

INFS 519 – Fall 2015

Program Design and Data Structures

Lecture 4

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Today

- Last Class
 - Linked Lists, Queues, Stacks, Recursion
- Today
 - Merge & Quick Sort
 - JavaDocs
 - Testing
 - Trees

Linked List

Big-O

Operation Implementation	get set	add remove last	insert remove front	insert remove middle	search
Singly Linked List	N	1 or N*	1	N	N
Doubly Linked List	N	1	1	N	N

* Add is 1 for double-ended, remove last is still N

- Only use memory proportional to N
- “Locality of reference” poor compared to array based

Last Class: Node Data Structure

- What **fields** do we need to store a node?
- What **methods** do we need?

Last Class: List Data Structure

- What **fields** do we need?
- What **methods** do we need?

List Methods

- **Minimum Required** to Emulate an Array
 - Get / Set element
 - Append
 - Get size
- **Cool new things**
 - Remove
 - Insert
 - Append can't (usually) run out of space (just like a dynamic list)



Questions?

Last Class: Stacks & Queues

- Stacks
 - _____ in _____ out?
 - 4 common methods?
- Queues
 - _____ in _____ out?
 - 4 common methods?

Last Class: Stacks & Queues

- Stacks
 - last in first out (LIFO)
 - push(), pop(), peek(), isEmpty()
- Queues
 - first in first out (FIFO)
 - BUT THERE ARE OTHER TYPES!
 - enqueue(), dequeue(), peek(), isEmpty()

Stack

Big-O

Operation Implementation	push	pop	peek	isEmpty	size
Dynamic Array	1	1	1	1	1
Doubly Linked List	1	1	1	1	1

Easiest to implement Dynamic Array, constants are better

Queue

Big-O

Operation Implementation	enqueue	dequeue	peek	isEmpty	size
Dynamic Array	1	1	1	1	1
Doubly Linked List	1	1	1	1	1

For Queue, Dynamic Array is called a “Circular Queue”
Easiest to implement Singly (double-ended) Linked List

Last Class: Priority Queues

- _____ comes out first
- Implementations
 - _____
 - _____

Last Class: Priority Queues

- Highest priority comes out first
 - Not FIFO or LIFO
 - Priority could be max or min
- Implementations
 - Multiple queues
 - Single queue
- Naive Approaches
 - Unordered Array
 - Ordered Array



Questions?

Last Class: Recursion

- Idea: keep doing the _____ thing, _____ the problem
- Key components
 - _____ case (when to _____)
 - _____ case (when to _____)

Last Class: Recursion

- Idea: keep doing the **same** thing, **reducing** the problem
 - smaller **subset** of the problem
 - **one step** closer to the answer
- Key components
 - **recursive** case (when to **keep going**)
 - **base** case (when to **stop**)



Questions?

Last Class: Binary Search

- Requires?
- Method?
- Big-O?
 - worst case?
 - best case?

Last Class: Binary Search

- Requires: **sorted** list
- Method, see:
 - http://en.wikipedia.org/wiki/Binary_search_algorithm
- Big-O
 - worst case: **$O(\log n)$**
 - why?
 - best case: **$O(1)$**
 - why?

Sorted Array Binary Search Code

```
public static int search( Comparable findItem, Comparable[] items )
{
    int index = -1;
    int lo = 0;
    int hi = items.length-1;
    while( lo <= hi )
    {
        // Find half way position between begin and end
        int mid = lo + (hi - lo) / 2;
        Comparable midItem = items[mid];

        if(      findItem.compareTo(midItem) < 0 ) hi = mid-1; // Move left
        else if( findItem.compareTo(midItem) > 0 ) lo = mid+1; // Move right
        else
        {
            // Must be equal, narrowed to one index, exit loop
            index = mid;
            break;
        }
    }
    return index;
}
```

Divide and Conquer

- **Divide** the problem
 - in half or some smaller portion
- Keep doing that (based on recursion)
 - until the problem is small enough to solve (**conquer**)
- If needed, use the smaller solved problems to solve the big one (**conquer**)
- Traditionally, “divide-and-conquer” algorithms have two or more recursive calls.

Recursion Alternatives

- Recursion can make a seemingly difficult problem easy to solve, is often the most elegant approach, and generally easier to verify
- However, it has non-trivial overhead to add activation records to the stack for each method call
- Factorials and triangle problems use recursion only for teaching, in practice use iterative approach
- Common recursion alternatives (Look at Triangle.java)
 - Iterative, not always possible
 - Explicit stack (call stack is implicit), used to be important, now compilers quite efficient and approach is less often useful



Questions?

