

Routing Protocols for Wireless Sensor Networks

Routing protocol survey

- **Traditional technique**
 - Flooding
 - Gossiping
- Routing technique
 - Flat-routing
 - Hierarchical-routing
 - Location-based routing

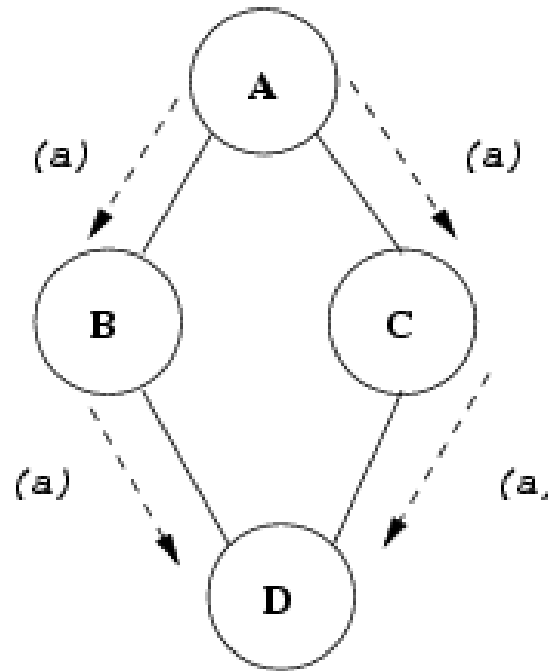
Flooding(1/2)

- A classical mechanisms to relay data in sensor networks without the need for any routing algorithms and topology maintenance.
- drawbacks:
 - Implosion
 - Overlap
 - Resource blindness

A Simple Solution: Flooding

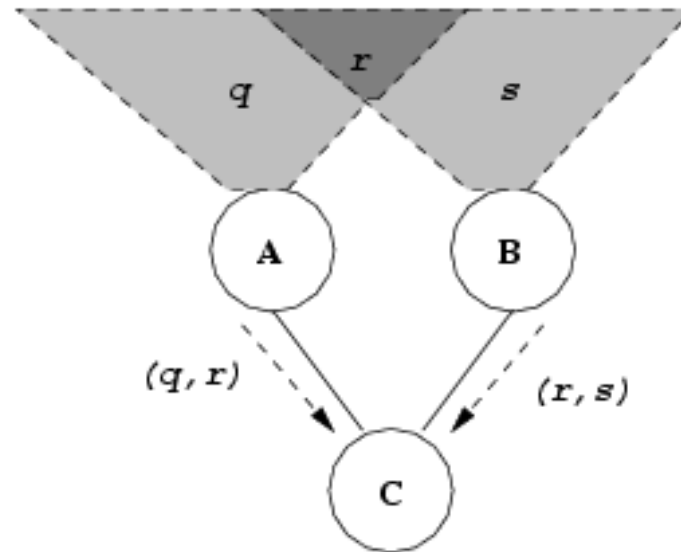
Suffers from 3
deficiencies:

- Implosion: Resources wasted by sending duplicate copies of the same data.



Flooding (Problems)

- Overlap: Common pieces of data are gathered and sent.
- Resource Blindness: Flooding does not modify its activities based on amount of energy available.



Gossiping

- A slightly enhanced version of flooding where the receiving node sends the packet to a randomly selected neighbor which picks another neighbor to forward the packet to and so on.
- Advantage: avoid the implosion
- Drawback: Transmission delay

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- **Current routing technique[1]**
 - **Flat-routing**
 - Hierarchical-routing
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[1]JAMAL N. AL-KARAKI, AHMED E. KAMAL," ROUTING TECHNIQUES IN WIRELESS SENSOR NETWORKS: A SURVEY",
IEEE Wireless Communications • December 2004

Flat-routing

- **SPIN (Sensor Protocols for Information via Negotiation)**
- DD (Directed diffusion)
- Rumor routing

SPIN's Approach

2 key innovations introduced:

- Negotiation

To solve implosion and resource overlap, SPIN nodes *negotiate* before sending data.

Use of descriptors called *meta-data* to name/describe data.

- Resource Adaptation

Each node has a *resource manager* that keeps track of resources at the node. Probed by applications.

