Matrix Multiplication with compressed gadget

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0.1 Matrix multiplications

There are several methdos for performing matrix-matrix multiplication on quantum computers. One method involves duplicating the ancilla qubits relying on block-encoding [1, 2]. The number of ancilla qubits and two-qubit swap gates required for matrix multiplication increases linearly with the number of matrix multiplications (see Appendix **??**). For example, computing E^p, the power p of matrix E, using the duplicate ancilla qubits technique requires additional $O(p)$ ancilla qubits and $O(p)$ swap gates. Each swap gate is decomposed into 3 $CNOT$ gates. Consequently, the duplicate ancilla qubits technique is inefficient for multiplications of non-unitary matrices. Another method for multiplication of the non-unitary matrices is called compression gadget [3]. The number of extra ancilla qubits needed for computing matrix multiplication increases logarithmically with the number of multiplications. For instance, E^p can be computed with $O(log_2(p))$ ancilla qubits using the compression gadget. First, we explain the compression gadget in a lemma, and then we describe the quantum circuit for the compression gadget.

Lemma 1(compression gadget for the matrix multiplications) [4]. Suppose we are given unitaries $U_1, U_2, \ldots, U_L$, each of which is a $(\alpha_l', m_l', 0)$-block encoding of A_l. Then the $(\alpha_{\text{comp}}, m_{\text{comp}}, 0)$-block encoding of $A_L \cdots A_2 A_1$ is constructed, where

$$\alpha_{\text{comp}} = \alpha_1' \alpha_2' \cdots \alpha_L', \quad m_{\text{comp}} = \max_l(m_l') + \lceil \log_2(L) \rceil + 1 \tag{1}$$

using one application of each U_l. In Eq. 1, m_l' is the number of qubits needed for block encoding of A_l. Fig. 1 shows the quantum circuit for the compression gadget for multiplication of $A_L ... A_2 A_1$. The counter register in Fig. 1 contains $log_2(L) + 1$ qubits. It is used to keep track of how many block encoding of A_l have been applied successfully. The ADD operation in Fig. 1 is a unitary operation. It maps $ADD\,|0\rangle =$

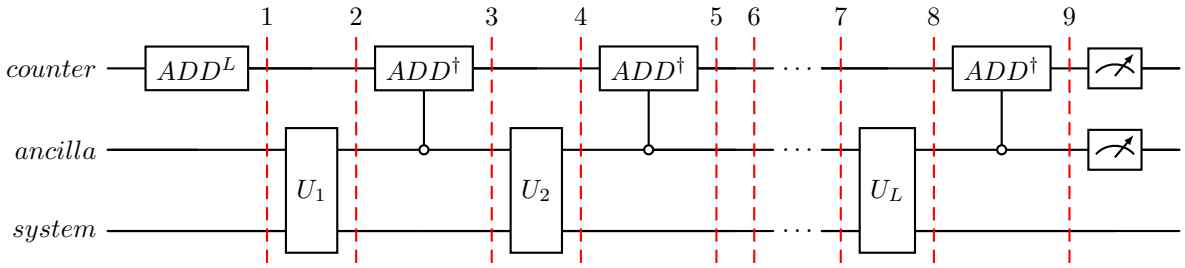

Fig. 1 Quantum Circuit for computation of matrix multiplication with compression gadget. ADD is applied L times on the counter register to maps $ADD^L|0\rangle = |L\rangle$ (step 1). At the next step, unitary block encoding U_1 is applied to the ancilla qubits and system qubits. This operation generates a superposition between $|0\rangle^{\otimes a}$ and $|1\rangle^{\otimes a}$ (step 2). The next step is to utilize the information in the ancilla register to change the information in the counter register. The controlled $ADD^\dagger$ operation, conditioned on a-ancilla qubits all in the state $|0\rangle^{\otimes a}$ is applied to change the information in the counter registers (step 3). Steps 2 and 3 are repeated until the quantum state of the counter register is brought back to $|0\rangle^{\otimes c}$. Finally, the counter register and ancilla registers are measured. The output after postselection is the matrix $A_L...A_2A_1$.