Sorting

- Often desirable to have output sorted
- More importantly, sorting input can make algorithms much more efficient (e.g. binary search arrays)
- Because it is a building block for other algorithms, must understand performance
- Three sub-quadratic algorithms (many more)
 - Shell Sort (will not cover, easy to implement)
 - Merge Sort, invented by Von Neumann
 - Quick Sort, invented by Hoare

Stable Sorting

- Relative ordering for duplicate values is maintained, important for objects

<u>Original</u>		
<u>Sorted by Name</u>		
Arda	8	PA1
Arda	7	PA2
Arda	10	PA3
James	8	PA1
James	3	PA2
James	8	PA3
Mary	6	PA1
Mary	6	PA2
Mary	6	PA3

<u>Not Stable</u>		
<u>Sorted by Grade</u>		
James	3	PA2
Mary	6	PA3
Mary	6	PA1
Mary	6	PA2
Arda	7	PA2
James	8	PA3
Arda	8	PA1
James	8	PA1
Arda	10	PA3

<u>Stable</u>		
<u>Sorted by Grade</u>		
James	3	PA2
Mary	6	PA1
Mary	6	PA2
Mary	6	PA3
Arda	7	PA2
Arda	8	PA1
James	8	PA1
James	8	PA3
Arda	10	PA3

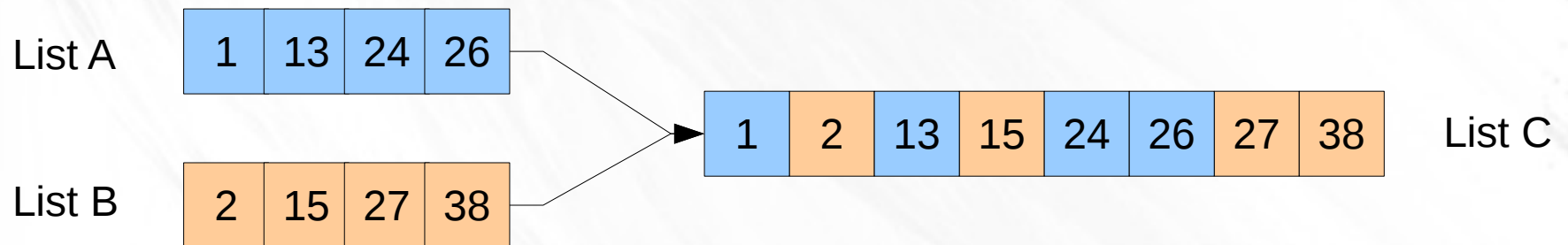
Merge Sort

- If the problem is too big
... find a smaller problem.
- Demo
- Resource with **animations** if you forget this:
 - http://en.wikipedia.org/wiki/Merge_sort
 - <http://www.sorting-algorithms.com/merge-sort>

Merge Operation

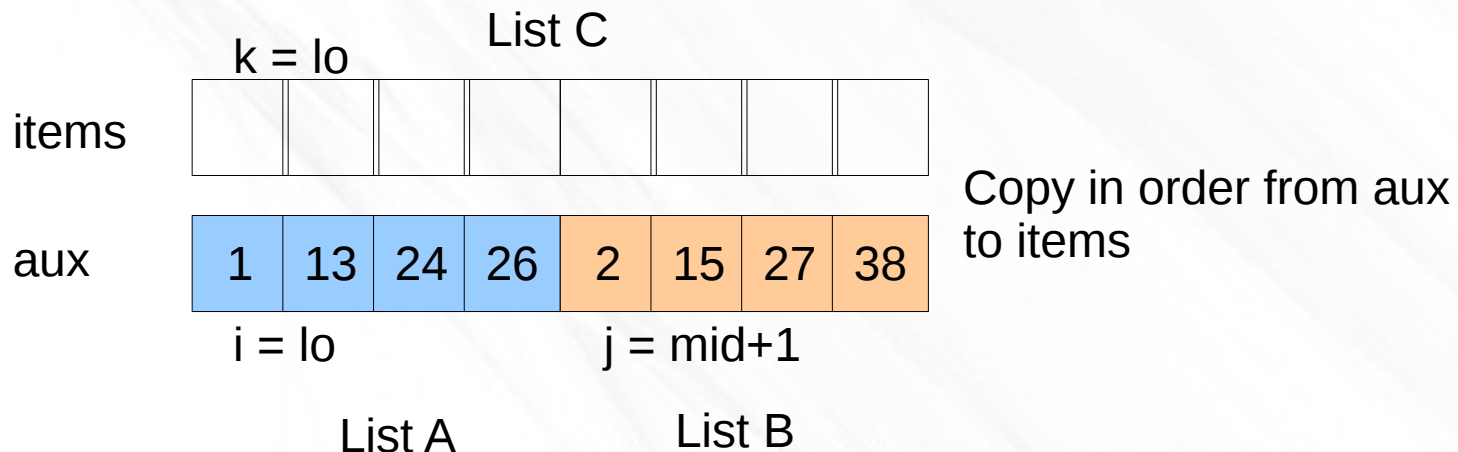
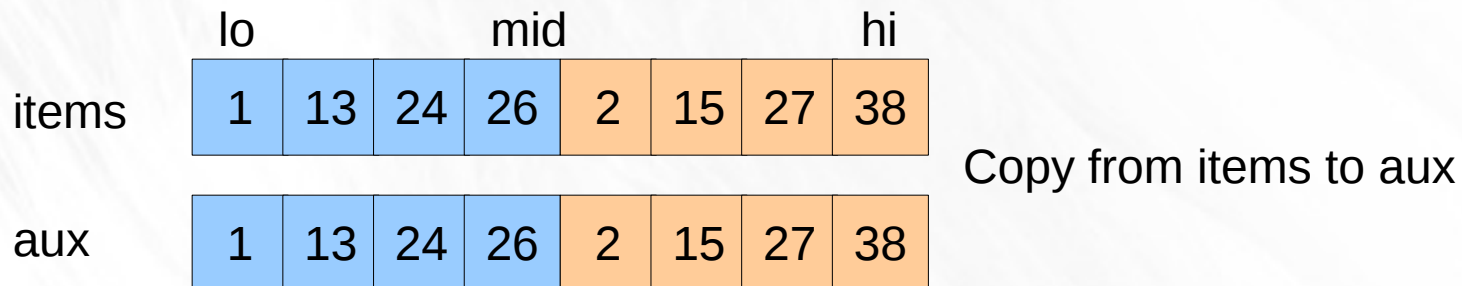
Weiss 8.5.1

- Given two sorted lists, merge into one sorted list in linear time $O(n)$



Merge Operation Memory

- Memory to copy from and memory to copy to.