4 variants of SPIN

1. A wireless Point to Point Network with no queuing delays or dropping of packets.
 - SPIN-PP and SPIN-EC

2. Realistic network with unreliable broadcast channel.
 - SPIN-BC and SPIN-RL

SPIN Overview

2 basic ideas:

- Exchanging sensor data might be expensive but exchanging data *about* sensor data need not be.
- Nodes must monitor and adapt to changes in their own energy resources.

SPIN features

- Application level control

Pursues the principles of Application Level Framing (ALF). **Common data naming followed** between transmission protocol and application.

Propose to make *routing decisions* in application controlled and application specific ways using knowledge of topology, resources, data layout.

SPIN features (contd.)

- Meta-Data

Used to describe data collected.

Must be **shorter than the data**, and if two pieces of data are the same then they should have the **same meta-data**.

Application-specific: Example a camera sensor could take (x,y,O) as metadata where x,y are coordinates and O is orientation.

- SPIN Resource Management

SPIN applications poll their resources to find out how much energy is available.

Make informed decisions by obtaining this information from a specified interface.

SPIN Message types

- ADV – new data advertisement. Contains the meta-data.
- REQ – request for data. Sent when a node wants a data.
- DATA – data message. Contains the actual data.

SPIN Protocols

1. SPIN – PP

Optimized for a point-to-point link where A and B can talk to each other without interfering with other nodes. To send to n-neighbors, cost is n times that of sending to one neighbor.

3 stages : ADV-REQ-DATA

2. SPIN-EC

When a node observes that its energy is below a threshold it reduces its participation. Only starts a stage of the protocol that it can finish.

