

Operating Systems (CS3000)

Lecture – 12 (exec(), wait() System Call)



INDIAN INSTITUTE OF INFORMATION TECHNOLOGY,
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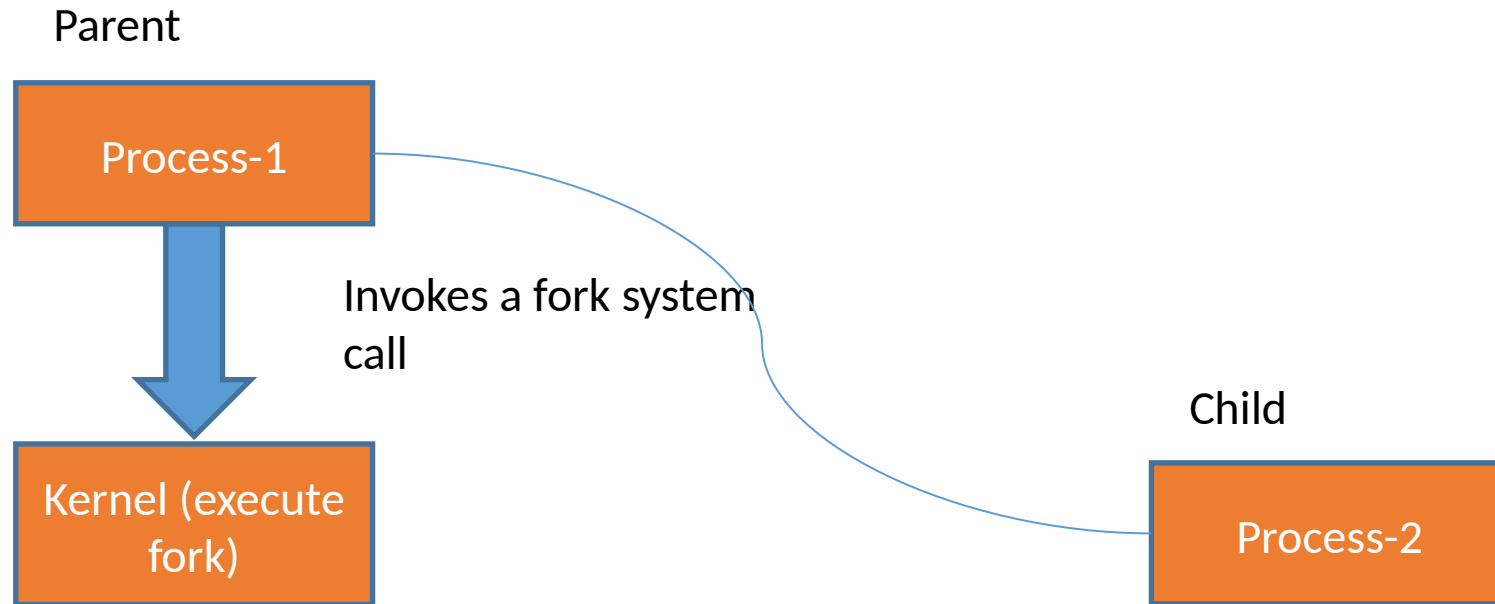
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Creating a Process by Cloning