Merge Operation Outline

“Top-down Mergesort”

```
public static void merge( Comparable[] items, Comparable[] aux,
                          int lo, int mid, int hi )
{
    // Copy items from lo to hi to the aux array

    // Now copy back into item, file out in order
    // Consider [lo,mid] list a, [mid+1,hi] list b
    int i = lo;    // start index for list a
    int j = mid+1; // start index for list b
    for( int k = lo; k <= hi; k++ ) // k index into items
    {
        // Handle cases where a list is empty
        //      If a is empty,    take from b
        // Else If b is empty,    take from a
        // Else If b is smaller, take from b
        // Else a is smaller or equal, take from a

        // Note order, if equal, get from a,
        // comes first to keep stable
    }
}
```

Merge Sort Outline

```
// Note: not real code...  
list mergeSort(list)  
{  
    if(list is empty or contains 1 element)  
        return list  
    list1 = mergeSort(first half of list)  
    list2 = mergeSort(second half of list)  
    return merge(list1, list2)  
}
```

Merge Sort Tree

0	1	2	3	4	5	6	7
26	1	13	24	38	15	2	27

26	1	13	24	38	15	2	27
----	---	----	----	----	----	---	----

26	1	13	24	38	15	2	27
----	---	----	----	----	----	---	----

26	1	13	24	38	15	2	27
----	---	----	----	----	----	---	----

Merge Sort Trace

lo=0 hi=7

lo=0 hi=3

lo=0 hi=1

lo=0 hi=0

lo=1 hi=1

Merge: [0, 1]

lo=2 hi=3

lo=2 hi=2

lo=3 hi=3

Merge: [2, 3]

Merge: [0, 3]

lo=4 hi=7

lo=4 hi=5

lo=4 hi=4

lo=5 hi=5 Merge: [4, 5]

lo=6 hi=7

lo=6 hi=6

lo=7 hi=7

Merge: [6, 7]

Merge: [4, 7]

Merge: [0, 7]

Merge Sort Heuristics 1/2

Sedgewick/Wayne 2.2

- Cutoff value below which simple sorting is used (e.g. InsertionSort)
 - Recursive call overhead to sort just a few items is relatively significant
 - Common technique for other recursive sorting algorithms
 - For small n (e.g. 8), simple sorts are faster
- Avoid copy to auxiliary array
 - Tricky, essentially reverse roles of the items and auxiliary arrays on recursive calls

Merge Sort Heuristics 2/2

Sedgewick/Wayne 2.2

- Check if two lists are already in sorted order to avoid the merge $O(N)$
 - Left list max value at `items[mid]`
 - Right list min value at `items[mid+1]`
- Use iterative approach (“bottom-up mergesort”) which avoids recursion altogether

Properties of Merge Sort

- Not in-place
 - requires $O(n)$ additional memory space
- Stable
 - relative order of equal elements preserved

Operation Implementation	worst	average	best	in place	stable	remarks
Selection Sort	N^2	N^2	N^2	yes	no	
Insertion Sort	N^2	N^2	N	yes	yes	
Merge Sort	$N \lg N$	$N \lg N$	$N \lg N$	no	yes	



Questions?