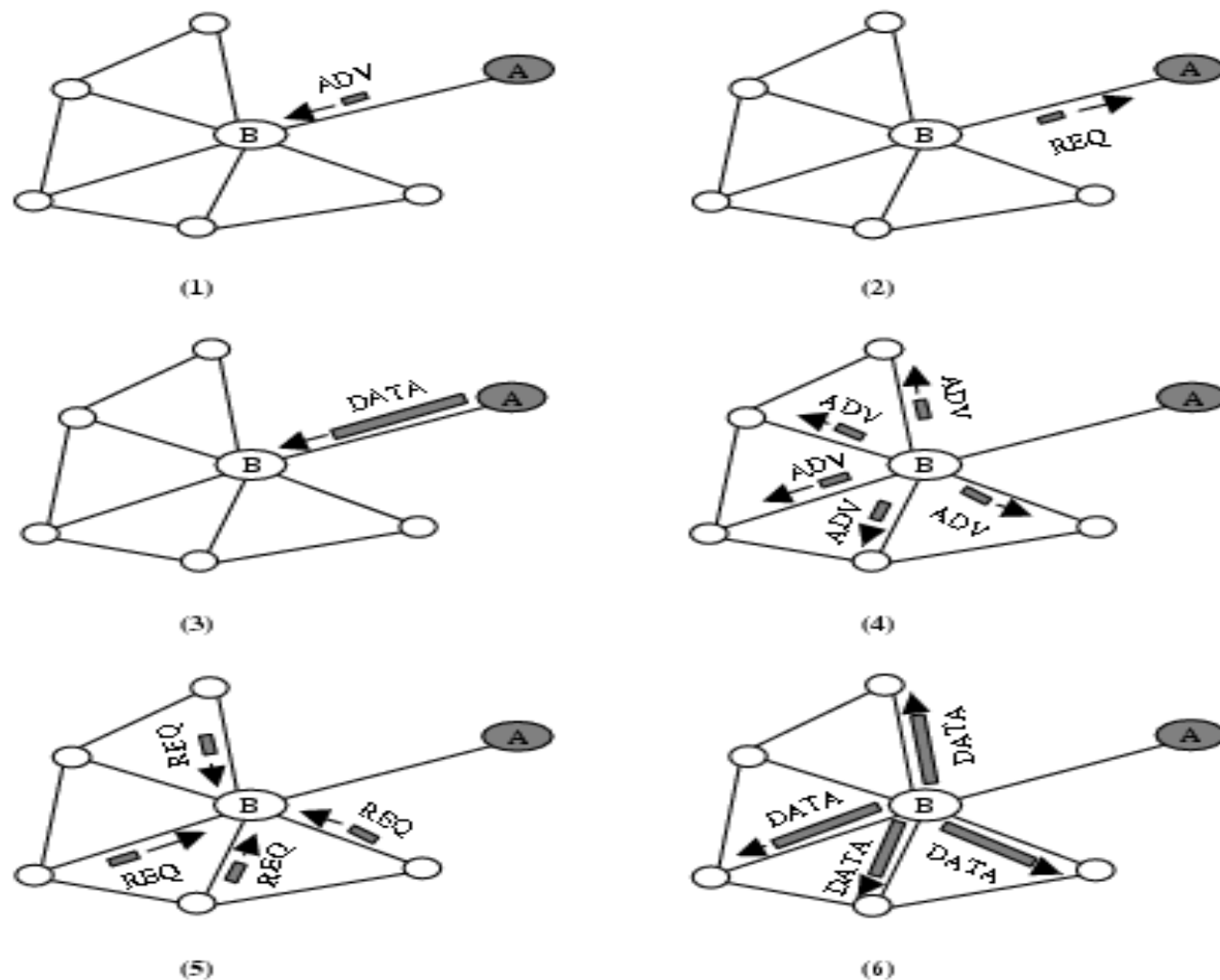


Figure 3: The SPIN-PP Protocol. Node A starts by advertising its data to node B (1). Node B responds by sending a request to node A (2). After receiving the requested data (3), node B then sends out advertisements to its neighbors (4), who in turn send requests back to B (5,6).

SPIN

- Topological changes are **localized**
- provides **more energy savings** than flooding, and metadata negotiation almost **halves the redundant data**.
- Drawback: SPIN's data advertisement mechanism cannot guarantee delivery of data.

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Hierarchical-routing

- *LEACH (Low Energy Adaptive Clustering Hierarchy)*
- PEGASIS (Power-Efficient Gathering in Sensor Information Systems)
- TEEN(APTEEN) (Threshold-Sensitive Energy Efficient Protocols)

LEACH[1]

- LEACH is a cluster-based protocol
- Setup phase
- Steady state phase

[1]. Heinzelman, A. Chandrakasan and H. Balakrishnan, “Energy-Efficient Communication Protocol for Wireless Microsensor Networks,” *Proc. 33rd Hawaii Int’l. Conf. Sys. Sci.*

Assumptions about sensor nodes

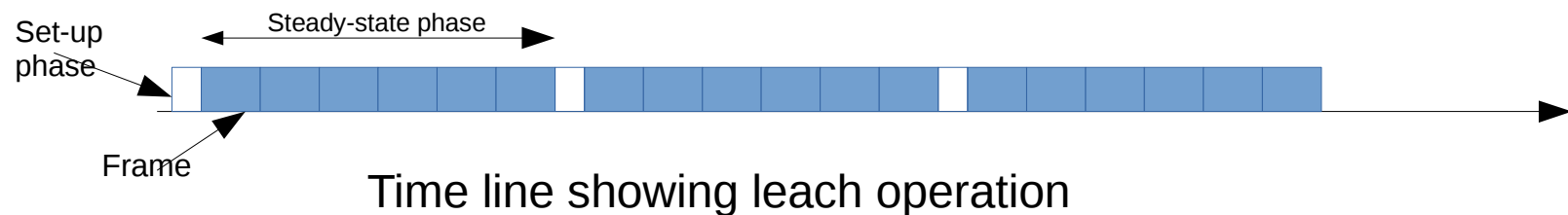
- All nodes can transmit with enough power to reach the BS.
- All nodes have the computational power to support different MAC protocols.

LEACH

- Nodes organized into local clusters.
- One node acting as the cluster head.
- All non-cluster head nodes transmit their data to the cluster head
- Cluster head receive data from all the members perform data aggregation and transmit data to remote BS.
- Cluster head node needs more energy compared to cluster members.
- To avoid draining the battery of any node cluster heads are randomly chosen on rotation basis.