$|1\rangle$, $ADD|1\rangle = |2\rangle$,..., $ADD|L-1\rangle = |L\rangle$ and $ADD|L\rangle = |0\rangle$ and the $ADD^\dagger$ is the hermitian conjugate of ADD. It maps $ADD^\dagger|0\rangle = |L\rangle$, $ADD^\dagger|L\rangle = |L-1\rangle$, ..., $ADD^\dagger|2\rangle = |1\rangle$, and $ADD^\dagger|1\rangle = |0\rangle$.

Let's consider that the initial state is $|0\rangle^{\otimes(c=log(L)+1)}|0\rangle^{\otimes(a=\frac{m'_l+1}{2})}|0\rangle^{\otimes(s=\frac{m'_l-1}{2})}$. First ADD operator applies L times on the counter register c to map $|0\rangle^{\otimes c}$ to $|L\rangle^{\otimes c}$. The quantum state after step 1 is $|L\rangle^{\otimes c}|0\rangle^{\otimes a}|0\rangle^{\otimes s}$. Then the first unitary block encoding U_1 is applied on $|L\rangle^{\otimes c}|0\rangle^{\otimes a}|0\rangle^{\otimes s}$. This operation generate a superposition between $|0\rangle^{\otimes a}$ and $|1\rangle^{\otimes a}$. The quantum state is $|L\rangle^{\otimes c}(|0\rangle^{\otimes a}A_1|0\rangle^{\otimes s} + |1\rangle^{\otimes a}|*\rangle^{\otimes s})$. $|*\rangle$ is called a junk state. The next step is to utilize the ancilla register to change the information in the counter register. This is achieved by applying a controlled $ADD^\dagger$ operation, conditioned on a-control qubits all in the state $|0\rangle^{\otimes a}$. The quantum state after this operation (step 3) is $|L-1\rangle^{\otimes c}|0\rangle^{\otimes a}A_1|0\rangle^{\otimes s} + |L\rangle^{\otimes c}|1\rangle^{\otimes a}|*\rangle^{\otimes s}$. After step 3, another unitary block encoding U_2 is applied. The quantum state after applying U_2 operation is $|L-1\rangle^{\otimes c}(|0\rangle^{\otimes a}A_2A_1|0\rangle^{\otimes s} + |1\rangle^{\otimes a}|*\rangle^{\otimes s}) + |L\rangle^{\otimes c}(|0\rangle^{\otimes a}|*\rangle^{\otimes s} + |1\rangle^{\otimes a}|*\rangle^{\otimes s})$ (step 4). Then another controlled $ADD^\dagger$ by considering the control qubit $|0\rangle^{\otimes a}$ is applied and a new counter register $|L-2\rangle^{\otimes c}$ is generated. The new quantum state is $|L-2\rangle^{\otimes c}|0\rangle^{\otimes a}A_2A_1|0\rangle^{\otimes s} + |L-1\rangle^{\otimes c}|1\rangle^{\otimes a}|*\rangle^{\otimes s} + |L-1\rangle^{\otimes c}|0\rangle^{\otimes a}|*\rangle^{\otimes s} + |L\rangle|1\rangle^{\otimes a}|*\rangle^{\otimes s}$ (step 5). These steps must be repeated until all the unitary block encoding operations are applied and the counter register is brought back to its initial state $|0\rangle^{\otimes c}$. The quantum state before measuring the counter qubit and the ancilla qubits is $|0\rangle^{\otimes c}|0\rangle^{\otimes a}A_L...A_2A_1|0\rangle^{\otimes s} + \sum_{j>0}|j\rangle|*\rangle|*\rangle$ (step 9). Measuring the counter qubits and the ancilla qubits on the basis $|0\rangle^{\otimes c}|0\rangle^{\otimes a}$ in combination with post-selection guarantee that the multiplication $A_L...A_2A_1$ is applied successfully (step 10). The code for matrix multiplication with the compressed gadget technique is available on GitHub.

References

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