1	Primal feasible	Dual feasible			
	Maximize $z = 2x_1 + x_2$,	Minimize $v = 4y_1 + 2y_2$,			
	subject to:	subject to:			
	$x_1 + x_2 \le 4,$ $x_1 - x_2 \le 2,$ $x_1 \ge 0, x_2 \ge 0.$	$y_1 + y_2 \ge 2,$ $y_1 - y_2 \ge 1,$ $y_1 \ge 0, y_2 \ge 0.$			
			2	Primal feasible and unbounded	Dual infeasible
				Maximize $z = 2x_1 + x_2$,	Minimize $v = 4y_1 + 2y_2$,
subject to:	subject to:				
$x_1 - x_2 \le 4,$	$y_1 + y_2 \ge 2$,				
$x_1 - x_2 \le 2,$	$-y_1 - y_2 \ge 1$,				
$x_1 \geq 0, x_2 \geq 0.$	$y_1 \ge 0, y_2 \ge 0.$				
3	Primal infeasible	Dual feasible and unbounded			
	$Maximize z = 2x_1 + x_2,$	$Minimize v = -4y_1 + 2y_2,$			
	subject to:	subject to:			
	$-x_1-x_2\leq -4,$	$-y_1 + y_2 \ge 2,$			
	$x_1 + x_2 \le 2,$	$-y_1 + y_2 \ge 1,$			
	$x_1 \ge 0, x_2 \ge 0.$	$y_1 \ge 0, y_2 \ge 0.$			
4	Primal infeasible	Dual infeasible			
	$Maximize z = 2x_1 + x_2,$	$Minimize v = -4y_1 + 2y_2,$			
	subject to:	subject to:			
	$-x_1+x_2\leq -4,$	$-y_1 + y_2 \ge 2$,			
	$x_1 - x_2 \le 2$	$v_1 - v_2 \ge 1$.			

We can show that the complementary-slackness conditions follow directly from the strong duality property just presented. Recall that, in demonstrating the weak duality property, we used the fact that:

$$\sum_{j=1}^{n} c_{j} \hat{x}_{j} \leq \sum_{i=1}^{m} \sum_{j=1}^{n} a_{ij} \hat{x}_{j} \hat{y}_{i} \leq \sum_{i=1}^{m} b_{i} \hat{y}_{i}$$
(18)

for any \hat{x}_j , j = 1, 2, ..., n, and \hat{y}_i , i = 1, 2, ..., m, feasible to the primal and dual problems, respectively. Now, since these solutions are not only feasible but optimal to these problems, equality must hold throughout. Hence, considering the righthand relationship in (18), we have:

$$\sum_{i=1}^{m} \sum_{j=1}^{n} a_{ij} \hat{x}_{j} \hat{y}_{i} = \sum_{i=1}^{m} b_{i} \hat{y}_{i},$$

which implies:

$$\sum_{i=1}^m \left[\sum_{j=1}^n a_{ij} \hat{x}_j - b_i \right] \hat{y}_i = 0.$$