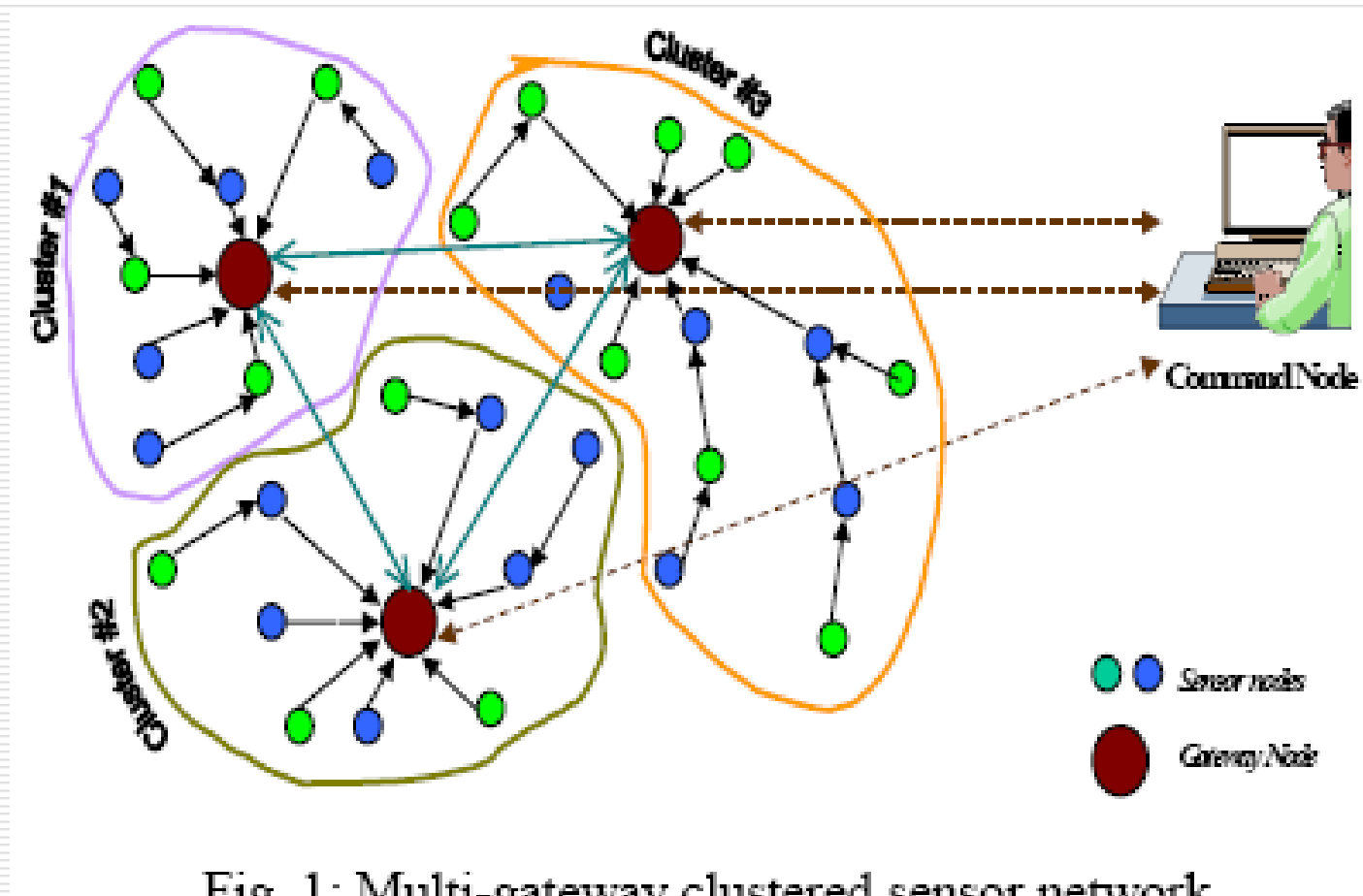
Leach operation

- Set-up phase: Clusters are organized
- Steady-state phase: Data are transferred from nodes to the cluster head and on to the BS.