The Quick Sort Algorithm

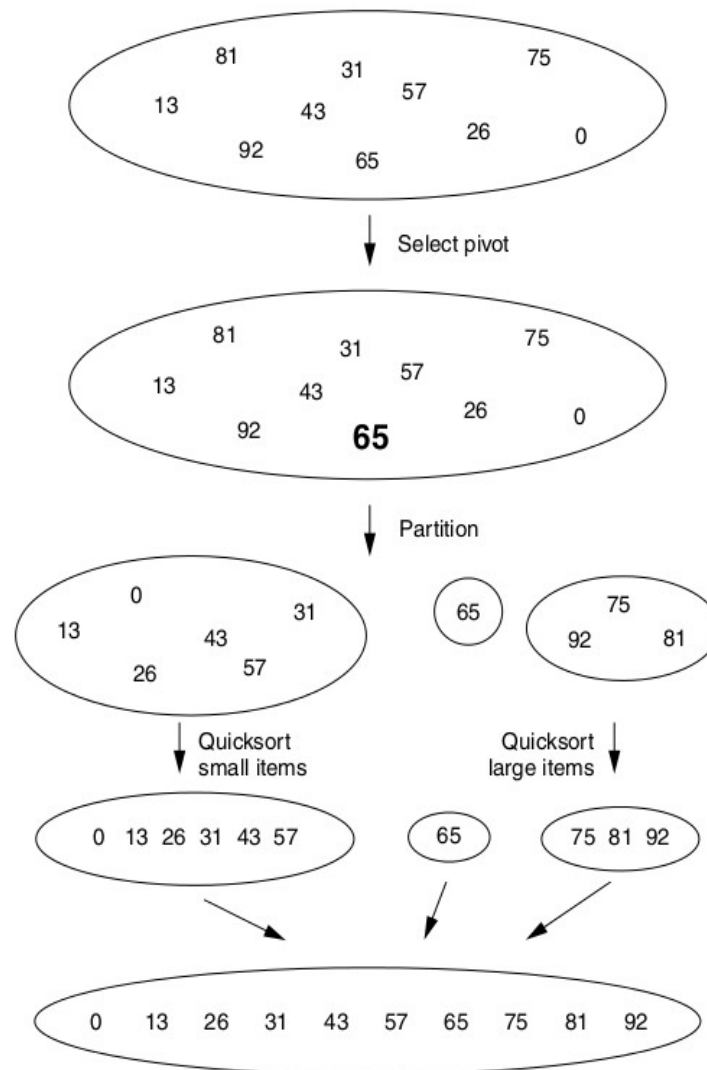
Weiss 8.6.1

The basic algorithm Quicksort(S) consists of the following four steps.

1. If the number of elements in S is 0 or 1, then return.
2. Pick any element v in S . It is called the pivot.
3. Partition $S - \{v\}$ (the remaining elements in S) into two disjoint groups:
 $L = \{ x \in S - \{ v \} \mid x \leq v \}$ and
 $R = \{ x \in S - \{ v \} \mid x \geq v \}$.
4. Return the result of Quicksort(L) followed by v followed by Quicksort(R).

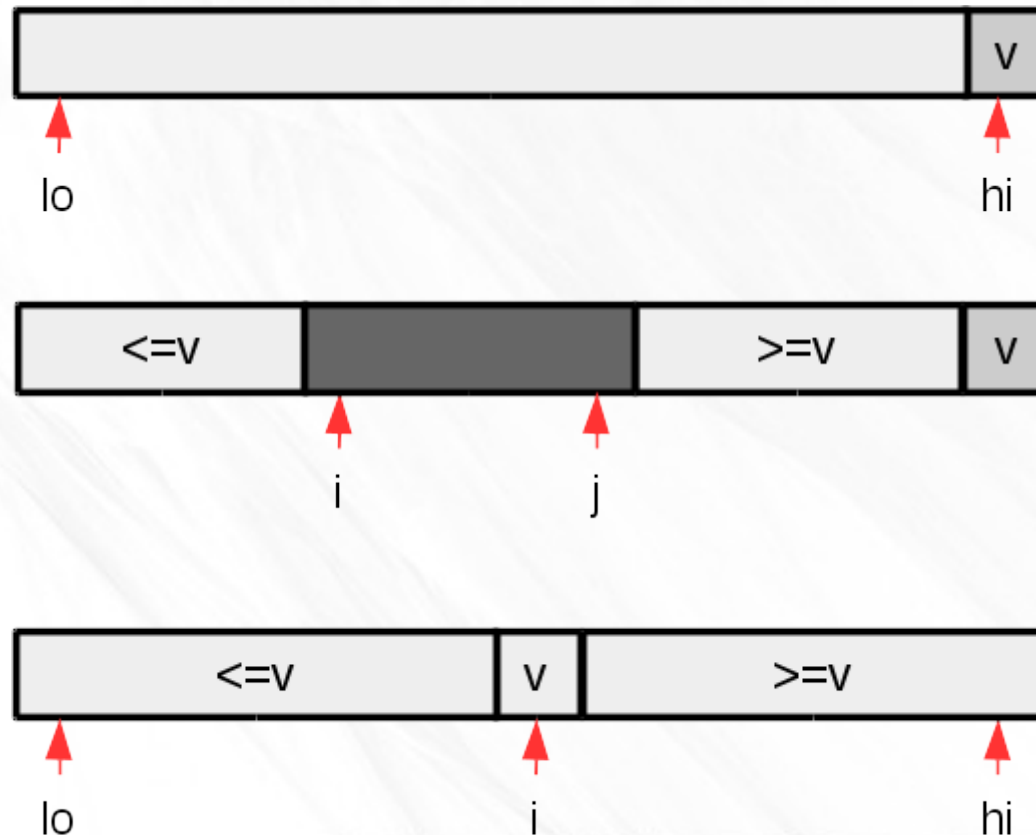
Steps of Quick Sort

Weiss Figure 8.11



The Quick Sort Partition 1/2

Invariant: For a given pivot, no larger values to the left of the pivot and no smaller values to the right of the pivot.



