the vertices with edges. Such a method ensures the graph becomes quadrangle closed. For example:



And also note that if the graph does not fulfill the criteria for having a quad. closed supergraph (i.e. does not have  $a_0a_1a_2a_3$ ), then it is not considered. Hence we have proved that the set A is not empty.

Now we find the smallest element of A. Let  $x_i \in A$ . Consider the intersection of all the graphs in A:

 $H = x_1 \cap x_2 \cap x_3 \dots \cap x_n$ 

H is a supergraph of G because it is an intersection of supergraphs.

H is quadrangle closed:

H contains the path  $a_0a_1a_2a_3$ , as do all the elements of A

So  $a_0a_3$  is also in H

And  $H \subseteq x \in A$ 

Now we show that H is unique:

Say another quad. closed supergraph of G exists, called H'

So  $H' \subseteq x \in A$ 

This implies  $H' \subseteq H$ 

But because H is quad. closed,  $H' \in A$ 

This implies:  $H' \subseteq H$ 

Therefore H = H' and is unique.

## 9.1.2

Want to show if G is bipartite, then the quadrangle closure of G is also bipartite. Proof by induction:

Let n represent the next edge added to graph G

Base case: the starting graph  $G^0 = G$  is bipartite

Inductive hypothesis: if  $G^n$  is bipartite  $\to G^{n+1}$  is bipartite

Let  $a_0a_1a_2a_3$  be a path in  $G^n$ 

Without loss of generality, we can start with saying  $a_0 \in S$ . To maintain a bipartite property,  $a_1 \in T$ ,  $a_2 \in S$  and  $a_3 \in T$ 

Now the edge  $a_0a_3$  maintains the bipartite property because the two vertices are separated into S and T

 $\rightarrow G^{n+1}$ 

## 9.2

There  $\exists$  cycle C such that  $C \subseteq G$ :

We know that there are exactly two distinct simple paths that don't share edges, between any 2 vertices

Let's say  $\exists u, v \in E_G$ , there must be another simple path connecting them, which means there is a cycle.

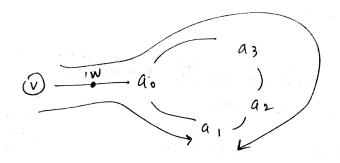
Now we show that  $C \supseteq G$ 

Let  $V_c$  be the vertices of  $C = \{a_0, a_1, a_2, ... a_n\}$  where  $a_0 a_n$  and  $a_i a_{i+1}$  are edges Proof:

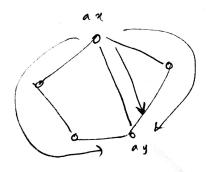
1) First we show that any vertex in G must be a vertex in C

Suppose  $\exists v : v \in V_G \land v \notin V_C$ 

Then if there exists a simple path from v to  $a_i$ , then there must exist another simple path, which contains the same edge. This leads to a contradiction.



2) Now we show every edge in G is an edge in C Suppose  $\exists$  edge  $a_xa_y\in E_G\land\not\in E_C$  But any simple path from  $a_x$  to  $a_y$  would also imply the existence of other simple paths, which leads to a contradiction



So now we have 3 paths and a contradiction

an, an+1 -- ay is a patu an, an-1 -- ay+1 ay is a patu

As a result of these two cases, G = C and is a cycle.

9.3