Cluster-Head Selection Algorithm

LEACH



LEACH: Operation

- Periodic process
- Three phases per round:
 - Advertisement
 - Election and membership
 - Setup
 - Schedule creation
 - Steady-State
 - Data transmission

LEACH: Advertisement

- Cluster head self-election
 - Status advertised to nearby nodes
- Non-cluster heads must listen to the medium
 - Choose membership based on signal strength
 - RSSI
 - E_b/N_0

LEACH: Setup

- Nodes broadcast membership status
 - CSMA-CA
- Cluster heads must listen to the medium
- TDMA schedule created
 - Dynamic number of time slots

LEACH: Data Transmission

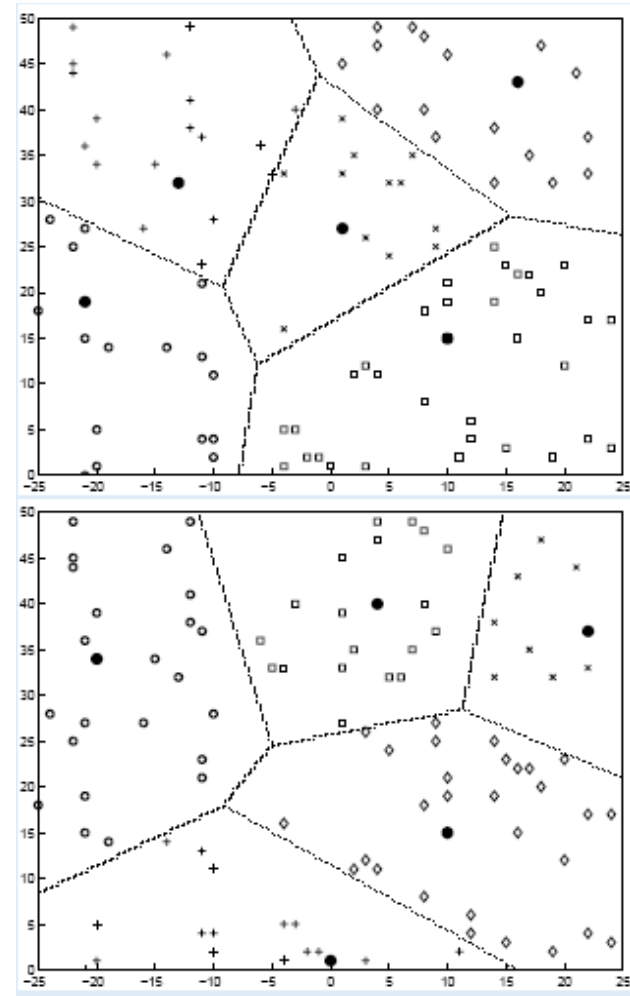
- Nodes sleep until its time slots
- Cluster heads must listen to each slot
- Cluster heads aggregate/compress and transmit to Sink
- Phase continues until the end of the round

Low-Energy Adaptive Clustering Hierarchy (LEACH)

- Adaptive Clustering
 - Distributed
- Randomized Rotation
 - Biased to balance energy loss
- Heads perform compression
 - Also aggregation
- In-cluster TDMA