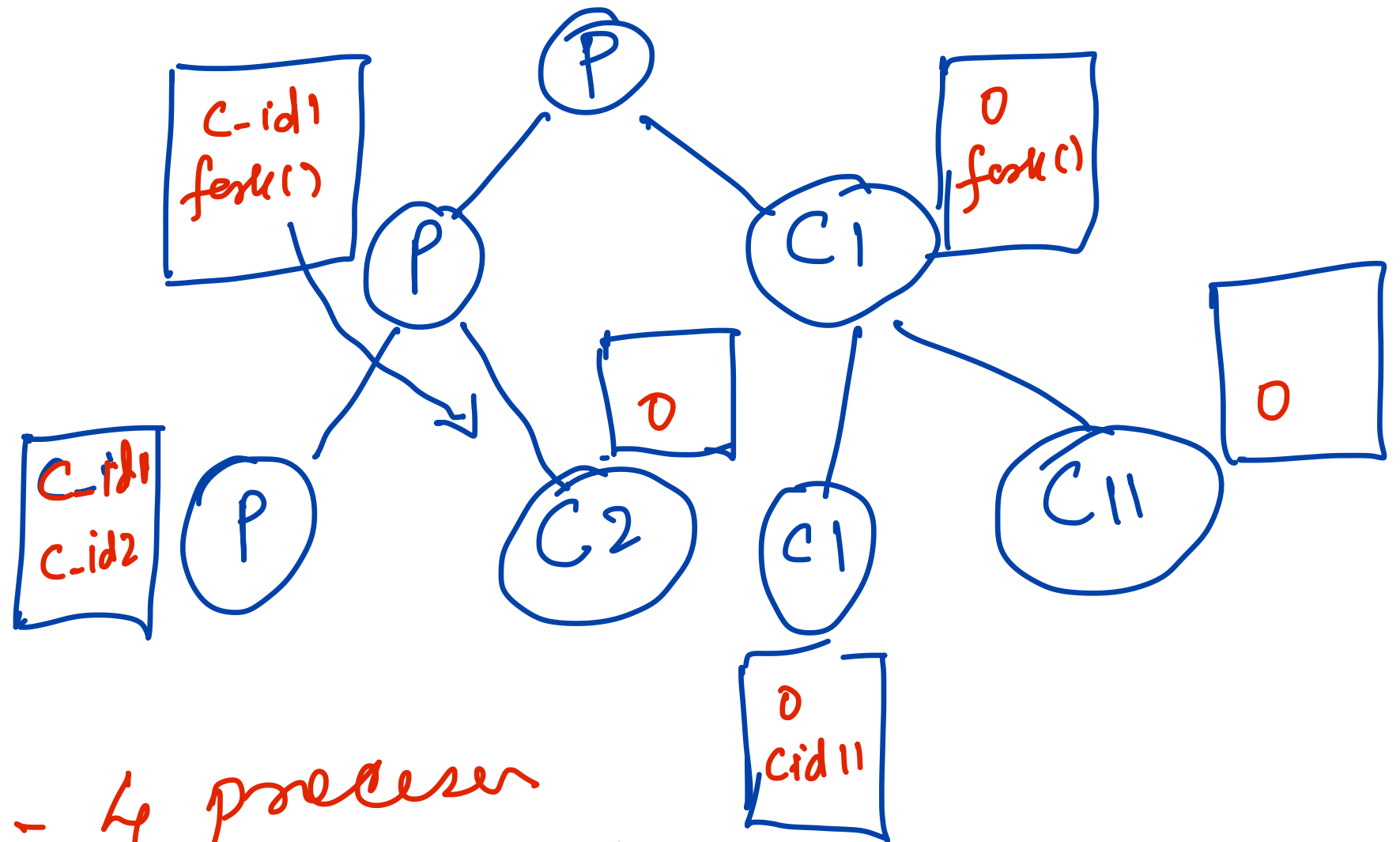
- fork()
 - Child Process is duplicate of parent process
 - PID \Rightarrow Parent process is Child's PID
 - PID \Rightarrow Child process is 0



{


fork();
fork();