The Quick Sort Partition 2/2

Weiss 8.6.4

8	1	4	9	0	3	6	2	7	5
---	---	---	---	---	---	---	---	---	---

lo

hi

Step0: Pick pivot (6)

8	1	4	9	0	3	5	2	7	6
---	---	---	---	---	---	---	---	---	---

Step1: Move out of way

i = lo

j = hi-1

8	1	4	9	0	3	5	2	7	6
---	---	---	---	---	---	---	---	---	---

i

j

Step2: Small elements to left of array and large elements to right of array

2	1	4	9	0	3	5	8	7	6
---	---	---	---	---	---	---	---	---	---

Swap 8 and 2

i

j

2	1	4	9	0	3	5	8	7	6
---	---	---	---	---	---	---	---	---	---

i

j

2	1	4	5	0	3	9	8	7	6
---	---	---	---	---	---	---	---	---	---

Swap 9 and 5

i

j

2	1	4	5	0	3	9	8	7	6
---	---	---	---	---	---	---	---	---	---

j

i

2	1	4	5	0	3	6	8	7	9
---	---	---	---	---	---	---	---	---	---

Step3: Swap 6 and 9

j

i

Quick Sort

- There is a good chance that
 - ... randomness is your friend
- Quick sort does work (partition) then recursion
- Merge sort does recursion then work (merge)
- Demo
- Resource with [animations](#) if you forget this:
 - <http://en.wikipedia.org/wiki/Quicksort>
 - <http://www.sorting-algorithms.com/quick-sort>

Quick Sort

```
// Note: not real code...
int quickSort(list)
{
    if(list is empty or contains 1 element)
        return list
    int pivot = some item in the list
    for(each item in the list)
    {
        if(item smaller than pivot)
            put in first "section" of list
        if(item larger than pivot)
            put in last "section" of list
    }
    put pivot in between two sections
    quickSort(first "section" of list)
    quickSort(last "section" of list)
    return list
}
```


Quick Sort Variants

Weiss 8.6 and 9.4

- Pick middle
- Estimate median (3 samples)
- Randomly shuffle input before sorting
 - Statistically guarantees average performance

```
// Similar to Knuth / Fisher-Yates shuffling algorithm
// If generator independent, uniform, so is output
public static void shuffle(Object[] items)
{
    for(int j = 1; j < items.length; j++)
    {
        swap( items, j, random.nextInt(j) ); // [0,j)
    }
} // returns some permutation of items, adds O( N )
```

Space Complexity: Quicksort

- Space complexity
 - $O(\log n)$ or $O(n)$
- Why?
- Just add more memory?
 - Is this stack or heap memory?
- Other variants to reduce memory to $\log(N)$
 - Tail recursion / Iteration
- Duplicate values, can degrade to $O(N^2)$!!!
 - Quick 3-Way solves, $[< v] [= v] [> v]$

Properties of Quick Sort

- Space complexity?
- In-place/Not in-place
- Unstable/Stable

Operation Implementation	worst	average	best	in place $O(1)$	stable	remarks
Selection Sort	N^2	N^2	N^2	yes	no	
Insertion Sort	N^2	N^2	N	yes	yes	
Merge Sort	$N \lg N$	$N \lg N$	$N \lg N$	no	yes	
Quick Sort	N^2	$N \lg N$	$N \lg N$	yes*	no	fast practice
???	$N \lg N$	$N \lg N$	N	yes	yes	Unknown

* Depending on variant, will assume $O(\lg(N)) \sim O(1)$

Sorting Summary

- Merge sort (Java objects, C++/Python stable)
 - Stable and efficient but requires extra memory
- Quick sort (Java primitives, Python, Matlab)
 - Little extra memory and efficient, not stable
- Which one to use?
 - If memory is limited (in-place), use quicksort
 - If stability is required, use mergesort
- Example: Java system programmers assume if using primitives then memory is limited and if using objects then stability is more desirable



Questions?

JavaDocs

- Use **javadoc** style for this class
 - Always use **@param** and **@return** for all methods with parameters and return types
 - Use these on **your assignments!**
- <http://www.oracle.com/technetwork/java/javase/documentation/index-137868.html>

Example Class Comment

```
/**  
 * Model of a banana  
 */  
public class Banana  
{  
  
}
```

Javadoc Annotations

```
/**  
 * Model of a banana  
 *  
 * @author Your Name Here  
 * @version 0.1-alpha  
 */  
public class Banana  
{  
  
}
```


HTML!

```
/**
 * <p>Model of a banana.</p>
 *
 * <p>Another paragraph about the awesome
 * banana class</p>
 *
 * @author Your Name Here
 * @version 0.1-alpha
 */
public class Banana
{

}
```

Example Method Comment

```
/**  
 * myMethod provides users with some  
 * cool functionality.  
 */  
public void myMethod()  
{  
  
}
```

Return Annotation

```
/**  
 * Counts how many bananas we have.  
 *  
 * @return number of bananas in the system  
 */  
public int countBananas()  
{  
  
}
```

Param Annotation

```
/**
 * Divides two numbers.
 *
 * @param a the numerator
 * @param b the denominator
 * @return the quotient of a and b
 */
public double div(int a, int b)
{
}
}
```

Param Annotation

```
/**
 * Divides two numbers.
 *
 * @param a the denominator
 * @param b the numerator
 * @return the quotient of a and b
 */
public double div(int a, int b)
{
}
}
```

Throws/Exception Annotation

```
/**
 * My method throws that exceptions...
 *
 * @throws UnsupportedOperationException
 *         Always throws this
 */
public void myMethod()
{

}
```

Member Comment

```
/**
 * Model of a banana
 */
public class Banana
{
    /**
     * Keeps count of number of bananas
     * in existence
     */
    public static int count = 0;
}
```

The background of the slide features a grayscale photograph of a rugged mountain range. A large, semi-transparent white rectangle is centered over the image, creating a clean space for the text.