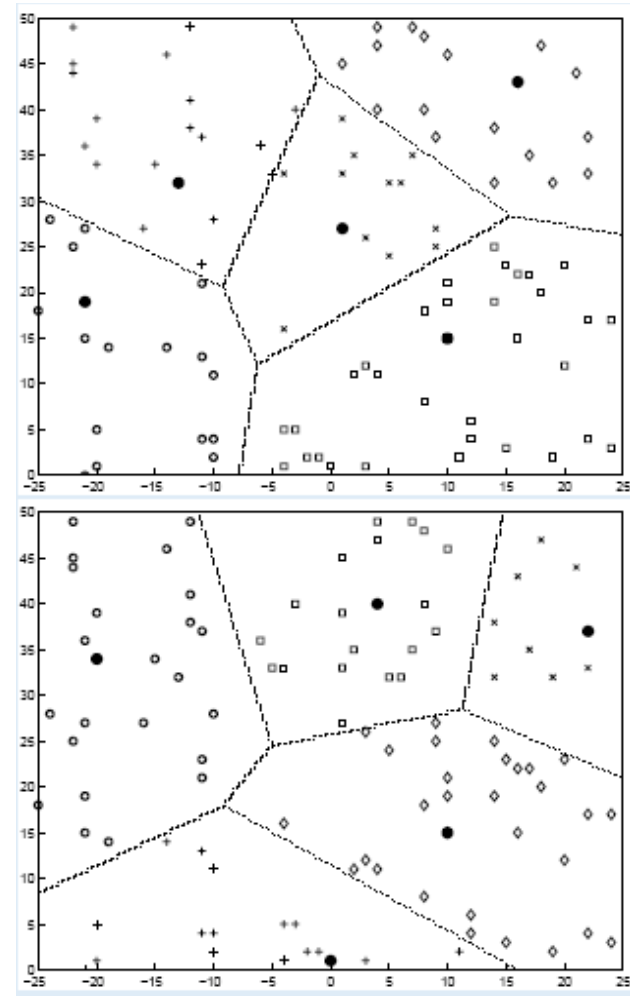
LEACH: Adaptive Clustering

- Periodic independent self-election
 - Probabilistic
- CSMA CA used to advertise
- Nodes select advertisement with strongest signal strength
- Dynamic TDMA time slots



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LEACH: Randomized Rotation

- Cluster heads elected every round
 - Recent cluster heads disqualified
 - Optimal number not guaranteed

- Residual energy considered

- P = Desired cluster head percentage

- r = Current Round

- G = Set of nodes which have not been cluster heads in $1/P$ rounds

$$T(n) = \begin{cases} \frac{P}{1 - P * (r \bmod \frac{1}{P})} & \text{if } n \in G \\ 0 & \text{otherwise} \end{cases}$$

LEACH: Hierarchical Clustering

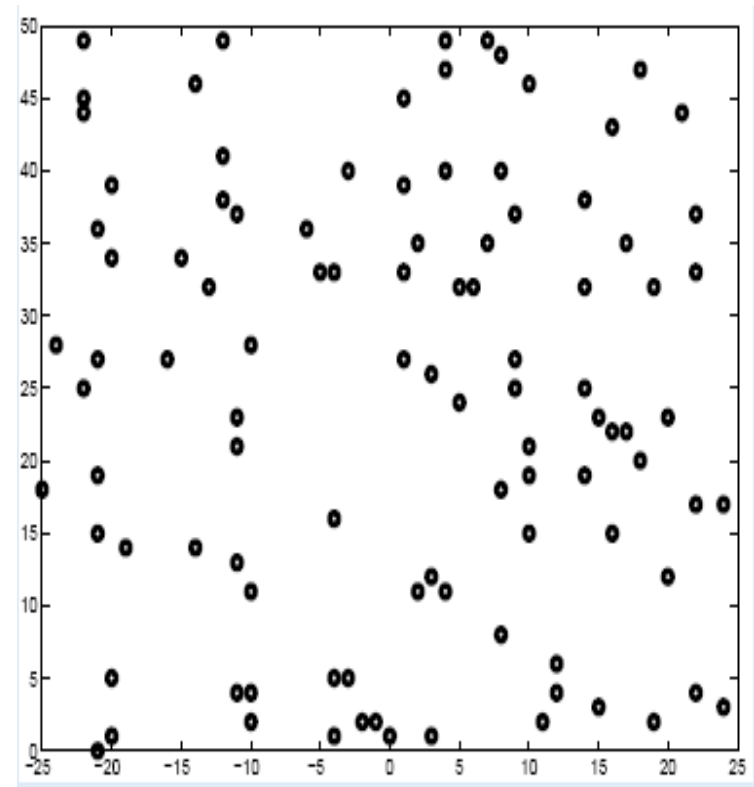
- Not currently implemented
- Efficient when network diameters are large

LEACH

- Drawbacks
 - It is not applicable to networks deployed in large regions
 - The idea of dynamic clustering brings extra overhead
 - The protocol assumes that all nodes begin with the same amount of energy capacity in each election round, assuming that being a CH consumes approximately the same amount of energy for each node

Performance: Parameters

- MATLAB Simulator
- 100-node random network
- $E_{elec} = 50\text{nj/bit}$
- $\varepsilon_{amp} = 100\text{pJ/bit/m}^2$
- $k = 2000$ bits



Comparison between SPIN LEACH and directed diffusion[1]

	SPIN	LEACH	Directed diffusion
Optimal route	No	No	Yes
Network lifetime	Good	Very good	Good
Resource awareness	Yes	Yes	Yes
Use of meta-data	Yes	No	Yes

[1]W. Heinzelman, A. Chandrakasan and H. Balakrishnan, "Energy-Efficient Communication Protocol for Wireless Microsensor Networks," *Proc. 33rd Hawaii Int'l. Conf. Sys. Sci.*, Jan. 2000.