}



Total = 4 processes
1 parent + 3 child

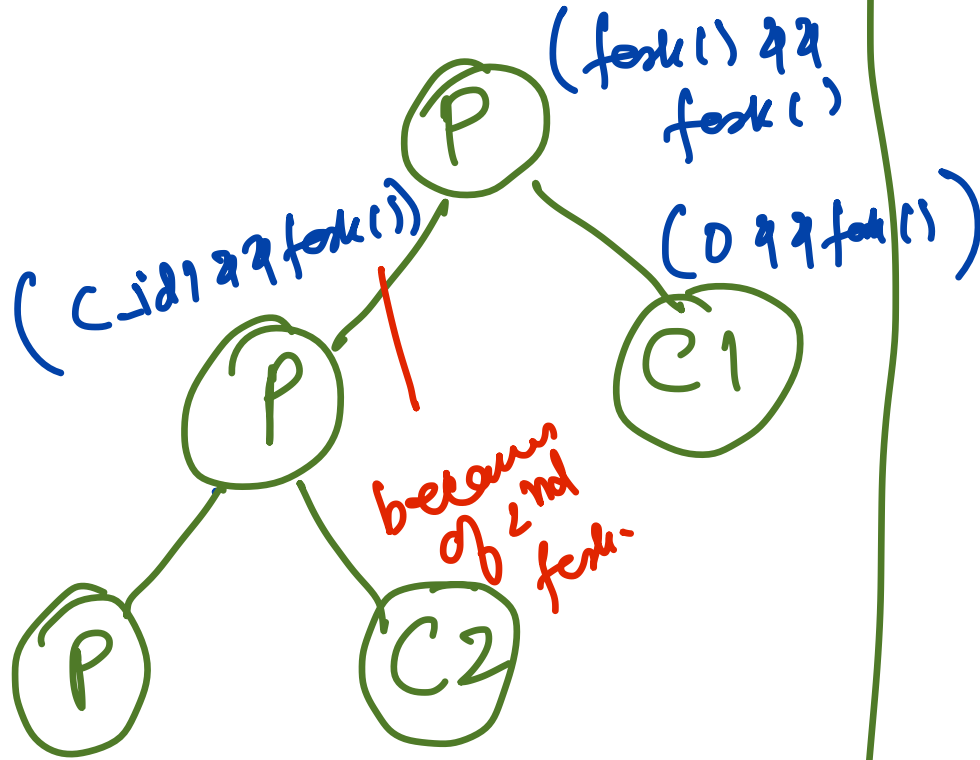
The classic fork() Bomb!

- Overall  processes in Main Memory as indicated at the leaf level nodes count
- In general n fork class (non conditional) will result in _____processes

```
int main()
{
fork();
fork();
fork();
printf("Magic of fork()\n");
return 0;
}
```

if (fork() && fork())

else { }



if (fork() || fork())