Questions?

Getting the Bugs Out

- Most **development time** is spent **testing** and **debugging** code
- How do you know when your **code works**?
- How do you **catch bugs**?
- How do you **fix bugs**?
 - If you fix a bug, how do you know it **stays fixed**?

Software Engineering: Testing

- **Quality Assurance** (QA) process – manner in which a company ensures its product is acceptable
- **Code Reviews** – meeting in which several people examine a design document or section of code
- For medium and large systems, testing must be a carefully **managed process**
 - Many organizations have a separate QA department to lead testing efforts

Unit and Integration Testing

Liskov and Guttag 10.7

- Unit testing checks each individual module for correctness in isolation. Typically software engineers are responsible for unit testing.
- Integration testing checks entire program correctness when all the modules are put together. Usually harder than unit testing and may be carried out by software or system engineers.

Automated Tests

- **Test Case** – set of inputs and actions coupled with expected results
- **Test Suite** – test cases and/or testing framework which is stored and reused as needed
- **Defect Testing** – execution of test cases to find errors
- **Regression Testing** – running previous test suites to ensure new errors have not been introduced

How is Testing Done

- Testing can mean many different things:
 - running a program on various inputs
 - human or computer assessment of quality
 - evaluations before writing code
- The **earlier** we find an problem, the **easier** and **cheaper** it is to fix
 - Some software engineering approaches advocate writing test cases **BEFORE** you even write any code (Termed “**Test Driven Development (TDD)**”)
 - Why?

Unit Testing 1/2

- Main test method (some languages do not allow more than one main). Simple, black box view.
- Java assert. Typically used to maintain internal invariants, not black box view.
- Testing framework. More complicated but more powerful black box view. Most common in industry.

Unit Testing 2/2

- Commonly used frameworks to test code
 - JUnit (<http://junit.org>) for java
 - NUnit for c#
- Idea:
 - Write code & test method
 - **@Test** (Java Annotation, not Javadoc)
 - Test methods do something & check if answer is correct, using assert-type methods
 - assertTrue(..), assertFalse(..), assertEquals(x,y)

Example: PA3

```
/**
 * Check HeapSort, if takes longer than timeout
 * test case will fail
 */
@Test(timeout=2000)
public void heapSort()
{
    .. .. .
    int[] sortedList = PA2.heapsort(list);
    list = insertionSort(list);
    assertEquals(sortedList, list);
}
```


Example: PA3

```
/**
 * Expected exception, test case succeeds if and
 * only if the exception is thrown
 */
@Test(expected=RuntimeException.class,
        timeout=2000)
public void myMethod()
{
}
}
```

Assert Things!

```
import org.junit.*; //junit
```

```
import static org.junit.Assert.*; //assert
```

- junit provides convenience functions for testing
 - Assert.[assertEquals\(\)](#)
 - 2 params: expected value & actual value
 - lot of versions of this function
 - Assert.[assertTrue\(\)](#) & Assert.[assertFalse\(\)](#)
 - Assert.[assertNull\(\)](#) & Assert.[assertNotNull\(\)](#)
- <http://junit.sourceforge.net/javadoc/org/junit/Assert.html>

JUnit Assert vs Java Assert

- Java has a keyword “assert”
 - assert <boolean>, e.g. assert isSorted(array);
 - If false, AssertionError is thrown
 - Have to enable “java -ea MergeSort” (default disabled)
- JUnit also has similar behavior
 - <http://junit.org/apidocs/org/junit/Assert.html>
 - assertTrue(...)
 - assertEquals(...)
 - Throws same AssertionError
- JUnit is considered more flexible and black box view of code.
- Either approach no running cost for production code
 - Java assert in compiled code, JUnit is not

Java Annotations

<http://docs.oracle.com/javase/tutorial/java/annotations/basics.html>

- Information for the **compiler**
 - Like comments but the compiler may not completely ignore them
- **Metadata** that summarizes the intent of code

Java Annotation Examples

- Java Annotations -
http://en.wikipedia.org/wiki/Java_annotation
- Examples
 - **@Test** This code tests other code (compiler may just ignore)
 - **@Deprecated** This code is old, unsupported, may disappear
 - **@Override** Warn if not overriding parent method

What Testing Is...

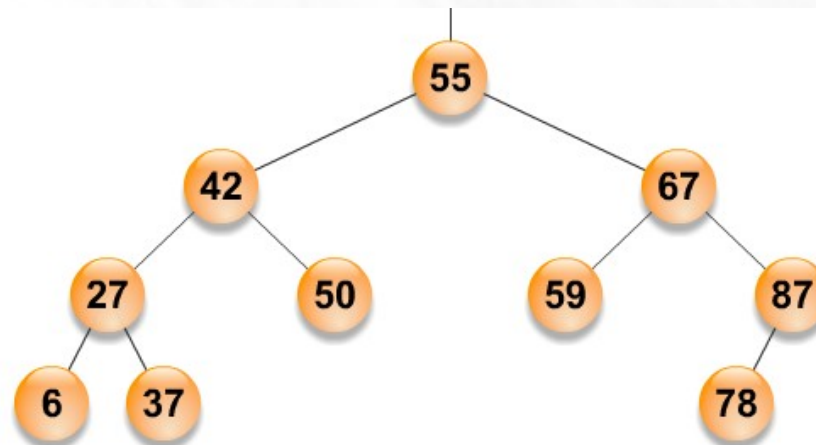
- Unit tests are just **more code**
 - Intended to test other code
 - Can have **bugs** (what's testing the tests?)
 - Can miss critical stuff
- Very common **job in industry**
- To really **learn** testing, **write** some tests!