{ }

else { }

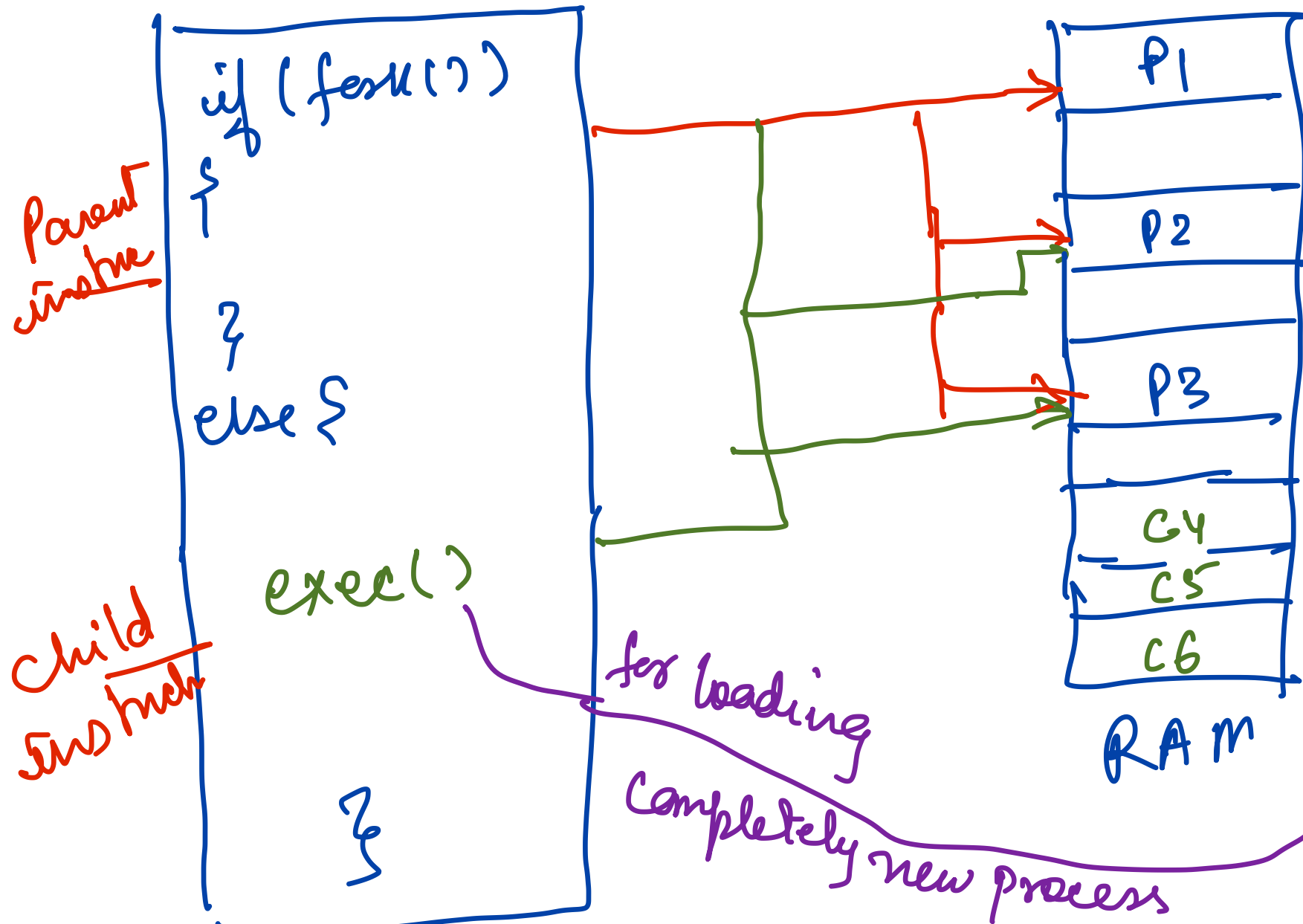
Draw
Process
Tree?

How to create a new definition of process?

- So far fork example we did, child process carried the same image as the parent process.
- Practicality requires child to have new definition.
- Is it possible?
 - Yes

exec() system call

- **fork()** system call is used to create a **SEPARATE, DUPLICATE** process with **non-shared pages** from which process (**parent**) it is called.
- Using **Copy-on-Write (COW)**-**SEPARATE, DUPLICATE** process with **Shared pages**
 - The new **child** process will have different PID.
- When **exec()** system call is invoked from (p1), the program specified in the parameters of exec() will **replace the entire process**.
 - exec() takes one parameter, which is another program(p2)
 - p1 will be replaced by p2.
 - Replace one process with another process (PID don't change)
 - As we are not creating new process
 - We are replacing an existing one
 - **Same PID with different content**



After `exec()` system call, based on new process, new pages gets uploaded in the RAM.

Enter new pages in pagetable

Copy on Write

- * Initially parent and child processes share the same pages.
- * If any of the parent or child wants to update any page, a new copy of that page will be created for that particular page.

wait() system call

- called in parent process
- `int wait(arg);`
- Parent goes to block state
 - Until one of its children terminates
 - -1: if no child is executing or exists
- When the child process exits i.e `exit(0)`, it would cause the parent process to wake up
- the wait function returns the child process's pid
- The parent waits for the child process using this system call.

Why wait() system call ?

- * Waiting for the completion of child process.
- * Sleep(t) system call can also be used but it will ^{not} guarantee that ^{once} child completes its execution, then only parent will start its execution.





Thank You

Any Questions?