Questions?

Trees

- Data structure which looks like an upside down **tree** (or the root system of a tree)
 - Nodes have **parents** and **children**
 - **No loops**



Trees

- Collection of **nodes** and **edges**
 - These nodes are **different!**
 - Any shape, but can't have a **loop**
 - **Acyclic** means “no cycles” (i.e. no loops)
- Nodes have:
 - **data**
 - (possibly) a “**key**” to sort/search by
 - (possibly) pointer to **children**
 - (possibly) pointer to **parent**

Types of Nodes

- By Relationship
 - Parent/Child Nodes
 - Ancestor/Descendent Nodes
- By Tree Location
 - Inner/Branch/Internal Nodes
 - Outer/Leaf/External/Terminal Nodes
 - Root Node (one 1!)
- Null Links

Tree Definitions 1/2

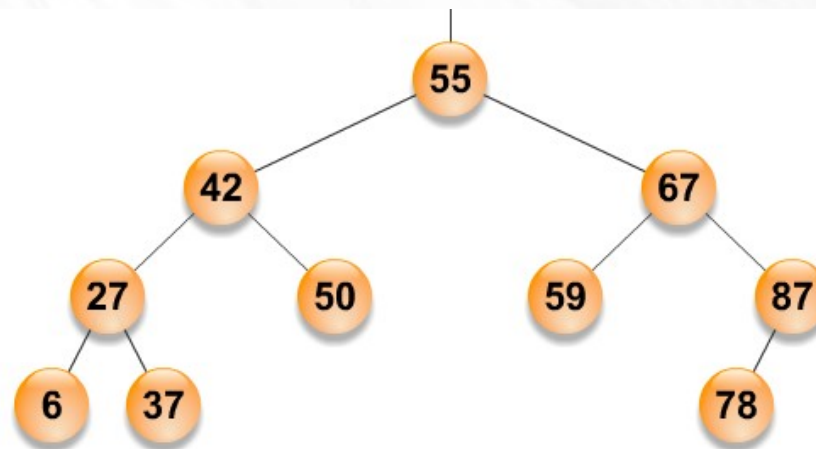
- The **descendants** of a node are all the nodes below it and includes itself
- The **ancestors** of a node are the nodes on the path from the node to the **root**
- The **leaf** nodes have no children, the **root** node has no parent
- Nodes are **siblings** if they have the same parent

Tree Definitions 2/2

- The **depth** of a node is the length (number of edges) of the path from the node to the root
- The **tree height** is the maximum **depth** of any node in the tree
- An **inner** node has at least one child (i.e. not a leaf)

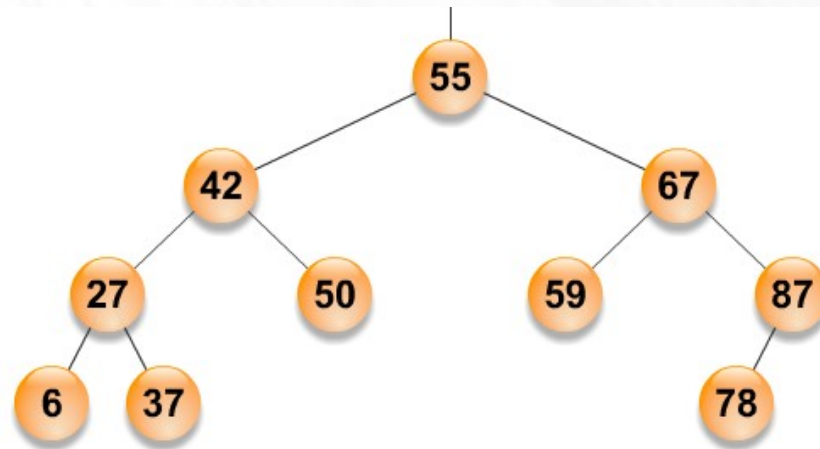
Examples

- What is the **parent** of 27?
- What are the **children** of 67?
- What are the **ancestors** of 59?
- What are the **descendants** of 55?



Examples

- What is the **root**?
- Which nodes are **leaf nodes**?
- Which nodes are **inner nodes**?
- Where are the **null links**?



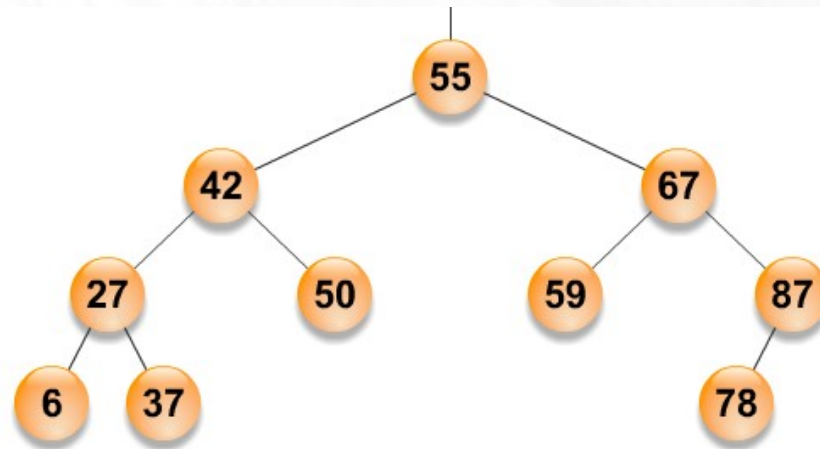
Types of Trees: Fullness

Note: Ambiguous Definitions

- **full**/perfect tree
 - every node other than the leaves has the max number of children
- perfect/**complete** tree
 - all leaves have the same depth
 - every node other than the leaves has the max number of children
- almost complete or **nearly complete** tree
 - last level is not completely filled

Examples

- Is this tree **full**?
- Is this tree **complete**?
- Is this tree **nearly complete**?
- What is the **tree height**?

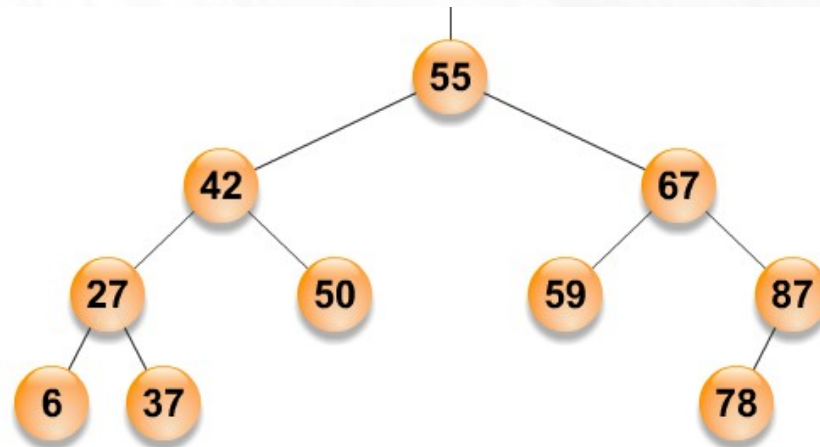


Types of Trees: Arrangement

- **balanced** tree
 - **height** of the left and right sub trees of every node differ by 1 or less
- **degenerate** tree
 - each parent node has only one associated child node
 - equiv. to **linked list**, maximum height?

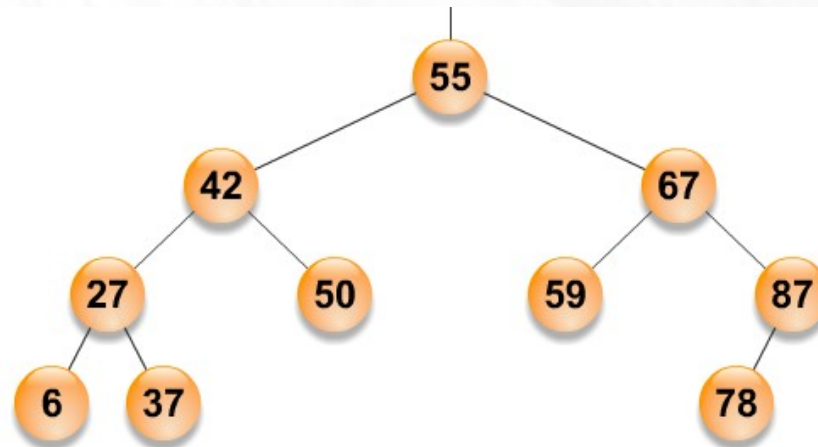
Examples

- Is this tree **balanced**?
- Is this tree **degenerate**?



Examples

- Is this tree **balanced**?
- Is this tree **degenerate**?
 - what nodes could we remove to make it degenerate?





Questions?

Common Tree Operations

- **Iterating (a.k.a. Enumerating)** = mention things one by one
 - all the items
 - a section of a tree
- **Searching** for an item
- **Adding/Deleting** items
- **Pruning/Grafting**
- **Balancing**

Tree Data Structures

- Arrays
 - Need to know **where** each item is
 - How? Need to **limit** number of **children**
 - Most common for **balanced binary trees**
 - Fast **memory** access (compared to linked list)
- Linked Data Structures
 - **easy** to add, remove, and swap around parts of the tree



Questions?

Assignments: PA2 / PA3 / PA4

- PA2: Grades posted
- PA3: Due Tomorrow
- PA4
 - Practice recursive techniques
 - Implement a sorting algorithm

Assignments: PA2

- Generics
- Array List (Java Collections)
- Compiling
- Exceptions (ranges negative) throw vs try/catch
- Edge case – reducing size to zero
- Insert
- Null checking
- Submission format
- Generics



Let's Look at the Files

The background of the slide is a grayscale photograph of a vast, snow-covered mountain range. The peaks are rugged and partially obscured by soft, white snow. The sky is a pale, uniform gray. A large, semi-transparent white rectangle is centered over the image, serving as a backdrop for the text.

Free Question Time!